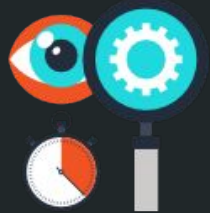
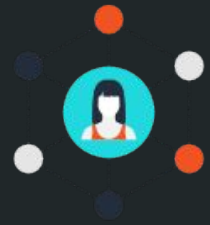


Blockchain

Nepal



Symmetric Key Ciphers

XOR

- XOR

$$0+0 = 0;$$

$$0+1 = 1;$$

$$1+0 = 1;$$

$$1+1 = 0$$

- XOR of bit vectors (words)

>> we do XOR for each corresponding bit pairs

e.g., $0011 + 1010 = 1001$

One Time Pad (OTP)

- Key: sequence of random bytes: $k_1k_2\dots k_l$
- Encryption:
 - message is represented as sequence of bytes : $m_1m_2\dots m_l$
 - cipher text = message XOR key
$$m_1m_2\dots m_l + k_1k_2\dots k_l = c_1c_2\dots c_l$$
where $c_i = m_i + k_i$ for all $i = 1, \dots, l$
- Decryption:
 - ciphertext represented as sequence of bytes : $c_1c_2\dots c_l$
 - plaintext = ciphertext XOR key
$$c_1c_2\dots c_l + k_1k_2\dots k_l = m_1m_2\dots m_l$$
 - Proof: $c_i + k_i = m_i + k_i + k_i = m_i$

Properties of OTP

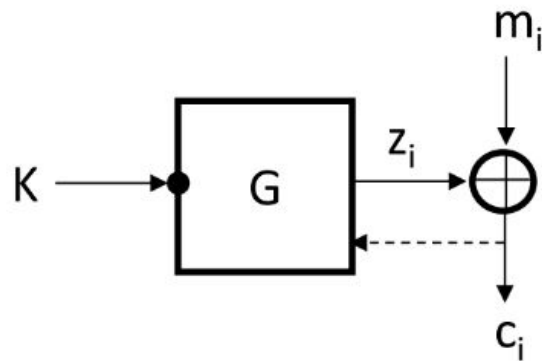
- Perfect secrecy
 - >> Looking at the encrypted messages provides no information about the original message.
- Large key size
 - >> Need a truly random key that has the same length as the message
- impractical in many applications
 - >> how to send the key in a secure way to the recipient?
 - >> key management is a pain

Stream Ciphers

Idea: simulate the truly random key stream of the one-time pad with a pseudo random sequence generated from a random seed.

Terminology:

- m_i - plaintext character
- c_i - ciphertext character
- z_i - key-stream character
- K - key(seed)
- G - key stream generator

