

Chapter 4: outline

4.1 introduction

4.2 virtual circuit and datagram networks

4.3 what's inside a router

4.4 IP: Internet Protocol

- datagram format
- IPv4 addressing
- ICMP
- IPv6
- NAT

4.5 routing algorithms

- link state
- distance vector
- hierarchical routing

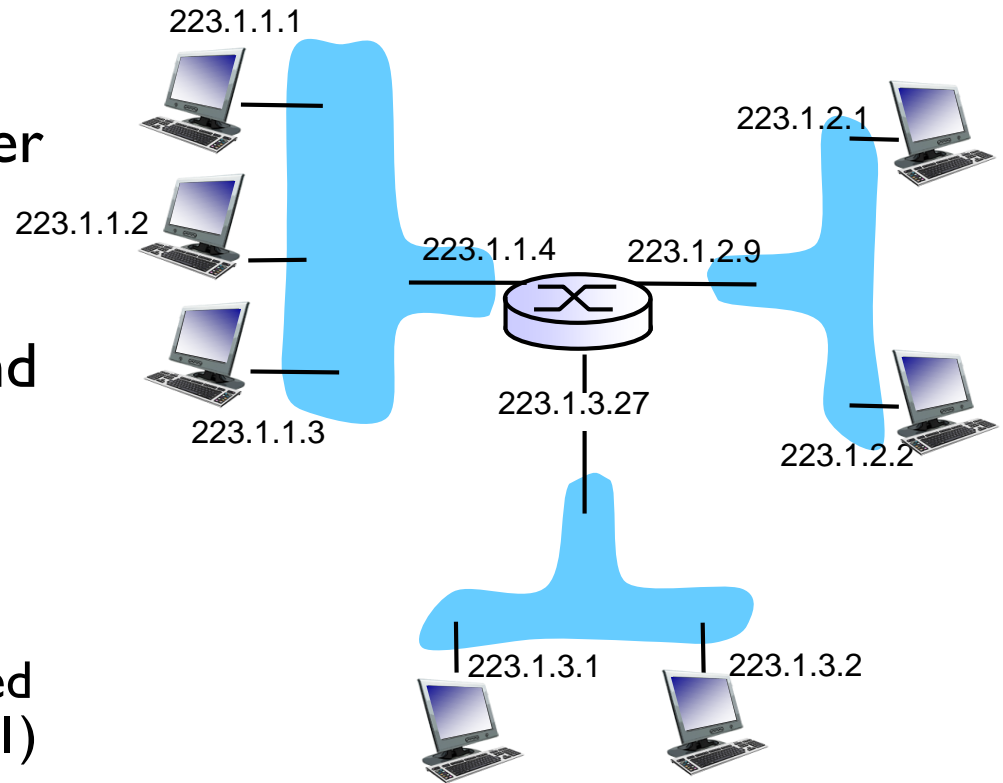
4.6 routing in the Internet

- RIP
- OSPF
- BGP

4.7 broadcast and multicast routing

IP addressing: introduction

- ❖ **IP address:** 32-bit identifier for host, router interface
- ❖ **interface:** connection between host/router and physical link
 - router's typically have multiple interfaces
 - host typically has one or two interfaces (e.g., wired Ethernet, wireless 802.11)
- ❖ **IP addresses associated with each interface**



$$223.1.1.1 = \underbrace{11011111}_{223} \underbrace{00000001}_1 \underbrace{00000001}_1 \underbrace{00000001}_1$$

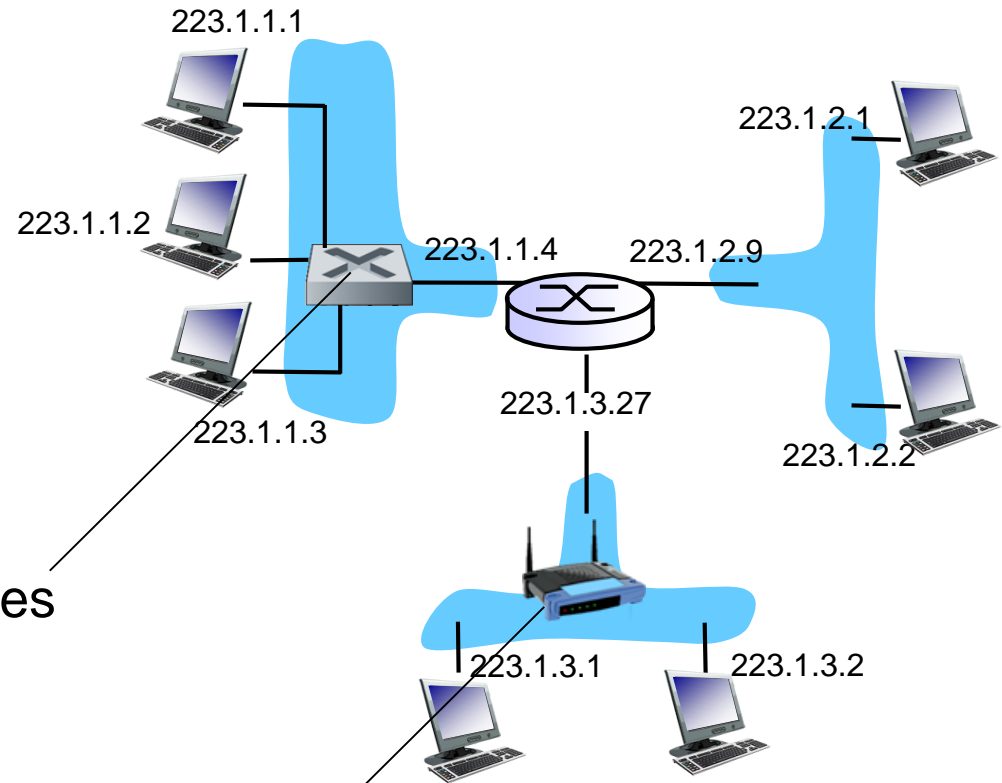
IP addressing: introduction

Q: how are interfaces actually connected?

