

A DEFENSE OF TEMPERATE EPISTEMIC TRANSPARENCY

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FEW 2011 - USC

EPISTEMIC TRANSPARENCY

- If S knows that p , S knows that she knows that p :
 - *KK Principle:* $Kp \rightarrow KKp$
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- *Knowledge reflexivity*
 - *Positive introspection*
 - *Self-knowledge*
 - *Transparency*
 - *Luminosity*

GOAL

- A defense of a moderate version of *KK*

RISE AND FALL OF KK

- ◉ 1960s: Dogma

Then....

- ◉ Externalism - e.g.: Reliabilism
- ◉ Williamson (2000)

STRATEGY

- ◉ (A) Why do we want transparency?
- ◉ (B) Indirect argument

WHY DO WE CARE ABOUT TRANSPARENCY?

- ◉ Ideal agents
 - Ideally *rational*?

WHY DO WE CARE ABOUT TRANSPARENCY?

- Knowledge and responsibility
 - How?

WHY DO WE CARE ABOUT TRANSPARENCY?

- Responsibility demands us to be in an appropriate reflective state.
- *What* reflective state?
 - Epistemic responsibility entails “ratifiability”.

A MODAL FRAMEWORK

- $\mathcal{F} = \langle W, R, P_{prior} \rangle$
- $K\varphi = \{w \in W : \forall x \in W (wRx \rightarrow x \in \varphi)\}$
- $R(w) = \{x \in W : wRx\}$

WILLIAMSON: IMPROBABLE KNOWING

- $P_w(\varphi)$: the evidential probability of φ in w .
- $P_w(\varphi) = P_{prior}(\varphi \mid R(w)) = P_{prior}(\varphi \cap R(w)) / P_{prior}(R(w))$
- $P_w(R(w)) = 1$
- $[P(\varphi) = r] =_{\text{def.}} \{w \in W: P_w(\varphi) = r\}$

IMPROBABLE KNOWING

- “The *KK* principle is equivalent to the principle that if the evidential probability of p is 1, then the evidential probability that the evidential probability of p is 1 is itself 1” (Williamson, p. 8).
- We can build a model in which $P_w([P(R(w)) = 1])$ is as low as we want.

PROBLEMS

*Why should we say that the evidential basis
is always $R(w)$?*

PROBLEMS

Recall that:

- $[P(\varphi) = r] = \{w \in W : P_w(\varphi) = r\}$

- $[P_w(\varphi) = r]$?

PROPOSAL (FIRST VERSION)

- We'll have a sequence of languages $L^0, L^1, \dots L^n$...with probability operators $P^0, \dots P^n \dots$
- We'll have a sequence of functions $P^1_w \dots P^n_w \dots$ on sentences φ^i of language L^i
- $P^i_w: L^{i-1} \rightarrow \mathbb{R}$

PROPOSAL (FIRST VERSION)

- Expressions of the form $P_{prior}(\varphi)$ or $P^i_w(\varphi)$ do not belong to any language of the sequence $L^0, L^1 \dots L^n$
- “ $P^i(\varphi)=r$ ” is true in w iff $P^i_w(\varphi)=r$.

PROPOSAL (FIRST VERSION)

- How should we conditionalize?

PROPOSAL (FIRST VERSION)

- ◉ For $P^1_w(\varphi)$, the relevant evidence is $R(w)$.
- ◉ For $P^2_w(\underline{P^1(\varphi)=r})$, the relevant evidence is $KR(w)$.

CONDITIONALIZATION (FIRST VERSION)

- C* rule:

For $i \geq 1$: $P_w^i (\underline{P^{i-1}(\dots P(\varphi)=r\dots)}) =$

$P_{\text{prior}} (\underline{P^{i-1}(\dots P(\varphi)=r\dots)} \mid \underline{K^{i-1}\dots KR(w)})$

where K^{i-1} is the same \underline{K} -operator iterated $i-1$ times

DIFFICULTIES

- C^* divorces probability 1 from knowledge.

A MODEL FOR MODERATE TRANSPARENCY (SECOND VERSION)

- $M = \langle W, R^1, \dots, R^n, P_{prior}, v \rangle$
- New operators $\underline{K^0} \dots \underline{K^n} \dots$, in addition to $\underline{P^0}, \dots, \underline{P^n} \dots$
- We define a sequence of relations $R^1 \dots R^n$ which correspond to the different Ks.
 - The R s are nested: $R^i \subseteq R^{i-1} \dots \subseteq R^1$
 - R^i is a reflexive relation over W , for all i , and transitive for $i > 1$.

