Structure and Interpretation of Computer Programs

COMP200

SUMMARY

What have we learned?

- type: a set of values
 - every value has a type
- procedure types include



- number of arguments required
- type of each argument
- type of result of the procedure
- types: a mathematical theory for reasoning efficiently about programs
 - useful for preventing certain common types of errors
 - basis for many analysis and optimization algorithms

```
(* 2 2)
(* 57 57)
(* k k)
```

```
(*22)
(* 57 57)
(* k k)
(lambda (x) (* x x))
     parameter
                 actual
                pattern
```

```
(* 2 2)
(* 57 57)
(* k k)
(lambda (x) (* x x))
(define square (lambda (x) (* x x)))
```

```
(*22)
(*5757)
(* k k)
(lambda (x) (* x x))
(define square (lambda (x) (* x x)))
            number — number
```

$$1 + 2 + \ldots + 100 = (100 * 101)/2$$

$$1 + 4 + 9 + \dots + 100^2 = (100 * 101 * 201)/6$$

$$1 + 1/3^2 + 1/5^2 + ... + 1/101^2 = \pi^2/8$$

$$\sum_{k=1}^{100} k$$

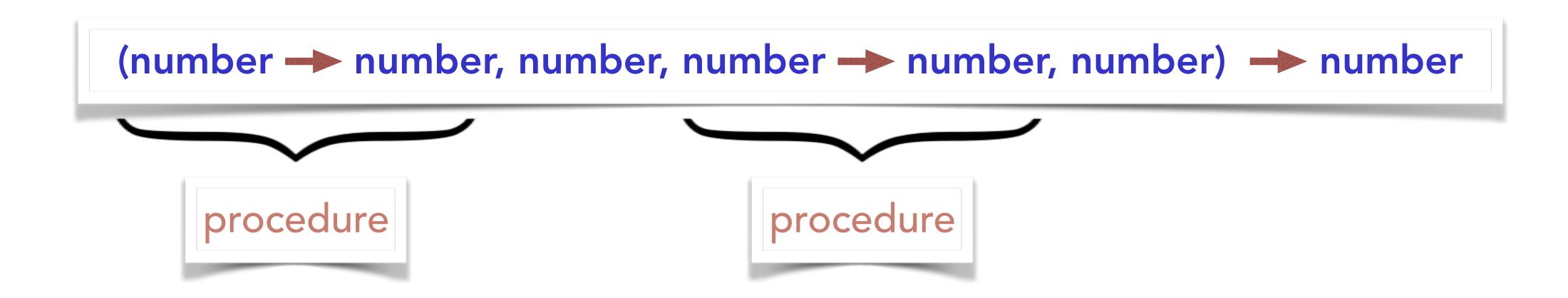
$$\sum_{k=1}^{100} k^{k}$$

$$\sum_{k=1}^{101} k^{-2}$$

```
(define (sum-integers a b)
\sum_{k=1}^{100} k
         (if (> a b)
              (+ a (sum-integers (+ 1 a) b))))
       (define (sum-squares a b)
\sum^{100} k^2
         (if (> a b)
k=1
               (+ (square a)
                  (sum-squares (+ 1 a) b))))
       (define (pi-sum a b)
         (if (> a b)
              (+ (/ 1 (square a))
                  (pi-sum (+ a 2) b))))
```

```
(define (sum-integers a b)
\sum_{k=1}^{100} k
         (if (> a b)
              (+ a (sum-integers (+ 1 a) b))))
       (define (sum-squares a b)
\sum^{100} k^2
         (if (> a b)
k=1
              (+ (square a)
                  (sum-squares (+ 1 a) b))))
       (define (pi-sum a b)
         (if (> a b)
              (+ (/ 1 (square a))
                  (pi-sum (+ a 2) b))))
```

```
(define (sum-integers a b)
100
         (if (> a b)
\sum_{k=1}^{\infty} k
              (+ a (sum-integers (+ 1 a) b))))
       (define (sum-squares a b)
100
         (if (> a b)
k=1
