

# Chapter 4: outline

## 4.1 introduction

## 4.2 virtual circuit and datagram networks

## 4.3 what's inside a router

## 4.4 IP: Internet Protocol

- datagram format
- IPv4 addressing
- ICMP
- IPv6

## 4.5 routing algorithms

- link state
- distance vector
- hierarchical routing

## 4.6 routing in the Internet

- RIP
- OSPF
- BGP

## 4.7 broadcast and multicast routing

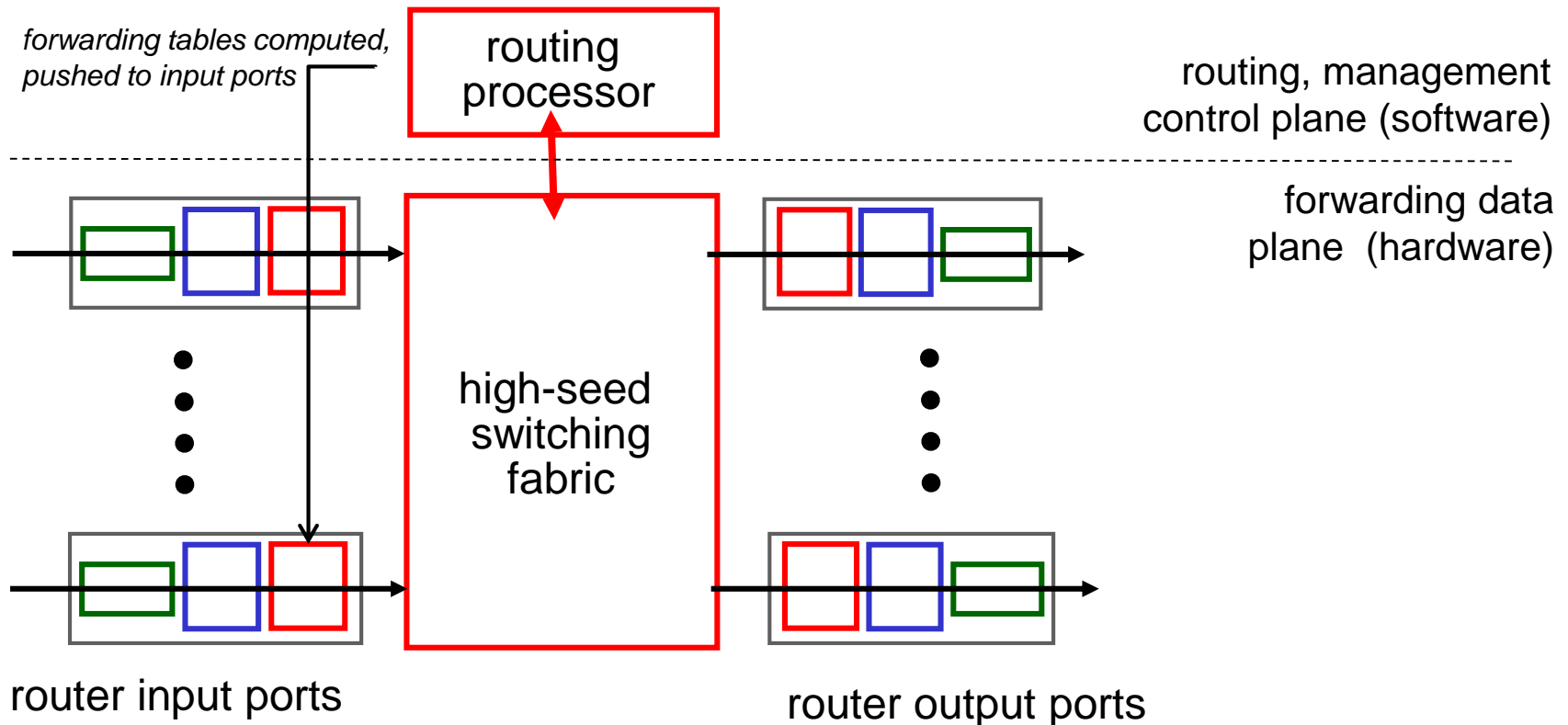
# What is a router?