A: we'll learn about that in chapter 5, 6.

A: wired Ethernet interfaces connected by Ethernet switches

For now: don't need to worry about how one interface is connected to another (with no intervening router)



A: wireless WiFi interfaces connected by WiFi base station

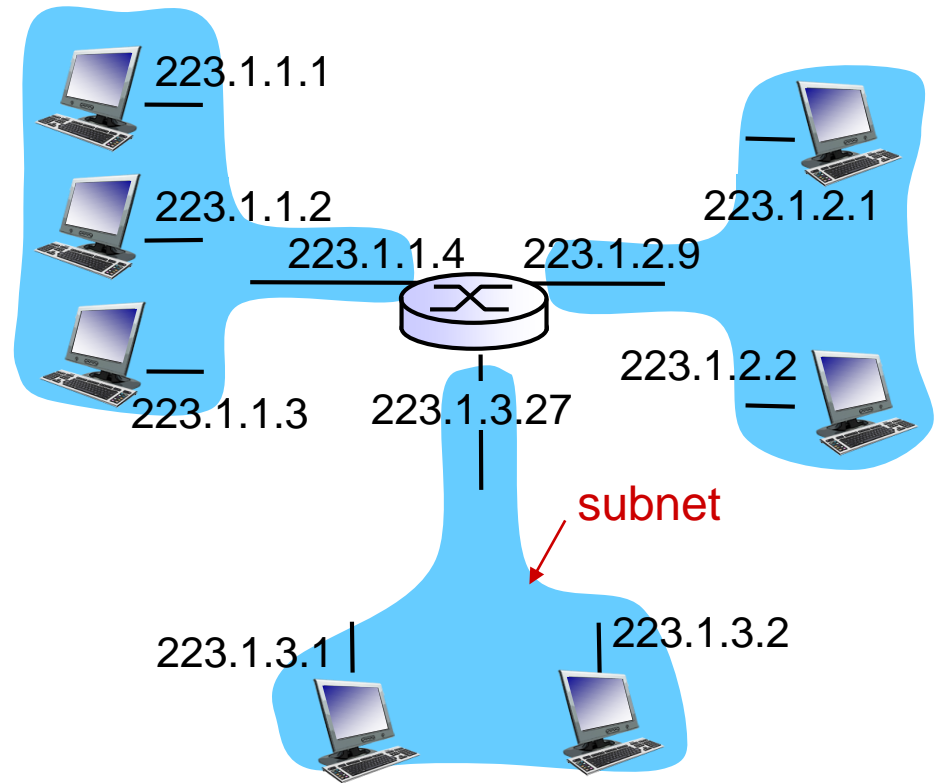
Subnets

❖ IP address:

- subnet part - high order bits
- host part - low order bits

❖ *what's a subnet ?*

- device interfaces with same subnet part of IP address
- can physically reach each other *without intervening router*

