

Experiments on automation of formal verification of devices at the binary level

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Abstract

With the advent of virtualization, more and more work is put into the verification of hypervisors. Being low level softwares, such verification should preferably be performed at binary level. Binary analysis platforms are being developed to help perform these proofs, but a lot of the work has to be carried out manually.

In this thesis, we focus on the formal verification of a Network Interface Controller (NIC), more specifically we look at how to automate and reduce the boilerplate work from an existing proof. We base our work on the HolBA platform, its hardware-independent intermediate representation language BIR and supporting tools, and we experiment on how to perform this proof by leveraging existing tools.

We first replaced the existing NIC model written in HOL4 to an equivalent one written using BIR, enabling the use of HolBA tools. Secondly, we developed some visualization tools to help navigate and gain some insight in the existing proof and its structure. Thirdly, we experimented with the use of Hoare triples in conjunction with an SMT solver to perform contract verification. Finally, we proved a simple contract written in terms of the formal NIC model on the BIR implementation of this model, unlocking the way of performing more complex proofs using the HolBA platform.

Keywords: binary analysis, formal verification, proof producing analysis, theorem proving

Résumé

Avec la démocratisation de la virtualisation, de plus en plus d'efforts sont consacrés à la vérification des hyperviseurs. S'agissant de logiciels de bas niveau, une telle vérification devrait de préférence être effectuée au niveau binaire. Des plates-formes d'analyse binaire sont en cours de développement pour aider à réaliser ces preuves, mais une grande partie du travail doit encore être effectuée manuellement.

Dans cette thèse, nous nous concentrons sur la vérification formelle d'un Contrôleur d'Interface Réseau (NIC), plus spécifiquement sur la manière d'automatiser et de réduire le travail répétitif d'une preuve existante. Nous nous basons sur la plate-forme HolBA, son langage de représentation intermédiaire indépendant du matériel, BIR et ses outils de support, et nous nous intéressons à la manière de réaliser cette preuve en utilisant des outils existants.

Nous avons d'abord remplacé le modèle NIC existant écrit en HOL4 par un modèle équivalent écrit en BIR, permettant ainsi l'utilisation des outils de HolBA. Deuxièmement, nous avons développé des outils de visualisation pour nous aider à naviguer et à mieux comprendre la preuve existante et sa structure. Troisièmement, nous avons expérimenté l'utilisation des triplets de Hoare en conjonction avec un solveur SMT pour effectuer une vérification par contrat. Enfin, nous avons prouvé un contrat simple écrit en termes du modèle formel du NIC sur l'implémentation de ce modèle en BIR, ouvrant la voie à la réalisation de preuves plus complexes avec la plate-forme HolBA.

Mot-clés : binary analysis, formal verification, proof producing analysis, theorem proving

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Glossary

BIR HolBA’s machine-independent binary intermediate representation. 3, 9, 31, 36

BSL BIR Simple Language. It is also a word play with DSL which means Domain Specific Language. 42

CFG Control Flow Graph. 16, 25, 26, 41

CMU BAP Carnegie Mellon University’s Binary Analysis Platform. 3, 8

CPU Control Processing Unit. This is the core processing unit of a computer. 5

DMA Direct Memory Access. 5

FreeRTOS FreeRTOS is a free and open-source market leading real time operating system for microcontrollers and small microprocessors. Website: <https://freertos.org/>. 1

HolBA HOL4 Binary Analysis Platform. vii, 2, 3, 5, 7, 9, 31, 36, 53

IoT Internet Of Things. This is the idea of connecting every devices, such that temperature sensors, lamps, video cameras or wearable devices, to the Internet. 1

IR Intermediate Representation. 7, 18

ISA Instruction Set Architecture. 3, 8, 10

MMU Memory Management Unit. 5

NIC Network Interface Controller. 2–4, 11, 33, 36, 41, 43, 49

proof-producing A proof-producing analysis or procedure is a procedure that yields its result in the form of a formal proof. Therefore, those procedures represent the highest degree possible of trustworthiness. However, great care must

be taken to ensure that the definitions used to phrase the proof are sound and actually represent the properties to be proved. Moreover, implementing proof-producing code is costly and thus should be used with discernment. 3, 9, 18, 24, 31, 33, 37, 46, 49

PROSPER Provably Secure Execution Platforms for Embedded Systems. 1, 6, 11

SAT Boolean SATisfiability Problem. 30

SML Standard ML. 9, 12, 24

SMT Satisfiability Modulo Theories. 29, 30, 33, 53, 55

transpilation Process carried out by a **transpiler**. 3

transpiler A transpiler, or source-to-source compiler, is a compiler that takes source code written in a programming language A, and producing an equivalent source code in a programming language B. A and B can be different languages or the same. They differ from traditional compilers in that compilers generally produce machine code. viii, 3

Chapter 1

Introduction

This chapter serves as an introduction to the degree project and presents the background of the work along with this thesis objective. Delimitations to the project and the choice of methodology are also discussed.

1.1 Background

Embedded systems are becoming more and more common with the current advent of **IoT** and mobile computing platforms such as smartphones. Those systems are fully-fledged computers with powerful hardware, complete operating systems and access to Internet. Such systems can run security-critical services, such as a building security system or automatic toll gates, or carry valuable information as it is the case for personal smartphones. Therefore, these two characteristics make them targets of choice for attackers.

The **Provably Secure Execution Platforms for Embedded Systems (PROSPER)** project [1] aims to develop a secure and formally verified hypervisor for embedded systems. Hypervisors are thin layers running directly on top of hardware providing the ability to run virtualized applications, such that operating systems or realtime control systems. Those virtualized applications then don't have privileged access to the hardware and have to go through the hypervisor. This allows different applications to share the same hardware while providing strong isolation between them, thus ensuring confidentiality and security. Moreover, security not only means protection from external attacks, but also resilience to bugs. If multiple critical systems are running on the same hardware, bugs or crashes in some systems shouldn't affect the others from behaving correctly. Figure 1.1 shows a system running two isolated Linux on top of a hypervisor.

Previous work in the **PROSPER** project achieved [2] to formally verify a simple separation kernel [3], which later resulted into an implementation of a working hypervisor. Then, they achieved to run both Linux and **FreeRTOS** on top of it. Finally, they for-

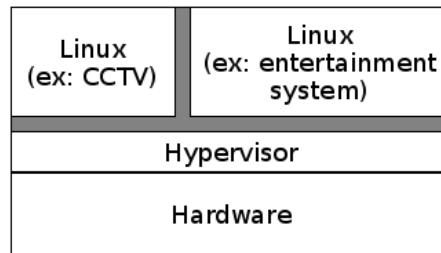


Figure 1.1: Two Linux on top of an hypervisor. They run isolated from each other and from the hypervisor.

mally verified memory isolation for virtualized applications [4]. Now, among other projects, the PROSPER team is working on device virtualization, allowing to give access to hardware devices to virtualized applications. An interesting example are **Network Interface Controller (NIC)** devices, which enable network communication and give the ability to communicate through the Internet.

A formal model of a **NIC** device has already been produced, on which some security theorems have been proved [5]. These high-level proofs relying on a layer of lower-level lemmas. This layer provides an abstraction over the raw formal model. This is illustrated in the left-hand side of Figure 1.2.

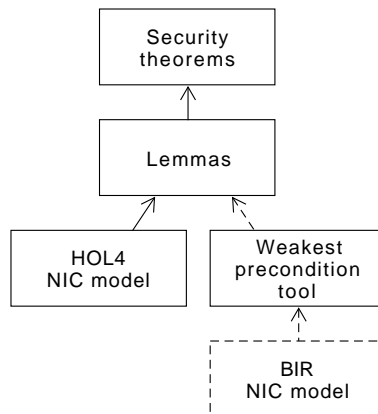


Figure 1.2: HOL4 v. BIR NIC models. The left hand side already exists. This project would consist in the dashed elements. The dotted lines represent the work to be done during this project.

The team is now developing a new framework for performing binary analysis in HOL4, an interactive theorem prover, named **HOL4 Binary Analysis Platform (HolBA)** [6]. This framework is based on two papers written in the team. The first one in-

troduces sound **transpilation** from binary to machine-independent code ¹[7]. The second paper, “TrABin: Trustworthy Analyses of Binaries” [8], lays the foundations of the **HolBA** platform: it formally models **BIR**, introduces various supporting tools, implements two **proof-producing transpilers** (ARMv8 and Cortex-M0) and a proof-producing weakest precondition generator for loop-free programs.

While this kind of **transpilers** and **proof-producing** weakest precondition tools already exist ², the novelty in this work is that the transpiler is proof-producing, i.e. it produces a formal proof that both binary representations are equivalent, under the simulation theory, with respect to the **Instruction Set Architecture (ISA)** model. With this method, you no longer need to trust the transpiler. Figure 1.3 gives an overview of the TrABin framework³.

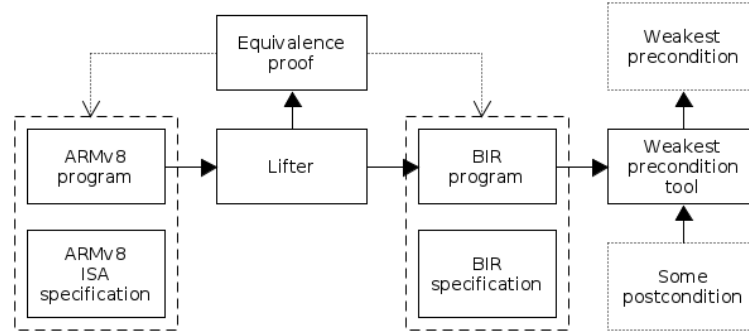


Figure 1.3: The HolBA framework. The lifter generates a BIR program and an equivalence proof from an ARMv8 program. The equivalence proof establishes a simulation property between the ARMv8 binary program and the generated BIR binary program, showing that they have the same behaviour with respect to the ARMv8 ISA specification and the BIR specification. HolBA also support the Cortex-M0 ISA.

The idea of this work is to translate the formal **NIC** model of [5] using **BIR**, then use HolBA’s proof-producing weakest precondition tool to prove the same lower-level lemmas than the formal model. With all the lemmas proved, the security properties are implied. Figure 1.2 gives an overview of this idea: using the proof-producing weakest precondition tool to bind together a newly written BIR NIC model and the work done on the formal model.

¹The machine-independent language used in the work is an implementation of **Carnegie Mellon University’s Binary Analysis Platform (CMU BAP)**’s BIL [9]. This implementation will evolve later in the TrABin paper into BIR that HolBA uses.

²See related discussion in [8].

³TrABin works with both ARMv8 and Cortex-M0 binary programs. Only ARMv8 is showed in Figure ?? to save some space.

1.2 Intended readers

In this thesis, formal verification is the central topic. The thesis presents how model a hardware device using a binary analysis platform and presents some formal verification techniques. A reader interested in this topic may find the results presented useful for further work. A casual reader will be presented with a light introduction to the concepts that are essential in order to understand the problem that this thesis aims to explore. Those concepts should also be useful in order to understand software verification in general. For a more coherent reading, the concepts are introduced throughout the document. However, the reader is expected to have a background in Computer Science in general, and knowledge in formal verification will make the thesis easier to digest.

1.3 Thesis objective

The primary goal of this thesis project is to explore verification techniques in order to automate parts, if not all, of the verification process of hardware devices using the HolBA platform. The formal **NIC** model of [5] is used as working example.

The ultimate goal would be to obtain a fully automatic pipeline for performing such verifications. However, it is evident that goal isn't reachable in such a small amount of time, or even at all. Thus, this thesis focuses instead on exploring what toolkit is needed in order to facilitate this work.

1.4 Delimitations

TODO: Use this section to explain the time limits that have made it impossible to perform certain experiments.

1.5 Choice of methodology

This work has been carried out step-by-step toward an ideal goal, i.e. re-establishing all the security properties. On the road, needs have been identified and tools have been implemented in order to tackle them. This approach made sense in this particular work because the needs weren't known in advance, and therefore needed to be identified. This thesis presents the steps taken during this work, the motivations of each tool that have been implemented, and discusses their limitations and future work in the conclusion.

Chapter 2

Related work

*This chapter will present the related work in the domain of secure execution platforms and binary analysis. After briefly introducing memory sharing between CPU and devices, it will explain the need of secure execution platforms, present existing binary analysis platforms and explain the novelty of **HolBA**, the platform used in this work.*

2.1 Memory sharing between CPU and devices

Figure 2.1 represents how the CPU and devices can share the main memory using **Direct Memory Access (DMA)**¹. To read from/write to the main memory, the CPU passes through the **Memory Management Unit (MMU)**, which is responsible for virtual/physical address translation. In a nutshell, memory addresses that the CPU uses are mapped on a virtual address space—each running process having its own—that are mapped to the actual physical addresses by the MMU. This enables virtualization and abstraction between processes, which can each have their own memory spaces in apparent isolation.