- ❖ Modern routers have varying amounts of input ports and output ports.
- ❖ Home-grade, SMB “routers” typically have the following internal components:
  - Switch
  - Router
  - Firewall
  - Wireless Radio
- ❖ Advanced devices often have:
  - VPN
  - Port Forwarding
  - QoS (Quality of Service)
  - Separate control over each port

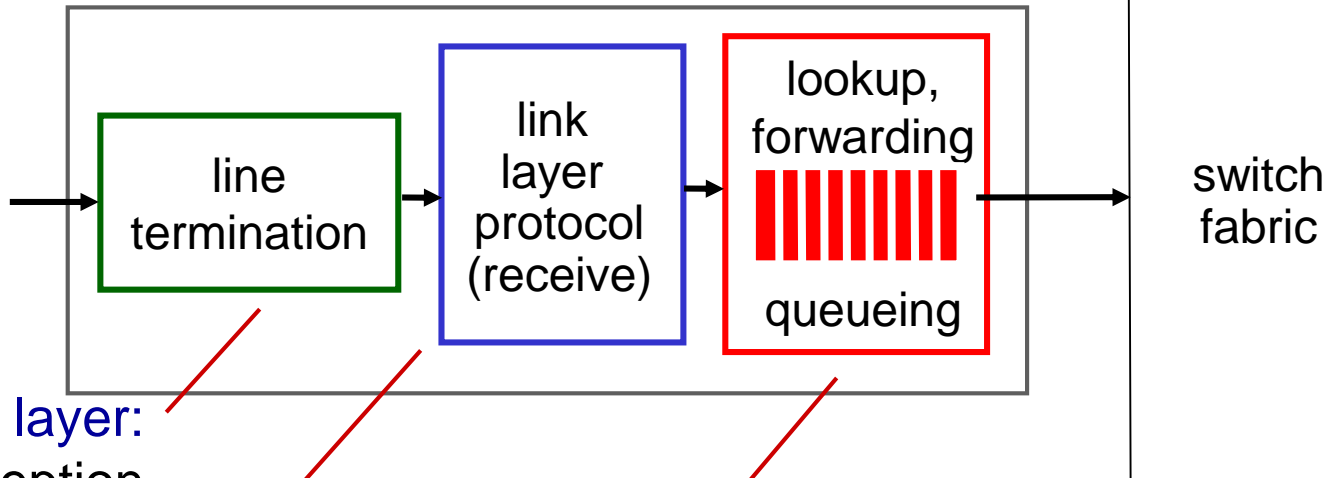
# Router (only) architecture overview

two key router functions:

- ❖ run routing algorithms/protocol (RIP, OSPF, BGP)
- ❖ *forwarding* datagrams from incoming to outgoing link



# Input port functions



physical layer:  
bit-level reception

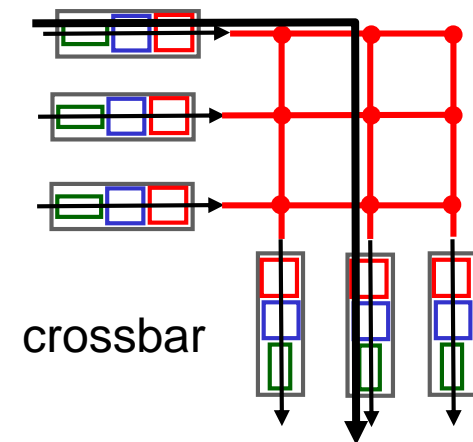
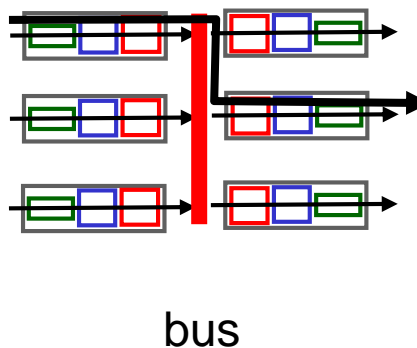
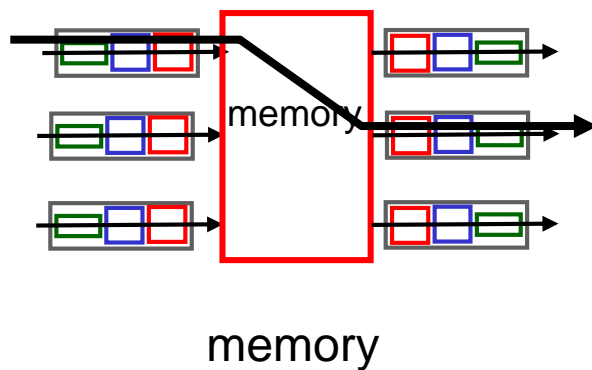
data link layer:  
e.g., Ethernet  
see chapter 5