              (+ (square a)
                  (sum-squares (+ 1 a) b))))
       (define (pi-sum a b)
         (if (> a b)
              (+ (/ 1 (square a))
                  (pi-sum (+ a 2) b))))
```



```
(define (sum term a next b)

(if (> a b)

(+ (term a)

(sum term (next a) next b))))

(define (sum term a next b)

(if (> a b)

(sum term (next a) next b))))
```

```
(define (sum-integers1 a b)
  (sum (lambda (x) x) a (lambda (x) (+ x 1)) b))

(define (sum-squares1 a b)
  (sum square a (lambda (x) (+ x 1)) b))

(define (pi-sum1 a b)
  (sum (lambda (x) (/ 1 (square x))) a (lambda (x) (+ x 2)) b))
```

```
(define (sum-integers a b)
(if (> a b))
(+ a (sum-integers (+ 1 a) b))))
(define (sum term a next b)
(if (> a b))
(+ (term a))
(sum term (next a) next b))))
```

```
(define (sum-integers1 a b)
  (sum (lambda (x) x) a (lambda (x) (+ x 1)) b))
```

```
(define (sum-squares a b)

(if (> a b)

(+ (square a)

(sum-squares (+ 1 a) b))))
(define (sum term a next b)

(if (> a b)

(+ (term a)

(sum term (next a) next b))))
```

```
(define (sum-squares1 a b)
  (sum square a (lambda (x) (+ x 1)) b))
```

```
\sum_{k=1 \text{ odd}}^{101} k^{-2} \text{ (if (> a b)} \\ \text{ (+ (/ 1 (square a))} \\ \text{ (pi-sum (+ a 2) b))))}
```

```
(define (pi-sum1 a b) (sum (lambda (x) (/ 1 (square x))) a (lambda (x) (+ x 2)) b))
```

$$f: x \mapsto x^2$$

$$Df: x^2 \mapsto 2x$$

$$f: x \mapsto x^3$$

$$Df: x^3 \mapsto 3x^2$$

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$$f: x \mapsto x^3$$

$$Df: x^3 \mapsto 3x^2$$

$$Df \approx \frac{f(x+\epsilon) - f(x)}{\epsilon}$$

Computing Derivatives

$$Df \approx \frac{f(x+\epsilon) - f(x)}{\epsilon}$$

(define deriv
 (lambda (f)

$$Df \approx \frac{f(x+\epsilon) - f(x)}{\epsilon}$$

Computing Derivatives

$$Df \approx \frac{f(x+\epsilon) - f(x)}{\epsilon}$$

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$$Df \approx \frac{f(x+\epsilon) - f(x)}{\epsilon}$$

```
(number → number) → (number → number)
```

(define epsilon 0.001)

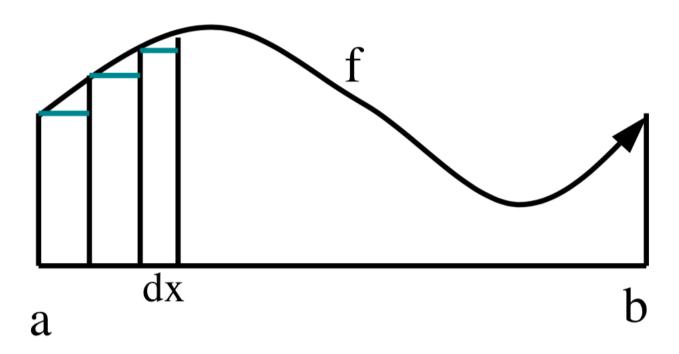
((deriv square) 5)