Similarly, devices can directly access the main memory using DMA. This technique enables to offload the CPU from copying each byte of memory and instead setting the DMA controller to do so. DMA controllers also give devices direct access to the main memory. While fast and convenient, DMA creates a whole new range of vulnerabilities. Indeed, if misconfigured, the DMA controller can give complete access to the main memory to devices, like kernel private memory, page table or executable memory [10].

2.2 The necessity of secure execution platforms

TODO: Explain the “necessity”.

¹This schema doesn't represent caches.

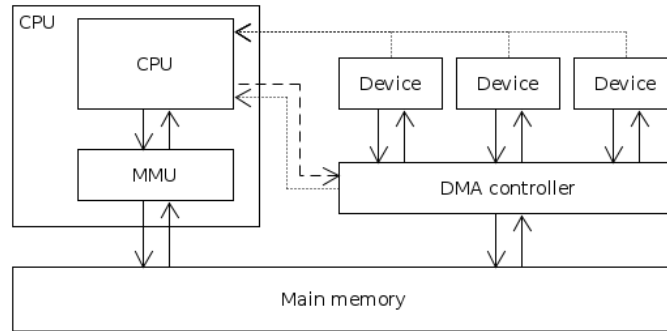


Figure 2.1: CPU schema from a memory point of view. The dashed line between the CPU and the DMA controller represent the capability of the CPU to send commands to the DMA controller through register writes. Dotted lines to the CPU represent the ability to raise interrupts.

The **PROSPER** project isn’t the only project focused on high-security execution platforms. Platforms such that seL4, Microsoft Hyper-V and INTEGRITY Multivisor are examples of platforms already used in production and providing strong security properties.

seL4 is a recent L4-based microkernel created in 2006 with the goal to produce a completely formally verified implementation of a L4 microkernel. This has been achieved in 2009 [11]. At this time, seL4 consisted of 8700 lines of C code and 600 lines of assembler. The implementation of seL4 has been formally proved from its abstract specification down to its C implementation. Later, the validity of the generated assembly code has been proved, removing the need of trusting the compiler [12].

Microsoft Hyper-V is Microsoft’s hypervisor, widely used today within the Microsoft Azure cloud platform. It has been released in 2008. Hyper-V is a huge code-base, as we can read on VCC’s website²: “Hyper-V consists of about 60 thousand lines of operating system-level C and x64 assembly code, it is therefore not a trivial target”. Microsoft has put a lot of work in formal verification³ of Hyper-V down to machine code [13]. However, they don’t appear to include device drivers in their formal verification.

INTEGRITY Multivisor is a commercial real-time operating system developed by Green Hills Software. Although not much information seems to be publicly available,

²<https://www.microsoft.com/en-us/research/project/vcc-a-verifier-for-concurrent-c/>

³Microsoft has several formal verification projects, many of which are freely available for non-commercial use: <https://github.com/Microsoft?q=verifier>

Green Hills Software has done considerable formal verification work [14]. Multivisor has several certifications, including, for example, ISO 26262 ASIL D automotive electronics, NSA-certified secure mobile phones or FAA DO-178B Level A-certified avionics controlling life-critical functions on passenger and military aircraft ⁴.

2.3 Binary analysis platforms

For this project, we will use the **HolBA** framework. However, several other binary analysis platforms have been created for various purposes, such as formal verification or static analysis. A common characteristic of these platforms is to use an **Intermediate Representation (IR)**. IRs are designed to be simpler to use for each platform's end purpose. As an example, the HolBA platform has BIR as its intermediate representation (BIR is presented in Section 4.1).

Microsoft Boogie is Microsoft's intermediate verification language. Boogie is the IR for multiple Microsoft tools, including VCC. Boogie as a tool can infer some invariants on the given Boogie program and then generate verification conditions that are passed to an SMT solver ⁵.

Valgrind is a framework for building program supervision tools, such as memory checkers, cache profilers or data-race detectors [15]. As its core, Valgrind is a JIT x86-to-x86 compiler, translating binary programs into its IR called UCode. Then, *skins*—the tools built on the Valgrind framework—are free to transform and work in the IR to perform several analysis.

LLVM is a compiler infrastructure which supports a unique multi-stage optimization system [16]. LLVM is built around its IR, LLVM Virtual Instruction Set, which can be described as a strict RISC architecture with high-level type information. This IR was what made LLVM successful because it is a pragmatic IR suitable for optimizations at multiple stages (link-, post-link and run-time) and supporting a wide variety of transformations. Leveraging this IR, an ecosystem grew around LLVM, providing tools such as symbolic execution (LLVM KLEE), benchmarking environments or static and dynamic analysers.

Mayhem is a system for automatically finding exploitable bugs in binary programs and generating working exploits as proof of the discovered vulnerabilities [17]. It leverages BAP, the Binary Analysis Platform from Carnegie Mellon University (CMU

⁴https://ghs.com/products/rtos/integrity_virtualization.html

⁵<https://www.microsoft.com/en-us/research/project/boogie-an-intermediate-verification-language/>

BAP) [18], as its IR. It proceeds by first JIT-ing each instruction to the BAP intermediate language (IL) and then performing a custom symbolic execution.

There are several other tools, platforms and frameworks, such as **CMU BAP** [9] or Angr. An interesting note though is that BIR's design is based upon CMU BAP's intermediate language.

The novelty introduced in HolBA is that proofs are performed directly on the generated assembly code, not at the source code level. Therefore, proofs can be performed on programs without needing their source code, and regardless on the programming language used as long as it can be compiled in assembly code in an **ISA** supported by the platform.

Chapter 3

Overview of the formal verification of the NIC device

*This chapter intends to introduce the concepts that are essential to the reader in order to understand the problem that this degree project aims to explore. This includes a presentation of Interactive Theorem Proving and formal proofs in HOL4, an introduction to **BIR**, completing what has already been said about the **HolBA** framework in the Introduction, and a brief overview of the formal NIC model and proof of [5].*
TODO TODO

3.1 Interactive Theorem Proving and HOL4

Interactive theorem provers are software producing formal proofs, in an interactive fashion, i.e. a human can step through the proof interactively while the proof assistant provides some automation (like rewriting of terms, arithmetic evaluation, integration with external tools like SMT solvers, ...). Coq, HOL4 or Isabelle are such tools.

HOL4 [19] stands for Higher-Order Logic. It is a programming environment deeply embedded into the **Standard ML (SML)** programming language enabling to prove theorems and write **proof-producing** programs. Since its first version in 1988, HOL has been focused on hardware verification, and has been successfully used in this domain, as shows its `examples/` directory [20].

HOL4 is built around a very small kernel: 3 axioms, 10 inference rules, two predefined constants and two types ¹[21]. A theorem can only be built using this core, every other higher-level theory must be built from lower-level ones. Furthermore, the kernel is consist of very few lines of code, leaving less space for bugs. Hence, for this reason, the trust needed is far inferior to standard systems.

¹This is true for the Hol Light Kernel, however the HOL4 kernel has slightly deviated from this simplicity for historical and performance reasons.

Several models of different ARM ISA have been realized, such as ARMv3, ARMv4, ARMv7 [20, 22] and very recently ARMv8 [23]. These models enable making proofs directly at the ISA level, avoiding the necessity to trust the compiler and enabling to formally verify program behaviours according to the CPU specification. However, this process is tedious and cannot actually be performed by hand. To respond to this problem, [7] introduced a proof-producing lifter that moves the verification from raw assembly code to an higher-level intermediate representation, and introduced later some supporting tools [8] in order to automate some of the work. This constitutes the foundation of HolBA.

3.2 Overview of the formal NIC model

The formal NIC model of [5] has been designed from reading the hardware specification of the device, because no model nor device driver is freely available.

The NIC model is designed as a transition system with four types of transitions: register read, register write, autonomous, and memory read request reply transitions. They are described by four functions which constitute its public interface: `read_nic_register`, `write_nic_register`, `nic_transition_autonomous` and `memory_read_reply`. Each of these functions update the given initial NIC state, simulating the real behaviour of the NIC. To make the model sound, the state is marked *dead* if the model is asked to describe any transition or operation that is not described by the specification. Dead states represent undefined states and cannot be left with further transitions.

The model is composed of five finite state automata describing the inner transitions of the whole system, each automata describing a part of the NIC. The automata are: initialisation, transmission, transmission teardown, reception and reception teardown. The function `nic_execute` performs one autonomous step of one of the five automata. The automaton that takes a step is decided by a scheduler, depending on the NIC inner state. If more than one automata are in a ready state, their order is not deterministic². Each finite state automata is defined as a set of transitions working on NIC states. Each transition takes a NIC state and returns an update one, possibly dead.

Being designed as a transition system, each of the four functions described earlier are loop-free, aspect that is crucial in order to apply contract based verification as described in Section 5. Each of the transition is implemented as a HOL4 function using exclusively **if-then-else** and state modification statements. There are no recursive definition. Hence, every function can be seen as some kind of decision tree in which some of the nodes “mutate” the state.

²To reason about non determinism, the notion of oracles has been introduced in [5]. To put it simply, an oracle can be seen as a list of steps to perform, but where each element of the list is undefined. As we will not directly reason about the oracle, we will not explain how this works in more details. The reader can read a thorough explanation in [5].

The five automata are of different complexity, as shown in Table 3.1, with the initialization automaton having only 4 transitions of which only one is autonomous, and the reception automaton being the most complex with 20 transitions. Calling the `nic_transition_autonomous` function makes one of the automata take an autonomous step. Non autonomous transitions are performed by the three other functions.

Automaton	# of (autonomous) transitions	LoC (w/o comments)
Initialization	4 (1)	21
Transmission	7 (5)	182
Transmission teardown	5 (4)	67
Reception	20 (20)	280
Reception teardown	7 (6)	93

Table 3.1: Statistics on each of the automata in the formal NIC model.

The NIC state is defined as a nested *datatype*—a record³ of records of words, booleans and enumerations. However, there is one exception: the NIC works on a memory data structure, called `CPPI_RAM`, which is represented in the formal model as a function from addresses to values.

The structure of the implementation of `read_nic_register` and `write_nic_register` is similar to the implementation of `nic_transition_autonomous` in terms of the statements they use. `memory_read_reply` is implemented as a non autonomous transition of the transmission automaton.

3.3 Overview of the formal proof of the NIC model

It has been formally proved with HOL4 that with **PROSPER**'s hypervisor, augmented with a monitor and in the presence of a **NIC**, only signed Linux code is executed. This security policy has been defined as invariant of the NIC state, then the proof consists of verifying that each transition of the NIC transition system preserves the invariant. This invariant states that the NIC isn't in an undefined state (represented by the dead state), that the transmission and reception queue are well-behaving (non overlapping, finite) and that the NIC is in a valid state (restrict the values of the NIC state variables to valid configurations). This invariant is stated as a conjunction of the invariants of each automaton plus a not-dead conjunct:

$$I_{NIC} \stackrel{def}{=} \neg NIC.dead \wedge I_{init} \wedge I_{tx} \wedge I_{td} \wedge I_{rx} \wedge I_{rd} \quad (3.1)$$

Each sub-invariant is proved individually on its respective part of the NIC model, then they are composed together in order to form the final one in Equation 3.1.

³Records in HOL4 are analogous to structures in imperative languages.

The majority of the lemmas are phrased in the same way, that we shall see in Chapter 5 correspond to the equation of Hoare Triples:

$$I \text{ nic} \wedge \text{nic}' = \text{transition nic} \implies I' \text{ nic}' \quad (3.2)$$

3.3.1 Visualizing proof dependencies

The model of the NIC consists of 1500 lines of HOL4 code, and required around three man-months of work. The NIC invariant consists of 650 lines of HOL4 code and the proof consists of approximately 55 000 lines of HOL4 code, including comments. Identifying the invariant and implementing the proof in HOL4 required around one man-year of work [24].

The proof being consequent, it is composed of a multitude of lemmas spread out in a multiple layers. The ideal objective of this work being to prove again the base lemmas on a different but equivalent BIR program, we therefore need to identify the relevant lemmas to be replaced. The proof is divided in two major parts, reception and transmission, and lemmas are scattered across a few files, each file being for a different part or abstraction level. Optimally, identifying the fringe of the dependency graph of the lemmas and theorems would give exactly the set of lemmas that all the other lemmas and theorems rely on, and therefore the smallest set of lemmas that are enough to prove in order to imply the security properties already proved in [24] by using the existing proofs.

