

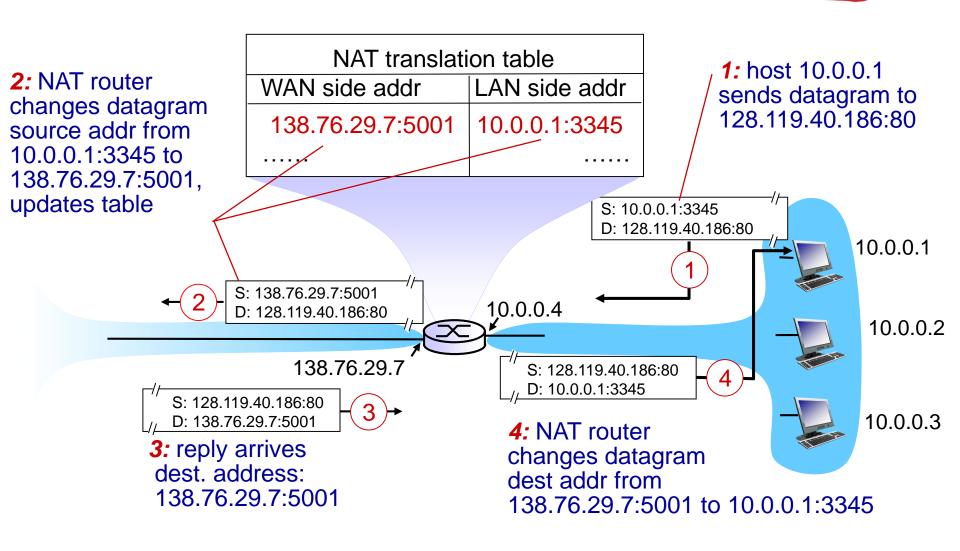
all datagrams leaving local network have same single source NAT IP address: 138.76.29.7, different source port numbers datagrams with source or destination in this network have 10.0.0/24 address for source, destination (as usual)

motivation: local network uses just one IP address as far as outside world is concerned:

- range of addresses not needed from ISP: just one IP address for all devices
- can change addresses of devices in local network without notifying outside world
- can change ISP without changing addresses of devices in local network
- devices inside local net not explicitly addressable, visible by outside world (a security plus)

implementation: NAT router must:

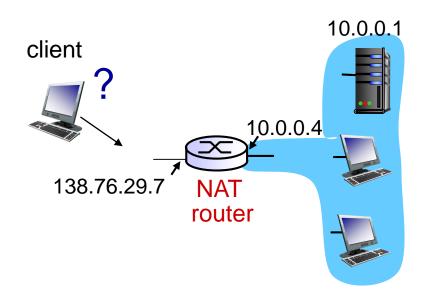
- outgoing datagrams: replace (source IP address, port #) of every outgoing datagram to (NAT IP address, new port #)
 - ... remote clients/servers will respond using (NAT IP address, new port #) as destination addr
- remember (in NAT translation table) every (source IP address, port #) to (NAT IP address, new port #) translation pair
- incoming datagrams: replace (NAT IP address, new port #) in dest fields of every incoming datagram with corresponding (source IP address, port #) stored in NAT table



- 16-bit port-number field:
 - 60,000 simultaneous connections with a single LAN-side address!
- NAT is controversial:
 - routers should only process up to layer 3 (not layer 4, network, where TCP ports are defined)
 - violates end-to-end argument
 - NAT possibility must be taken into account by app designers, e.g., P2P applications
 - address shortage should instead be solved by IPv6

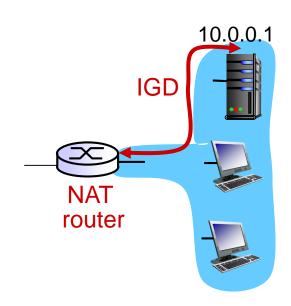
NAT traversal problem

- client wants to connect to server with address 10.0.0.1
 - server address I0.0.0.I local to LAN (client can't use it as destination addr)
 - only one externally visible NATed address: 138.76.29.7
- solution I: statically configure NAT to forward incoming connection requests at given port to server at given port
 - e.g., (123.76.29.7, port 2500)
 always forwarded to 10.0.0.1 port 25000
 - Use different ports for different services



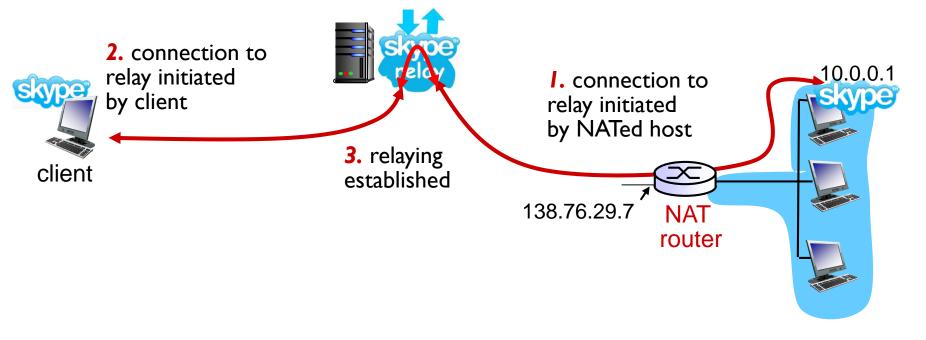
NAT traversal problem

- solution 2: Universal Plug and Play (UPnP) Internet Gateway Device (IGD) Protocol. Allows NATed host to:
 - learn public IP address (138.76.29.7)
 - add/remove port mappings (with lease times)
 - i.e., automate static NAT port map configuration
- Holy buckets is that insecure!
 - Not used on commercial routers