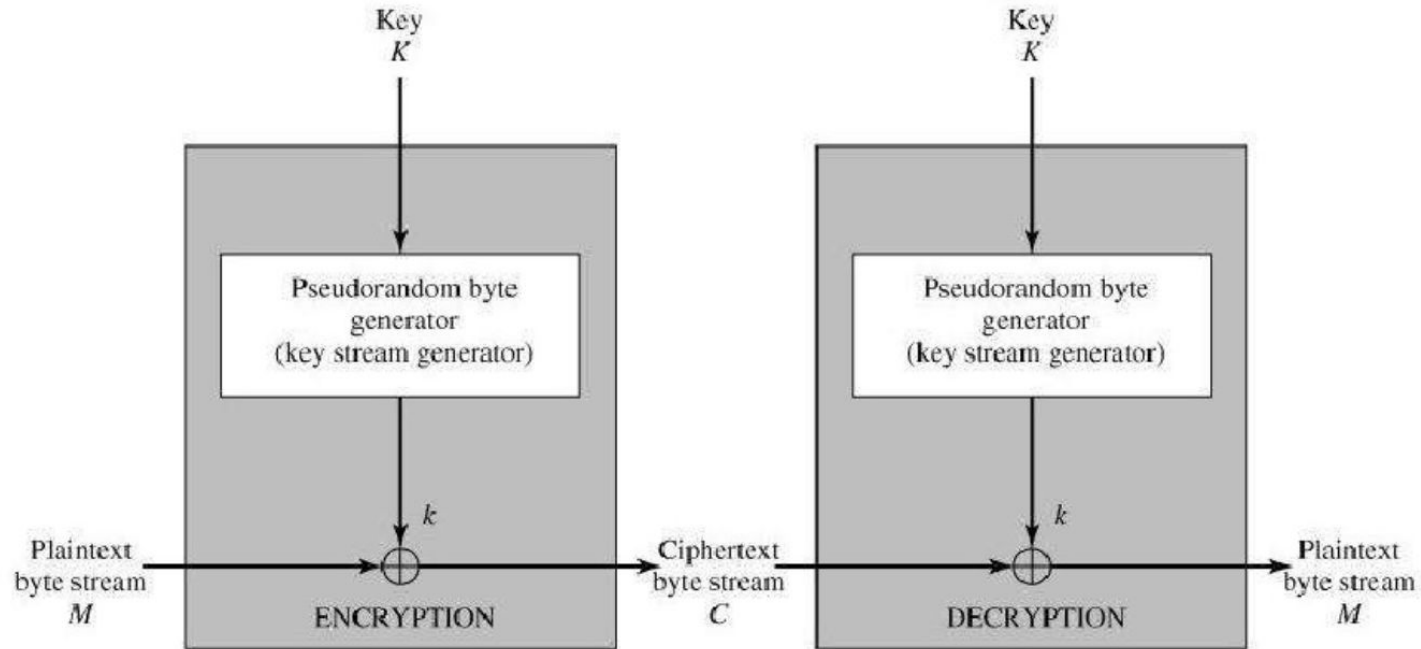


Examples:- LSFR, RC4

Properties of Stream Ciphers

- Usually very efficient
 - >> Fast (hardware impl)
 - >> Requires less memory
- Ciphertext always has the same length as the plaintext

RC4



RC4

1. Initialize an array of 256 bytes
2. Run the Key Scheduling Algorithm (KSA) on them
3. Run the Pseudo-random Generation Algorithm (PRGA) on the (KSA) output to generate Key stream
4. XOR the data with a key stream

RC4- Algorithm

- initialization:

```
for i = 0 to 255 do
  S[i] = i
end

j = 0
for i = 0 to 255 do
  j = j+S[i]+K[i mod len(K)] mod 256
  swap(S, i, j)
end

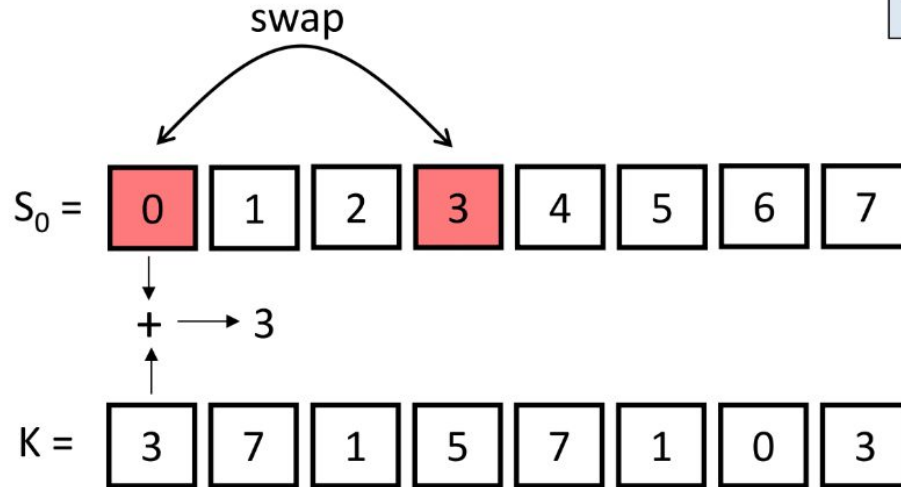
i = 0
j = 0
```

- generation:

```
i = i+1 mod 256
j = j+S[i] mod 256
swap(S, i, j)
return S[ S[i]+S[j] mod 256 ]
```