As an attempt to visualize this fringe, we developed a tool, called DepGraph⁴, that can extract the dependency structure of HOL4 proofs in the form of a dependency graph. DepGraph can extract dependencies between HOL theories, i.e. compiled **SML** files containing proofs of lemmas and theorems, and between definitions, theorems and lemmas. Figure 3.2 shows the dependencies between theories about the transmission part of the NIC proof, and Figure 3.3 shows the dependencies between definitions, theorems and lemmas of the same part of the proof. However, as can be seen on Figures 3.2 and 3.3, this tool presents some critical shortcomings:

- The theory dependencies exporter uses files generated by Holmake, HOL4 compile system, in order to get the dependencies between theories. However, those files don't really represent dependencies but the files to be loaded before this script can be loaded, in a recursive fashion. Therefore, they represent the transitive reduction of the dependency graph. Figure 3.1 presents transitive reduction. Because of this fact, precious knowledge is lost and cannot be recovered by using this method: edges representing direct dependencies can be removed if the remaining edges still account for this dependency. Therefore, we are still able to

⁴DepGraph's design is presented in Appendix B.

tell which nodes depend on some node n , but we cannot identify the aforementioned fringe. In order to solve this problem, different approaches exist, such as implementing simplified a SML parser that looks only at dependencies, or injecting code inside the dependency resolution of an existing SML compiler. However, this would involve too much work that isn't the direct focus of this thesis.

- The definition, theorem and lemma dependencies exporter uses word-based heuristics in order to extract dependencies, and is as such not quite reliable and cannot give any guarantee. As above, there exist similar solutions in order to get multiple levels of guarantees, such that implementing a SML parser or injecting code inside HOL4 theory and definitions handling, but this would also require too much work. Destructuring HOL4 theories doesn't work because of how HOL4 has been designed, i.e. no trace is kept on disk about how to prove a given saved theorem, except of course with the corresponding script file. Moreover, such dependency graphs become quickly big, making them unusable in practice, and some additional work would be needed in order to represent them in a convenient way. Therefore, as above, no further work has been put into this exporter.

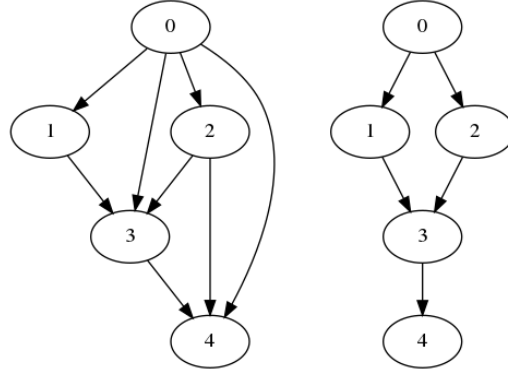


Figure 3.1: The right-hand side graph represent the transitive reduction of the left-hand side graph. Edges $(0 \rightarrow 3)$, $(0 \rightarrow 4)$ and $(2 \rightarrow 4)$ have been deleted because the transitive dependency relationship is preserved in the remaining edges. For example, the dependency $(0 \rightarrow 3)$ is represented in paths $(0 \rightarrow 1 \rightarrow 3)$ and $(0 \rightarrow 2 \rightarrow 3)$.



Figure 3.2: Dependency graph between theories about the transmission part of the formal NIC proof. An edge from a node A to a node B means that A depends on B . On this figure, edges are directed from left to right. Therefore, on this figure, higher-level files are on the left and lower-level ones on the right. Yellow nodes represent theories that depend on $txTheory$, i.e. the theory describing the transmission automaton presented in Section 3.2. Red nodes represent the leafs of the dependency graph, i.e. files without dependencies. Each node's edges are of different colors.

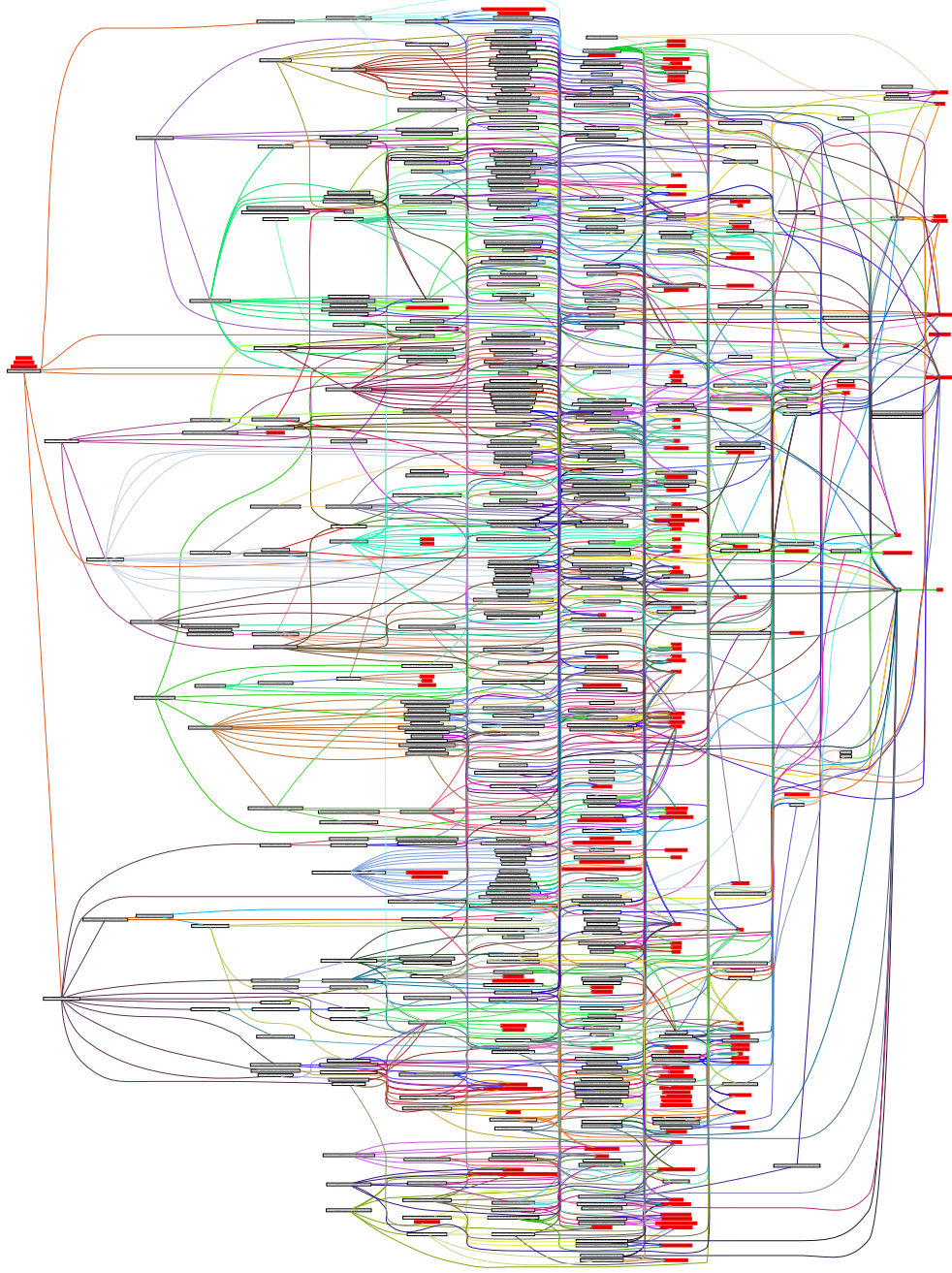


Figure 3.3: Dependency graph between definitions, theorems and lemmas of the transmission part of the formal NIC proof. An edge from a node A to a node B means that A depends on B . Therefore, on this figure, higher-level files are on the left and lower-level ones on the right. Each node's edges are of different colors. Red nodes represent leaf nodes, i.e. theorems and definitions without dependencies. The red nodes on the left-hand side of the Figure are nodes for which the heuristic failed to capture the dependencies.

Chapter 4

The BIR NIC model

This chapter intends to discuss about the different ways of implementing an equivalent BIR model from the formal NIC model. After an introduction of BIR, it will present different possible ways of doing this translation and present three approaches. Then, the final BIR model will be introduced along with the new tools that have been implemented in order to build it.

4.1 HolBA’s Binary Intermediate Representation (BIR)

HolBA’s Binary Intermediate Representation (BIR) [8], introduced in the Introduction, is a machine independent binary representation. It aims to be the simplest possible while still being able to represent all possible binary programs but self-modifying programs. It does so by having a limited syntax—introduced in Table 4.1—and forbidding implicit side-effects. A statement can only have explicit state changes and can only affect one variable.

This representation allows to produce proofs more easily than with classical binary representations, whose design are focused on execution speed rather than offline analysis. Moreover, BIR is fully specified and doesn’t have unspecified behaviour.

BIR is implemented as a set of HOL4 *datatypes*, and possesses a completely defined semantic. Section 5.5 contains a more thorough discussion of the BIR semantic. Section 4.2.3 implements a toy BIR program and presents the concrete BIR syntax.

Among its supporting tools, HolBA features a tool to visualize the **Control Flow Graph (CFG)** of BIR programs.

$$\begin{aligned}
prog &:= block^* \\
block &:= (string \mid integer, stmt^*, estmt) \\
stmt &:= \mathbf{assign} (string, exp) \mid \mathbf{assert} (exp) \\
estmt &:= \mathbf{jmp} (exp) \mid \mathbf{cjmp} (exp, exp, exp) \mid \mathbf{halt} (exp) \\
exp &:= integer \mid \mathbf{var} string \\
&\quad \mid \mathbf{if-then-else} (exp, exp, exp) \\
&\quad \mid \diamond_u exp \mid exp \diamond_b exp \\
&\quad \mid \mathbf{load} (exp, exp, \tau) \mid \mathbf{store} (exp, exp, exp, \tau)
\end{aligned}$$

Table 4.1: BIR’s syntax. Valid BIR programs must be well-typed. *integers* represent bounded N-bit integers. \diamond_u and \diamond_b represent respectively unary and binary operators. BIR blocks are tuples, with the first element being its label, the second a list of statements and the third the end statement. BIR syntax contains some other statements that won’t be used in this work and that have been omitted. For more information, see [8].

4.2 Exploring multiple approaches of translating the formal NIC model

Multiple approaches of translation of the formal NIC model described in Section 3.2 to an equivalent BIR program have been considered:

- **handwritten BIR program:** The most straightforward way would be to directly write the BIR program by hand by looking at the actual formal model and directly converting it to BIR. This is easily feasible because of the limited set of statement used in the implementation of the NIC model which are all representable in the BIR syntax. However, this does not seem optimal if more than one such translation need to be performed. One remedy of this problem would be to implement a set of tools that facilitate the implementation of device models and reduce the boiler plate. This idea is explored in the following Sections 4.2.3 to 4.2.5.
- **lifted C program:** An alternative would be to implement the NIC model in a more convenient higher-level programming language, compile it to assembly code and leverage the existing lifter to generate a BIR program. One advantage of this method is the possibility to use existing frameworks and tools that already exist in other languages, in which the ease of development is higher. However, the gained ease of development is maybe not worth the complexity introduced

by those steps. The following Section 4.2.2 experiments this idea.

- **device model specific IR:** This idea rejoins the previous one in the sense that it uses an additional intermediate representation (C in the previous idea) to implement the model, then uses some compiler to generate the BIR program (standard C compiler and HolBA's lifter). This idea would have the advantage that it introduces a new higher-level IR that would be common to possible later other device modelisations. Tools could also be developed for this IR. However, this approach would require to develop a new IR and, if we later want to develop a **proof-producing** lifter (compiler) the implementation would be costly, especially since this language would be more complex. The following Section 4.2.1 tries to use visual flowchart language as IR.

4.2.1 Using flowcharts as Intermediate Representation

As already discussed in the previous Section, the structure of the formal model looks like a tree. Therefore, using flowcharts can be a convenient way to represent such structures. An attempt has been made to design a flowchart representation. Figures 4.1, 4.2 and 4.3 show respectively a preview of the scheduler, transmission automaton and particular transition of this automaton.

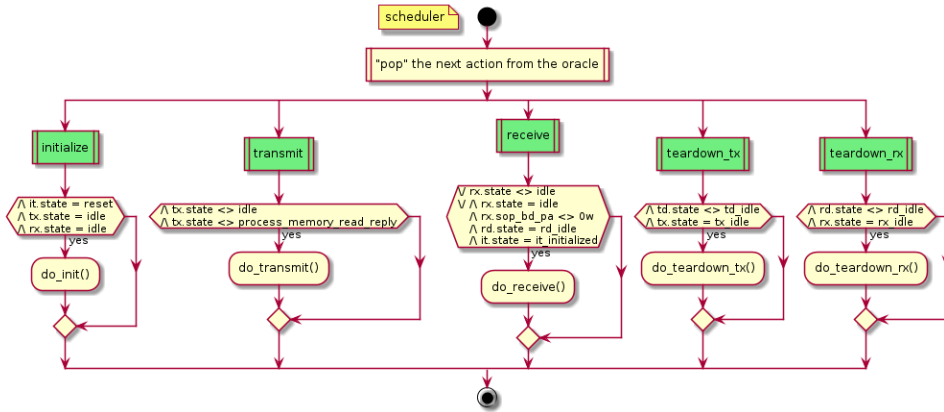


Figure 4.1: Flowchart of the scheduler of the NIC model. Green nodes represent condition statements, here they represent the value of the popped action from the oracle. The full dot represent the entry point, and the other point the exit point.

