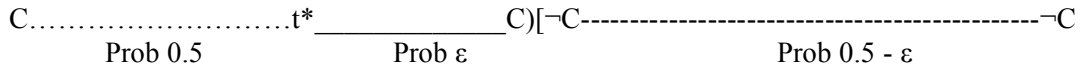


## Why Epistemology Can't be Operationalized

- (L) Condition C is luminous if and only if in every case in which C obtains, one is in a position to know that C obtains.
- (O0) C is luminous. Assumption
- (O1<sub>i</sub>) If at  $t_i$  one is in a position to know that C obtains, then at  $t_{i+1}$  C obtains. Assumption
- (O2<sub>i</sub>) If at  $t_i$  C obtains, then at  $t_i$  one is in a position to know that C obtains. O0, LP
- (O3<sub>i</sub>) At  $t_i$  C obtains. Assumption
- (O4<sub>i</sub>) At  $t_i$  one is in a position to know that C obtains. O2<sub>i</sub>, O3<sub>i</sub>, MP
- (O3<sub>i+1</sub>) At  $t_{i+1}$  C obtains. O1<sub>i</sub>, O4<sub>i</sub>, MP
- (O3<sub>0</sub>) At  $t_0$  C obtains. Assumption
- (O3<sub>n</sub>) At  $t_n$  C obtains.  $O3_i \rightarrow O3_{i+1}, i=0, \dots, n-1$
- (O5) At  $t_n$  C does not obtain. Assumption
- (LP1) Condition C is *luminous in probability 1* if and only if in every case in which C obtains, the probability that C obtains is 1.
- (0) C is luminous in probability 1. Assumption
- (1<sub>i</sub>) If at  $t_i$  the probability that C obtains is 1, then at  $t_{i+1}$  C obtains. Assumption
- (2<sub>i</sub>) If at  $t_i$  C obtains, then at  $t_i$  the probability that C obtains is 1. 0, LP
- (3<sub>i</sub>) At  $t_i$  C obtains. Assumption
- (4<sub>i</sub>) At  $t_i$  the probability that C obtains is 1. 2<sub>i</sub>, 3<sub>i</sub>, MP
- (3<sub>i+1</sub>) At  $t_{i+1}$  C obtains. 1<sub>i</sub>, 4<sub>i</sub>, MP
- (3<sub>0</sub>) At  $t_0$  C obtains. Assumption
- (3<sub>n</sub>) At  $t_n$  C obtains.  $3_i \rightarrow 3_{i+1}, i=0, \dots, n-1$
- (5) At  $t_n$  C does not obtain. Assumption

- (0) C is luminous in probability 1. Assumption
- (1) If C obtains at no time in a nonempty interval  $(t, t^{**})$ , then at no time in the interval  $(t-\epsilon, t^{**}+\epsilon)$  is the probability that C obtains 1. Assumption
- (2) If at no time in the interval  $(t-\epsilon, t^{**}+\epsilon)$  is the probability that C obtains 1, then C obtains at no time in the interval  $(t-\epsilon, t^{**}+\epsilon)$ . 0, LP
- (6) If C obtains at no time in a nonempty interval  $(t, t^{**})$ , then C obtains at no time in the interval  $(t-\epsilon, t^{**}+\epsilon)$ . 1, 2
- (7<sub>0</sub>) C obtains at no time in the interval  $(t-\delta, t)$ . Assumption
- (7<sub>i</sub>) C obtains at no time in the interval  $(t-\delta-i\epsilon, t+i\epsilon)$ . (6), (7<sub>0</sub>)
- (7) C obtains at no time in the interval  $[t_0, t_n]$ . (7<sub>i</sub>)
- (3<sub>0</sub>) At  $t_0$  C obtains. Assumption

(LP<sub>x</sub>) Condition C is *luminous in probability x* if and only if in every case in which C obtains, the probability that C obtains is x.



(LP<sub>>x</sub>) Condition C is *luminous in probability > x* if and only if in every case in which C obtains, the probability that C obtains is more than x.

In every case in which C obtains, the probability that C obtains is more than 0.5.

Expectation (finite case)  $E(X) = \sum_{x \in I} xP(X = x)$

Example. Time is discrete. At each time  $t$ :  $P(T = t-1) = P(T = t) = P(T = t+1) = 1/3$ . Thus always  $E(T) = T$ .

There are countably many atomic variables  $X, Y, Z, \dots$  (informally, denoting real numbers; they correspond to random variables).

For each rational number  $c$  there is an atomic constant  $[c]$  (informally, denoting  $c$ ).

If  $T$  and  $U$  are terms then  $T+U$  is a term ( $+$  is informally read as ‘plus’).

If  $A$  is a formula then  $\neg A$  is a formula ( $\neg$  is informally read as ‘it is not the case that’).

If  $A$  and  $B$  are formulas then  $A \& B$  is a formula ( $\&$  is informally read as ‘and’).

If  $A$  is a formula then  $V(A)$  is a term ( $V$  is informally read as ‘the truth-value of’).