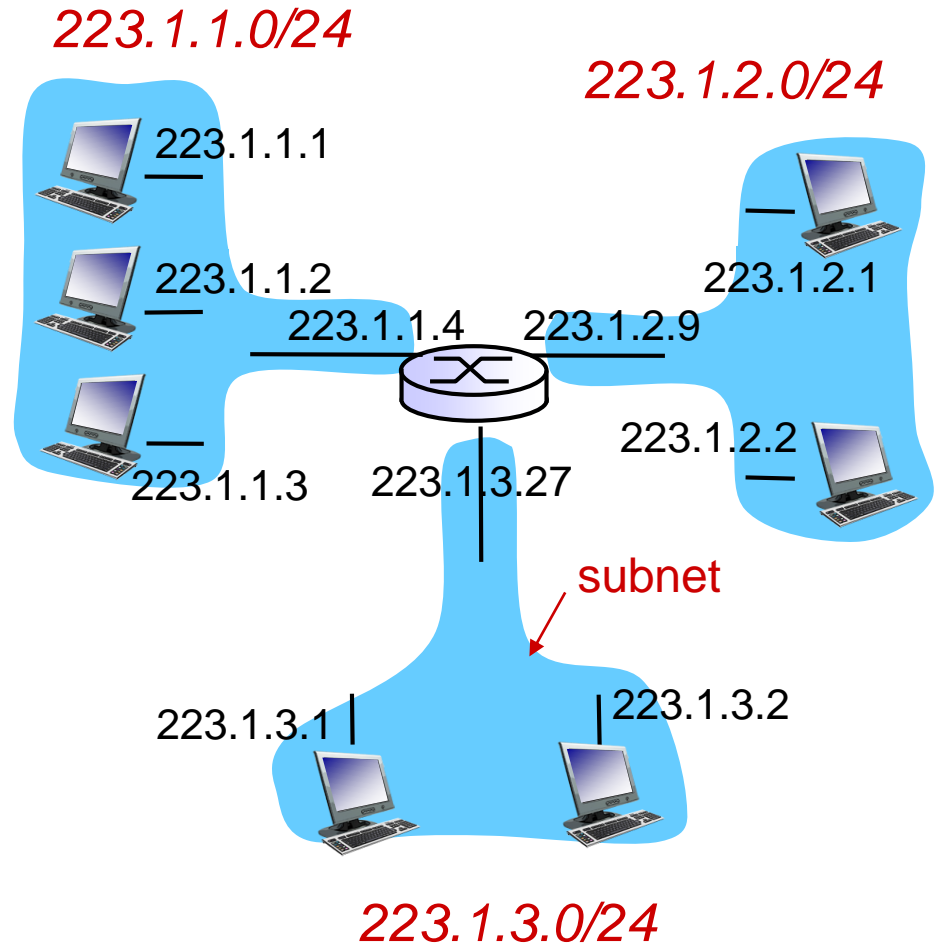


network consisting of 3 subnets

Subnets

recipe

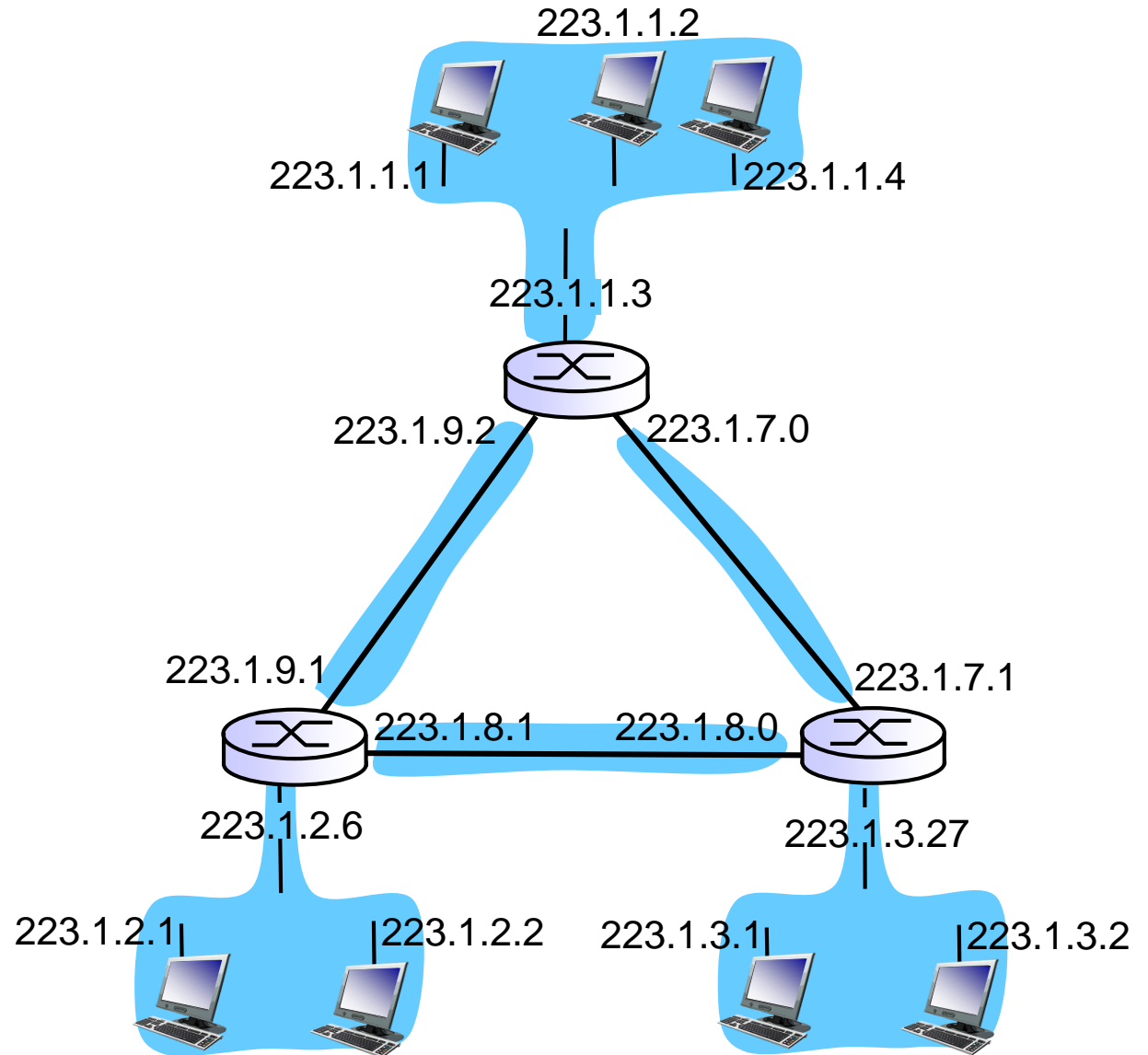
- ❖ to determine the subnets, detach each interface from its host or router, creating islands of isolated networks
- ❖ each isolated network is called a *subnet*