Example: RC4 Initialization

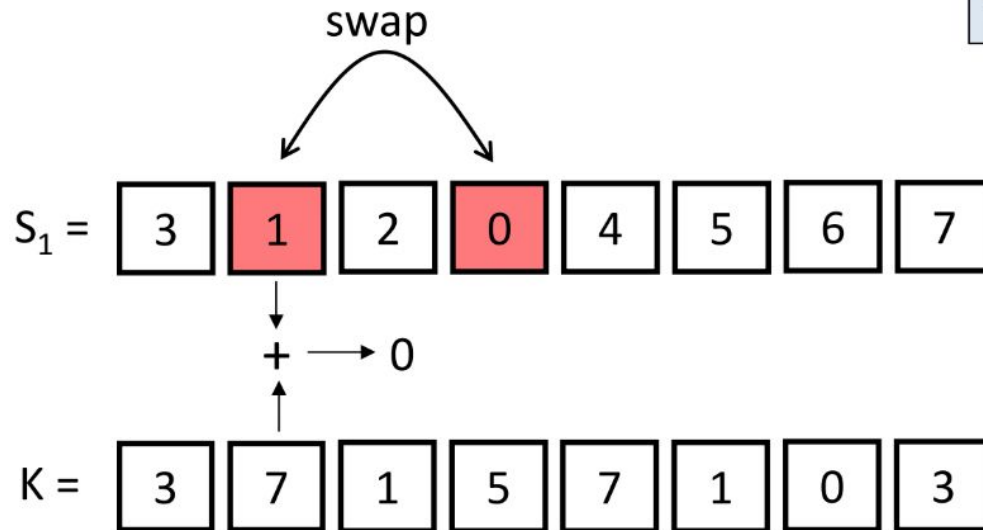
```
for i = 0 to 7 do  
  j = j + S[i] + K[i] mod 8  
  swap (S, i, j)  
end
```



i = 0 j = 0
 j₀ = 3

Example: RC4 Initialization

```
for i = 0 to 7 do
  j = j + S[i] + K[i] mod 8
  swap (S, i, j)
end
```