While this visual representation was useful to get to know the formal NIC model, we encountered several shortcomings:

- Flowcharts of each transition rapidly grew in size with the complexity of its formal counterpart. Possible orkarounds include the use of nested diagrams, as it is the case of Figure 4.3 representing one node of Figure 4.2, namely

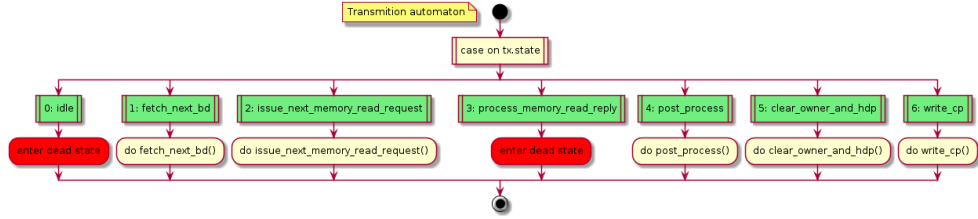


Figure 4.2: Flowchart of the transmission automaton of the NIC model. Green nodes are similar to the ones of Figure 4.1. Red nodes represent non autonomous transitions leading to dead states.

`fetch_next_bd`, or usage of shorter ways of representing common patterns, as it is the case for representing dead transitions on Figure 4.3.

- It is hard to define a coherent visual language able to represent the full set of features needed in order to realize device models. Additionally, this language must be compatible or easily translatable to BIR.
- It is hard to design a textual representation of this visual language other than conventional programming languages, so using such representation would require a substantial implementation effort in order to implement all the tools needed to use it. Developing visualization tools for a conventional language appears to be a more reasonable approach than developing a visual language.

For those reasons, it has been decided to not go further with this visual representation, and to focus instead of existing tools of the HolBA platform. However, as we shall see in the rest of this thesis, we will explore other visualization tools, and flowcharts will be used as a visual help while designing the model using other methods.

4.2.2 Writing the model in C

One-to-one translation of the transmission automaton, scheduler, state and *CPPI_RAM* related definitions of the formal model has been realized in C. This translation has been quite easy to perform and has been completed in roughly one day, and no difficulty arose during this translation. Regarding the NIC state, C has all types needed to represent it. An array has been used to represent the function type of *CPPI_RAM*. Enumerations have been used in order to represent HOL4 enumerations. While not optimal because of their non-restricted usage allowed by the C language—they can be freely modified as integers without checking their validity in C, while HOL4 and functional languages in general are restrictive and force to use correct values—enumerations are enough if used correctly while being well-aware of those shortcomings.

When studying the compiled assembly code and lifted BIR program, we noticed that

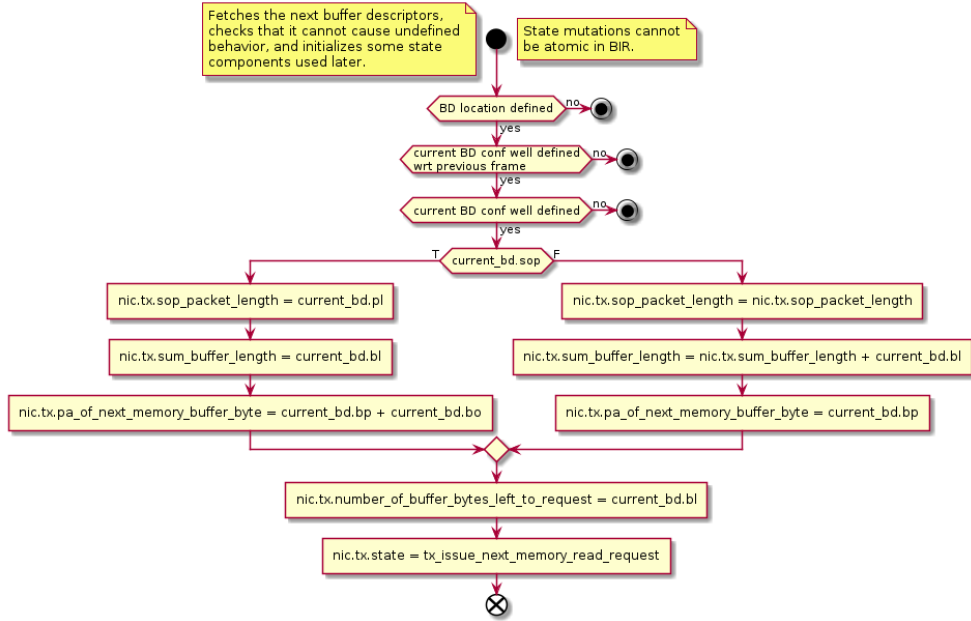


Figure 4.3: Flowchart of the `fetch_next_bd` transition of the transmission automaton of the NIC model. The full dot represent the entry point, \otimes represents the exit point and the other dots are shorthands to represent dead transitions (the symbols have been changed because of technical limitations of the tool used to draw the diagram).

all the convenient naming that we can use in the formal or the C model is lost and replaced with abundant usage of the stack. While this is completely normal behaviour for a C compiler, this is not convenient when performing later proofs on the model: we would first need to rename them in the proof using some definitions, and this process would be lengthy and cumbersome (if not automated), resulting in more code than if the model had been directly written in BIR in the first place. This experiment made us realize that writing the model is a rapid operation and that we should rather focus on making the verification step as smooth as possible, because it is the most difficult one to perform.

For those reasons, it has been decided to try using BIR directly to write the NIC model. The following Section 4.2.3 presents a first prototype that has been realized in order to visualize the shape of a BIR program representing the NIC model, and identify the tools that would be needed in order to reduce the boilerplate of such implementation in BIR.

4.2.3 Implementing a toy BIR model

Before writing the whole NIC model by hand, we shall identify the structure of the model and develop tools that facilitate its implementation. Using well-designed tools can reduce the boilerplate work of the implementation, helping to focus only on the meaningful content of the implementation, and can also reduce the chance of introducing bugs as the code is factored and mechanically shorter.

It has been decided to implement a transition system similar to the one of the NIC model, with two inner independent automata: Alice and Bob. The two automata feature a simple linear transition system, and each of them have a non-autonomous transition, that are performed by two external functions `bus_arrived` and `taxi_arrived` that represent memory accesses from the CPU in the NIC model. The function `autonomous_step` performs one autonomous transition via the scheduler. Figure 4.4 gives an overview of the two automata, Figure 4.6 shows the role of the three functions and Figure 4.7 describes in details Alice's automaton. (Flowcharts from the previous Section 4.2.1 have been used in this section in order to facilitate designing the toy model. However, no effort has been made to come up with a coherent and formalized language.)

From the design presented in Figures 4.4, 4.6 and 4.7, writing the BIR program is a repetitive but straightforward step. The resulting BIR program is 450 lines of code long. The following issues have been identified:

- BIR, as a HOL4 embedded language, is very verbose. Simple operations like additions or assignments require long constructions, as shown in Listing 4.1. Section 4.2.4 presents BSL, a less verbose way of writing BIR programs.

Listing 4.1: BIR

```
(* A1. Eating *)
<|bb_label := BL_Label "alice_automaton.step.1";
  bb_statements := [
    (* state.alice.hunger -= 20 *)
    BStmt_Assign (BVar "state.alice.hunger" (BType_Imm Bit32))
      (BExp_BinExp BExp_Minus
        (BExp_Den (BVar "state.alice.hunger" (BType_Imm Bit32)))
        (BExp_Const (Imm32 20w)));
    (* state.alice.step := 2 *)
    BStmt_Assign (BVar "state.alice.step" (BType_Imm Bit32))
      (BExp_Const (Imm32 2w))
  ];
  bb_last_statement := BStmt_Jmp (BLE_Label (BL_Label "alice_automaton.end"))
|>
```

- BIR features only one conditional statement controlling the control flow of a program: conditional jumps. Hence, BIR is not convenient for representing **if-then-else** statements with more than two branches. Section 4.2.5 presents some helper functions that have been implemented in order to be able to abstract the raw BIR code and work at a higher-level.

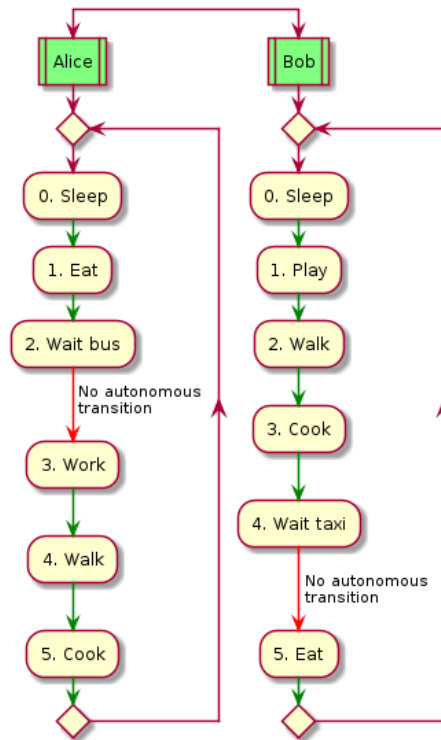


Figure 4.4: Overview of the two independent automata. The green arrow represent the autonomous transitions and the red ones the non-autonomous. As in Section 4.2.1, the green nodes represent conditional statements depending on the scheduler's oracle, as shown in Figure 4.5.

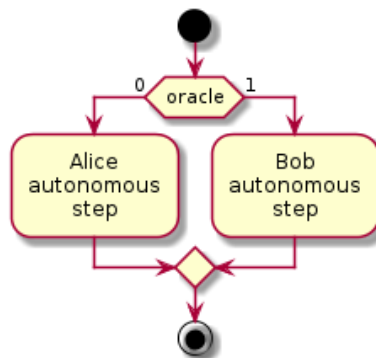


Figure 4.5: Scheduler of the toy BIR model. We can see how the oracle decides which automaton takes a step. If an automaton is in a state whose transition isn't autonomous, the automaton state is returned unchanged.

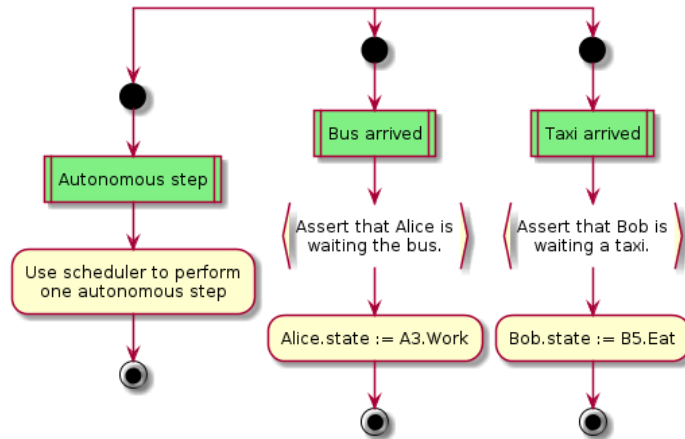


Figure 4.6: The three entrypoints of the toy program: `autonomous_step`, `bus_arrived` and `taxi_arrived`.

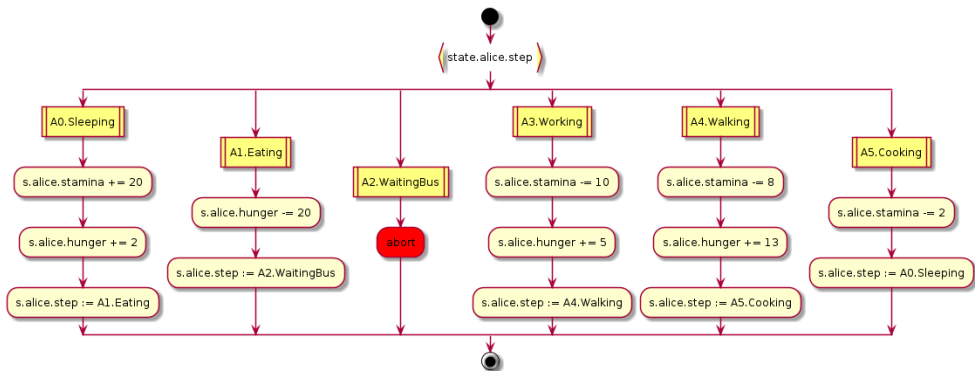


Figure 4.7: The six transitions of Alice's automaton. Each transition only updates the state. Conditional statements hasn't been added in this toy model because it already features conditional statements in the scheduler. The third transition is a non-autonomous transition; `bus_arrived` is in charge of updating Alice's state from *WaitingBus* to *Working*.

