Branden Fitelson Philosophy 148 Lecture

Philosophy 148 — Announcements & Such

- Overall, people did very well on the mid-term ($\mu = 90, \sigma = 16$).
- HW #2 graded will be posted very soon. Raul won't be able to give back the HW's until next week, and I will collect HW #3's for him today.
- Recall: we're using a straight grading scale for this course. Here it is:
 - A+ > 97, A (94,97], A- (90,94], B+ (87,90], B (84,87], B- (80,84], C+ (77,80], C (74,77], C- (70,74], D [50,70], F < 50.
- I've posted solutions for HW #1 (HW #2 solutions are coming soon).
- HW #3 is due today. HW #4 has been posted. We will have a discussion devoted to it in two weeks time. Next week, I'm in Europe so no office hours for me. We'll have guest lecturers so the class will meet!
- Today: More on "Logical" Probability and Inductive Logic
 - Review of Carnapian theory, and some new stuff on Carnap.
 - Another way of thinking about Inductive Logic.

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Review of Carnapian "Logical" Probability 2

- Exactly how $\mathfrak c$ (i.e., \Pr_T) should be applicable to epistemology is not completely clear. But, Carnap says things which *constrain* applicability.
 - It should be possible for logically independent claims to be *correlated* with each other. Presumably, because it is possible for logically independent claims to *evidentially support* each other.
 - Specifically, Carnap thinks Ga and Gb should be correlated " $a\ priori$ ":

 $Pr(Gb \mid Ga) > Pr(Gb)$ [this is called "instantial relevance"]

- This *rules-out* m[†]. And, this leads Carnap to adopt m* instead. But, m* has "inapplicability problems" of its own. Carnap later came to think we should also be able to have the following chain of inequalities:

 $Pr(Gb \mid Ga) > Pr(Gb \mid Ga \& Fa \& \sim Fb) > Pr(Gb)$

- But (HW #3!), *neither* \mathfrak{m}^{\dagger} *nor* \mathfrak{m}^{*} are compatible with this. This lead Carnap to continue his search. He moved to a more complex $\mathfrak{m}^{\lambda,y}$.

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Review of Carnapian "Logical" Probability 1

- Carnap endorses the *logicality desideratum* (\mathcal{D}_2) for \mathfrak{c} , and since Carnap defines \mathfrak{c} in terms of Pr, Carnap concludes this mandates a "logical" kind (theory/interpretation) of *probability* the search for "logical Pr" is on.
- Carnap assumes that $c(C, P) = Pr(C \mid P)$. The principle of indifference (PI) follows from this assumption, since then "K does not *favor* any s_i over any s_j " \Rightarrow "K *confirms* each s_i to the same degree as K *confirms* each s_j " $\Rightarrow c(s_i, K) = c(s_i, K) \Rightarrow Pr(s_i \mid K) = Pr(s_i \mid K)$, for *all* i and j, *ergo* (PI).
- This leads Carnap, initially, to endorse the Wittgensteinian function \mathfrak{m}^{\dagger} , which assigns equal probability to each state description s_i of \mathcal{L} .
- But, \mathfrak{m}^{\dagger} implies that there can be no correlations between logically independent propositions. Carnap thinks this renders \mathfrak{m}^{\dagger} *inapplicable*.
- Here, Carnap is presupposing what I will call an applicability desideratum:
- (\mathcal{D}_3) Inductive logic (*i.e.*, the theory of confirmation) should be *applicable* (presumably, to *epistemology*) in *some substantive way*.

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Carnap's Final Theory of "Logical" Probability $(m, n = 2, \lambda = \frac{1}{2})$