NAT traversal problem

- solution 3: relaying (used in Skype)
 - NATed client establishes connection to relay
 - external client connects to relay
 - relay bridges packets between to connections



Chapter 4: outline

- 4.1 introduction
- 4.2 virtual circuit and datagram networks
- 4.3 what's inside a router
- 4.4 IP: Internet Protocol
 - datagram format
 - IPv4 addressing
 - ICMP
 - IPv6

- 4.5 routing algorithms
 - link state
 - distance vector
 - hierarchical routing
- 4.6 routing in the Internet
 - RIP
 - OSPF
 - BGP
- 4.7 broadcast and multicast routing

ICMP: internet control message protocol

*	used by hosts & routers to communicate network-	Type	Code	description
		0	0	echo reply (ping)
	level information	3	0	dest. network unreachable
	error reporting:	3	1	dest host unreachable
	unreachable host, network,	3	2	dest protocol unreachable
	port, protocol	3	3	dest port unreachable
	echo request/reply (used by	3	6	dest network unknown
	ping)	3	7	dest host unknown
*	network-layer "above" IP:	4	0	source quench (congestion
	ICMP msgs carried in IP			control - not used)
	datagrams	8	0	echo request (ping)
	•	9	0	route advertisement
**	ICMP message: type, code	10	0	router discovery
	plus first 8 bytes of IP	11	0	TTL expired
	datagram causing error	12	0	bad IP header

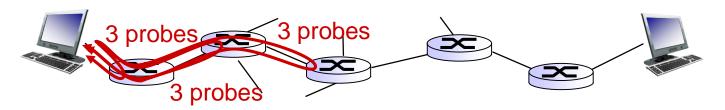
Traceroute and ICMP

- source sends series of UDP segments to dest
 - first set has TTL = I
 - second set has TTL=2, etc.
 - unlikely port number
- when nth set of datagrams arrives to nth router:
 - router discards datagrams
 - and sends source ICMP messages (type II, code 0)
 - ICMP messages includes name of router & IP address

 when ICMP messages arrives, source records RTTs

stopping criteria:

- UDP segment eventually arrives at destination host
- destination returns ICMP "port unreachable" message (type 3, code 3)
- source stops



IPv6: motivation

- initial motivation: 32-bit address space is completely allocated.
- additional motivation:
 - header format helps speed processing/forwarding
 - header changes to facilitate QoS

IPv6: motivation

- IPv6 uses 128-bit addresses
 - $2^{128} > 2^{32}$
 - \bullet 3.4 * $10^{38} > 4.2 * 10^9$
 - This is 50 octillion addresses for each of 6.5 billion people on earth

50 octillion addys*:

50,000,000,000,000,000,000,000,000

IPv6: datagram format

IPv6 datagram format:

- fixed-length 40 byte header
- no fragmentation allowed

IPv6 datagram format

priority: identify priority among datagrams in flow flow Label: identify datagrams in same "flow." (concept of flow not well defined).

next header: identify upper layer protocol for data (i.e.TCP)

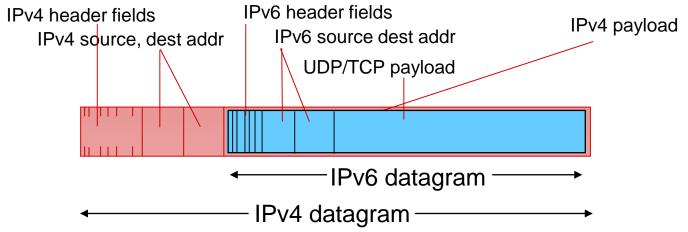
ver	pri	flow label					
Ķ	payload	llen	next hdr	hop limit			
source address (128 bits)							
destination address (128 bits)							
data							
◆ 32 bits							

Other changes from IPv4

- checksum: removed entirely to reduce processing time at each hop
- options: allowed, but outside of header, indicated by "Next Header" field
- * ICMPv6: new version of ICMP
 - additional message types, e.g. "Packet Too Big"
 - multicast group management functions

Transition from IPv4 to IPv6

- not all routers can be upgraded simultaneously
 - no "flag days"
 - how will network operate with mixed IPv4 and IPv6 routers?
- tunneling: IPv6 datagram carried as payload in IPv4 datagram among IPv4 routers



Tunneling

IPv4 tunnel connecting IPv6 routers logical view: IPv6 IPv6 IPv6 IPv6 Ε Α В physical view: IPv6 IPv6 IPv6 IPv4 IPv4 IPv6

Tunneling

IPv4 tunnel В Ε connecting IPv6 routers logical view: IPv6 IPv6 IPv6 IPv6 Α В Ε F physical view: IPv6 IPv6 IPv6 IPv6 IPv4 IPv4 src:B flow: X src:B src:B flow: X src: A src: A dest: E dest: E dest: E dest: F dest: F Flow: X Flow: X Flow: X Src: A Src: A Src: A Dest: F Dest: F data Dest: F data data data data A-to-B: E-to-F: B-to-C: C-to-D: D-to-F: IPv6 IPv6 IPv6 inside IPv6 insideIPv6 inside IPv4 IPv4 IPv4

Network Layer 4-19

IPv6: adoption

- US National Institutes of Standards estimate [2013]:
 - ~3% of industry IP routers
 - ~II% of US gov't routers
- Long (long!) time for deployment, use
 - 20 years and counting!
 - think of application-level changes in last 20 years: WWW, Facebook, ...
 - Why?

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