Listing 4.2: Example of prelude of a SML script that defines functions tailor-made to its use. The function *o* is used for function combination.

```
val bvarstate = bvarimm32
val bdenstate = (bden o bvarstate)
val bstateval = bconst32
val bjmplabel_str = (bjmp o belabel_str)
```

4.2.4 BSL: BIR Simple Language

Because of its extreme verbosity, BIR is not a language convenient to manually use. Its primary use has been as an output from the **proof-producing** lifter and as a machine-independent intermediate represent with a very limited feature-set convenient to reason about. A less verbose and more convenient method for producing BIR code is needed in order to open this language to other uses. Therefore, a library has been written that offers the same expressiveness than BIR but with shorter constructs. This library has been kept simple and will serve as the base layer of possible later abstractions. As such, it has been decided that no feature other than pure syntactic construct, such that type inference, would be included. This library has been named BIR Simple Language (BSL) and is implemented in the file *bslSyntax.sml*¹.

BSL is composed of a set of functions with short names and a coherent interface. Since this library will be directly included in every SML script file in order to avoid usage of the lengthy `bslSyntax.function_name` SML construction, we decided to prefix every function with the letter `b`. As BSL doesn't feature type inference, types must be explicitly given when necessary. For every such function, such as `bconst`, implementations for every possible types have been added, in addition to a general functions taking the size as argument. Thanks to partial function application in SML, composing functions is easy and scripts can create the set of functions that they need. However, in practice, a BIR program uses a very limited set of types and it is often enough in scripts to statically define the needed set of sized functions (as shown in Listing 4.2). Finally, BSL has been designed to be fully interoperable with HOL4 terms; it exists for every BIR operation a related BSL function taking a HOL4 term as parameter, and every BSL function return a HOL4 term. Scripts can therefore not be limited by the restricted feature-set of BSL.

¹This file has been named in order to follow the HOL4 conventions: **Syntax.sml* files contain functions that create, destruct and check HOL4 terms. *bslSyntax.sml* is a library for creating only BIR terms, and should therefore follow this convention.

Listing 4.3: Call to the `gen_state_helpers` function that generates a record containing the list of the states, the list of the autonomous states, a function from state name to state id, a function from state name to boolean telling if the state is autonomous and a function that generates BIR code of the state id from the state name. `gen_state_helpers` takes as parameters the name of the automaton and a list of its states with an id and a boolean telling if the state is autonomous.

```
val tx_state = gen_state_helpers "tx" [
  ("tx1_idle", (1, false)),
  ("tx2_fetch_next_bd", (2, true)),
  ("tx3_issue_next_memory_read_request", (3, true)),
  ("tx4_process_memory_read_reply", (4, false)),
  ("tx5_post_process", (5, true)),
  ("tx6_clear_owner_and_hdp", (6, true)),
  ("tx7_write_cp", (7, true))
]
```

4.2.5 Implementing the real model

From knowledge gained in the previous section 4.2.3, a set of helper functions has been implemented, mainly in order to facilitate reasoning about the state machine. Listings 4.3 and 4.4 show the SML code respectively used to represent the transmission automaton and to generate the BIR blocks that jumps to the right label of the program depending on the current step of the automaton (analogous to a *case* or *switch* construct in higher-level languages).

Because of time constraints, not every transitions of the formal NIC model have been implemented in this new model. Instead, we decided to focus on only some transitions in two automata: initialization and transmission. This focus is pragmatic for two reasons: (a) we decided to write the model along with the proof, only needed transitions for the current proof, in order to grow the model at a reasonable pace and to no clutter it with unneeded features or at the contrary missing critical aspects during the first sketch (b) as the verification was performed at the same time, we decided to start with easier, but not obvious, transitions, hence the choice to postpone verification of the more complex reception automaton to later in the work.

When visualizing the CFG of the BIR NIC model (Figure 4.8), we notice that it has the shape of a tree, rejoining the idea introduced in Section 3.2: each node of the tree is a BIR block, each edge a jump. Since every jump in the NIC model go to a static target, a jump can only have one or two destinations. Hence, nodes can only have one or two outgoing edges. In the model, each node represent either an empty jump block, possibly conditional, or a mutation of the state.

The static and loop-free control-flow of this BIR program is very suitable for use with HolBA's weakest precondition tool in order to perform contract based verification, as

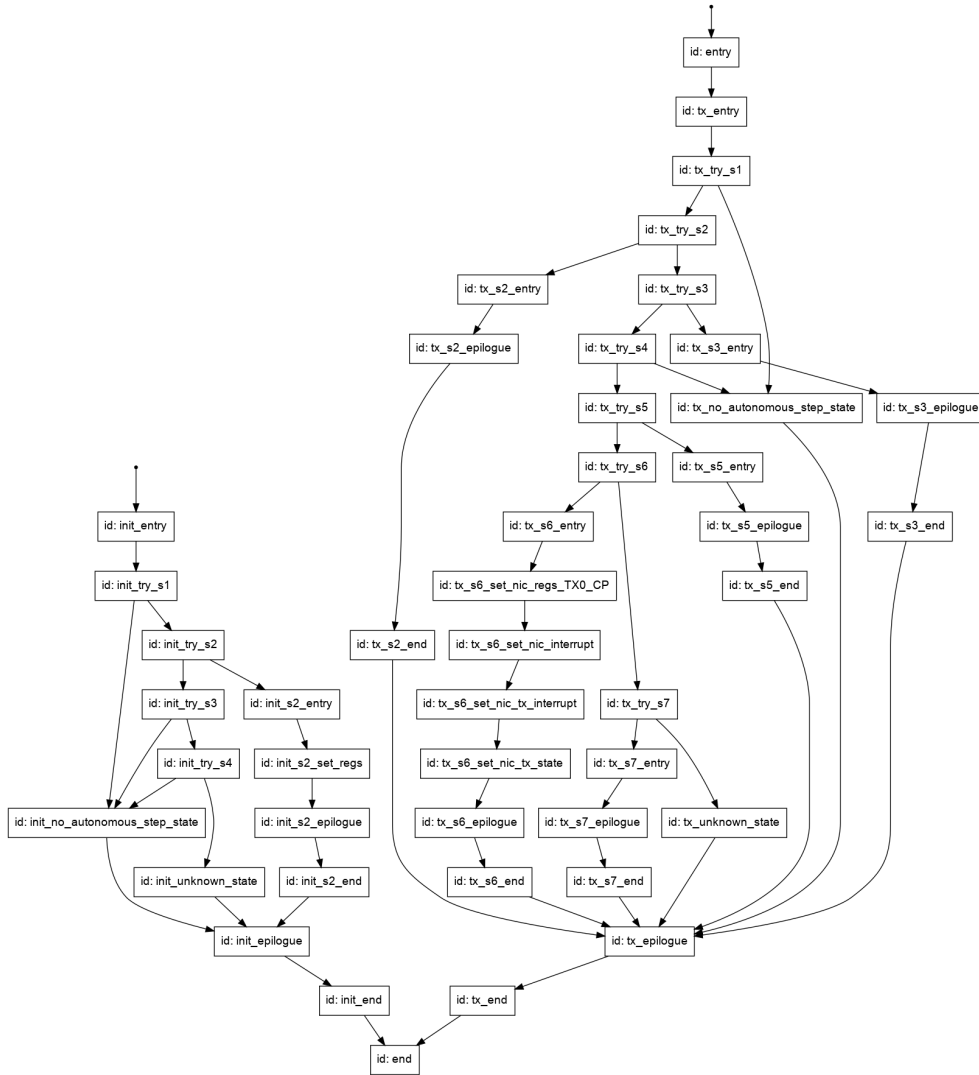


Figure 4.8: **CFG** of the implemented transitions of the NIC BIR model. Each node represents one BIR block and each edge one jump target. Because the scheduler hasn't been implemented in this model, the **entry** block is jumping arbitrarily on *tx_entry*; all contract verification procedures take labels in order to define entry and end points of the section of a program that the tools should generate the weakest precondition from, so this isn't a problem in practice as long as no attempt to perform proofs on more than one automaton are made. The CFG of the full program would resemble a tree with one node at the top, *entry*, and the tree getting larger with the depth because of all the conditional jumps that go to the transition being performed.

Listing 4.4: Autonomous transition jump table used for the transmission automaton. This table uses the names of each state and transition to generate the BIR blocks that will perform the conditional jumps depending on the value of the current transmission state. The parameters of the function are: a tuple of the name of the variable containing the current state, the label of the block to jump to when the current state isn't handled and a function that gives the state number from state name, and a list of jumps, each jump being a tuple of the label of the conditional BIR block, the value of the state that leads to this block and the name of the label to go to.

```
(* Autonomous transition jump *)
@ bstate_cases ("nic_tx_state", "tx_unknown_state", bstateval_tx) [
  ("tx_try_s1", "tx1_idle", "tx_no_autonomous_step_state"),
  ("tx_try_s2", "tx2_fetch_next_bd", "tx_s2_entry"),
  ("tx_try_s3", "tx3_issue_next_memory_read_request", "tx_s3_entry"),
  ("tx_try_s4", "tx4_process_memory_read_reply", "tx_no_autonomous_step_state"),
  ("tx_try_s5", "tx5_post_process", "tx_s5_entry"),
  ("tx_try_s6", "tx6_clear_owner_and_hdp", "tx_s6_entry"),
  ("tx_try_s7", "tx7_write_cp", "tx_s7_entry")
]
```

we shall see in the next chapter.

Chapter 5

Contract based verification

This chapter will present contract based verification, explaining the underlying theory: Hoare Triples, weakest precondition derivation and the use of SMT solvers. It will also present the current status of contract based verification in HolBA, and present some of the work needed in order to reason about BIR memories with SMT solvers.

5.1 Hoare triples

Contract based verification is a powerful approach for verifying programs. For a given program *prog* consisting of a list of instructions and two predicates *P* and *Q* called respectively pre- and postcondition, a Hoare triple $\{P\} \text{ prog } \{Q\}$ states that when executing the program *prog* from a state *S* terminates in a state *S'*, if *P* holds in *S* then *Q* will hold in *S'* (Equation 5.1). Hereafter, we assume programs and states to be well-typed.

$$\{P\} \text{ prog } \{Q\} \triangleq S' = \text{exec}(S, \text{prog}) \implies P(S) \implies Q(S') \quad (5.1)$$

For example, $\{P\} \emptyset \{P\}$ holds because an empty program doesn't change the state of the execution. $\{n = 1\} n := n + 1 \{ \text{even}(n) \}$, with $n \in \mathbb{N}$, holds because $1 + 1 = 2$, which is even.

In order to perform the verification, the Hoare logic introduces a set of axioms describing the effect of each instruction of a given language over the execution state [25]. For an assignment $x := f$ where *x* is a variable identifier and *f* an expression without side-effects, Equation 5.2 defines the axiom of assignment, where $P[f/x]$ denotes the substitution of all occurrences of *x* by *f* in *P*.

$$\{P[f/x]\} x := f \{P\} \quad (5.2)$$

5.2 Weakest precondition derivation

While Hoare logic introduces sufficient preconditions, Dijkstra introduced the concept of necessary and sufficient preconditions, called “weakest” preconditions. Such weakest preconditions can be automatically derived from a program $prog$ and a postcondition Q . Let’s call $WP(prog, Q)$ such a weakest precondition. Then, from Equation 5.1 follows:

$$\forall(prog, Q), \{ WP(prog, Q) \} prog \{ Q \} \quad (5.3)$$

For the program $n := n + 1$ mentioned in the previous section, we can generate the weakest precondition for the postcondition $even(n)$. First, we can rewrite $even(n)$ as $n \text{ MOD } 2 = 0$ with MOD denoting the arithmetic modulo. Then, we derive the weakest precondition of the statement $n := n + 1$ by transforming the predicate $n \text{ MOD } 2 = 0$ by substituting all occurrences of n by $n + 1$:

$$WP(“n := n + 1”, n \text{ MOD } 2 = 0) = (n + 1 \text{ MOD } 2 = 0) \quad (5.4)$$

From the properties of the modulo, we can simplify $n+1 \text{ MOD } 2 = 0$ to $n \text{ MOD } 2 = 1$ or $odd(n)$. Therefore, $\{ odd(n) \} n := n + 1 \{ even(n) \}$, i.e. incrementing the value of an odd integer variable by one makes it even.

While the triple $\{ n = 1 \} n := n + 1 \{ even(n) \}$ uses a sufficient precondition for establishing its postcondition, the triple $\{ odd(n) \} n := n + 1 \{ even(n) \}$ uses the weakest precondition. The later being the weakest precondition of the former, the two contracts are in relation:

$$n = 1 \implies odd(n) \quad (5.5)$$

More generally, for a triple $\{ P \} prog \{ Q \}$ to hold, P must be stronger than the weakest precondition, i.e. we need to prove that $P \implies WP(prog, Q)$:

TODO: Here and above, should we mention/explain termination?

$$(P \implies WP(prog, Q)) \implies \{ P \} prog \{ Q \} \quad (5.6)$$

5.3 Using SMT solvers to prove contracts

From Equation 5.6 we see that, in order to prove that a triple $\{ P \} prog \{ Q \}$, we need to prove $P \implies WP(prog, Q)$. While multiple methods exist to perform such proofs, **SMT** solvers offer a convenient and automatic solution.