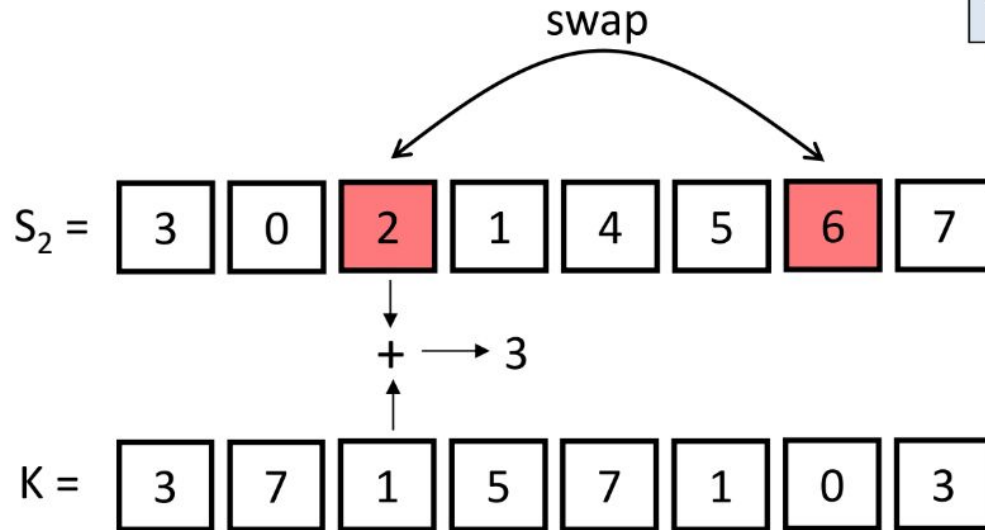
$j = 0$

$i = 0$ $j_0 = 3$

$i = 1$ $j_1 = 3$

Example: RC4 Initialization

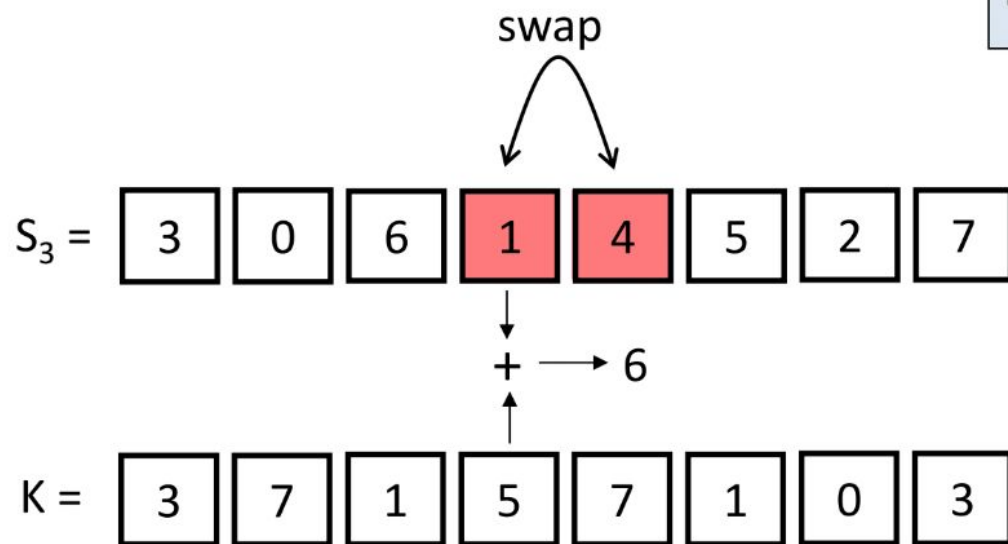
```
for i = 0 to 7 do  
  j = j + S[i] + K[i] mod 8  
  swap (S, i, j)  
end
```



$j = 0$
 $i = 0 \quad j_0 = 3$
 $i = 1 \quad j_1 = 3$
 $i = 2 \quad j_2 = 6$

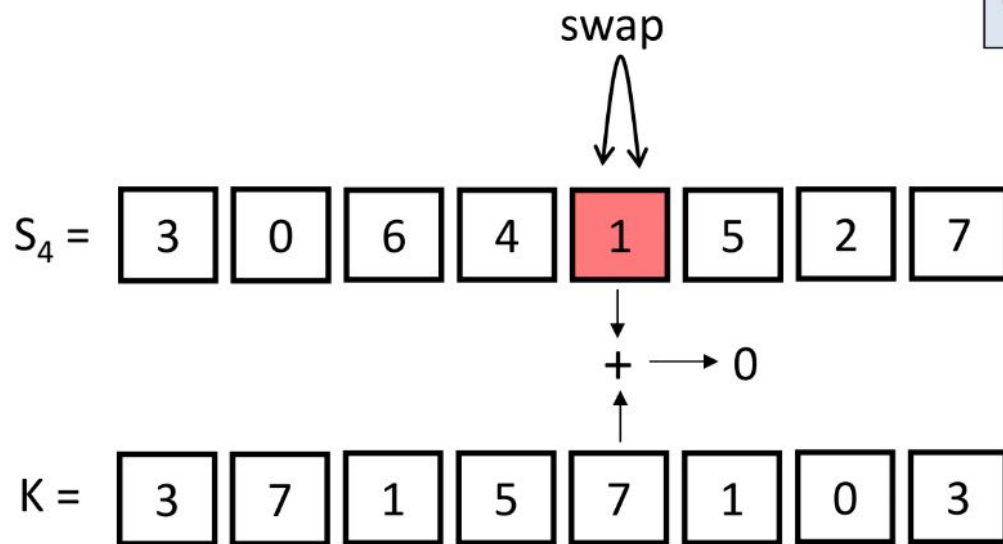
Example: RC4 Initialization

```
for i = 0 to 7 do  
  j = j + S[i] + K[i] mod 8  
  swap (S, i, j)  
end
```



Example: RC4 Initialization

```
for i = 0 to 7 do  
  j = j + S[i] + K[i] mod 8  
  swap (S, i, j)  
end
```

