

The simplest Lewis-style triviality proof yet?

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In his celebrated ‘Probabilities of conditionals and conditional probabilities’ David Lewis showed that the identification of the probability of conditionals $a \Rightarrow c$ with the conditional probability of consequent given antecedent when that antecedent has non-zero probability,

$$\text{i.e. for all } a \text{ and } c, P(a \Rightarrow c) = P(c|a), \text{ when } P(a) > 0,$$

trivializes the probability distribution in question. Lewis presented three triviality results:

- LEWIS 1 If $P(a \& c) > 0$ and $P(a \& \sim c) > 0$ then $P(c|a) = P(c)$;
- LEWIS 2 P assigns non-zero probabilities to at most two of any set of pairwise inconsistent propositions;
- LEWIS 3 P takes at most four values.

The premisses for Lewis’s results are

- (i) the probability expansion rules,

$$\begin{aligned} P(a \Rightarrow c) &= P((a \Rightarrow c) \& b) + P((a \Rightarrow c) \& \sim b) \\ &= P(a \Rightarrow c|b)P(b) + P(a \Rightarrow c|\sim b)P(\sim b), \end{aligned}$$

the first holding generally, the second when $0 < P(b) < 1$;

- (ii) the family of probability distributions satisfying the identification of conditional probability and probability of conditional is closed under conditionalization, so that

$$P(a \Rightarrow c|b) = P(c|a \& b) \text{ when } P(a \& b) > 0.$$

Lewis used the second expansion rule, expanding the probability of the conditional with respect to its consequent, i.e. setting $b = c$ in (i), then applied (ii) to obtain LEWIS 1, from which he derived LEWIS 2 and, from it in turn, LEWIS 3.

It’s more fun to expand with respect to the corresponding material conditional. To make that pay off we need two facts.

- (1) It was, it seems, Karl Popper – see Dorn 1992/93 – who first remarked that the probability of the material conditional, $a \supset c$, is never less than the conditional probability $P(c|a)$ and but for exceptional cases to be noted exceeds it.

When $P(a) > 0$,

$$\begin{aligned} P(a \supset c) &= P(\sim a \vee (a \& c)) = P(\sim a) + P(a \& c) \\ &= P(\sim a) + P(c|a)P(a) \end{aligned}$$

$$\geq P(c|a)P(\sim a) + P(c|a)P(a) = P(c|a),$$

with equality if, and only if, $P(\sim a) = 0$ or $P(c|a) = 1$.

- (2) $a \Rightarrow c$ being a conditional, an obvious question to ask is, how probable is $a \Rightarrow c$ given the corresponding material conditional?
The surprising answer:

$$\begin{aligned} \text{when } P(a \& c) > 0, P(a \Rightarrow c | a \supset c) &= P(c | a \& (a \supset c)) \\ &= P(c | a \& c) = 1. \end{aligned}$$

Expanding $P(a \Rightarrow c)$ we find:

$$\begin{aligned} P(a \Rightarrow c) &= P((a \Rightarrow c) \& (a \supset c)) + P((a \Rightarrow c) \& \sim(a \supset c)) \\ &\geq P((a \Rightarrow c) \& (a \supset c)) \\ &= P(a \Rightarrow c | a \supset c)P(a \supset c) = P(a \supset c), \end{aligned}$$

when $P(a \& c) > 0$.

Putting (1) and (2) together yields a new variation on the triviality theme:

$$P(a \Rightarrow c) = P(a \supset c) \text{ when } P(a \& c) > 0.$$

From (1) we know that with $P(a) > 0$, $P(a \Rightarrow c) = P(a \supset c)$ when, and only when, $P(a) = 1$ or $P(c|a) = 1$. So,

$$P(a) = 1 \text{ or } P(c|a) = 1 \text{ when } P(a \& c) > 0.$$

Put another way, if $0 < P(a) < 1$ and $0 < P(a \& c)$ then $P(c|a) = 1$. Since $P(c|a) = 0$ when $P(a) > 0$ and $P(a \& c) = 0$, we have our *Basic Triviality Result (BTR)*:

the function $P(.|a)$ is two-valued, i.e. takes only the values 0 and 1, when $0 < P(a) < 1$.

BTR tells us that if one learns/comes to believe something – anything – about which one is not currently certain one way or the other, i.e. $0 < P(a) < 1$, then, updating by conditionalization, one will be certain, one way or the other, about everything!

BTR affords easy proofs of LEWIS 1, LEWIS 2 and LEWIS 3:

LEWIS 1 If $P(a \& c) > 0$ and $P(a \& \sim c) > 0$, then $P(a) > 0$ and $0 < P(c|a) = P(a \& c)/[P(a \& c) + P(a \& \sim c)] < 1$.

As $P(.|a)$ is not two-valued, $P(a) = 1$ so $P(c|a) = P(c)$.

This proof gives us extra information about the independence under P of a and c when $P(a \& c) > 0$ and $P(a \& \sim c) > 0$. It shows it to be essentially trivial, resulting from the fact that

$$P(a) = 1 \text{ when } P(a \& c) > 0 \text{ and } P(a \& \sim c) > 0.$$

What is in essentials the same proof works for

LEWIS 2 If $P(a) > 0$, $P(b) > 0$, and a and b are inconsistent, $P(a \vee b) = P(a) + P(b) > 0$ and $0 < P(a|a \vee b) = P(a)/[P(a) + P(b)] < 1$.

As $P(.|a \vee b)$ is not two-valued, $P(a \vee b) = 1$. Consequently, if c is inconsistent with both a and b and so inconsistent with $a \vee b$, $P(c) = 0$.

LEWIS 3 If, for some a , $0 < P(a) < 1$ then $P(.|a)$ and $P(.|\sim a)$ are both two-valued. But then, for any b , as

$$P(b) = P(b|a)P(a) + P(b|\sim a)P(\sim a),$$

$$P(b) \in \{0, P(a), P(\sim a), 1\}.$$

BTR, LEWIS 1 and LEWIS 2 are equivalent. As mentioned above, Lewis derived LEWIS 2 from LEWIS 1. It remains to show that LEWIS 2 entails BTR:

Suppose that $0 < P(a) < 1$ and that $0 < P(c|a) < 1$. Then $P(a \& c) > 0$, $P(a \& \sim c) > 0$, and $P(\sim a) > 0$. But $a \& c$, $a \& \sim c$, and $\sim a$ are pairwise inconsistent.

LEWIS 3, perhaps the most eye-catching of the original triviality results, is weaker. LEWIS 3 does not entail LEWIS 2: consider a fair, three-ticket lottery.

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