## decentralized switching:

- ❖ given datagram dest., lookup output port using forwarding table in input port memory (*“match plus action”*)
- ❖ goal: complete input port processing at ‘line speed’
- ❖ queuing: if datagrams arrive faster than forwarding rate into switch fabric

# Switching fabrics

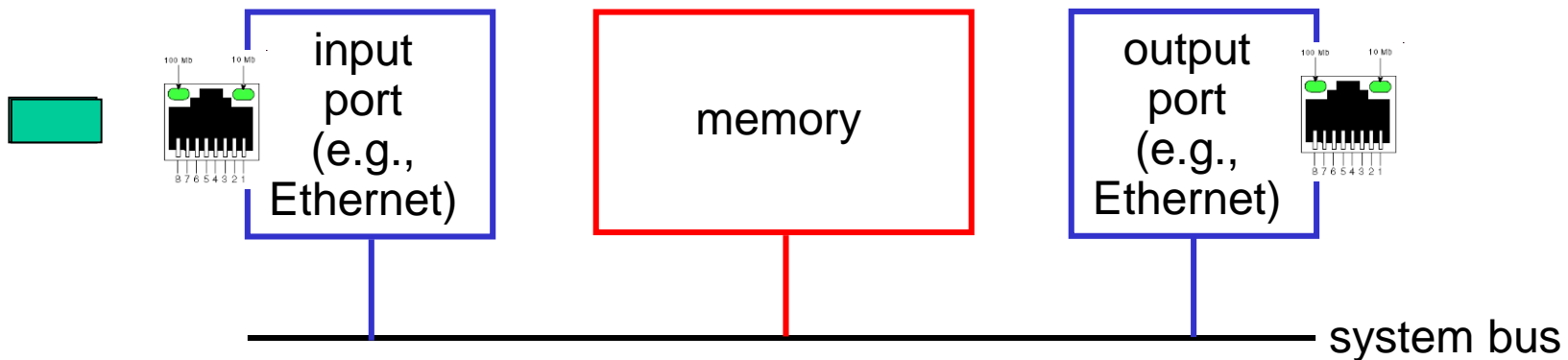
- ❖ transfer packet from input buffer to appropriate output buffer
- ❖ switching rate: rate at which packets can be transfer from inputs to outputs
  - often measured as multiple of input/output line rate
  - N inputs: switching rate N times line rate desirable
- ❖ three types of switching fabrics



# Switching via memory

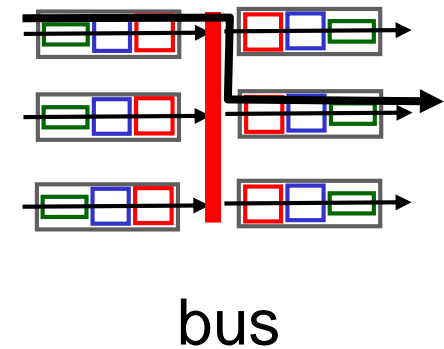
## *first generation routers:*

- ❖ traditional computers with switching under direct control of CPU
- ❖ packet copied to **system's** memory
- ❖ speed limited by memory bandwidth (2 bus crossings per datagram)



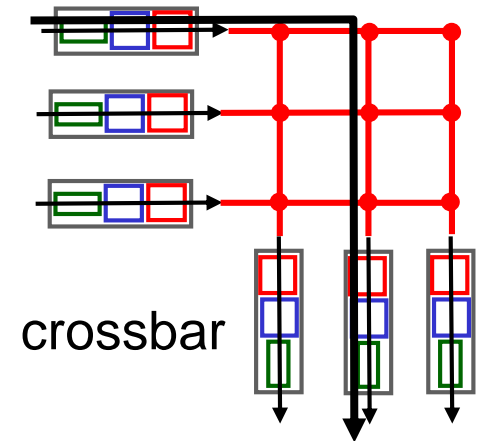
# Switching via a bus

- ❖ datagram from input port memory copied to output port memory via a shared bus
- ❖ *bus contention*: switching speed limited by bus bandwidth
- ❖ 32 Gbps bus, Cisco 5600: sufficient speed for access and enterprise routers