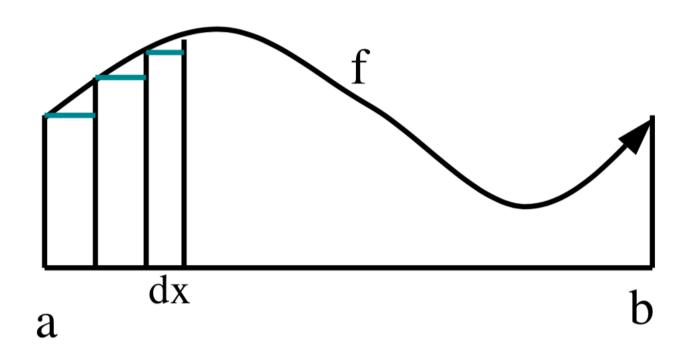
```
(define deriv
                        HIGHER ORDER PROCEDURES
 (lambda (f)
   (lambda (x)
                                   Using Derivatives
     (/ (- (f (+ x epsilon)) (f x))
       epsilon))))
                                  ((lambda (x)
                                     (/ (- ((lambda (y) (* y y))(+ x epsilon))
      (define square
                                            ((lambda (y) (* y y)) x)))
        (lambda (y) (* y y)))
                                     epsilon) 5)
      (define epsilon 0.001)
                                  (lambda (x)
                                     (/ (- ((lambda (y) (* y y))(+ 5 epsilon))
      ((deriv square) 5)
```

epsilon)

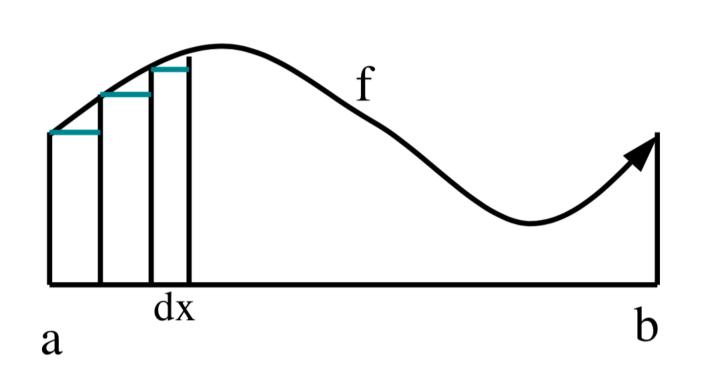
((lambda (y) (* y y)) 5)))

Mhhs





$$dx (f(a) + f(a + dx) + f(a + 2dx) + \cdots + f(b))$$



$$dx (f(a) + f(a + dx) + f(a + 2dx) + \cdots + f(b))$$

```
(define my-atan
  (lambda (a)
        (integral (lambda (x) (/ 1 (+ 1 (square x)))) 0 a 1000)))
(my-atan 2.0)
(my-atan 3.0)
```

$$\sqrt{x} = x/\sqrt{x}$$

$$f: y \mapsto x/y$$

Fixed Points

$$\sqrt{x} = x/\sqrt{x}$$

$$f: y \mapsto x/y$$

if we can find a
$$y = \sqrt{x}$$

then
$$f(y) = y$$

such a y is called a fixed point of f

- Given a guess x_1 , let new guess by $f(x_1)$
- Keep computing f of last guess, until close enough

- Given a guess x_1 , let new guess by f(x)
- Keep computing f of last guess, until close enough

```
(fixed-point (lambda (x) (+ 1 (/ 1 x))) 1)
1.6180
```

$$x = 1 + (1/x)$$
 when $x = (1 + \sqrt{5})/2 = 1.6180$

```
(define (my-sqrt x)
  (fixed-point
    (lambda (y) (/ x y)) 1))
```

Fixed Points

```
(define (my-sqrt x)
  (fixed-point
    (lambda (y) (/ x y)) 1))
```

What happens if we try: (my-sqrt 2)

