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| **ITU-T** | **X.609.7** | |
| TELECOMMUNICATION STANDARDIZATION SECTOR OF ITU | | (12/2018) |
|  | SERIES X: DATA NETWORKS, OPEN SYSTEM COMMUNICATIONS AND SECURITY  OSI networking and system aspects – Networking | | | |
|  | **Managed P2P communications: Content distribution peer protocol** | | | |
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**Recommendation ITU-T X.609.7 (ex.** **X.mp2p-cdpp)**

**Managed P2P communications: Content distribution peer protocol**

**Summary**

Recommendation ITU-T X.609.7 specifies a content distribution peer protocol (CDPP) that runs on the interface among entities of managed P2P communications. CDPP is used to distribute one or more contents to a number of peers. Content distribution over traditional P2P communications has incurred various issues such as distribution of illegal content, uncontrollable participation, and synchronized distribution of the updated contents. Different from the content distribution over the traditional P2P communications which is not capable of providing manageability, the content distribution over managed P2P communications can be managed by content provider or service provider. In the content distribution over managed P2P communications, as an example, participation in an overlay network can be controlled so that only predefined peers can join the overlay network and distribute contents each other. In addition, content to be distributed over an overlay network can be updated anytime, and every update will be applied to all peers in the overlay network. The protocol is capable of managing content distribution under control of content provider or service provider. This Recommendation provides protocol operations, and message formats for content distribution over managed P2P network.

**Keywords**

<Managed P2P, content distribution, peer protocol>

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**Draft new Recommendation ITU-T X.609.7**

**Managed P2P communications: Content distribution peer protocol**

**Summary**

Recommendation ITU-T X.609.7 specifies a content distribution peer protocol (CDPP) that runs on the interface among entities of managed P2P communications. CDPP is used to distribute one or more contents to a number of peers. Content distribution over traditional P2P communications has incurred various issues such as distribution of illegal content, uncontrollable participation, and synchronized distribution of the updated contents. Different from the content distribution over the traditional P2P communications which is not capable of providing manageability, the content distribution over managed P2P communications can be managed by content provider or service provider. In the content distribution over managed P2P communications, as an example, participation in an overlay network can be controlled so that only predefined peers can join the overlay network and distribute contents each other. In addition, content to be distributed over an overlay network can be updated anytime, and every update will be applied to all peers in the overlay network. The protocol is capable of managing content distribution under control of content provider or service provider. This Recommendation provides protocol operations, and message formats for content distribution over managed P2P network.

**Keywords**

<Managed P2P, content distribution, peer protocol>

# Scope

This Recommendation specifies a content distribution protocol over managed P2P communications. It describes the following details:

* overview of content delivery peer protocol;
* protocol messages and its parameters;
* protocol behaviours

# References

The following ITU-T Recommendations and other references contain provisions, which, through reference in this text, constitute provisions of this Recommendation. At the time of publication, the editions indicated were valid. All Recommendations and other references are subject to revision; users of this Recommendation are therefore encouraged to investigate the possibility of applying the most recent edition of the Recommendations and other references listed below. A list of the currently valid ITU-T Recommendations is regularly published. The reference to a document within this Recommendation does not give it, as a stand-alone document, the status of a Recommendation.

[ITU-T X.609] Recommendation ITU-T X.609 (2015), *Managed P2P communications: Functional architecture*

[ITU-T X.609.6] Recommendation ITU-T X.609.6 (2018), *Managed P2P communications: Content delivery signalling requirements*

[IETF RFC 7159] IETF RFC 7159 (2014), *The JavaScript Object Notation (JSON) Data Interchange Format*

# Definitions

## Terms defined elsewhere

This Recommendation uses the following terms defined elsewhere:

**3.1.1 Overlay network** [ITU-T X.1162]: An overlay network is a virtual network that runs on top of another network. Like any other network, the overlay network comprises a set of nodes and links between them. Because the links are logical ones, they may correspond to many physical links of the underlying network.

**3.1.2 Peer** [ITU-T X.1161]: Communication node on P2P network that functions simultaneously as both "client" and "server" to the other nodes on the network.

**3.1.3 Peer-to-peer (P2P)** [ITU-T Y.2206]: A system is considered to be P2P if the nodes of the system share their resources in order to provide the service the system supports. The nodes in the system both provide services to other nodes and request services from other nodes.