# Switching via interconnection network

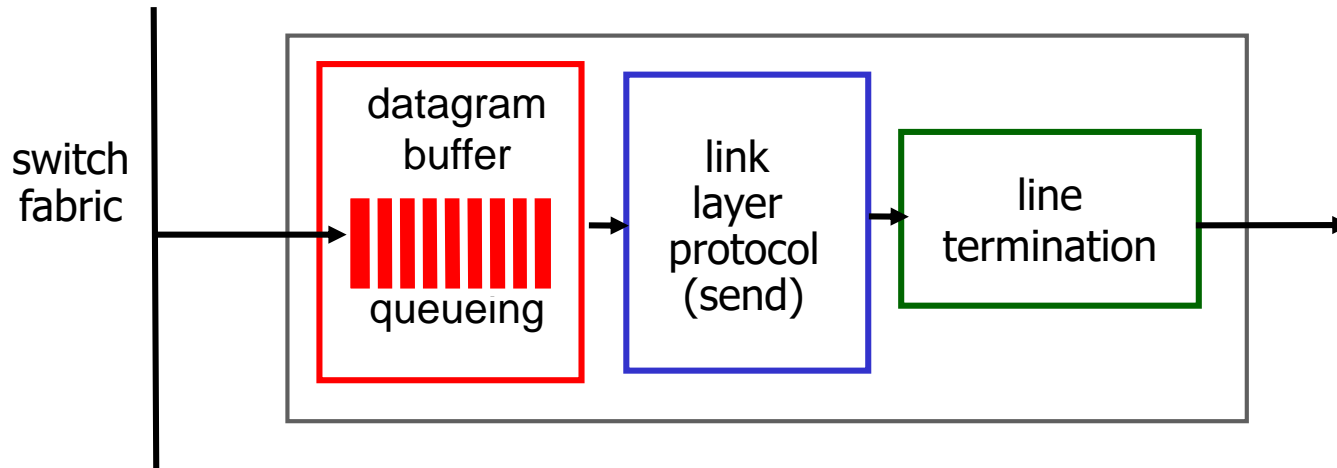
- ❖ overcome bus bandwidth limitations
- ❖ banyan networks, crossbar, other interconnection nets initially developed to connect processors in multiprocessor
- ❖ advanced design: fragmenting datagram into fixed length cells, switch cells through the fabric.
- ❖ Cisco I2000: switches 60 Gbps through the interconnection network





# Output ports

*This slide is HUGELY important!*



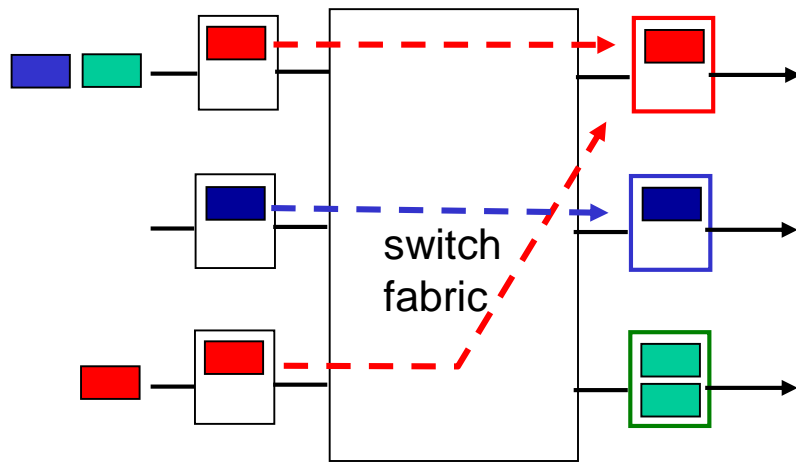
- ❖ *buffering* required from fabric faster rate