	Fa	Ga	Fb	Gb	Carnap's $\mathfrak{m}^{\frac{1}{2},\gamma}(s_i)$ [where $\gamma_F, \gamma_G \in (0,1)$]
	Т	Т	Т	Т	$\frac{1}{9}\gamma_F\gamma_G\left(\gamma_G+\gamma_F\left(5\gamma_G+1\right)+2\right)$
	Т	Т	Т	F	$-\frac{1}{9}\gamma_F\left(5\gamma_F+1\right)\left(\gamma_G-1\right)\gamma_G$
	Т	Т	F	Т	$-\frac{1}{9}\left(\gamma_F-1\right)\gamma_F\gamma_G\left(5\gamma_G+1\right)$
	Т	Т	F	F	$\frac{5}{9}\left(\gamma_F-1\right)\gamma_F\left(\gamma_G-1\right)\gamma_G$
-	Т	F	Т	Т	$-\frac{1}{9}\gamma_F\left(5\gamma_F+1\right)\left(\gamma_G-1\right)\gamma_G$
	Т	F	Т	F	$\frac{1}{9}\gamma_F(\gamma_G-1)(\gamma_G+\gamma_F(5\gamma_G-6)-3)$
_	Т	F	F	Т	$\frac{5}{9} (\gamma_F - 1) \gamma_F (\gamma_G - 1) \gamma_G$
-	Т	F	F	F	$-\frac{1}{9}\left(\gamma_F-1\right)\gamma_F\left(\gamma_G-1\right)\left(5\gamma_G-6\right)$
	F	Т	Т	Т	$-\frac{1}{9}\left(\gamma_F-1\right)\gamma_F\gamma_G\left(5\gamma_G+1\right)$
	F	Т	Т	F	$\frac{5}{9} (\gamma_F - 1) \gamma_F (\gamma_G - 1) \gamma_G$
-	F	Т	F	Т	$\frac{1}{9}\left(\gamma_F-1\right)\gamma_G\left(\gamma_F+\left(5\gamma_F-6\right)\gamma_G-3\right)$
	F	Т	F	F	$-\frac{1}{9}(\gamma_F - 1)(5\gamma_F - 6)(\gamma_G - 1)\gamma_G$
_	F	F	Т	Т	$\frac{5}{9} (\gamma_F - 1) \gamma_F (\gamma_G - 1) \gamma_G$
-	F	F	Т	F	$-\frac{1}{9}(\gamma_F - 1)\gamma_F(\gamma_G - 1)(5\gamma_G - 6)$
	F	F	F	Т	$-\frac{1}{9}(\gamma_F - 1)(5\gamma_F - 6)(\gamma_G - 1)\gamma_G$
	F	F	F	F	$\frac{1}{9}(\gamma_F - 1)(\gamma_G - 1)(-6\gamma_G + \gamma_F(5\gamma_G - 6) + 9)$

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Carnapian Logical Probability: Analogy and Similarity 1

• More generally: if a has properties $P_1 \dots P_n$, and b has $P_1 \dots P_{n-2}$, but lacks P_{n-1} , that should be of *some* relevance to b's having P_n ; n=3 case: $\Pr(Hb \mid Ha) > \Pr(Hb \mid Ha \& Fa \& Ga \& Fb \& \sim Gb)$

$$> \Pr(Hb \mid Ha \& Fa \& Ga \& \sim Fb \& \sim Gb) > \Pr(Hb)$$

- *I.e.*, Differing on 2 properties should be worse than 1, but neither should completely undermine instantial relevance. Pr[†] and Pr* violate this.
- The "analogical" idea here involves certain judgments of "similarity" of objects, where "similarity" is measured by "counting shared predicates."
- The slogan seems to be: The more properties *a* and *b* "share", the more this *enhances* instantial relevance; and, the more properties they are known to *differ* on, the more this *undermines* instantial relevance.
- Unfortunately, any theory that satisfies this analogical principle will be *language-variant*, so long as the languages contain 3 or more predicates.

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Carnapian Logical Probability: Analogy and Similarity 3

- That is, "the degree of similarity of a and b" depends sensitively on the *syntax* of the language one uses to *describe* a and b. [Note: if n = 2, there is no language-variance it requires 3 or more predicates.] Here's why.
- The ABCD language consists of four predicates A, B, C, and D. And, the XYZU language also has four predicates X, Y, Z, and U such that Xx = ||Ax| = |Bx|, Yx = ||Ax| = |Bx|, |Ax| = |Bx|, |Ax| = |Bx|, |Ax| = |Bx|, |Ax| = |Ax|, and |Ax| = |Ax|.
- ABCD and XYZU are (extra-systematically) **expressively equivalent**.