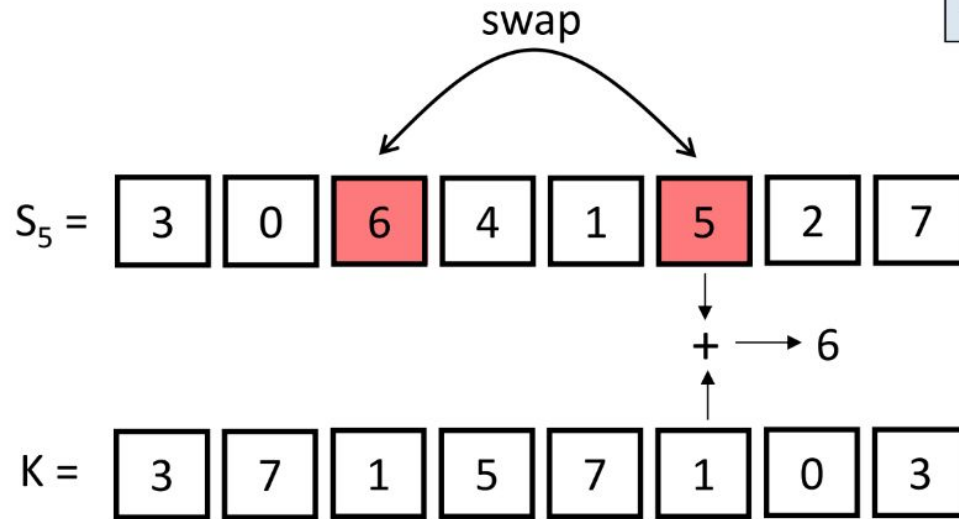


$j = 0$

$i = 0$	$j_0 = 3$
$i = 1$	$j_1 = 3$
$i = 2$	$j_2 = 6$
$i = 3$	$j_3 = 4$
$i = 4$	$j_4 = 4$

Example: RC4 Initialization

```
for i = 0 to 7 do  
  j = j + S[i] + K[i] mod 8  
  swap (S, i, j)  
end
```

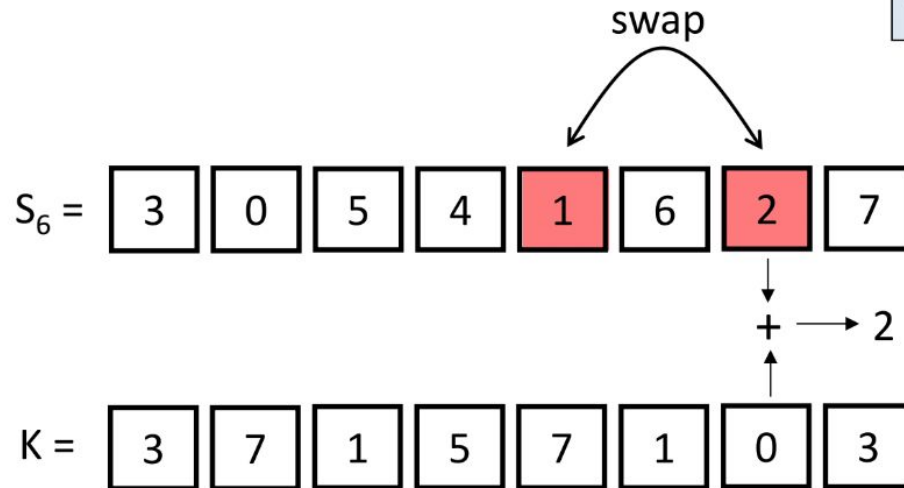


$j = 0$

$i = 0$	$j_0 = 3$
$i = 1$	$j_1 = 3$
$i = 2$	$j_2 = 6$
$i = 3$	$j_3 = 4$
$i = 4$	$j_4 = 4$
$i = 5$	$j_5 = 2$

Example: RC4 Initialization

```
for i = 0 to 7 do  
  j = j + S[i] + K[i] mod 8  
  swap (S, i, j)  
end
```