Datagram (packets) can be lost due to congestion, lack of buffers

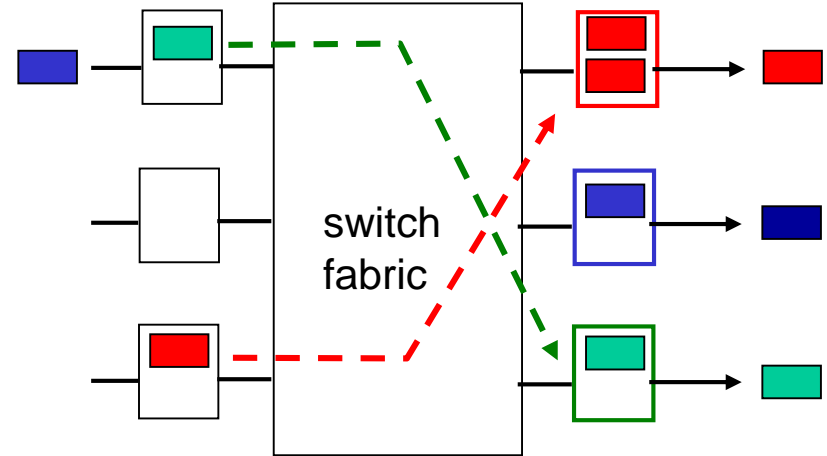
- ❖ *scheduling* datagrams

Priority scheduling – who gets best performance, network neutrality

# Output port queueing



at time  $t$ , packets move  
from input to output



one packet time later

- ❖ buffering when arrival rate via switch exceeds output line speed
- ❖ *queueing (delay) and loss due to output port buffer overflow!*

# How much buffering?

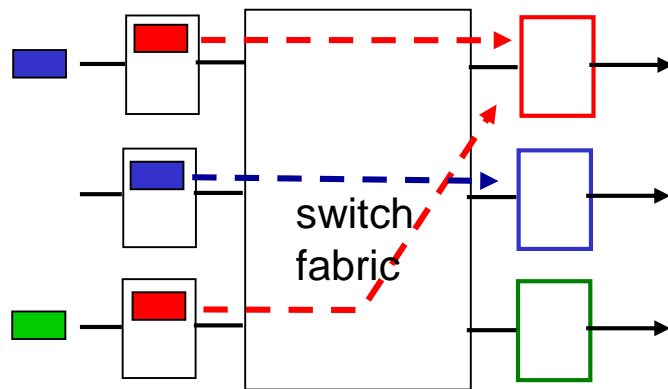
- ❖ RFC 3439 rule of thumb: average buffering equal to "typical" RTT (say 250 msec) times link capacity  $C$ 
  - e.g.,  $250\text{ms} * 10 \text{ Gbps} \Rightarrow 2.5 \text{ Gbit buffer}$
- ❖ recent recommendation: with  $N$  TCP "flows", buffering equal to 
$$\frac{\text{RTT} \cdot C}{\sqrt{N}}$$

thus, with  $\text{RTT} = 0.25\text{s}$ ,  $C = 10 \text{ Gbps}$ , and 10 flows:

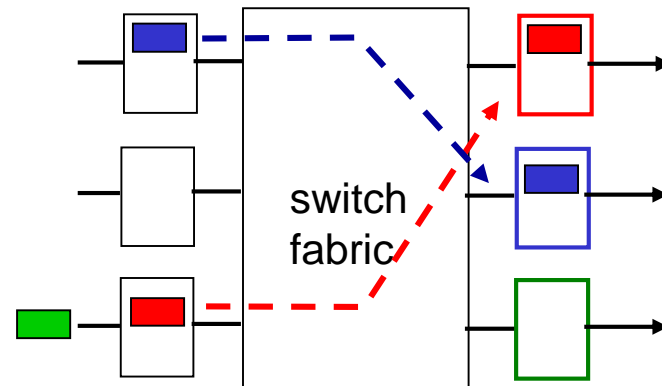
- e.g.,  $(0.25\text{s} * 10 \text{ Gbps}) / \text{sqrt}(10) \Rightarrow 0.79 \text{ Gbit buffer}$

# Input port queuing

- ❖ fabric slower than input ports combined -> queueing may occur at input queues,
  - *queueing delay and loss due to input buffer overflow!*
- ❖ **Head-of-the-Line (HOL) blocking:** queued datagram at front of queue prevents others in queue from moving forward



output port contention:  
Assume only one **red** datagram  
can be transferred per time  $t$ .  
*lower red packet is blocked*



one packet time later: **green**  
packet experiences HOL  
blocking - can't be sent, even  
though it's desired output is  
not busy!