subnet mask: /24

Subnets

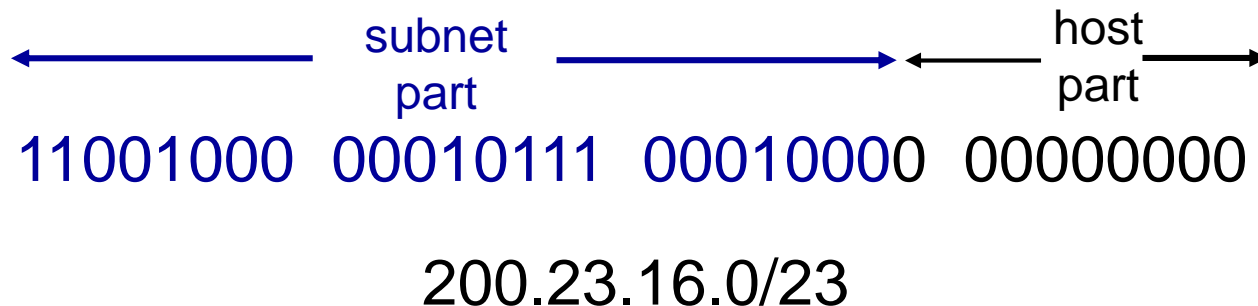
how many?



IP addressing: CIDR

CIDR: Classless InterDomain Routing

- Replaces older “class” based system (A, B, C, D)
- subnet portion of address of *arbitrary length*
- address format: **a.b.c.d/x**, where x is # bits in subnet portion of address



CIDR Subnets and masks

192.168.1.0/24

- ❖ For this particular address, the first 24 bits are the network prefix used for routing:
 - 192.168.1 =
11000000 10101000 00000001
- ❖ We write this prefix like this:
 - 192.168.1.0 =
11000000 10101000 00000001 00000000
- ❖ The remaining 8 bits are used for *host addressing*:
 - 192.168.1.1 - 192.168.1.254
 - .0 and .255 are special reserved addresses.

CIDR Subnets and masks

- ❖ The subnet mask is a binary number that, when applied by a bitwise AND operation to *any* IP address in the network, yields the routing prefix.
- ❖ 192.168.1.0/24 Network prefix has the corresponding subnet mask of 255.255.255.0.
- ❖ Example:

- 192 . 168 . 1 . 0
11000000 10101000 00000001 00001010

- 255 . 255 . 255 . 0
11111111 11111111 11111111 00000000

- AND
11000000 10101000 00000001 00001010
11111111 11111111 11111111 00000000

11000000 10101000 00000001 00000000 = 192.168.1.0

Subnets and masks

255.255.255.255	11111111.11111111.11111111.11111111	/32	Host (single addr)
255.255.255.254	11111111.11111111.11111111.11111110	/31	Unuseable
255.255.255.252	11111111.11111111.11111111.11111100	/30	2 useable
255.255.255.248	11111111.11111111.11111111.11111000	/29	6 useable
255.255.255.240	11111111.11111111.11111111.11110000	/28	14 useable
255.255.255.224	11111111.11111111.11111111.11100000	/27	30 useable
255.255.255.192	11111111.11111111.11111111.11000000	/26	62 useable
255.255.255.128	11111111.11111111.11111111.10000000	/25	126 useable
255.255.255.0	11111111.11111111.11111111.00000000	/24	"Class C" 254 useable
255.255.254.0	11111111.11111111.11111110.00000000	/23	2 Class C's
255.255.252.0	11111111.11111111.11111100.00000000	/22	4 Class C's
255.255.248.0	11111111.11111111.11111000.00000000	/21	8 Class C's
255.255.240.0	11111111.11111111.11110000.00000000	/20	16 Class C's
255.255.224.0	11111111.11111111.11100000.00000000	/19	32 Class C's
255.255.192.0	11111111.11111111.11000000.00000000	/18	64 Class C's
255.255.128.0	11111111.11111111.10000000.00000000	/17	128 Class C's
255.255.0.0	11111111.11111111.00000000.00000000	/16	"Class B"
255.254.0.0	11111111.11111110.00000000.00000000	/15	2 Class B's
255.252.0.0	11111111.11111100.00000000.00000000	/14	4 Class B's
255.248.0.0	11111111.11111000.00000000.00000000	/13	8 Class B's
255.240.0.0	11111111.11110000.00000000.00000000	/12	16 Class B's
255.224.0.0	11111111.11100000.00000000.00000000	/11	32 Class B's
255.192.0.0	11111111.11000000.00000000.00000000	/10	64 Class B's
255.128.0.0	11111111.10000000.00000000.00000000	/9	128 Class B's
255.0.0.0	11111111.00000000.00000000.00000000	/8	"Class A"
254.0.0.0	11111110.00000000.00000000.00000000	/7	
252.0.0.0	11111100.00000000.00000000.00000000	/6	
248.0.0.0	11111000.00000000.00000000.00000000	/5	
240.0.0.0	11110000.00000000.00000000.00000000	/4	
224.0.0.0	11100000.00000000.00000000.00000000	/3	
192.0.0.0	11000000.00000000.00000000.00000000	/2	
128.0.0.0	10000000.00000000.00000000.00000000	/1	
0.0.0.0	00000000.00000000.00000000.00000000	/0	IP space

Network Layer

IP addresses: how to get one?

Q: How does a *host* get IP address?