 Anything that can be said in ABCD can be said in XYZU, and conversely intuitively, there is no semantic difference between the two languages.
 - Now, consider two objects *a* and *b* such that:

Aa & Ba & Ca & Da Ab & ~Bb & Cb & Db

• Question: How similar are a and b in our "predicate-sharing" sense?

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Carnapian Logical Probability: Analogy and Similarity 2

- Let $x_1, ..., x_n$ be the objects that fall under the predicate X [*i.e.*, those in Ext(X)]. And, let $\mathfrak{s}(x_1, ..., x_n)$ be some measure of "the degree to which the objects falling under X are similar to each other".
- Carnap doesn't offer much in the way of a *theory* of $\mathfrak{s}(x_1, ..., x_n)$. But, his discussion suggests the following account of $\mathfrak{s}(x_1, ..., x_n)$:
- Let $\mathcal{P}(x)$ be the set of predicates that x falls under. Then, define:

$$\mathfrak{s}(x_1,\ldots,x_n) = \left|\bigcap_i \mathcal{P}(x_i)\right|$$

- That is, $\mathfrak{s}(x_1,\ldots,x_n)$ is the size (cardinality) of the intersection of all the $\mathcal{P}(x_i)$. This is "the size of the set of shared predicates of the x_i ".
- There is a problem with this idea. Next, I will present an argument which shows that this measure of similarity is *language variant*.

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Carnapian Logical Probability: Analogy and Similarity 4

• Answer: That depends on which of our expressively equivalent languages we use to describe *a* and *b*! To see this, note that in *XYZU* we have:

$$Xa \& Ya \& Za \& Ua$$

 $\sim Xb \& \sim Yb \& Zb \& Ub$

• Therefore, in *ABCD*, *a* and *b* share three predicates. But, in *XYZU*, *a* and *b* share only two predicates. Or, to use a modified notation, we have:

$$\mathfrak{s}_{ABCD}(a,b) = 3 \neq 2 = \mathfrak{s}_{XYZU}(a,b)$$

- On the other hand, *probabilities* should *not* be language-variant. It shouldn't matter which language you use to describe the world equivalent statements should be *probabilistically indistinguishable*.
- One consequence is that if p = ||x| = q, x = ||x| = y and z = ||x| = u, then we shouldn't have both $\Pr(p \mid x) > \Pr(p \mid u)$ and $\Pr(q \mid y) < \Pr(q \mid z)$. Carnapian principles of analogy and similarity *contradict* this requirement $(n \ge 3)$.

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Carnapian Logical Probability: Analogy and Similarity 5

- Here's a concise way of stating the general Carnapian analogical principle:
 - (*A*) If n > m, then $Pr(Xa | Xb \& \mathfrak{s}(a, b) = n) > Pr(Xa | Xb \& \mathfrak{s}(a, b) = m)$
- Applying this principle to our example yields *both* of the following:
- $(1) \Pr(Da \mid Db \& Aa \& Ba \& Ca \& Ab \& \sim Bb \& Cb) > \Pr(Da \mid Db \& Aa \& Ba \& Ca \& Ab \& \sim Bb \& \sim Cb)$
- $(2) \ \Pr(Ua \mid Ub \& Xa \& Ya \& Za \& \sim Xb \& Yb \& Zb) > \Pr(Ua \mid Ub \& Xa \& Ya \& Za \& \sim Xb \& \sim Yb \& Zb)$
- Now, let $p \stackrel{\text{def}}{=} Da$, $q \stackrel{\text{def}}{=} Ua$, and

$x \triangleq Db \& Aa \& Ba \& Ca \& Ab \& \sim Bb \& Cb$	$y \triangleq Ub \& Xa \& Ya \& Za \& \sim Xb \& \sim Yb \& Zb$
$z \triangleq Ub \& Xa \& Ya \& Za \& \sim Xb \& Yb \& Zb$	$u \triangleq Db \& Aa \& Ba \& Ca \& Ab \& \sim Bb \& \sim Cb$

- Then, p = || || q, x = || || y and z = || || u, but the Carnapian principle (A) implies $both(1) \Pr(p \mid x) > \Pr(p \mid u)$, and (2) $\Pr(q \mid y) < \Pr(q \mid z)$. Bad.
- It seems that principle (*A*) must go. Otherwise, some restriction on the choice of language is required to block inferring both (1) and (2) from it.