NOTE– Peer is the node in a P2P system.

**3.1.4 Managed P2P** [ISO/IEC TR 20002]: P2P with manageability features to manage the P2P-based service and P2P network by the P2P participants such as P2P service provider, ISP, and peer.

**3.1.5 Buffermap** [ITU-T X.609]: A map showing downloading status of fragments comprising a shared content.

**3.16 Fragment** [ITU-T X.609]: A piece of the shared content.

**3.1.7 Fragmentation** [ITU-T X.609]: A process that divides the shared content into multiple fragments for sharing the content in a distributed manner.

**3.1.8 Source peer** [ITU-T X.609.3]:A peer that streams the multimedia contents to the overlay network. The peer only provides content data to other peers and does not receive it. This peer generates fragments using the multimedia data received from the content source.

**3.1.9 Client peer** [ITU-T X.609.3]:A peer that sends fragments received from other peers to other peers, and does not generate its own fragments.

## Terms defined in this Recommendation

This Recommendation defines the following terms:

None

# Abbreviations and acronyms

This Recommendation uses the following abbreviations:

|  |  |
| --- | --- |
| CDPP | Content distribution peer protocol |
| CP | Completed point |
| CS | Completed Section |
| DP | Downloading point |
| DS | Downloading Section |
| IXS | Index server |
| MP2P | Managed peer-to-peer |
| SP | Start point |
| TCP | Transmission control protocol |
| UDP | User Datagram Protocol |

# Conventions

None

# Overview

Figure 6-1 shows the framework and reference points of managed peer-to-peer (MP2P) communications defined in [ITU-T X.609].