- ❖ Statically assigned - hard-coded by system admin:
 - Windows: control-panel->network->configuration->tcp/ip->properties
 - UNIX: /etc/rc.config
- ❖ Dynamically assigned:
 - **DHCP**: **D**ynamic **H**ost **C**onfiguration **P**rotocol:
dynamically get address from a server
 - "plug-and-play"

DHCP: Dynamic Host Configuration Protocol

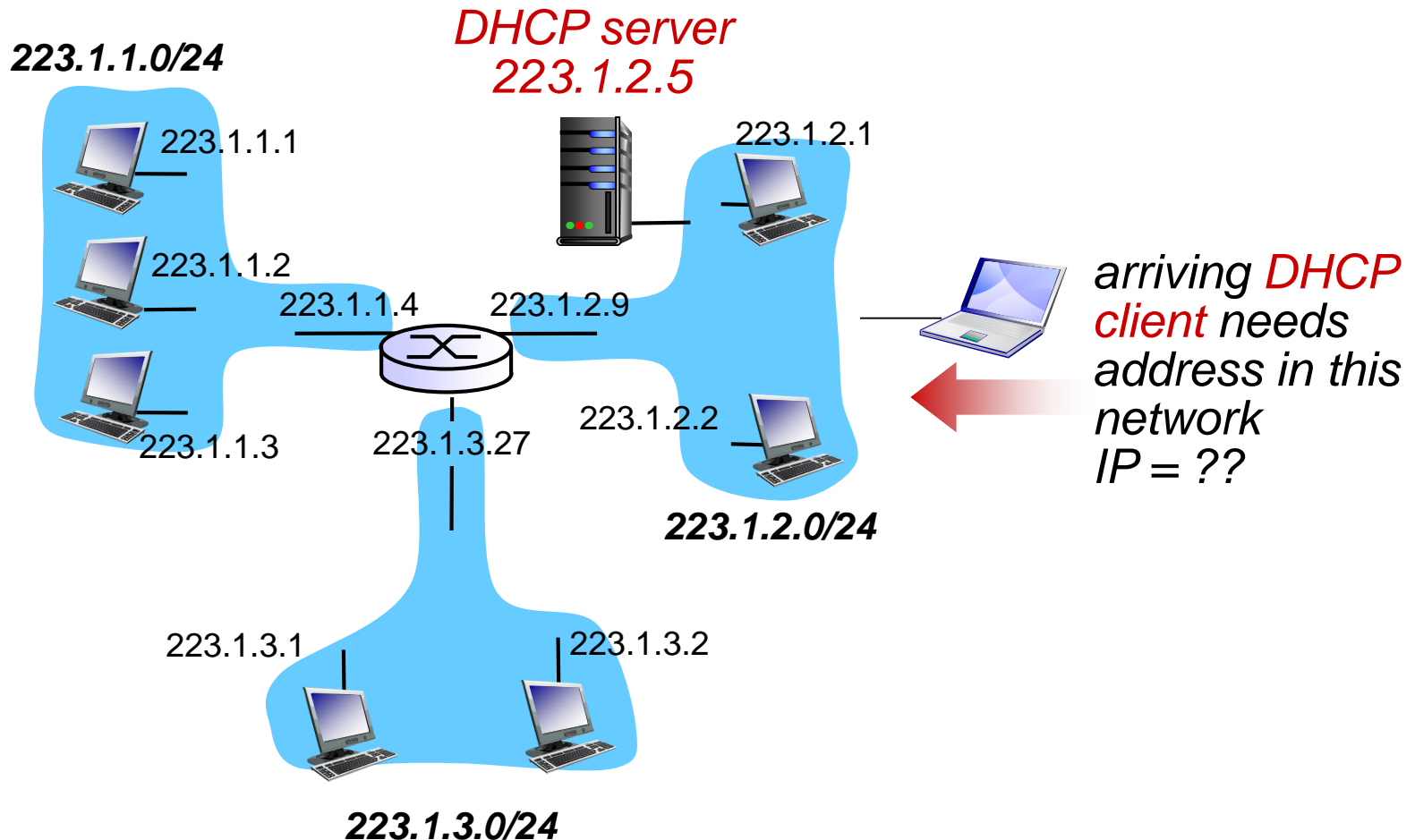
goal: allow host to *dynamically* obtain its IP address from network server when it joins network, and keep it for a time (a “lease”)

- can renew its lease on address in use
- allows reuse of addresses (only hold address while connected or “on”)
- support for mobile users who want to join network (more shortly)

DHCP overview:

- host broadcasts “DHCP discover” msg [optional]
- DHCP server responds with “DHCP offer” msg [optional]
- host requests IP address: “DHCP request” msg
- DHCP server sends address: “DHCP ack” msg

DHCP client-server scenario



DHCP client-server scenario

0.0.0.0 =
"I don't have an address"

255.255.255.255 =
"Send to everybody"

DHCP server: 223.1.2.5

DHCP discover

src : 0.0.0.0:68
dest.: 255.255.255.255:67
yiaddr: 0.0.0.0
transaction ID: 654

arriving
client

Broadcast: is there a
DHCP server out there?

DHCP offer

src: 223.1.2.5:67
dest: 255.255.255.255:68
yiaddr: 223.1.2.4
transaction ID: 654
lifetime: 3600 secs

Broadcast: I'm a DHCP
server! Here's an IP
address you can use

DHCP request

src: 0.0.0.0:68
dest: 255.255.255.255, 67
yiaddr: 223.1.2.4
transaction ID: 655
lifetime: 3600 secs

Broadcast: OK. I'll take
that IP address!

DHCP ACK

src: 223.1.2.5: 67
dest: 255.255.255.255:68
yiaddr: 223.1.2.4
transaction ID: 655
lifetime: 3600 secs

Broadcast: OK. You've
got that IP address!