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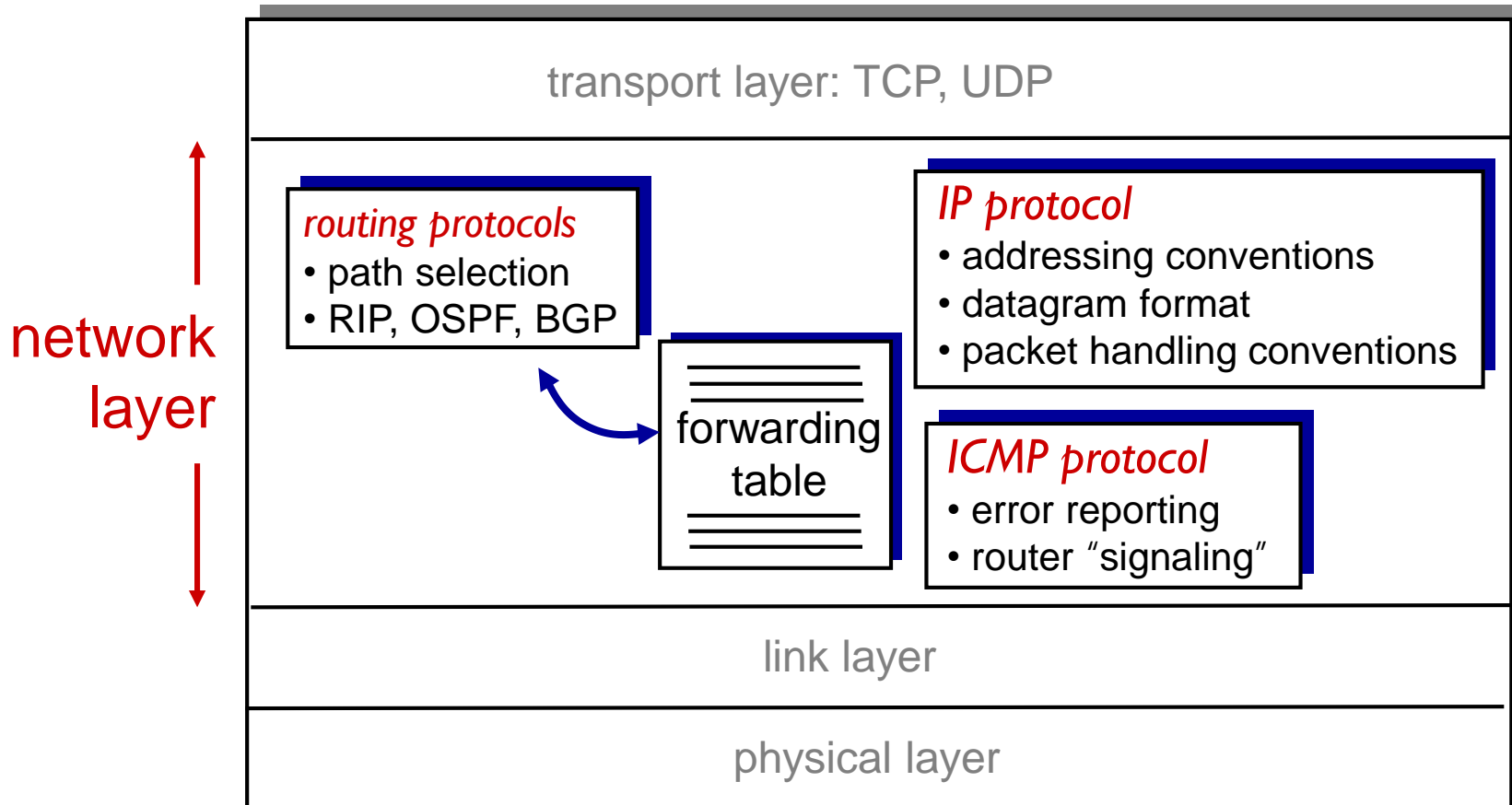
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