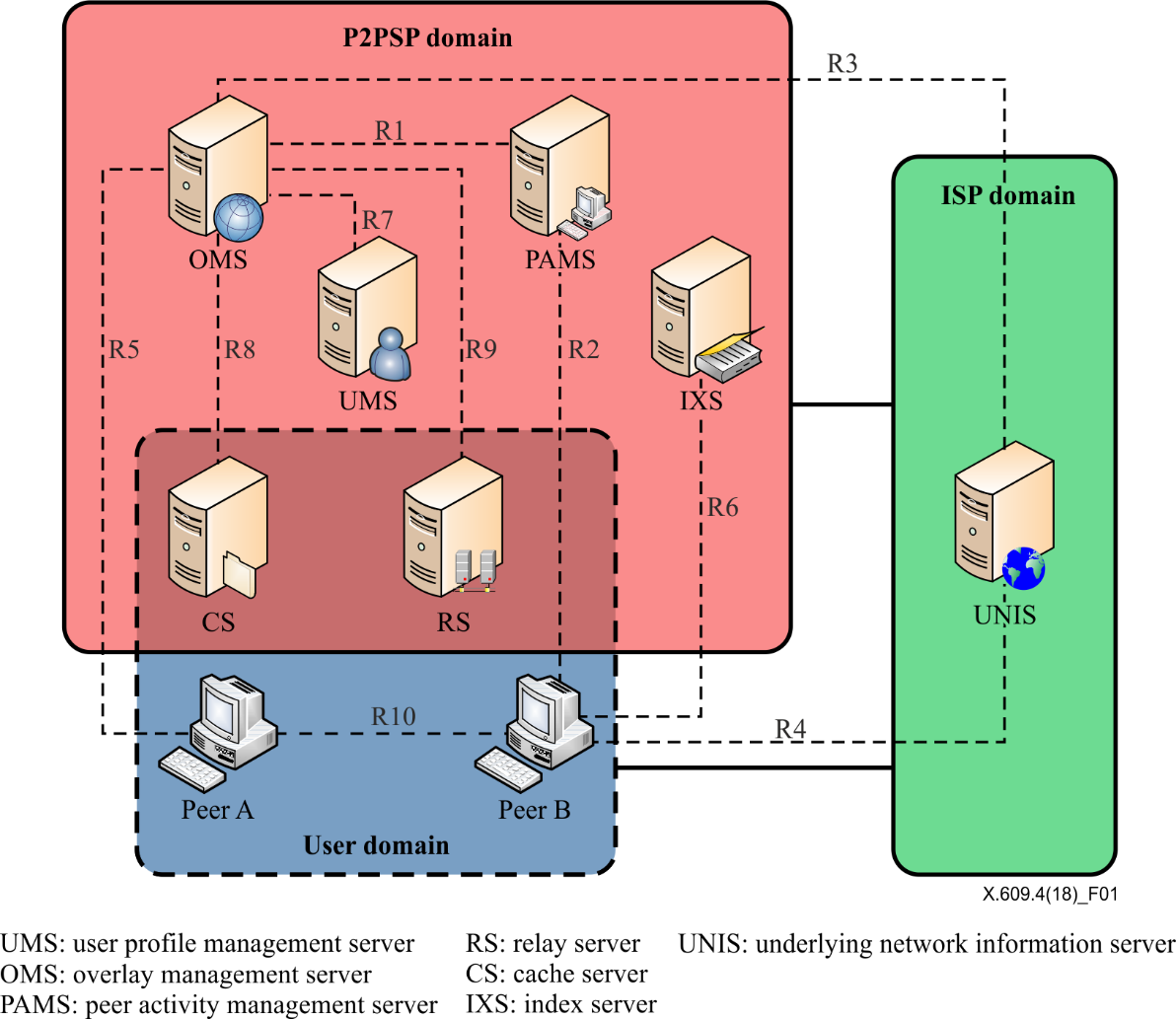


Figure 6-1 – Framework and reference points of managed P2P [ITU-T X.609]

Based on the functional architecture defined in [ITU-T X.609], [ITU-T X.609.6] describes signalling requirements for content distribution services. Figure 6-2 shows an architectural overview of content distribution service over MP2P communications defined in [ITU-T X.609.7]. This Recommendation defines a protocol for the reference point R10.



Figure 6-2 – Functional architecture of MP2P communications for content distribution [ITU-T X.609.7]

# 7 Messages and behaviours

# Messages and behaviours

This clause describes messages and procedures for content distribution service based on MP2P communications defined in [ITU-T X.609.7]. This Recommendation defines the peer protocol for reference point R10 of the functional architecture depicted in Figure 6-2.

* 1. **Buffer and buffermap**

A peer generates a buffer when it starts content distribution. A source peer willing to distribute a content has the buffer fully filled with entire fragments of the content. Client peer starts with an empty buffer, and the buffer is gradually filled with the fragments by communication with the source peer and other client peers. For the communication, each peer maintains the information of own buffer and shares the information with other peers in order to find which fragments it needs to get and which peer can send the fragments. The information is called as buffermap.

### 7.1.1 Structure of a buffer

Figure 7-1 shows the structure of a buffer.

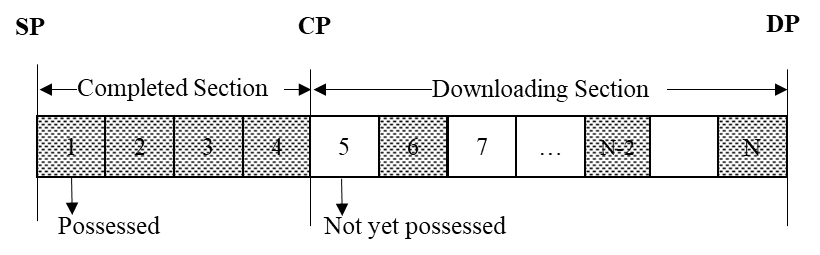


Figure 7-1 – Structure of buffer

A buffer is composed of the following sections;

* **Completed section (CS)**: CS contains the contiguous fragments received from other peer(s). Thus a peer can requests any fragment in the CS to other peer;
* **Downloading section (DS):** DS includes zero or more fragments received from other peer(s). Thus a peer needs a buffermap describing which fragments are in DS;

The following points are used to represent the start or end of each section in a buffer:

* **Completed point (CP)**: CP points to the end of the completed section in the buffermap;
* **Downloading point (DP):** DP points to the end of the downloading section in the buffermap;
* **Start point (SP)**: SP points to the identifier at the beginning of the buffermap, and is always set to zero.

### 7.1.2 Structure of a buffermap

A buffermap of a buffer represents the possession status of each fragment of a content within a downloading section. The possessed fragment is represented as one, and the fragment not yet possessed is represented as zero. The buffermap corresponding to the buffer shown in Figure 7-1 is depicted in Figure 7-2. A buffermap is used as a payload of HELLO message, PELLO message, and BUFFERMAP message.