DHCP: more than IP addresses

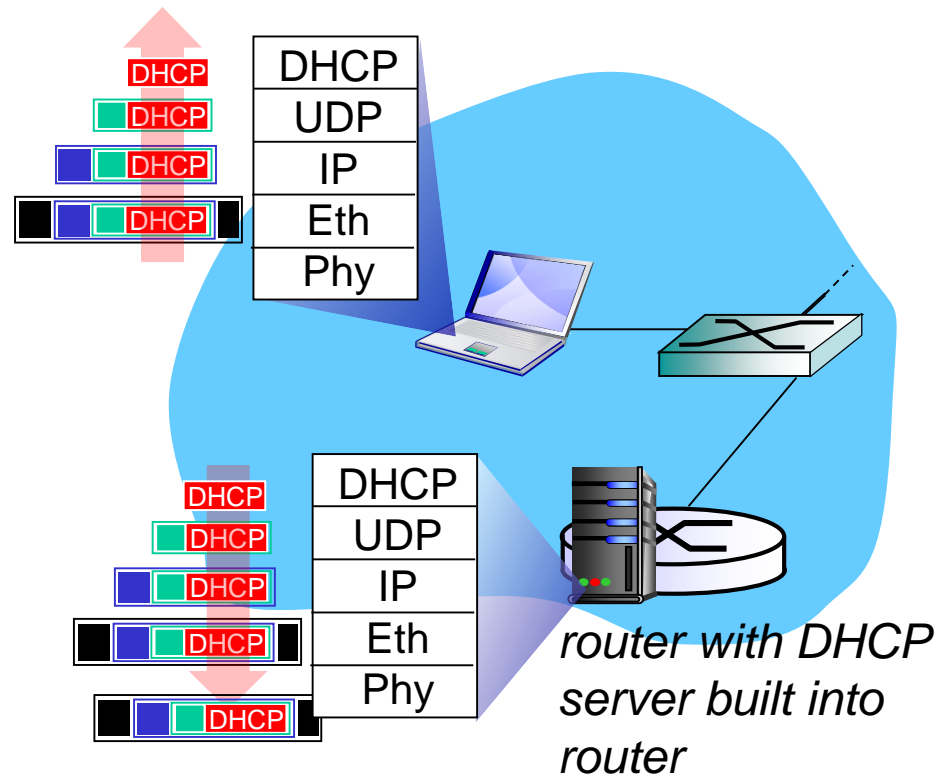
DHCP can return more than just allocated IPv4 address on subnet:

- address of first-hop router for client (default gateway)
- name and IP address of DNS server
- network mask (indicating network versus host portion of address)
- Windows Domain name
- IPv6 addressing information



- ❖ connecting laptop needs its IP address, addr of first-hop router, addr of DNS server: use DHCP
- ❖ DHCP request encapsulated in UDP, encapsulated in IP, encapsulated in 802.1 Ethernet
- ❖ Ethernet frame broadcast (dest: FFFFFFFFFF) on LAN, received at router running DHCP server
- ❖ Ethernet demuxed to IP demuxed, UDP demuxed to DHCP

DHCP: example



- ❖ DCP server formulates DHCP ACK containing client's IP address, IP address of first-hop router for client, name & IP address of DNS server
- ❖ encapsulation of DHCP server, frame forwarded to client, demuxing up to DHCP at client
- ❖ client now knows its IP address, name and IP address of DSN server, IP address of its first-hop router

DHCP: Wireshark output (home LAN)

Message type: **Boot Request (1)**

Hardware type: Ethernet

Hardware address length: 6

Hops: 0

Transaction ID: 0x6b3a11b7

Seconds elapsed: 0

Bootp flags: 0x0000 (Unicast)

Client IP address: 0.0.0.0 (0.0.0.0)

Your (client) IP address: 0.0.0.0 (0.0.0.0)

Next server IP address: 0.0.0.0 (0.0.0.0)

Relay agent IP address: 0.0.0.0 (0.0.0.0)

Client MAC address: Wistron_23:68:8a (00:16:d3:23:68:8a)

Server host name not given

Boot file name not given

Magic cookie: (OK)

Option: (t=53,l=1) **DHCP Message Type = DHCP Request**

Option: (61) Client identifier

Length: 7; Value: 010016D323688A;

Hardware type: Ethernet

Client MAC address: Wistron_23:68:8a (00:16:d3:23:68:8a)

Option: (t=50,l=4) Requested IP Address = 192.168.1.101

Option: (t=12,l=5) Host Name = "nomad"

Option: (55) Parameter Request List

Length: 11; Value: 010F03062C2E2F1F21F92B

1 = Subnet Mask; 15 = Domain Name

3 = Router; 6 = Domain Name Server

44 = NetBIOS over TCP/IP Name Server

.....

request

Message type: **Boot Reply (2)**

Hardware type: Ethernet

Hardware address length: 6

Hops: 0

Transaction ID: 0x6b3a11b7

Seconds elapsed: 0

Bootp flags: 0x0000 (Unicast)

Client IP address: 192.168.1.101 (192.168.1.101)

Your (client) IP address: 0.0.0.0 (0.0.0.0)

Next server IP address: 192.168.1.1 (192.168.1.1)

Relay agent IP address: 0.0.0.0 (0.0.0.0)

Client MAC address: Wistron_23:68:8a (00:16:d3:23:68:8a)

Server host name not given

Boot file name not given

Magic cookie: (OK)

Option: (t=53,l=1) DHCP Message Type = DHCP ACK

Option: (t=54,l=4) Server Identifier = 192.168.1.1

Option: (t=1,l=4) Subnet Mask = 255.255.255.0

Option: (t=3,l=4) Router = 192.168.1.1

Option: (6) Domain Name Server

Length: 12; Value: 445747E2445749F244574092;

IP Address: 68.87.71.226;

IP Address: 68.87.73.242;

IP Address: 68.87.64.146

Option: (t=15,l=20) Domain Name = "hsd1.ma.comcast.net."

reply

IP addresses: how to get one?