Satisfiability Modulo Theories (SMT) problem is a decision problem for logical formulas with respect to combinations of background theories such as arithmetic, bit-vectors, arrays, and uninterpreted functions [26]. **Satisfiability Modulo Theories (SMT)** problem is a generalization of **Boolean SATisfiability Problem (SAT)** problem supporting more theories. When given a formula, a **SMT** solver decides if the formula is satisfiable, i.e. if there exist a valuation of its variables where the formula evaluates to true. As a **SMT** solver can fail to decide a given instance, there are three possible outputs: “satisfiable”, “unsatisfiable” and “unknown”. Another useful feature of some **SMT** solvers is the ability to ask for a satisfying model, which represents a counter-example of a false predicate.

A predicate P holds if it evaluates to true for all possible values of its variables. Alternatively, the negation of a predicate $\neg P$ holds if there exist no valuation of its variables where the predicate evaluates to true, i.e. if the instance is unsatisfiable. Therefore, if a **SMT** solver report that $\neg P$ is “unsatisfiable”, then P holds.

Another way of thinking about how to prove logical formulas with **SMT** solvers is by using De Morgan’s Laws: we know that $\neg(P \implies WP) \equiv (P \wedge \neg WP)$. Therefore, proving that $\neg(P \implies WP)$ is “unsatisfiable” using an **SMT** solver can be seen as proving that there exist no model where P holds and WP doesn’t.

5.3.1 The BitVector theory

To reason about fixed-size integers, **SMT** solvers often implement a “BitVector”, or “FixedSizeBitVectors”, theory. In order to understand its particularities, we can try to prove Equation 5.7. Hereafter, we will use Z3, a popular and efficient **SMT** solver implemented by Microsoft Research¹, and SMT-LIB 2.0, which is a standard format for **SMT** solvers [27]. Listing 5.1 shows the SMT-LIB 2.0 representation of this proof attempt.

$$\forall x. x + 1 > x, \text{ with } x \text{ an unsigned 32-bit integer} \quad (5.7)$$

Listing 5.1: SMT-LIB 2.0 representation of Equation 5.7.

```
(declare-const x (_ BitVec 32))
(assert (not
  (bvugt (bvadd x (_ bv1 32)) x)))
(check-sat)
(get-model)
```

When given Listing 5.7 as input, Z3 gives the following output:

¹Z3 is available on GitHub at <https://github.com/Z3Prover/z3/>.

Listing 5.2: Z3 output for Listing 5.1.

```
sat
(model (define-fun x () (_ BitVec 32) #xffffffff))
```

Z3 is telling us that Equation 5.7 is false, and gives a counterexample: $x = 2^{32} - 1$. Indeed, with this value of x , $x + 1$ wraps around and result in 0 which is smaller than $2^{32} - 1$. This behaviour is due to the bounded nature of fixed-size integers. The correct equation here would be:

$$\forall x. x \neq 2^{32} - 1 \implies x + 1 > x, \text{ with } x \text{ an unsigned 32-bit integer} \quad (5.8)$$

Listing 5.3 and 5.4 show the input and output of Z3 used to successfully prove Equation 5.8.

Listing 5.3: SMT-LIB 2.0 representation of Equation 5.8.

```
(declare-const x (_ BitVec 32))
(assert (not
  (bvugt (bvadd x (_ bv1 32)) x)))
(assert (not (= x #xffffffff)))
(check-sat)
```

Listing 5.4: Z3 output for Listing 5.3.

```
unsat
```

5.4 Contract based verification in HolBA

HolBA provides a **proof-producing** tool for automatically deriving weakest preconditions on loop-free **BIR** programs whose control flow can be statically identified [8]. This tool is proof-producing in that it proves Theorem 5.9 which is the instantiation of Definition 5.10, with $(p, \text{entry_l}, \text{end_ls})$ defining the program, wp the derived weakest precondition, $post$ the given postcondition.

$$\text{bir_exec_to_labels_triple } prog \text{ entry_l end_ls } \mathbf{wp} \text{ post} \quad (5.9)$$

$$\begin{aligned}
& \vdash \forall (prog : \alpha \text{ bir_program_t}) (entry_l : \text{bir_label_t}) (end_ls : \text{bir_label_t} \rightarrow \text{bool}) \\
& \quad (pre : \text{bir_exp_t}) (post : \text{bir_exp_t}). \\
& \quad \text{bir_exec_to_labels_triple } prog \text{ entry_l } end_ls \text{ pre } post \Leftrightarrow \\
& \quad \forall (s : \text{bir_state_t}) (r : \alpha \text{ bir_execution_result_t}). \\
& \quad \quad \text{bir_env_vars_are_initialised } s.\text{bst_environ } (\text{bir_vars_of_program } prog) \\
& \quad \Rightarrow s.\text{bst_pc.bpc_index} = 0 \wedge s.\text{bst_pc.bpc_label} = \text{entry_l} \\
& \quad \Rightarrow s.\text{bst_status} = \text{BST_Running} \\
& \quad \Rightarrow \text{bir_is_bool_exp_env } s.\text{bst_environ } pre \\
& \quad \Rightarrow \text{bir_eval_exp } pre \text{ } s.\text{bst_environ} = \text{bir_val_true} \\
& \quad \Rightarrow \text{bir_exec_to_labels } end_ls \text{ } prog \text{ } s = r \\
& \quad \Rightarrow \exists (obs : \alpha \text{ list}) (step_count : \text{num}) (pc_count : \text{num}) (s' : \text{bir_state_t}). \\
& \quad \quad r = \text{BER_Ended } obs \text{ step_count } pc_count \text{ } s' \\
& \quad \quad \wedge s'.\text{bst_status} = \text{BST_Running} \\
& \quad \quad \wedge \text{bir_is_bool_exp_env } s'.\text{bst_environ } post \\
& \quad \quad \wedge \text{bir_eval_exp } post \text{ } s'.\text{bst_environ} = \text{bir_val_true} \\
& \quad \quad \wedge s'.\text{bst_pc.bpc_index} = 0 \wedge s'.\text{bst_pc.bpc_label} \in end_ls
\end{aligned} \tag{5.10}$$

It is to be noted that this tool doesn't produce a theorem stating that the generated expression is actually *the* weakest precondition. However, this theorem isn't needed to perform contract-based verification if the generated “weakest” precondition is weak enough so that our precondition can imply it. However, without this theorem it is impossible to prove that a given precondition P isn't strong enough to establish the postcondition. We will still use the term “weakest precondition” as it is in practice how we are using this tool.

Definitions 5.10 introduces additional conditions about well-typedness and initialization that are needed in BIR today², as well as the notion of “Block Program Counter” for multi-statement blocks.

This tool doesn't provide a simple interface to compute weakest preconditions for a given program and postcondition, nor does it provide and support for proving the relation between the precondition and the weakest precondition. Then, in order to prove that the Hoare triple holds from this generated Theorem 5.9, we need to prove:

$$\text{bir_exec_to_labels_triple } p \text{ entry_l } end_ls \text{ pre } post \tag{5.11}$$

²Removal of the need of initialization is being discussed at the time of the writing, because actual hardware registers and memories are in facts always initialized: <https://github.com/kth-step/HolBA/issues/63>

Assuming well-typedness and initialization, after rewriting the definition of *bir_exec_to_labels_triple*, we have to show *bir_eval_exp wp s.bst_environ = bir_val_true* in order to prove our goal using the *modus ponens* with Theorem 5.9. This correspond to proving the following implication:

$$\begin{aligned} & \text{bir_eval_exp pre } s.\text{bst_environ} = \text{bir_val_true} \\ \implies & \text{bir_eval_exp wp } s.\text{bst_environ} = \text{bir_val_true} \end{aligned} \quad (5.12)$$

In Equation 5.12 we can recognize Equation 5.6 that we discussed how to prove using **SMT** solvers in Section 5.3. However, the expressions are expressed as BIR expressions. We then have to find a way to use an SMT solver. This is the focus of the following of this thesis. Section 6 will use a non **proof-producing** method for translating those BIR expressions into an equivalent formula that SMT solvers can work on, then focus on automating the whole verification process. Chapter 7 will complete this proof and use it to lift properties that have been proved on the BIR implementation to the **NIC** model. The following Section 5.5 will discuss how to make proofs about BIR memories using **SMT** solvers.

5.5 BIR memories and SMT solvers

TODO: Mention SMT-LIB logics?

TODO: Mention proof reconstruction?

HOL4 features a library for interfacing **SMT** solvers and HOL4, called *HolSmtLib*. This library supports Yices 1 and Z3 as external provers. Yices 1 being an abandoned project that doesn't support SMT-LIB 2.0, we will focus on Z3 and the standard format SMT-LIB 2.0 [27]. While *HolSmtLib* supports export for some SMT-LIB 2.0 theories, it doesn't support the *ArraysEx* theory and doesn't know about BIR. In Section 5.4, we discussed the translation from BIR expressions to *wordsTheory*. However, this theory doesn't contain anything about memories or arrays in general. Therefore, some modifications are needed.

BIR memories are semantically defined as functions from addresses to values.

There exist five types of BIR expressions operating directly on memories (cf. Listing 6.2 for the list of BIR expressions, and Section 6.1 for a more precise discussion of the BIR semantic):

- **BExp_Den**: this operation enables reading values from the environment. It is analogous to reading registers or the memory in assembly programs. This operation semantically equivalent to free variables that *HolSmtLib* already support.

- `BExp_MemEq`: this operation is the equality binary operation on BIR memories. This operation is semantically equivalent to equality between its operands. *HolSmtLib* already supports this operation.
- `BExp_Store`: this operation is used to represent memory writes. It is semantically defined as successive function updates of consecutive segments of the word being stored, because the length of the memory value-type can be less or equal than the length of values stored in the memory. A function update in *combinTheory* is defined with Definition 5.13. *HolSmtLib* cannot currently export function update operations.
- `BExp_Load`: this operation is used to read from memories. BIR memories are semantically defined as functions from addresses to values. A function application in *combinTheory* is defined with Definition 5.14. Then, a memory load operation is the concatenation of multiple function application of consecutive addresses. *HolSmtLib* supports function application of uninterpreted functions only.

$$\vdash \forall a b. a =+ b = (\lambda f c. \text{if } a = c \text{ then } b \text{ else } f c) \quad (5.13)$$

$$\vdash \forall x f. x := f = f x \quad (5.14)$$

We saw in the previous list that we need to implement the support for *combinTheory* function update and application in the *HolSmtLib* SMT-LIB 2.0 exporter. Since we only need two operations on the memory, load and store, the *ArraysEx* theory is a good fit.

SMT-LIB 2.0 [27] defines the *ArraysEx* theory using the three following axioms:

Listing 5.5: SMT-LIB 2.0 axioms of the *ArrayEx* theory.

```
(forall ((a (Array s1 s2)) (i s1) (e s2))
  (= (select (store a i e) i) e))

(forall ((a (Array s1 s2)) (i s1) (j s1) (e s2))
  (=> (distinct i j)
    (= (select (store a i e) j) (select a j))))

(forall ((a (Array s1 s2)) (b (Array s1 s2)))
  (=> (forall ((i s1)) (= (select a i) (select b i)))
    (= a b)))
```

If those axioms hold in *combinTheory* then the translation is sound. *combinTheory*'s *UPDATE_APPLY* theorem in Equation 5.15 is equivalent to the first two theorems,

the first and the second conjunct corresponding respectively to the first and the second axiom. The third axiom can be proved using both *combinTheory*'s *APP_def* theorem (Theorem 5.14) and *boolTheory*'s *EQ_EXT* theorem (Theorem 5.16)

$$\begin{aligned} \vdash \quad & \forall a \, x \, f. (a =+ x) \, f \, a = x \\ & \wedge \forall a \, b \, x \, f. a \neq b \implies ((a =+ x) \, f \, b = f \, b) \end{aligned} \quad (5.15)$$

$$\vdash \forall f \, g. (\forall x. f \, x = g \, x) \implies (f = g) \quad (5.16)$$

Since this translation is sound, it has been added in *HolSmtLib*. The translation is direct: from `:>` and `=+` to respectively `select` and `store`. Section 6.4.2 presents a test using BIR memories.