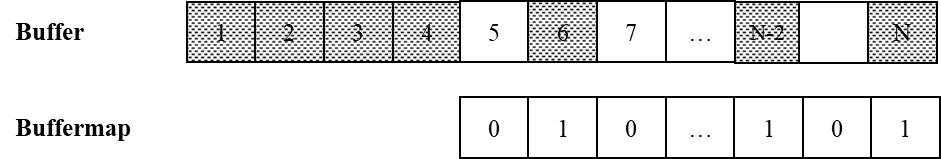


Figure 7-2 – A buffermap corresponding to an example buffer

* 1. **Messages**

This clause defines protocol messages for content distribution service, and the messages are represented as JSON [IETF RFC7159]. Since JSON itself cannot accommodate binary data, messages are converted to binary format such as BSON [b-BSON 1.1].

* + 1. **HELLO**

HELLO message is used for establishing a connection between two peers. After client peer A establishes a transmission control protocol (TCP) connection with client peer B, it sends a HELLO message, and then client peer B responds with a HELLO message. The syntax of the message is as follows;

{

"method":"HELLO",

"proto-version":integer,

"index-version":integer,

"peer-id":string,

"overlay-id":string,

"sp-index":integer,

"cp-length":integer,

"dp-index":integer,

"ds-length":integer,

"buffermap":binary

}

The semantic of each field in the message is as follows:

* *proto-version* is the version of the protocol;
* *index-version* is the version of the index file;
* *peer-id* is the identifier of the peer;
* *overlay-id* is the identifier of overlay network;
* *sp-index* indicates the starting index number of fragment represented by the buffermap;
* *cp-length* indicates the length of completed section within buffermap beginning *sp-index*;
* *dp-index* indicates the starting index number of downloading section;
* *ds-length* indicates the length of downloading section;
* *buffermap* represents the possession status of fragments within downloading section. If the peer possesses particular fragment, it set to ‘1’, and set to ‘0’, if not.
  + 1. **PELLO**

PELLO message is used to establish a connection with the peer behind a network address translation (NAT) or a firewall. If a source peer is located behind a NAT or a firewall, client peers cannot connect to the source peer. Client peer A receiving this message from client peer B acts as if it sent a HELLO message to client peer B and uses the same TCP connection in order to request and receive fragments. A source peer can actively establish TCP connections with client peers and sends this message to the client peers. The syntax of the message is as follows;

{

"method":"PELLO",

"proto-version":integer,

"index-version":integer,

"peer-id":string,

"overlay-id":string,

"sp-index":integer,

"cp-length":integer,

"dp-index":integer,

"ds-length":integer,

"buffermap":binary

}

The semantic of each field in the message is as follows:

* *proto-version* is the version of the protocol;
* *index-version* is the version of the index file;
* *peer-id* is the identifier of the peer;
* *overlay-id* is the identifier of overlay network;
* *sp-index* indicates the starting index number of fragment represented by the buffermap;
* *cp-length* indicates the length of completed section within buffermap beginning *sp-index*;
* *dp-index* indicates the starting index number of downloading section;
* *ds-length* indicates the length of downloading section;
* *buffermap* represents the possession status of fragments within downloading section. If the peer possesses particular fragment, it set to ‘1’, and set to ‘0’, if not.
  + 1. **REFRESH**

This message is used to request a peer to send the latest buffermap. When a client peer receives all fragments of another peer, which may a client peer or a source peer, it requests the peer to send the latest buffermap by sending this message. The syntax of the message is as follows;

{

"method": "REFRESH",

"piece-index"; integer,

"piece-number": integer

}

The semantic of each field in the message is as follows:

* *piece-index* is an index number of fragment that requesting peer is interested. In this Recommendation, this value is always set to ‘1’;
* *piece-number* indicates the number of fragments to be represented by buffermap. If this value is set to ‘0’, the peer sends whole buffermap after *piece-index*. In this Recommendation this value is set to ‘0’ as default, but other value can be used.
  + 1. **BUFFERMAP**

This message is used to deliver a buffermap of a peer. The syntax of the message is as follows;

{

"method": "BUFFERMAP"

"piece-index": integer,

"cp-length": integer,

"dp-index": integer,

"ds-length": integer,

"buffermap": binary

}

The semantic of each field in the message is as follows:

* *piece-index* is an index number requested by REFRESH message. this value is always set to ‘1’;
* *cp-length* indicates the length of completed section within buffermap;
* *dp-index* indicates the starting index number of downloading section;
* *ds-length* indicates the length of downloading section;
* *buffermap* represents the possession status of fragments within downloading section. If the peer possesses particular fragment, it set to ‘1’, and set to ‘0’, if not.
  + 1. **GET**

This message is used for request a particular fragment to a peer. The syntax of the message is as follows;

{

"method":"GET"

"piece-index":integer,

"offset":integer

}

The semantic of each field in the message is as follows:

* *piece-index* is the index number of fragment being requested;
* *offset* is set to 0(zero) if it needs whole data of the fragment. If it has part of the fragment, a requesting peer specify the offset value to this field, and then requested peer sends the fragment from the offset.

When a client peer requests an index file of another peer, the client peer sends a GET message with the following format;

{

"method":"GET"

"piece-index":0   
}

The semantic of each field in the message is as follows:

* *piece-index* is set to 0(‘zero’) for requesting index file.
  + 1. **BUSY**

This message is used to rejecting connection establishment with a requesting peer due to lack of resources of the requested peer. The syntax of the message is as follows;

{

"method":"BUSY"

"reason":string,

"index-file":binary,

}

The semantic of each field in the message is as follows:

* *reason* is set to convey reasons of denial. This can be “shortage of bandwidth, the number of concurrent connection has been exceeding, internal resources such as storage and CPU are busy”);
* *index-file* contains index file of the requested peer if the peer requesting a connection has the lower version of index file.
  + 1. **DATA**

This message is used for delivering fragments. A peer can use the hash value to verify integrity of the fragment from another peer. This message may contain signature that is signed by a private key of a source peer on the hash value to prevent from malicious manipulation. The syntax of the message is as follows;

{

"method":"DATA"

"piece-index":integer,

"offset":integer,

"data-size":integer,

"hash":string,

"signature":string,

"encrypted-hash":string,

"data":binary

}

The semantic of each field in the message is as follows:

* *piece-index* is an index number of fragment;
* *offset* indicates the offset for a particular fragment;
* *data-size* indicates the size of fragment;
* *timestamp* indicates the creation time of fragments with NTP timestamp format;
* *hash* contains SHA-1 hashing value for the fragment;
* *signature* contains the signed hash value by digital signature of source peer;
* *encrypted-hash* contains encrypted hash value by digital signature of source peer;
* *data* contains fragment.

When a peer needs to send its index file, it uses a DATA message with the following format;

{

"method":"DATA"

"piece-index":0,

"data-size":integer,

"data": binary

}

The semantic of each field in the message is as follows:

* *piece-index* is set to 0(‘zero’) for sending index file;
* *data-size* indicates the size of index file;
* *data* contains index file.
  + 1. **CANCEL**

This message is used to revoke a pending request on a particular fragment to a peer. The syntax of the message is as follows;

{

"method": "CANCEL",

"piece-index": integer,

"offset": integer

}

The semantic of each field in the message is as follows:

* *piece-index* is an index number of fragment;
* *offset* indicates the offset for a particular fragment.
  + 1. **BYE**

This message is used for the release a connection between two peers. The syntax of the message is as follows;

{

"method": "BYE"

}

## 7.3 Behaviour of peers

### 7.3.1 Overview

In this Recommendation, a peer is either a source peer or a client peer. A source peer is the origin of the content to be distributed and generates an overlay network for distributing the content. The source peer can also provide the latest index file if there is any change in the content. As a receiver of the content, a client peer participates in an overlay network corresponding to the content and receives the content. The client peer also shares the received fragments of the content with other client peers.

For distributing a content, peers need to conduct protocol behaviours, and the behaviour of a peer is dependent on operational phases and the type of the peer. [ITU-T X.609.7] defines three operational phases; provisioning phase, peering phase, and updating phase. This clause describes the peer behaviour in peering phase and in updating phase.

NOTE – A source peer establishes an overlay network for distributing a content in provisioning phase. However, the provisioning phase is not under scope of this Recommendation.

### 7.3.2 Peering phase

In the peering phase, peers exchange protocol messages with each other to establish a connection for sending and receiving fragments of a content. After establishing a connection between peers, they can exchange the fragments or an index file.