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A Closer Look at Inductive Logic and Applicability 2

- Popper discusses examples of the following kind (we've seen an example like this in the class before), which involve testing for a rare disease.
 - Let *E* report a positive test result for a very rare disease (for someone named John), and let *H* be the (null) hypothesis that John does *not* have the disease in question. We assume further that John knows (his *K* entails) the test is highly reliable, and that the disease is very rare.
- In such an example, it is plausible (and Carnap should agree) that (to the extent that Pr_T is *applicable* to modeling the *epistemic* relations here):
- (1) $Pr_T(H \mid E \& K)$ is very high.
- (2) But, $Pr_T(H \mid E \& K) < Pr_T(H \mid K)$.
- Because of (2), it would be odd to say that *E supports H* (for John) in this context. (2) suggests that *E* is (intuitively) evidence *against H* here.
- But, because of (1), Carnap's (RTE_C) implies that *E* supports *H* (for John) here. This looks like a counterexample to [the ← of] Carnap's (RTE_C).

A Closer Look at Inductive Logic and Applicability 1

- Carnap suggested various *bridge principles* for connecting inductive logic and inductive epistemology. The most well-known of these was:
 - **The Requirement of Total Evidence.** In the application of inductive logic to a given knowledge situation, the total evidence available must be taken as a basis for determining the degree of evidential support.
- A more precise way of putting this principle is:
- (RTE) E evidentially supports H for an agent S in an epistemic context $C \iff \mathfrak{c}(H, E \mid K) > r$, where K is S's total evidence in C.
- For Carnap, $c(H, E \mid K) = Pr_T(H \mid E \& K)$, where Pr_T is a suitable "logical" probability function. So, we can restate *Carnap's* (RTE) as follows:
- (RTE_C) E evidentially supports H for an agent S in an epistemic context $C \iff \Pr_{\mathsf{T}}(H \mid E \& K) > r$, where K is S's total evidence in C.
- Carnap's version of (RTE) faces a challenge (first articulated by Popper) involving *probabilistic relevance vs high conditional probability*.

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A Closer Look at Inductive Logic and Applicability 3

- This suggests the following refinement of Carnap's (RTE $_C$):
- (RTE'_C) E evidentially supports H for an agent S in an epistemic context C $\Longrightarrow \Pr_{\mathsf{T}}(H \mid E \& K) > \Pr_{\mathsf{T}}(H \mid K)$, where K is S's total evidence in C.
- In other words, (RTE'_C) says that *evidential support in (for S in C) implies probabilistic relevance, conditional upon K* (for a suitable Pr_T function).
- Note: this only states a *necessary* condition for evidential support.
- While (RTE'_C) avoids Popper's objection, it faces serious challenges of its own (e.g., Goodman's "Grue" example more on that later). Here's one:
- Consider any context in which *S* already knows (*with certainty*) that *E* is true. That is, *S*'s total evidence in the context *entails* $E(K \Vdash E)$.
- In such a case, Pr(H | E & K) = Pr(H | K), for *any* probability function Pr.
 Thus, (RTE'_C) implies that *E cannot support anything* (for *S*, in any such
 C). This shows that (RTE'_C) isn't a correct principle either. ["Old Evidence"]

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A Closer Look at Inductive Logic and Applicability 4

- I think this whole way of approaching inductive logic is wrongheaded.
- First, why must Pr *itself* be logical, if c (which is defined in terms of Pr) is to be logical? Analogy: must the truth-value assignment function v *itself* be logical, if ⊨ (which is defined in terms of v) is to be logical?
 - But: there is a crucial disanalogy here, which I will discuss below.
- Second, Carnap's proposal $c(H, E \mid K) = \Pr_T(H \mid E \& K)$ is suspect, because (as Popper pointed out) it is not sensitive to *probabilistic relevance*.
 - Note: this undermines Carnap's argument for the "logicality" of (PI).
- Third, the applicability desideratum (\mathcal{D}_3) may be fundamentally misguided. The search for logic/epistemology "bridge principles" is fraught with danger, even in the *deductive* case. And, since IL is supposed to *generalize* DL, it will also face these dangers *and new ones* (as above).
 - I think this is the true (but, surprisingly, un-appreciated) lesson of Goodman's "grue" example. I will explain why in the confirmation unit.

