

Cryptography: HW1

Due electronically on Friday 15 Jan

1. Alice is using one-time pad and notices that when her key is the all-zeroes string $k = 0^\lambda$, then $\text{Enc}(k, m) = m$ and her message is sent in the clear! To avoid this problem, she decides to modify KeyGen to choose a key uniformly from $\{0, 1\}^\lambda \setminus \{0^\lambda\}$. In this way, her plaintext is never sent in the clear.

Does the modified scheme still have one-time secrecy? Justify your answer.

2. Show that the following encryption scheme does **not** have one-time secrecy, by constructing a program that distinguishes the two relevant libraries from the one-time secrecy definition.

$\mathcal{K} = \mathbb{Z}_{10}$	<u>KeyGen:</u>	<u>Enc(k, m):</u>
$\mathcal{M} = \mathbb{Z}_{10}$	$k \leftarrow \mathbb{Z}_{10}$	return $k \times m \% 10$
$\mathcal{C} = \mathbb{Z}_{10}$	return k	

3. Consider the following encryption scheme:

$\mathcal{K} = \mathbb{Z}_n$	<u>KeyGen:</u>	<u>Enc(k, m):</u>	<u>Dec(k, c):</u>
$\mathcal{M} = \mathbb{Z}_n$	$k \leftarrow \mathbb{Z}_n$	return $(k + m) \% n$??
$\mathcal{C} = \mathbb{Z}_n$	return k		

- (a) Fill in the details of the Dec algorithm so that the scheme satisfies correctness.
 - (b) Prove that the scheme satisfies one-time secrecy.
4. Suppose there are 9 people on an important committee: Alice, Bob, Carol, David, Eve, Frank, Gina, Harold, & Irene. Alice, Bob & Carol form a subcommittee; David, Eve & Frank form another subcommittee; and Gina, Harold & Irene form another subcommittee. Suggest how a dealer can share a secret so that it can only be opened when a majority of each subcommittee is present. Describe why a 6-out-of-9 threshold secret-sharing scheme does **not** suffice.

You do not have to give a formal proof, but please give an informal argument about why every unauthorized set of users has no information about the secret.