#### 7.3.2.1 Source peer

A source peer in peering phase may exchange a HELLO message with client peers or send a PELLO message to client peers for establishing a connection between the client peers. After a connection is established, the source peer may receive a GET message requesting fragments of its content and then it can send a DATA message including the requested fragments.

##### 7.3.2.1.1 Connection establishment

On receiving a HELLO message from a client peer, a source peer can send a HELLO message including own buffermap composed of CS only. If there is no resource available, the source peer responds with a BUSY message and disconnects the TCP connection.

To establish a connection, a source peer can also send a PELLO message to a client peer instead of waiting a HELLO message.

##### 7.3.2.1.2 Fragment transmission

On receiving a GET message from a client peer, a source peer can send a DATA message including the fragment(s) as requested by the client peer. If the source peer receives a CANCEL message before sending the requested fragment(s), it cancels fragment transmission as requested.

##### 7.3.2.1.3 Connection termination

A source peer can terminate a connection with a client peer by sending a BYE message. The source peer may receive a BYE message from a client peer. In both cases, the source peer disconnects the TCP connection with the client peer.

If the connection with a certain peer is lost without receiving a BYE message, the source peer assumes that a BYE message had been received.

#### 7.3.2.2 Client peer

A client peer in peering phase may exchange a HELLO message with a source peer or other client peers. After a connection is established with another peer, the client peer can send a GET message in order to request fragments which it needs. The client peer may receive a GET message from other client peers and then it can send a DATA message including the requested fragments.

##### 7.3.2.1.1 Connection establishment

After sending a HELLO message to a source peer or to another client peer, a client peer will receive a HELLO message as a response. After exchanging a HELLO message with either a source peer or a client peer, the client peer can request the fragments which the peer in the established connection has.

When a client peer receives a BUSY message, it will disconnect the TCP connection.

On receiving a PELLO message from a source peer, a client peer compares its own buffermap with the buffermap included in the PELLO message in order to find any fragment which the client peer needs. The client peer that received the PELLO message behaves as if it receives a HELLO message from another peer without sending a HELLO message.

##### 7.3.2.1.2 Fragment reception

After establishing a connection, a client peer can request fragment which the peer connected with has. The client peer sends a GET message indicating a specific fragment which it needs. As a response, the client peer may receive a DATA message including the requested fragments.

##### 7.3.2.1.3 Fragment transmission

On receiving a GET message from another client peer, a client peer can send a DATA message including the fragment(s) as requested by the client peer. If the client peer receives a CANCEL message before sending the requested fragment(s), it cancels fragment transmission as requested.

##### 7.3.2.1.4 Connection termination

A client peer can terminate a connection with another peer by sending a BYE message. The client peer may receive a BYE message from another peer. In both cases, the client peer disconnects the TCP connection with the peer.

If the connection with a certain peer is lost without receiving a BYE message, the client peer assumes that a BYE message had been received.

### 7.3.3 Updating phase

Updating phase is happened when a content had been distributed has changes. A source peer needs to distribute a new version of index file and content. A client peer will try to get the new version of content after having the new version of index file.

#### 7.3.3.1 Source peer

The change of a content includes addition or removal of files organizing the content. After the content is changed, a source peer generates a new version of index file and distribute the index file and the new version of content.

##### 7.3.3.1.1 Index file distribution

The source peer with a new version of index file sends the index file to IXS, which is not under scope of this Recommendation. By exchanging HELLO messages, the source peer can also notify the new version of index file of client peers. Upon receiving a GET message requesting the new version of index file, the source peer sends a DATA message including the index file.

##### 7.3.3.1.2 Fragment transmission

Fragment transmission in updating phase is same as the fragment transmission in peering phase.

#### 7.3.3.2 Corresponding peer

By receiving a HELLO message or a PELLO message with a newer version, a client peer can aware of change in the content being distributed. Then the client peer tries to get the new version of index file and to apply the change in the content. If a client peer participates in an overlay network after a source peer sends the latest index file to IXS, the client peer can get the latest index file from IXS.

##### 7.3.3.2.1 New version awareness

A client peer may receive a HELLO message or a PELLO message from other client peers or a source peer, respectively. If the version information in the messages is higher than the version which the client peer knows, the client peer can aware of change in the content.

##### 7.3.3.2.2 Index file reception

In order to get the latest version of index file, a client peer sends a GET message to the peer sent the HELLO message or the PELLO message notifying change in the content. The requested index file will be delivered through a DATA message as a response of the GET message.