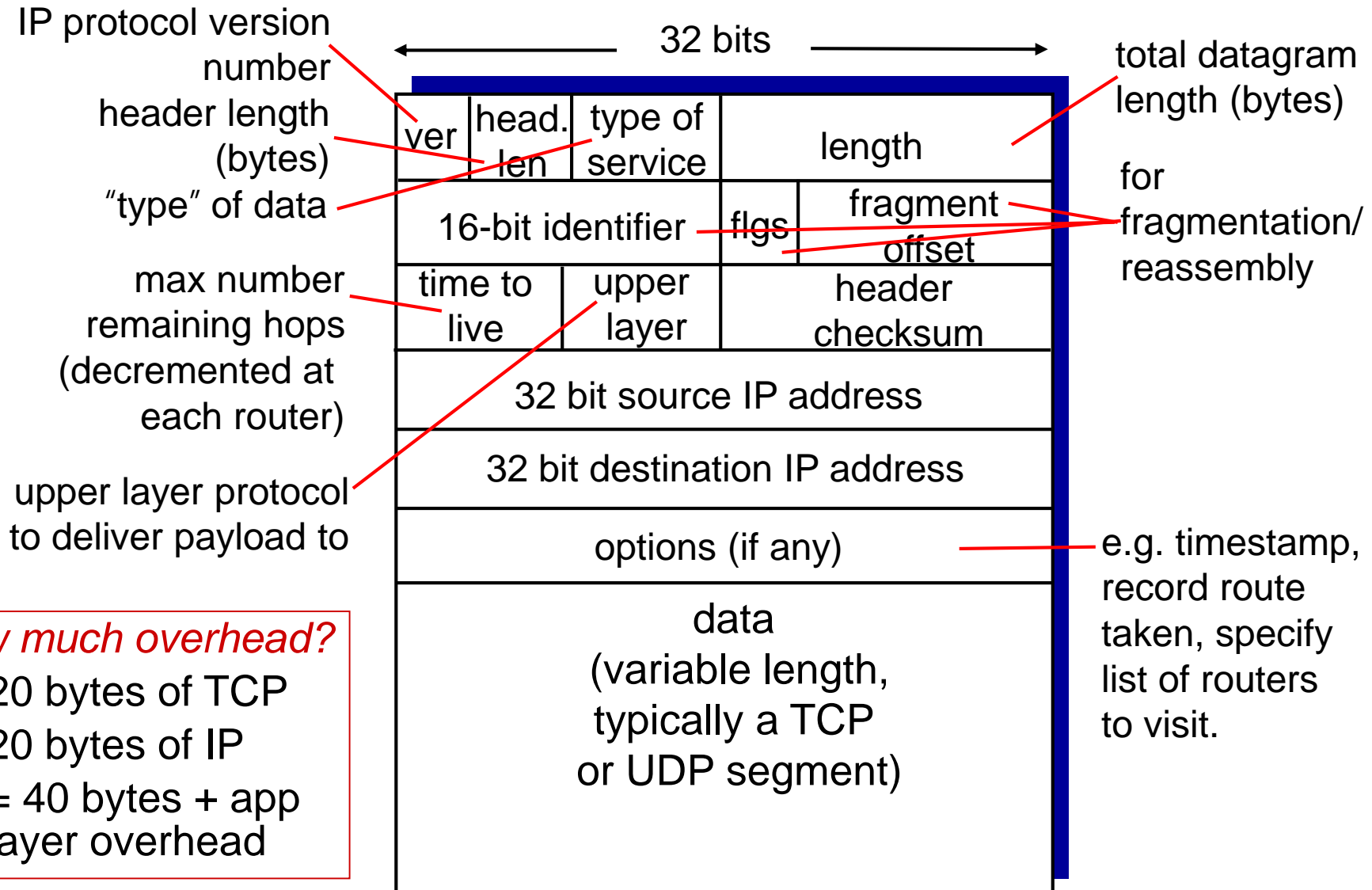
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# The Internet network layer