$j = 0$

$i = 0 \quad j_0 = 3$

$i = 1 \quad j_1 = 3$

$i = 2 \quad j_2 = 6$

$i = 3 \quad j_3 = 4$

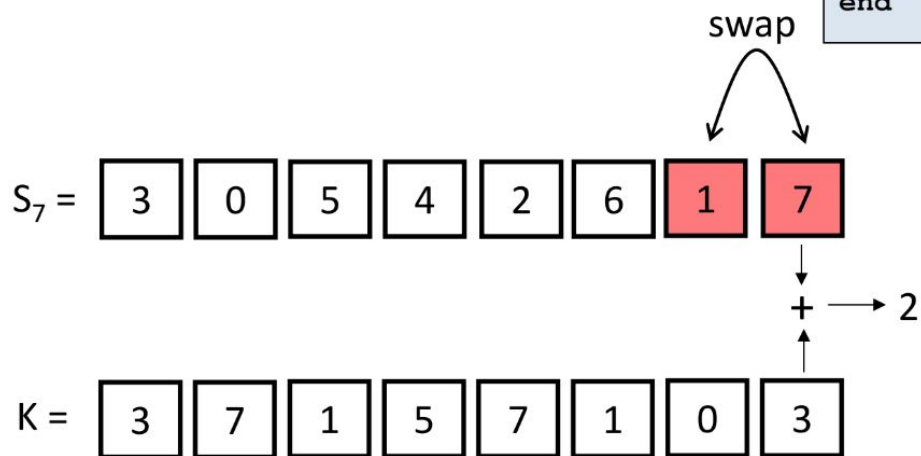
$i = 4 \quad j_4 = 4$

$i = 5 \quad j_5 = 2$

$i = 6 \quad j_6 = 4$

Example: RC4 Initialization

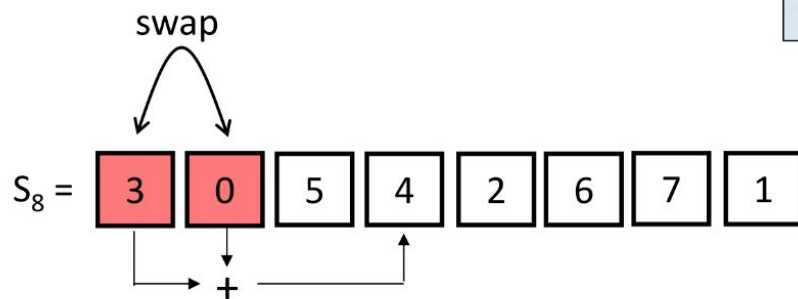
```
for i = 0 to 7 do  
  j = j + S[i] + K[i] mod 8  
  swap (S, i, j)  
end
```



$j = 0$	
$i = 0$	$j_0 = 3$
$i = 1$	$j_1 = 3$
$i = 2$	$j_2 = 6$
$i = 3$	$j_3 = 4$
$i = 4$	$j_4 = 4$
$i = 5$	$j_5 = 2$
$i = 6$	$j_6 = 4$
$i = 7$	$j_7 = 6$

Example: RC4 Generation

```
i = i+1 mod 8  
j = j + S[i] mod 8  
swap (S, i, j)  
return S[ S[i]+S[j] mod 8 ]
```



$i = 0$ $j = 0$

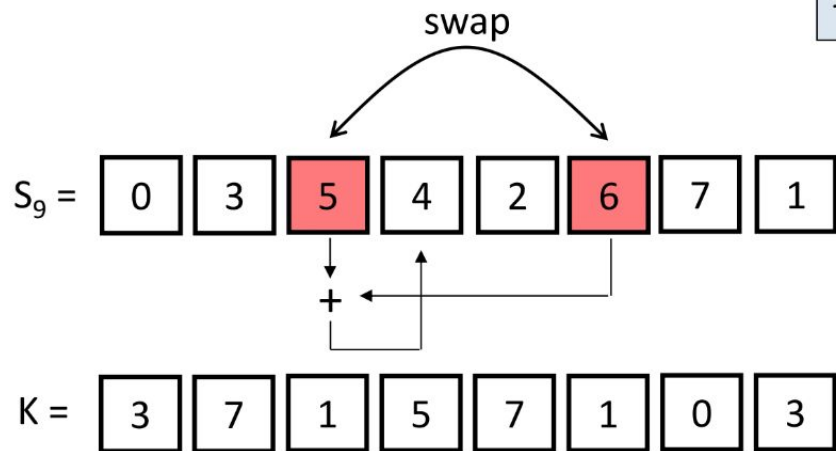
$i = 1$ $j_1 = 0 + S[1] = 0$

$K =$ [3] [7] [1] [5] [7] [1] [0] [3]

$\text{output}_1 = S[3] = 4$

Example: RC4 Generation

```
i = i+1 mod 8  
j = j + S[i] mod 8  
swap (S, i, j)  
return S[ S[i]+S[j] mod 8 ]
```



$i = 0$ $j = 0$
 $i = 1$ $j_1 = 0 + S[1] = 0$
 $i = 2$ $j_2 = 0 + S[2] = 5$

$$\text{output}_2 = S[5 + 6 \bmod 8] = S[3] = 4$$

Exercise 2

> Implement RC4 Stream cipher and implement proper unit tests for encryption and decryption.