Chapter 6

Non proof-producing automatic contract verification

*In the previous section, we learned about contract verification and the current status of **HolBA**'s implementation. To perform verification on the **NIC** model, we would like to automate the process as much as possible. **HolBA** currently offers tools for automatic weakest precondition generation, therefore we need to close the gap between BIR expression and SMT solvers, as well as to implement a convenient interface on top.*

6.1 Exporting BIR expressions to SMT solvers

As an intermediate language for formal verification, **BIR** possesses a precise semantic. The semantic of BIR expressions is expressed as a set of definitions describing what are the equivalent operations using *wordsTheory* and *combinTheory*. These theories contains definitions and theorems about “words”, i.e. bounded N -bit integers that are used to reason about integer types in programming languages and hardware memory in general, and function application and update used for BIR memories as already discussed in Section 5.5. For example, the semantic of binary operators in BIR is defined with the following theorems¹:

$$\begin{aligned} \vdash \text{bir_bin_exp_GET_OPER } BIEp_And &= \text{words_and} \\ \wedge \text{bir_bin_exp_GET_OPER } BIEp_Or &= \text{words_or} \end{aligned} \quad (6.1)$$

$$\begin{aligned} \vdash \forall (\text{bin_op} : \text{bir_bin_exp_t}) (\text{w1} : \text{word64}) (\text{w2} : \text{word64}). \\ \text{bir_bin_exp } \text{bin_op} (\text{Imm64 } \text{w1}) (\text{Imm64 } \text{w2}) \\ = \text{Imm64 } (\text{bir_bin_exp_GET_OPER } \text{bin_op } \text{w1 } \text{w2}) \end{aligned} \quad (6.2)$$

¹Theorems 6.1 and 6.2 have been reduced to only two operators and well-typed 64-bit expressions.

Similarly, a set of definitions and theorems describe the semantic of operations on BIR memories. However, correct handling of endianness, alignment and genericity over the size of memories cells and addresses, these definitions and theorems are pretty complicated to work with. The same is true for the semantic of operations on BIR variables, because of well-typedness and initialization.

For this reason—i.e. writing **proof-producing** code is costly—, we decided to write a non-proof producing function `bir_exp_to_words` that translates BIR expressions to the equivalent words expression. The obvious downside of such a function is that we now have to trust the translation to be sound, because we no longer get any guarantee from the theorem prover. However, development time is dramatically decreased and offers more time for experimenting. Moreover, this function can later be implemented in a proof-producing way for more trustful verification. Then, in order to have a high confidence of correctness, software engineering practices apply:

- write small and understandable pieces of code and compose them, and
- write a comprehensive suite of tests.

BIR expressions are defined as a HOL4 algebraic data type in Listing 6.1. Hence, in order to translate BIR expressions to words expressions, we need to handle every variant. This has been done² using an exhaustive `if-then-else` statement³. The code is mostly destructuring HOL4 terms and creating new *wordsTheory* and *combinTheory* terms. In order to obtain an easily reviewable code, a balance between expressivity and conciseness has to be carefully decided. Table 6.1 shows that the length of each variant is relatively small in terms of lines of codes, from 1 line for constants to 51 for memory store expressions. This achieves the first point of the previous list.

Testing of `bir_exp_to_words` has been done using a set of $(bir_exp, expected)$ couples with increasing complexity, where `bir_exp_to_words` is used to translate each *bir_exp* and the result is compared to *expected*. Then, BIR expression being defined as an algebraic data type, nesting of BIR expressions follow naturally. This achieves the second point of the previous list.

Listing 6.1: `bir_exp_t` definition.

```
Datatype 'bir_exp_t =
  BExp_Const      bir_imm_t
| BExp_Den        bir_var_t
```

²Code available here: <https://github.com/kth-step/HolBA/commit/2fcc54dcb04a20716e7697f64b5a4578f8a8af9>

³Pattern matching would have been optimal, but isn't possible because of how HOL4 is embedded in SML.

```

| BExp_Cast
    bir_cast_t bir_exp_t bir_immtype_t

| BExp_UnaryExp    bir_unary_exp_t bir_exp_t
| BExp_BinExp      bir_bin_exp_t bir_exp_t bir_exp_t
| BExp_BinPred     bir_bin_pred_t bir_exp_t bir_exp_t
| BExp_MemEq       bir_exp_t bir_exp_t

| BExp_IfThenElse bir_exp_t bir_exp_t bir_exp_t

| BExp_Load        bir_exp_t bir_exp_t bir_endian_t
    ↪ bir_immtype_t
| BExp_Store       bir_exp_t bir_exp_t bir_endian_t
    ↪ bir_exp_t `

```

bir_exp_t variant	Lines of code
BExp_Const	1
BExp_Den	21
BExp_Cast	not implemented
BExp_UnaryExp	8
BExp_BinExp	9
BExp_BinPred	10
BExp_MemEq	10
BExp_IfThenElse	9
BExp_Load	48
BExp_Store	51

Table 6.1: Length of each `bir_exp_t` variant in the implementation of `bir_exp_to_words`.

6.2 Pretty-printing to visualize BIR expressions

When working with complex constructs, the need of visualization techniques often arise. Generated weakest preconditions grow quickly with the number of statements in a program, linearly or exponentially depending on the type of statements—control flow statements produce exponential growth. While clever techniques can be implemented to keep their size reasonable [8], we often need to read and analyze them.

Printing of BIR terms in general is very verbose. For example, the expression 6.3 with a 64-bit x integer defined using the BSL code in Listing 6.2 yields the printing in Figure 6.1 using HOL4’s default printing capabilities.

$$\text{if } (x \leq 100) \vee (y + 1 > 10) \vee (x + y \leq 20) \text{ then } 2 \times x \text{ else } 3 \times y + 1 \quad (6.3)$$

Listing 6.2: BSL code

```
bite (
  borl [
    ble ((bden o bvarimm64) "x", bconst64 100),
    bnot (ble (bplus ((bden o bvarimm64) "y",
                      bconst64 1),
                      bconst64 10)),
    ble (bplus ((bden o bvarimm64) "x",
                (bden o bvarimm64) "y"),
          bconst64 20)
  ],
  bmult ((bden o bvarimm64) "x", bconst64 2),
  bplus (bmult ((bden o bvarimm64) "x", bconst64 3),
          bconst64 1))
```

```
BExp_IfThenElse
(BExp_BinExp BIEExp_Or
 (BExp_BinExp BIEExp_Or
  (BExp_BinPred BIEExp_LessOrEqual
   (BExp_Den (BVar "x" (BType_Imm Bit64))) (BExp_Const (Imm64 100w)))
  (BExp_UnaryExp BIEExp_Not
   (BExp_BinPred BIEExp_LessOrEqual
    (BExp_BinExp BIEExp_Plus (BExp_Den (BVar "y" (BType_Imm Bit64)))
                             (BExp_Const (Imm64 1w))) (BExp_Const (Imm64 10w))))))
 (BExp_BinPred BIEExp_LessOrEqual
  (BExp_BinExp BIEExp_Plus (BExp_Den (BVar "x" (BType_Imm Bit64)))
                           (BExp_Den (BVar "y" (BType_Imm Bit64))) (BExp_Const (Imm64 20w))))
 (BExp_BinExp BIEExp_Mult (BExp_Den (BVar "x" (BType_Imm Bit64)))
                          (BExp_Const (Imm64 2w)))
 (BExp_BinExp BIEExp_Plus
  (BExp_BinExp BIEExp_Mult (BExp_Den (BVar "x" (BType_Imm Bit64)))
                           (BExp_Const (Imm64 3w))) (BExp_Const (Imm64 1w))))
```

Figure 6.1: Default HOL4 printing

This expression is relatively small and yet the printed term is 17 lines long. Compared to the BSL expression that is 8 lines long⁴, that is a two time increase in size. Moreover, lines are long and verbose: for example, a “less-than” binary expression is written as “BExp_BinPred BIEExp_LessOrEqual e1 e2”. Comparatively, the math expression “ $e1 \leq e2$ ” and BSL expression “ble e1 e2” are shorter and arguably more readable.

To answer to these kinds of issues, HOL4 provides the ability to implement “pretty-

⁴8 lines correspond to the length in documents where line length is limited to 100 characters, instead of the 60 in the report. 100 characters is the usual setting for maximum line lengths and corresponds more closely to the reality.

printers”, which are custom printing functions for a given type. Four pretty-printers have been implemented to shorten the verbosity of the printed representation and to add colors to the output. Figure 6.2 shows the same expression printed with the pretty-printers enabled.

```
BExp_If
  (BExp_Or
    (BExp_LessOrEqual
      (BExp_Den (BVar "x" (BType_Imm Bit64))) (BExp_Const (Imm64 100w)))
    (BExp_Not
      (BExp_LessOrEqual
        (BExp_Plus
          (BExp_Den (BVar "y" (BType_Imm Bit64))) (BExp_Const (Imm64 1w)))
          (BExp_Const (Imm64 10w))))
      (BExp_LessOrEqual
        (BExp_Plus
          (BExp_Den (BVar "x" (BType_Imm Bit64)))
          (BExp_Den (BVar "y" (BType_Imm Bit64))))
          (BExp_Const (Imm64 20w))))
  )
  (BExp_Then
    (BExp_Mult (BExp_Den (BVar "x" (BType_Imm Bit64))) (BExp_Const (Imm64 2w)))
  )
  (BExp_Else
    (BExp_Plus
      (BExp_Mult
        (BExp_Den (BVar "x" (BType_Imm Bit64))) (BExp_Const (Imm64 3w)))
      (BExp_Const (Imm64 1w)))
    )
  )
)
```

Figure 6.2: Same expression printed with the pretty-printers enabled

The pretty-printers introduce a set of features:

- Simplification of verbose constructs as discussed before (e.g. `BExp_BinExp` `BExp_Or` is written as `BExp_Or`).
- Different representation of **if-then-else** statements, simplifying reading the expression when either the condition or the **then** expression are very long.
- Consistent breaking—new lines—of long expressions, because the default printer isn’t aware of the structure of printed expressions. In Figure 6.1, we can see inconsistent breaking in addition and multiplication binary operations.
- Highlighting of types, facilitating debugging when the expression isn’t well-typed.
- Highlighting of all strings, facilitating reading labels and variable names.
- Gathering of nested binary expressions of the same type on the same level. We can see this feature in Figure 6.2 with the two nested **or** binary operators, where the three operands are printed on the same level.
- Rainbow parenthesis, i.e. matching pairs of parenthesis are printed in the same

Listing 6.3: Ideal interface for “prove_contract”

```
fun prove_contract contract_name prog_def
  (precond_lbl, precond_bir_exp)
  postcond_lbl_and_bir_exp_list
```

Listing 6.4: Actual interface for “prove_contract”

```
fun prove_contract contract_name prog_def
  (precond_lbl, precond_bir_exp)
  (postcond_lbl_list, postcond_bir_exp)
```

color. This feature is really useful when reading long expression in order to quickly identify where a sub-expression ends.

6.3 Implementing a convenient interface

In order to perform a high number of proofs on the **NIC** model, we want to hide as much as possible the implementation details of the contract verification procedure. Ideally, we want a function “prove_contract” taking a program fragment, a pre- and a post-condition as parameters, and producing a proof about the Hoare triple if the contract holds or a comprehensive and useful error message if it doesn’t. Listing 6.3 shows the ideal interface that we would want, and Listing 6.4 shows the actual interface that have been implemented.

Interface in Listing 6.3 leverages the general idea of how the weakest precondition generation procedure works: it starts from end labels, setting the weakest precondition there to the postcondition, then propagate the weakest precondition of each node of the **CFG** to the previous nodes, and stops when it meets the entry label. Then, it is in theory possible to provide different postconditions to each end label, hence the last parameter being a list of $(end_label, postcond_exp)$ pairs. However, the current tool only supports using the same postcondition for the multiple end labels, therefore the interface has been constrained⁵.

When implementing this function, high attention has been paid to provide useful and comprehensive feedback in the case of failure. To that end, extensive use of exception wrapping has been made in order to give precise context to exceptions, and a logging library has been implemented (cf. Annex A).

⁵A proposal is being discussed at the time of writing this report about making the weakest precondition generation and the Hoare triple definition more general, and possibly allowing this feature.

Listing 6.5: Equivalent pseudocode of the *cjmp* test.

```

entry:
  x = 1;
  goto (if x=1 then assign_y_100 else assign_y_200)
assign_y_100:
  y = 100;
  goto end
assign_y_200:
  y = 200;
  goto end
end:

```

When using **BSL** to express pre- and post-conditions, this function provides an automatic solution to prove contracts. In the following sections, we will then see usage of this function, first to test it and then to perform proofs on the NIC model.

6.4 Testing the automatic proof procedure

Performing simple proofs is needed in order to test that the proof procedure works. The following examples introduce two of the tests that have been implemented, focusing on the critical parts of each of them. To this end, some liberties have been taken in order to reduce the complexity for the reader. Moreover, even if the test on conditional jumps has been the last one introduced in chronological order⁶, it will be presented first because of its relative simplicity.

6.4.1 Conditional jump test

Here we are interested in testing the `prove_contract` function in the presence of a conditional jump with its condition being just an equality test. Feature-wise, this test contains only jump, conditional jump and assignment statements. Listing 6.5 gives a pseudocode representation of this test program, the conditional jump being represented with the “**goto-if-then-else**” construct.

In this test, we want to check that the triple $\{\top\} \text{ prog } \{y = 100\}$ holds. Intuitively, this contract means “for every possible initial state S , executing the program will result in a state S' with $y = 100$ ”. It is interesting to note that the precondition \top means “for every initial state”, analogous to the universal quantifier \forall in logic. This comes from the fact that \top is the weakest precondition possible: $\forall x. x \implies \top$. Thus, for

⁶The test on conditional jumps has been introduced in order to fix a bug in the weakest precondition simplification library.