host, router network layer functions:

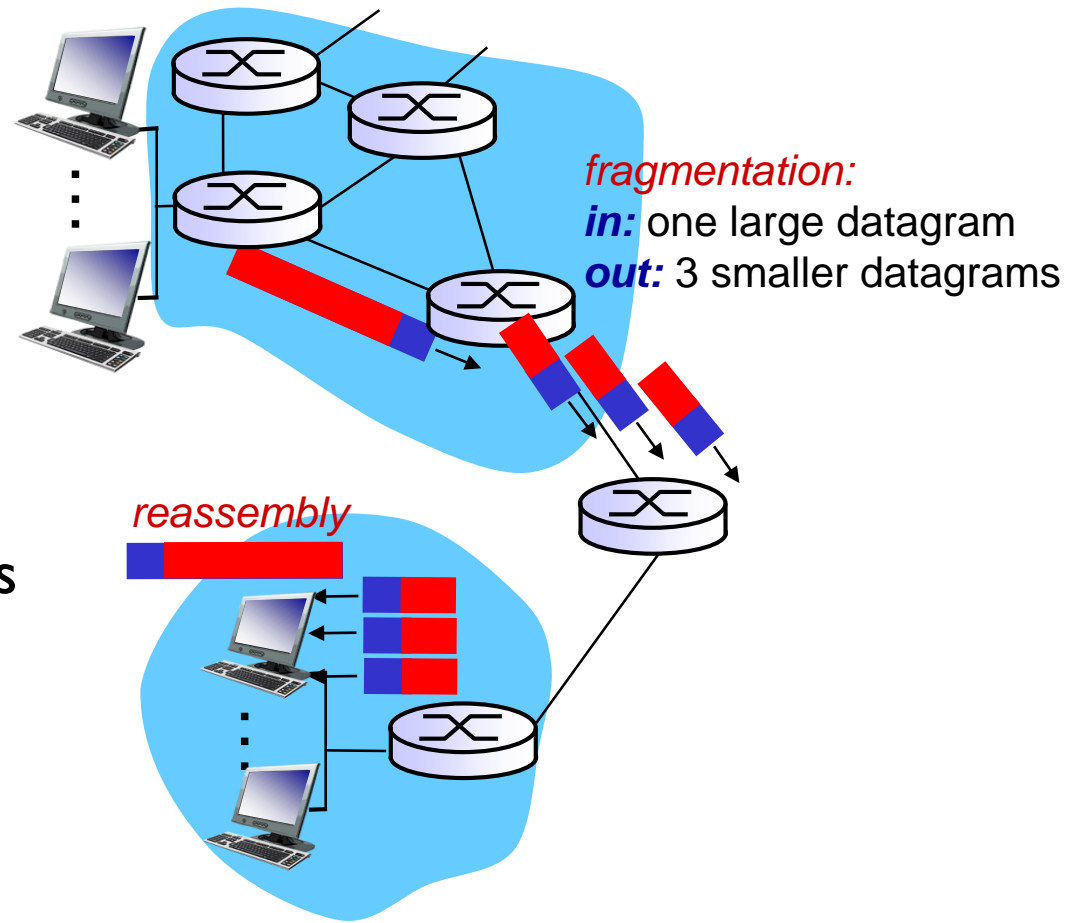


# IP datagram format



# IP fragmentation, reassembly

- ❖ network links have MTU (max.transfer size) - largest possible link-level frame
  - different link types, different MTUs
- ❖ large IP datagram divided ("fragmented") within net
  - one datagram becomes several datagrams
  - "reassembled" only at final destination
  - IP header bits used to identify, order related fragments





# IP fragmentation, reassembly

## *example:*

- ❖ 4000 byte datagram
- ❖ MTU = 1500 bytes

	length	ID	fragflag	offset	
	=4000	=x	=0	=0	

*one large datagram becomes  
several smaller datagrams*

1480 bytes in  
data field

	length	ID	fragflag	offset	
	=1500	=x	=1	=0	

	length	ID	fragflag	offset	
	=1500	=x	=1	=185	

	length	ID	fragflag	offset	
	=1040	=x	=0	=370	

offset =  
 $1480/8$

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