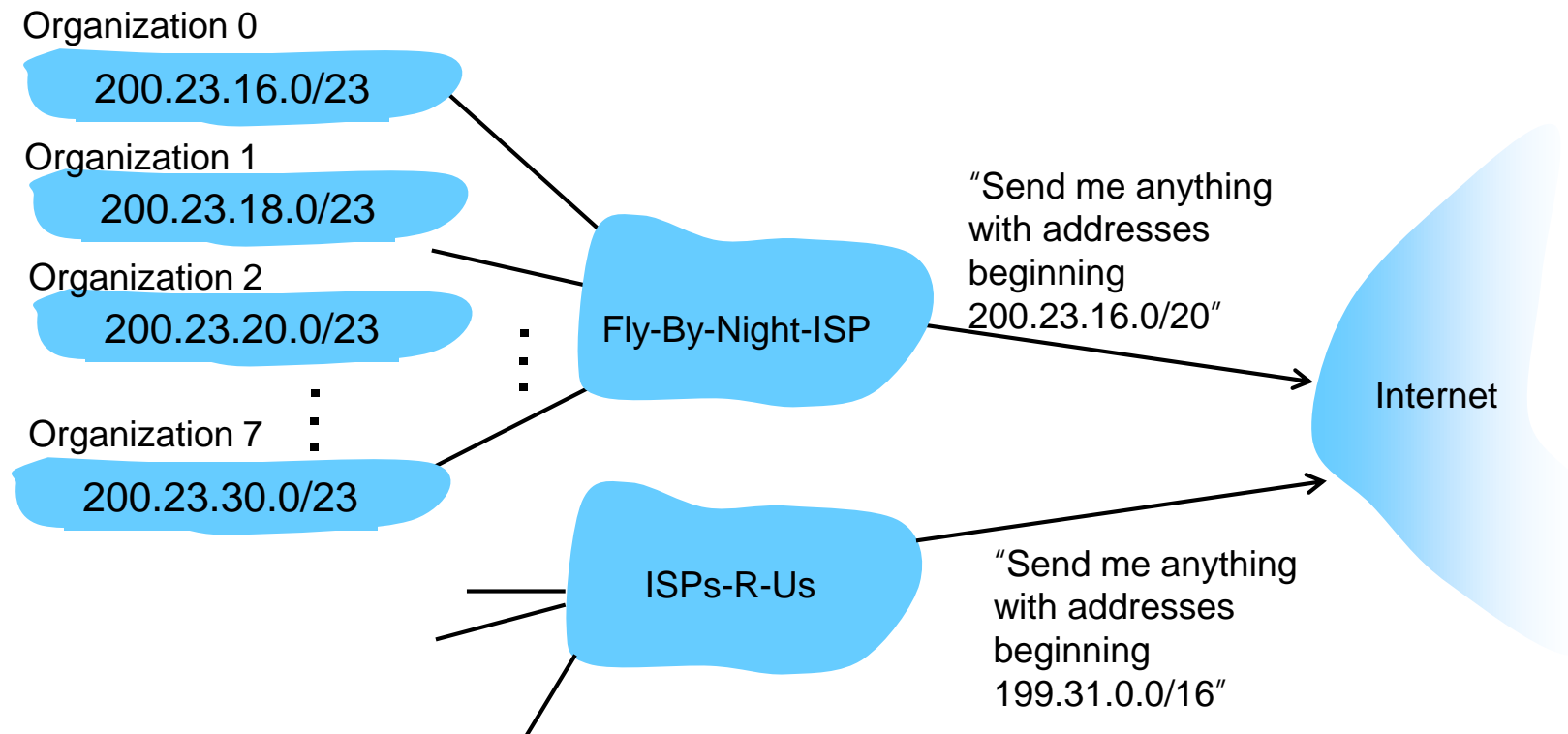
Q: how does *network* get subnet part of IP addr?

A: gets allocated portion of its provider ISP's address space

ISP's block	<u>11001000</u>	<u>00010111</u>	<u>00010000</u>	00000000	200.23.16.0/20
Organization 0	<u>11001000</u>	<u>00010111</u>	<u>00010000</u>	00000000	200.23.16.0/23
Organization 1	<u>11001000</u>	<u>00010111</u>	<u>00010010</u>	00000000	200.23.18.0/23
Organization 2	<u>11001000</u>	<u>00010111</u>	<u>00010100</u>	00000000	200.23.20.0/23
...
Organization 7	<u>11001000</u>	<u>00010111</u>	<u>00011110</u>	00000000	200.23.30.0/23