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An Alternative Conception of Inductive Logic 2

- Here's the idea. Confirmation is a *four*-place relation, between *E*, *H*, *K*, *and a probability function* Pr. The resulting relation is still *logical* in Carnap's sense, since, *given* a choice of Pr, c is logically (mathematically, if you prefer) determined, provided only that c is defined in terms of Pr.
- So, on this conception, desiderata (\mathcal{D}_1) and (\mathcal{D}_2) are satisfied.
- As usual, the subtle questions involve the applicability desideratum (\mathcal{D}_3).
- What do we say about that? Well, I think any naive "bridge principle" like (RTE or RTE') is doomed to failure. But, perhaps there is *some* connection.
- Thinking back to the deductive case, there may be *some* connection between deductive *logic* and deductive *inference*. But, what is it?
- This is notoriously difficult to say. The best hope seems to be that there is *some* connection between *knowledge* and entailment. [Note: connecting or bridging *justified belief* and *entailment* seems much more difficult.]

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An Alternative Conception of Inductive Logic 1

• In light of the above considerations, we might seek a measure c satisfying the following (provided that *E*, *H*, and *K* are *logically contingent*):

$$\mathsf{c}(H,E\,|\,K) \text{ should be} \begin{cases} \mathsf{Maximal}\;(>0,\,\mathsf{constant}) & \Leftarrow E\,\&\,K \vDash (\mathsf{or}\,\Vdash)\,H. \\ > 0\;(\mathsf{confirmation}\;\mathsf{rel.}\;\mathsf{to}\,K) & \Rightarrow \mathsf{Pr}(H\,|\,E\,\&\,K) > \mathsf{Pr}(H\,|\,K). \\ = 0\;(\mathsf{irrelevance}\;\mathsf{rel.}\;\mathsf{to}\,K) & \Rightarrow \mathsf{Pr}(H\,|\,E\,\&\,K) = \mathsf{Pr}(H\,|\,K). \\ < 0\;(\mathsf{disconfirmation}\;\mathsf{rel.}\;\mathsf{to}\,K) & \Rightarrow \mathsf{Pr}(H\,|\,E\,\&\,K) < \mathsf{Pr}(H\,|\,K). \\ \mathsf{Minimal}\;(<0,\,\mathsf{constant}) & \Leftarrow E\,\&\,K \vDash (\mathsf{or}\,\Vdash) \sim\!H. \end{cases}$$

- Carnap would add: "and Pr should be a 'logical' probability function Pr_T ". But, I suggested that this was a mistake. OK, but then what do *I* say about the Pr's above? There is an implicit quantifier over the Pr's above...
 - $-\exists$ is *too weak* a quantifier here, since there will *always* be *some* such Pr.
 - \forall is *too strong* a quantifier here, because that is *demonstrably false*!
 - What's the alternative? The alternative is that Pr is a *parameter* in c itself. That is, perhaps Pr is simply an *argument* of the function c.

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An Alternative Conception of Inductive Logic 3

- Many people think there is *some* connection between knowledge and entailment. But, simple/naive version of bridge principles don't work. Here is a progression of increasingly subtle "bridge principles":
 - 1. If *S* knows *p* and $p \models q$, then *S* knows *q*.
 - What if *S* doesn't *know* that $p \models q$?
 - 2. If *S* knows p and *S* knows $p \models q$, then *S* knows q.
 - What if *S* doesn't even *believe q*? [Discuss Dretske's Zebra Example.]
 - 3. If *S* knows *p* and *S* knows $p \models q$ and *S* believes *q*, then *S* knows *q*.
 - What if *S* believes *q* for reasons *other than p*?
 - 4. If *S* knows p and *S* knows $p \models q$ and *S* comes to believe q because of their (initial) belief that p, then *S* knows q.
 - What if *S no longer* believes *p*, *while/after* they infer *q*?
 - 5. If S knows p & S knows $p \models q \& S$ competently deduces q from p (thus coming to believe q) while maintaining their belief in p, then S knows q.