

**In the name of God**

**Electronic advanced algorithms-4012**

**Homework-1: Divide and Conquer**

**Professor: Dr. Farahani**

**TA: Sara Charmchi**

**Student: Mehdi Khaefi**

**Student-id: 401422072**

# 1. Divide-and-conquer multiplication

There is a faster way to multiply, though, called the *divide-and-conquer* approach.

## Algorithm

With divide-and-conquer multiplication, we split each of the numbers into two halves, each with  $n/2$  digits. I'll call the two numbers we're trying to multiply  $a$  and  $b$ , with the two halves of  $a$  being  $a_L$  (the left or upper half) and  $a_R$  (the right or lower half) and the two halves of  $b$  being  $b_L$  and  $b_R$ .

Basically, we can multiply these two numbers as follows.

$$\begin{aligned} ab &= (a_L 10^{n/2} + a_R) (b_L 10^{n/2} + b_R) \\ &= a_L b_L 10^n + a_L b_R 10^{n/2} + a_R b_L 10^{n/2} + a_R b_R \\ &= a_L b_L 10^n + (a_L b_R + a_R b_L) 10^{n/2} + a_R b_R \end{aligned}$$

We can reduce the number of  $n/2$ -digit multiplications from four to *three*!

The idea works as follows: We're trying to compute

$$a_L b_L 10^n + (a_L b_R + a_R b_L) 10^{n/2} + a_R b_R$$

What we'll do is compute the following three products using recursive calls.

$$\begin{aligned} x_1 &= a_L b_L \\ x_2 &= a_R b_R \\ x_3 &= (a_L + a_R) (b_L + b_R) \end{aligned}$$

These have all the information that we want, since the following is true.

$$\begin{aligned} &x_1 10^n + (x_3 - x_1 - x_2) 10^{n/2} + x_2 \\ &= a_L b_L 10^n + ((a_L b_L + a_L b_R + a_R b_L + a_R b_R) - a_L b_L - a_R b_R) 10^{n/2} + a_R b_R \\ &= a_L b_L 10^n + (a_L b_R + a_R b_L) 10^{n/2} + a_R b_R \end{aligned}$$

And we already reason that this last is equal to the product of  $a$  and  $b$ .

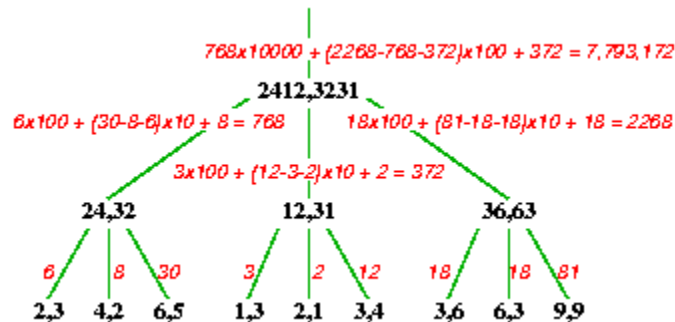
## Pseudocode: Divide Conquer Combine

```
BigInteger multiply(BigInteger a, BigInteger b) {
    int n = max(number of digits in a, number of digits in b)
    if(n == 1) {
        return a.intValue() * b.intValue();
    } else {
        BigInteger aR = bottom n/2 digits of a;
        BigInteger aL = top remaining digits of a;
        BigInteger bR = bottom n/2 digits of b;
        BigInteger bL = top remaining digits of b;
        BigInteger x1 = Multiply(aL, bL);
        BigInteger x2 = Multiply(aR, bR);
        BigInteger x3 = Multiply(aL + aR, bL + bR);
        return x1 * pow(10, n) + (x3 - x1 - x2) * pow(10, n / 2) + x2;
    }
}
```

## 2. The smallest sub-problem:

The smallest sub-problem in this method is multiplying two one-digit numbers.

Let's do an actual multiplication to illustrate how this works. I'm going to draw a recursion tree, labeling the edges with the final values computed by each node of the tree.



#### 4. Time complexity analysis

For  $n$  digit integer, we have to perform 3 multiplications of integers of size  $(n / 2)$ .  
Recurrence equation for this problem is given as,

$$T(n) = 3T(n/2), \text{ if } n > 1$$

$$T(n) = 1, \text{ if } n = 1$$

**Proof:**

$$T(n) = 3T(n/2) \dots \textbf{(3)}$$

Let us solve this recurrence by an iterative approach. Substitute  $n$  by  $n/2$  in Equation (3)

$$T(n/2) = 3T(n/4) \dots \textbf{(4)}$$

Put this value in Equation (3),

$$T(n) = 3(3T(n/4)) = 3^2T(n/2^2) \dots \textbf{(5)}$$

Substitute  $n$  by  $n/2$  in Equation (4)

$$T(n/4) = 3T(n/8)$$

Put this value in Equation (3)

$$T(n) = 3(3^2T(n/8)) = 3^3T(n/2^3) \dots \textbf{(6)}$$

.

.

.

.

After  $k$  iterations,

$$T(n) = 3^kT(n/2^k) \dots \textbf{(7)}$$

Every time number of digits in number reduces by factor 2, so it can go as deep as  $\log_2 n$ ,

So,  $k = \log_2 n \Rightarrow n = 2^k$

Thus from equation (5),  $T(n) = 3^k T(2^k / 2^k)$

$T(n) = n^{\log_2 3} \times T(1) (\because n^{\log_b a} = a^{\log_b n})$

$T(1) = 1$  (Only one multiplication is required to multiply two numbers of digits 1)

So,  $T(n) = n^{\log_2 3} = O(n^{1.58})$

Grade school method multiplies each digit of the multiplier with each digit of the multiplicand. So for each digit of the multiplier,  $n$  multiplications are performed with multiplicand.

This is done each of the  $n$  bits of the multiplier. So running time of that method was  $O(n^2)$ , whereas the divide and conquer approach reduces the running time to  $O(n^{1.58})$ . For large numbers, the difference becomes significant.

## References:

<http://www.cburch.com/csbsju/cs/160/notes/31/1.html>

<https://codecrucks.com/large-integer-multiplication-using-divide-and-conquer/>