##### 7.3.3.2.3 Index file transmission

A client peer with the latest version of index file may receive a GET message requesting the index file. Upon receiving the GET message, the client peer can send a DATA message including the requested index file to the peer.

Annex A  
Consideration on Web-based content distribution

(This annex forms an integral part of this Recommendation)

## A.1 Overview

Nowadays, content is usually consumed in the web platform such as web browser as well as native application. Hence, it is worth to consider the support of Web-based content distribution. For the Web-based content distribution, it is also considered that web environment cannot use legacy transport protocol, such as TCP and UDP. Instead, the modern web application uses WebRTC and WebSocket for communication with other application or a server. This Annex describes how X.609.7 can be used with WebRTC for Web-based content distribution

## A.2 Recommendation ITU-T X.609.7 over WebRTC

Figure A-1 shows the general procedure for Web-based content distribution.

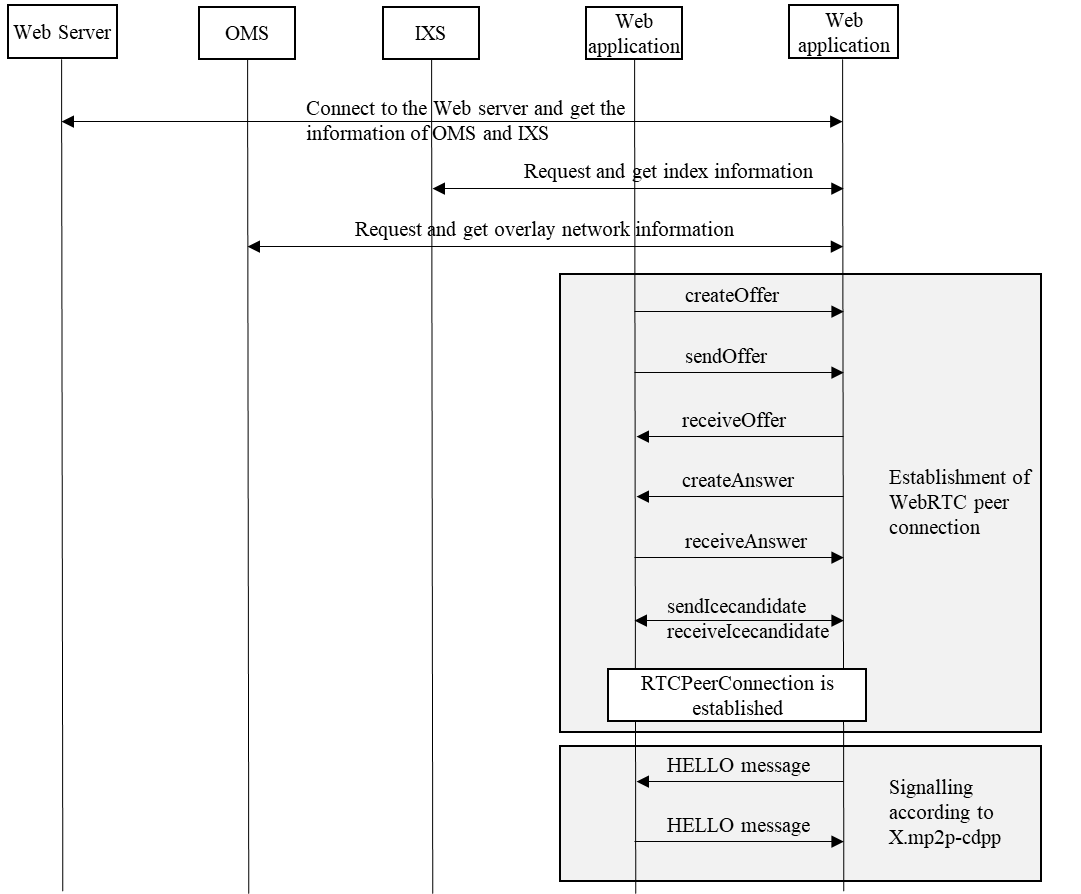


Figure A-1 – General procedure for Web-based content distribution

**Bibliography**

[b-BSON 1.1] *BSON – Binary JSON specification Version 1.1 (2013),* [*http://bsonspec.org*](http://bsonspec.org)

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