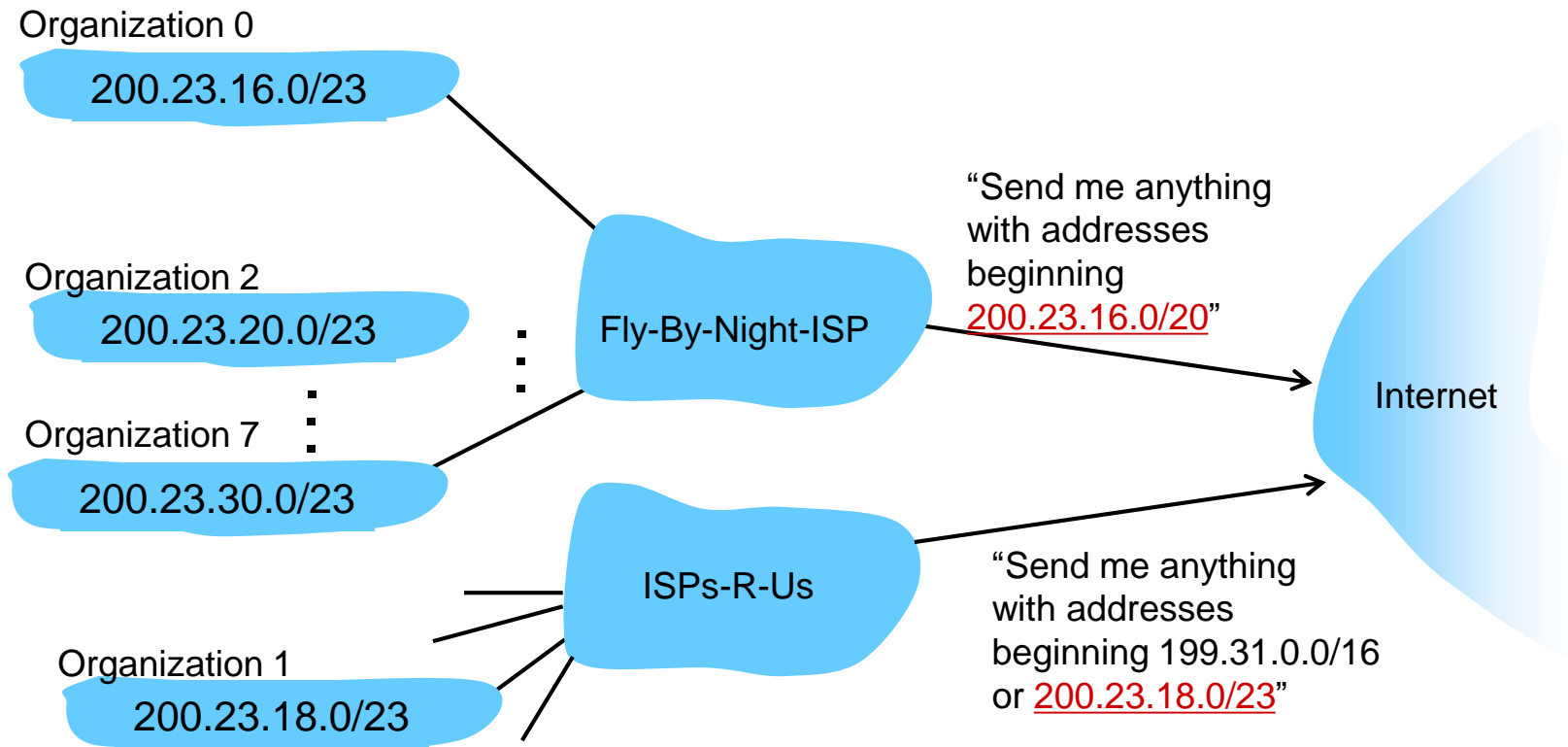
Hierarchical addressing: route aggregation

hierarchical addressing allows efficient advertisement of routing information:



Hierarchical addressing: more specific routes

ISPs-R-Us has a more specific route to Organization 1



IP addressing: the last word...

Q: how does an ISP get block of addresses?

A: ICANN: Internet Corporation for Assigned Names and Numbers <http://www.icann.org/>

- allocates addresses
- manages DNS
- assigns domain names, resolves disputes
- There are no more IPv4 addresses to give out!

Who's got them...

Companies/orgs with IPv4 /8 blocks from IANA – Internet Assigned Numbers Authority (a dept. of ICANN)

Owner	/8 Blocks	~IP addresses
US Military (Department of Defense etc.)	12	201 million
Level 3 Communications, Inc.	2	33 million
Hewlett-Packard	2	33 million
AT&T Bell Laboratories (Alcatel-Lucent)	1	16 million
AT&T Global Network Services	1	16 million
Bell-Northern Research (Nortel Networks)	1	16 million
Amateur Radio Digital Communications	1	16 million
Apple Computer Inc.	1	16 million
Cap Devis CCS (Mercedes-Benz)	1	16 million
Computer Sciences Corporation	1	16 million
Department of Social Security of UK	1	16 million
E.I. duPont de Nemours and Co., Inc.	1	16 million
Eli Lilly and Company	1	16 million
Ford Motor Company	1	16 million
General Electric Company	1	16 million
Halliburton Company	1	16 million
IBM	1	16 million
Interop Show Network	1	16 million
Merck and Co., Inc.	1	16 million
MERIT Computer Network	1	16 million
Massachusetts Institute of Technology	1	16 million
Performance Systems International (Cogent)	1	16 million
Prudential Equity Group, LLC	1	16 million
Société Internationale De Telecommunications Aero.	1	16 million
U.S. Postal Service	1	16 million
UK Ministry of Defence	1	16 million
Xerox Corporation	1	16 million

<http://royal.pingdom.com/2008/02/13/where-did-all-the-ip-numbers-go-the-us-department-of-defense-has-them/>