A MODEL FOR MODERATE TRANSPARENCY

- Our conditionalization rule now incorporates operators $\underline{K^1}, \dots \underline{K^n} \dots$ defined on the basis of relations $R^1, \dots R^n \dots$
- C** rule:

For $i \geq 1$: $P_w^i (\underline{P^{i-1}(\dots P(\varphi)=r\dots)}) =$

$P_{\text{prior}} (\underline{P^{i-1}(\dots P(\varphi)=r}) \mid \underline{K^{i-1} \dots KR(w)})$

where “ $\underline{K^{i-1} \dots KR(w)}$ ” includes $i-1$ higher-order \underline{K} -operators

A MODEL FOR MODERATE TRANSPARENCY

- Intended interpretation of the formalism:
 - $(^*)K^2p$ does not make sense!
 - A second-order evidential probability claim is the evidential probability of a probability statement.
 - *Mutatis mutandis* for higher-order levels and for *conditional* evidential probabilities.

SOME CONSEQUENCES

- Why should we demand such requirements for the R s?
They are not *ad hoc*!
- Higher-order probability requires increasingly complex probabilistic claims.
- For second-order evidential probability in w :
 - We conditionalize over $KR(w)$
 - Thus the second-order probability of $R(w)$ is 1
 - Thus the agent *knows* that $KR(w)$
 - $K^2KR(w)$ should be true in w

SOME CONSEQUENCES

- $\underline{K\varphi \rightarrow K^2K\varphi}$ **KK^2 Principle**
 - (if $[K^2KR(w)]$ is not empty, for any w)
- $\Diamond (\underline{K\varphi \rightarrow K^2K\varphi})$ **KK^\Diamond Principle**
 - (if $[K^2KR(w)]$ is not empty, for some w)
- $\underline{K^2K\varphi \rightarrow K^3K^2K\varphi}$ **KK^+ Principle**

SOME CONSEQUENCES

- A restricted version of possitive introspection holds:
→ *Quasi-transparency principles*
- KK^+ , KK^\Diamond and KK^2 result from conditionalizing over higher-order levels of evidence and from the attempt to adjust probability language and knowledge attribution in a progressively coherent way.

SOME CONSEQUENCES

- Links between lower- and higher-order probabilities.
 - If $P^1_w(\varphi) = r = 1$ or 0 , then $P^2_w(\underline{P^1(\varphi)=r}) = 1$.
 - If R^1 is an equivalence relation, $P^2_w(\underline{P^1(\varphi)=r})$ is either 1 or 0 .
 - Suppose $P^2_w(\underline{P(\varphi)=r}) = s$. If $0 \neq r \neq 1$ and R^1 is not transitive, then s need not be either 1 or 0 .

ON THE PROBABILISTIC REFLECTION PRINCIPLE (PRP)

- *PRP:*

$$P^2_w (\varphi \mid \underline{P^1(\varphi)=r}) = r \text{ (for } w \in W)$$

- *Iterated PRP:*

$$P^i_w (\varphi \mid \underline{P^{i-1}(\varphi \mid P^{i-2}(\varphi \mid \dots) \dots) = r}) = r$$

- Is PRP a theoretical truth of M ?

ON PROBABILISTIC REFLECTION

- Necessary and sufficient condition for Iterated *PRP*

$R^i = R^{i-1} \in M$ is an equivalence relation

iff

for all $w \in W$ and any $\varphi \in L^0$:

if $P^{i+1}_w(\cdot \mid \cdot)$ exists, then

$$P^{i+1}_w(\varphi \mid P^i(\varphi \mid P^{i-1}(\varphi \mid \dots) \dots) = r) = r$$

RELATION TO OTHER WORK

Paul Egré/ Jérôme Dokic

- Principal motivation: to deactivate Williamson's soritic argument on inexact knowledge
- Perceptual vs. reflective knowledge
- (KK') $K_\pi\varphi \rightarrow KK_\pi\varphi$
- Transparency failures do not generalize

RELATION TO OTHER WORK

Differences

- 1. Egré/ Dokic do not offer a probabilistic framework.
 - 2. They focus on reflection over perceptual knowledge, exclusively.
 - 3. KK' Principle is imposed “from the outside”.
- The present model for quasi-transparency can be seen as a refinement and extension of some aspects of the system suggested by Egré - Dokic.

CONCLUSIONS

- Once we clarify some conceptual aspects of higher-order probabilities...
- ...we obtain the vindication of a number of introspective principles, or *principles of quasi-transparency*.
- Quasi-transparency principles were not just assumed to hold, but they have been obtained as a result of implementing a number of natural constraints on the structure of the system.
→ Formally speaking, they behave quite differently from presuppositions of consistency or deductive closure.

CONCLUSIONS

- The framework vindicates the intuition that first- and second-order knowledge differ substantially:
 - Different attitudes about ignorance
 - Different attitudes toward “margin of error” principles
 - Second-order knowledge is concerned with the “ratification” of first-order attitudes.

CONCLUSIONS

- Quasi-transparency fully vindicates the normative link between self-knowledge and responsibility.

→ K^+Kp : “responsible knowledge” of p .

CONCLUSIONS

- Second-order knowledge, as a state of epistemic responsibility, is a desideratum we have *qua* agents.

Thank you!