If  $T$  and  $U$  are terms then  $T \leq U$  is a formula ( $\leq$  is informally read as ‘is less than or equal to’).

If  $T$  is a term then  $E(T)$  is a term ( $E$  is informally read as ‘the expectation of’).

‘ $T=U$ ’ abbreviates  $T \leq U \& U \leq T$ .

‘ $T < U$ ’ abbreviates  $T \leq U \& \neg U \leq T$

‘ $P(A)$ ’ abbreviates  $E(V(A))$  ( $P$  is informally read as ‘the probability of’).

A *model* is a triple  $\langle W, \text{Prob}, F \rangle$ , where  $W$  is a nonempty set,  $\text{Prob}$  is a function from members of  $W$  to probability distributions over  $W$  and  $F$  is a function from atomic terms to their intensions (functions from members of  $W$  to real numbers).

$\text{val}(w, A)$ : truth-value at  $w$  of  $A$

$\text{den}(w, T)$ : denotation at  $w$  of  $T$ ,  $\text{den}(w, T)$

If  $T$  is an atomic variable,  $\text{den}(w, T) = F(T)(w)$ .

If  $c$  is a rational number,  $\text{den}(w, [c]) = c$ .

If  $T$  and  $U$  are terms,  $\text{den}(w, T+U) = \text{den}(w, T) + \text{den}(w, U)$ .

If  $A$  is a formula,  $\text{val}(w, \neg A) = 1 - \text{val}(w, A)$ .

If  $A$  and  $B$  are formulas,  $\text{val}(w, A \& B) = \min\{\text{val}(w, A), \text{val}(w, B)\}$ .

If  $A$  is a formula,  $\text{den}(w, V(A)) = \text{val}(w, A)$ .

If  $T$  and  $U$  are terms,  $\text{val}(w, T \leq U) = 1$  if  $\text{den}(w, T) \leq \text{den}(w, U)$  and 0 otherwise.

If  $T$  is a term,  $\text{den}(w, E(T))$  is the expectation with respect to the probability distribution  $\text{Prob}(w)$  of the random variable whose value at each world  $x \in W$  is  $\text{den}(x, T)$ .

For example if  $W$  is finite, then  $\text{den}(w, E(T)) = \sum_{x \in W} \text{den}(x, T) \text{Prob}(w)(\{x\})$

Count  $\langle W, \text{Prob}, F \rangle$  as a model only if  $\text{den}(w, E(T))$  is well-defined for every  $w \in W$  and term  $T$ .

Example.  $W = \{u, v, w\}$ .

$\text{Prob}(u)(\{u\}) = \text{Prob}(u)(\{v\}) = 1/2$        $\text{Prob}(u)(\{w\}) = 0$

$\text{Prob}(v)(\{u\}) = \text{Prob}(v)(\{v\}) = \text{Prob}(v)(\{w\}) = 1/3$

$\text{Prob}(w)(\{u\}) = 0$        $\text{Prob}(w)(\{v\}) = \text{Prob}(w)(\{w\}) = 1/2$

$\text{den}(u, X) = F(X)(u) = 8$ ;  $\text{den}(v, X) = F(X)(v) = 4$ ;  $\text{den}(w, X) = F(X)(w) = 0$

$\text{den}(u, E(X)) = 6$ ,  $\text{den}(v, E(X)) = 4$ ,  $\text{den}(w, E(X)) = 2$

$\text{den}(u, E(E(X))) = 5$ ,  $\text{den}(v, E(E(X))) = 4$ ,  $\text{den}(w, E(E(X))) = 3$

$E(E(X)) = E(X)$ ?

$\text{Var}(X) = E((X - E(X))^2)$

Example. Where  $E(T) = T$ .

Wide-scope calculation:  $\text{Var}(T) = ((T-1-T)^2 + (T-T)^2 + (T+1-T)^2)/3 = 2/3$ .

Narrow-scope calculation: Since certainly  $E(T) = T$ , 'Var'(T) = 0.

**SORITES**      For any  $w, x \in W$ , there are  $w_0, \dots, w_n \in W$  such that  $w = w_0, x = w_n$  and for  $0 \leq i, j \leq n$ , if  $|i-j| \leq 1$  then  $\text{Prob}(w_i)(\{w_j\}) > 0$ .

**Proposition.** Let the set of worlds be finite and SORITES hold. Then for any random variable  $T$ , if  $E(T) = T$  at every world then  $T$  has the same value at every world.

**UNCMAX**      For some  $w \in W$ : for all  $x \in W$ ,  $\text{den}(x, T) \leq \text{den}(w, T)$ , but  $\text{Prob}(w)(\{x: \text{den}(x, T) < \text{den}(w, T)\}) > 0$ .

**Proposition.** Let the set of worlds be finite and SORITES hold. Then for any random variable  $T$ , the sequence  $T, E(T), E(E(T)), \dots$  converges to the same limit whichever world one is at.

TW 18.v.2006