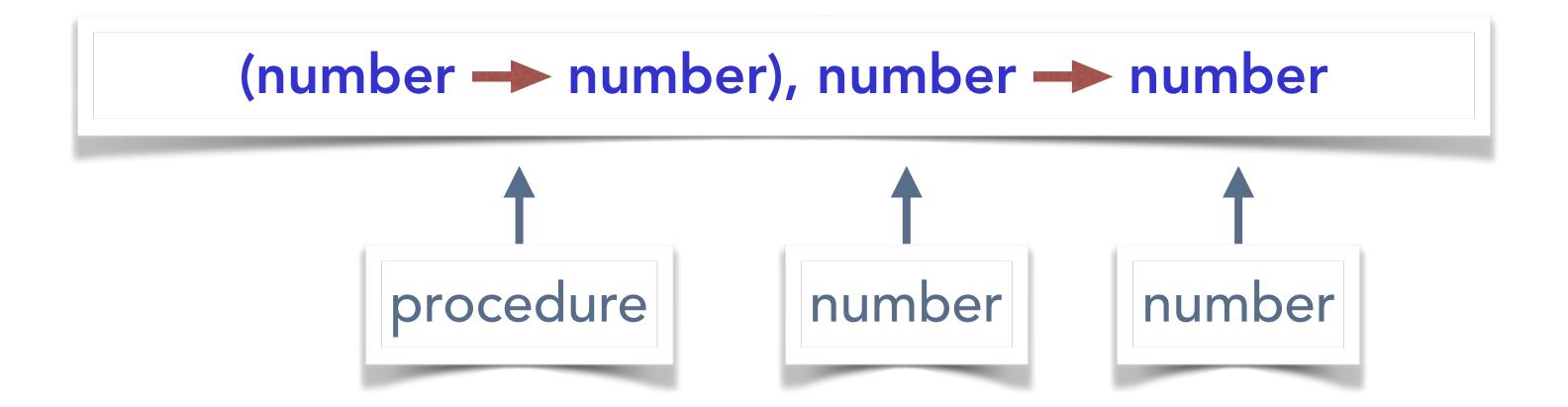
```
(define (average-damp f)
  (lambda (x)
          (average x (f x))))
```

```
(define (average-damp f)
  (lambda (x)
          (average x (f x))))
```

```
(number → number) → (number → number)
```

```
(define (my-sqrt x)
  (fixed-point
    (average-damp
       (lambda (y) (/ x y)))
    1))
```

```
(define hop1
  (lambda (f x)
    (+ 2 (f (+ x 1))))
(hop1 square 3)
(+ 2 (square (+ 3 1)))
(+ 2 (square 4))
(+ 2 (* 4 4))
(+216)
18
(hop1 (lambda (x) (* x x)) 3)
```



```
(define compose
  (lambda (f g x)
    (f (g x))))
(compose square double 3)
(square (double 3))
(square (* 3 2))
(square 6)
(*66)
```

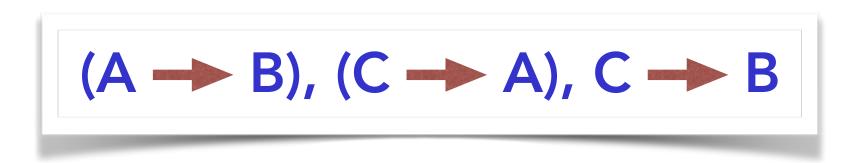
Example

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Example

(number → number), (number → number), number → number

```
(define compose
  (lambda (f g x)
    (f (g x))))
(compose
 (lambda (p)
                                   boolean -> string
   (if p "hi" "bye"))
 (lambda (x)
                                  number - boolean
  (> \times 0)
                                       number
-5)
                                       string
                            result:
```



- Meaning of type variables
 All places where a given type variable appears must match when you fill in the actual operand types.
- The constraints are:
 - F and G must be functions of one argument
 - the argument type of G matches the type of X
 - the argument type of F matches the result type of G
 - the result type of compose is the result type of F