Listing 6.6: Invocation of `prove_contract` for the *cjmp* test.

```
val thm = prove_contract "cjmp"
  cjmp_prog_def
  (* Precondition *) (blabel_str "entry", btrue)
  (* Postcondition *) (
    [blabel_str "end"],
    beq ((bden o bvarimm32) "y", bconst32 100)
  )
```

this Hoare Triple to hold, the generated weakest precondition must be \top . Listing 6.6 shows the invocation of `prove_contract`.

Figure 6.3 shows the auto-generated BIR $P \implies WP$ expression, and Equation 6.4 shows the same expression after translation to a *wordsTheory* expression. This expression can be trivially simplified to \top , which SMT solvers can very efficiently do. Hence, this invocation to `prove_contract` succeeds.

```
(BExp_Or
  (BExp_Not (BExp_True))
  (BExp_Not
    (BExp_Equal
      (BExp_Den (BVar "x_wp_0" (BType_Imm Bit32))) (BExp_Const (Imm32 1w))))
  (BExp_And
    (BExp_Or
      (BExp_Not
        (BExp_Equal
          (BExp_Den (BVar "x_wp_0" (BType_Imm Bit32)))
            (BExp_Const (Imm32 1w))))
        (BExp_Equal (BExp_Const (Imm32 100w)) (BExp_Const (Imm32 100w))))
      (BExp_Or
        (BExp_Equal
          (BExp_Den (BVar "x_wp_0" (BType_Imm Bit32)))
            (BExp_Const (Imm32 1w)))
          (BExp_Equal (BExp_Const (Imm32 200w)) (BExp_Const (Imm32 100w))))))
    (BExp_Or
      (BExp_Equal
        (BExp_Den (BVar "x_wp_0" (BType_Imm Bit32)))
          (BExp_Const (Imm32 1w)))
        (BExp_Equal (BExp_Const (Imm32 200w)) (BExp_Const (Imm32 100w))))
      (BExp_Equal (BExp_Const (Imm32 100w)) (BExp_Const (Imm32 100w))))))
```

Figure 6.3: Auto-generated $P \implies WP$ BIR expression for the *cjmp* test.

$$\top \vee (\neg(x = 1w) \vee ((\neg(x = 1w) \vee 100w = 100w) \wedge (x = 1w \vee 200w = 100w))) \quad (6.4)$$

6.4.2 Load and store test

The **NIC** manipulating a buffer descriptor queue, represented in BIR using memories, we need to ensure that `prove_contract` works with programs containing

memories. In this test, we will store a number N in memory at address A , then load a number into x from address B . We want to check the following Hoare Triple: $\{A = B\} \text{ prog } \{x = N\}$. Listing 6.7 shows the equivalent pseudocode of the test program, Figure 6.4 the auto-generated $P \implies WP$ expression, Figure 6.5 the same expression translated in *wordsTheory* and Listing 6.8 the auto-generated SMT-LIB 2.0 instance featuring *select* and *store* operations.

Listing 6.7: Equivalent pseudocode of the *load and store* test

```
MEM = store(MEM, ADDR1, 42)
x = load(MEM, ADDR2)
```

```
(BExp_Or
  (BExp_Not
    (BExp_Equal
      (BExp_Den (BVar "ADDR1" (BType_Imm Bit32)))
      (BExp_Den (BVar "ADDR2" (BType_Imm Bit32))))
    (BExp_Not
      (BExp_MemEq (BExp_Den (BVar "MEM_wp_0" (BType_Mem Bit32 Bit8)))
        (BExp_Store (BExp_Den (BVar "MEM" (BType_Mem Bit32 Bit8)))
          (BExp_Den (BVar "ADDR1" (BType_Imm Bit32))) BEnd_BigEndian
          (BExp_Const (Imm16 42w))))
      (BExp_Equal
        (BExp_Load (BExp_Den (BVar "MEM_wp_0" (BType_Mem Bit32 Bit8)))
          (BExp_Den (BVar "ADDR2" (BType_Imm Bit32))) BEnd_BigEndian Bit16)
        (BExp_Const (Imm16 42w))))
```

Figure 6.4: Auto-generated $P \implies WP$ BIR expression for the *load and store* test.

```
⊢ ¬(if ADDR1 = ADDR2 then 1w else 0w) ||
¬(if
  MEM_wp_0 =
    MEM |+ (ADDR1 + 1w, (15 >< 8) 42w) |+ (ADDR1 + 0w, (7 >< 0) 42w)
  then
    1w
  else 0w) ||
(if MEM_wp_0 ' (ADDR2 + 1w) @@ MEM_wp_0 ' (ADDR2 + 0w) = 42w then 1w
else 0w) =
1w
```

Figure 6.5: $P \implies WP$ expression translated in *wordsTheory* for the *load and store* test. $><$ is the bitwise extraction operation, $@@$ the word concatenation operation, $|+$ the memory update operation and $'$ the memory load operation. Bit extraction and concatenation is needed because we are working with 16-bit words in a 8-bit memory.

Listing 6.8: Autogenerated SMT instance for the *load and store* test

```
(set-info :source |Automatically generated from HOL4 by
  ↪ SmtLib.goal_to_SmtLib.
Copyright (c) 2011 Tjark Weber. All rights reserved.|)
(set-info :smt-lib-version 2.0)
(declare-fun v0_ADDR1 () (_ BitVec 32))
(declare-fun v1_ADDR2 () (_ BitVec 32))
(declare-fun v2_MEM_wp_0 ()
  (Array (_ BitVec 32) (_ BitVec 8)))
(declare-fun v3_MEM ()
  (Array (_ BitVec 32) (_ BitVec 8)))
(assert
  (not
    (=
      (bvor
        (bvnot (ite (= v0_ADDR1 v1_ADDR2)
                     (_ bv1 1) (_ bv0 1)))
        (bvor
          (bvnot
            (ite
              (= v2_MEM_wp_0
                (store
                  (store
                    v3_MEM
                    (bvadd v0_ADDR1 (_ bv1 32))
                    ((_ zero_extend 0) ((_ extract 15 8) (_ bv42
                      ↪ 16))))))
              (bvadd v0_ADDR1 (_ bv0 32))
              ((_ zero_extend 0) ((_ extract 7 0) (_ bv42 16))
                ↪ )))
            (_ bv1 1)
            (_ bv0 1)))
          (ite
            (=
              (concat
                (select v2_MEM_wp_0 (bvadd v1_ADDR2 (_ bv1 32)))
                (select v2_MEM_wp_0 (bvadd v1_ADDR2 (_ bv0 32))))
              (_ bv42 16))
            (_ bv1 1)
            (_ bv0 1))))
      (_ bv1 1))))
(check-sat)
(exit)
```

This expression is harder to prove manually. However, SMT solvers can report very efficiently that the negated expression is unsatisfiable, proving the contract. Therefore, we see that contracts involving BIR memories can be proved, thanks to the work of Section 5.5.

Other preconditions have been tested to verify that `prove_contract` doesn't prove false contracts and succeeds to prove true ones. With the current program, the precondition $load(MEM, B = N) \wedge B = A + 2$ can establish the postcondition. The second conjunct is important because N is stored in two consecutive 8-bit memory locations. Interestingly, the precondition \perp works for all contracts, because $\forall x. \perp \implies x$, and it works in this particular case.

6.5 Simple automatized proofs on the NIC model

Chapter 4 presented the implementation of parts of the NIC model using BIR. The previous sections of this chapter presented the non **proof-producing** contract verification library that have been implemented. This section will present some of the properties that have been proved on the NIC BIR model. However, as we shall see, `prove_contract` has some intrinsic limitations that prevent verifying some contracts.

Section 3.3 presents an overview of the formal HOL4 proof made in [5] on the formal NIC model. We saw that the properties are phrased in terms of invariants holding in each transition of the model. Such properties can be represented as Hoare Triple and therefore can be proved using this chapter's non proof-producing contract verification library as long as the pre- and post-conditions can be represented using BIR.

In this section, we will see that the non proof-producing function `prove_contract` is able to verify some invariants on the NIC BIR model. We will perform the verification of the Hoare Triples shown in Equations 6.5 and 6.6. Those two Hoare Triple respectively represent that, stating from a non-dead initial NIC state, performing one autonomous transition of the transmission automaton doesn't end in an undefined state (i.e. a dead state) and performing one non-autonomous transition of the initialization automaton does end in an undefined state.

$$\{\neg NIC.dead \wedge NIC.tx.state \in \{tx2, tx3, tx5, tx6, tx7\}\} \\ tx_automaton\{\neg NIC.dead\} \quad (6.5)$$

$$\{\neg NIC.dead \wedge NIC.init.state \in \{it1, it3, it4\}\} \\ init_automaton\{NIC.dead\} \quad (6.6)$$

Listing 6.9 shows the SML code used to prove the first contract. Code for the other contract is similar. This example shows that the library implemented in the chapter fulfills its goal of being simple to use and fully automatic when given a Hoare Triple and a program. Combined with BSL and custom helper functions, it becomes easy to express the pre- and post-conditions, and it is still possible to use arbitrary SML functions or HOL4 terms.

Listing 6.9: SML code used to prove contract 6.5.

```
(* TX automaton: NIC doesn't die in autonomous transition *)
val (_, _, p_imp_wp_word_thm) = prove_contract
  "tx_automaton_doesnt_die"
  nic_program_def
  (* Precondition *) (
    blabel_str "tx_entry",
    bandl [
      invariant_nic_not_dead,
      borl (List.map
        (fn s => beq (bdenstate "nic_tx_state",
          bstateval_tx s))
        (#autonomous_step_list tx_state))
    ])
  (* Postcondition *) (
    [blabel_str "tx_end"],
    invariant_nic_not_dead
  )
val _ = info "Successfully proved: tx_automaton_doesn't_die"
val _ = if !level_log >= logLib.INFO
  then (pprint_thm p_imp_wp_word_thm)
  else ();
```

6.5.1 Limitations of this approach

While being convenient and working well for simple contracts, the contract verification library implemented in this chapter suffers from two limitations:

- The first one is straightforward: this library isn't proof-producing. Using HOL4 requires a significant learning effort, and easier non-proof producing tools exist. Hence, when using HOL4, one usually wants to reason about proofs, and this library is not proof-producing and could therefore introduce subtle bugs or inconsistencies. Moreover, this library doesn't offer the possibility to compose Hoare Triples because it doesn't output any related theorem. If one wanted to do so, an oracle would have to be used in order to generate such theorems and enabling theorem compositions and further reasoning. Nevertheless, as already explained, proof-producing code is more costly is the desired level of trustworthiness should be carefully decided, and this library offers a pragmatic solution.

- Secondly, this library and its approach of contract based verification is limited by the expressiveness of BIR, since the pre- and post-conditions are BIR terms. BIR doesn't have quantifiers, and it is therefore impossible to prove contracts containing existential quantifiers. The existing formal proof on the NIC of [5] requires existential quantifiers in order to reason about the buffer descriptor queue in the *CPPI_RAM* memory of the NIC, and this part of the proof can therefore not be replaced by using `prove_contract`. Moreover, as explained in the first point of this list, composing verification produced with this library with raw HOL4 proofs is not straightforward. The following Chapter 7 explores a new proof-producing approach for performing contract based verification on the BIR model that answer to this issue.

Chapter 7

Trustful analysis on the NIC model

*In the previous chapter, we implemented an automated non **proof-producing** contract verification library, the non proof-producing part being the translation from BIR expressions to the equivalent wordsTheory and combinTheory expression. Moreover, this verification library can only produce contracts on BIR programs. In order to perform trustworthy verification on the **NIC** model, we need to make proofs directly on the **NIC** state. In this chapter, we will perform a proof on a BIR program and then lift it to the **NIC** model.*

In order to prove the feasibility of this approach, we will prove a simple property. Listing 7.1 contains the NIC state on which we want to prove a property, Equations 7.1, 7.2 and 7.3 present the property that we want to prove, and Listing 7.2 contains a pseudocode representation of the BIR program on which we will perform the verification. Figure 7.1 represents visually the structure of the verification and the steps that we will take during the proof.

Listing 7.1: NIC state used in this proof

```
Datatype 'nic_state = <|  
  dead : bool;  
  x : word32  
|> '
```

$$\vdash \forall nic. P_{NIC} nic \stackrel{def}{=} \neg nic.dead \wedge nic.x = 0w \quad (7.1)$$

$$\vdash \forall nic nic'. Q_{NIC} nic nic' \stackrel{def}{=} \neg nic'.dead \wedge nic'.x = nic.x + 1w \quad (7.2)$$

$$\vdash \forall nic nic'. P_{NIC} nic \wedge exec_prog nic bir_prog nic' \implies Q_{NIC} nic nic' \quad (7.3)$$

Listing 7.2: Pseudocode of the program used in this proof

```

nic.x := nic.x + 1
if nic.x > 10:
    nic.dead := true

```

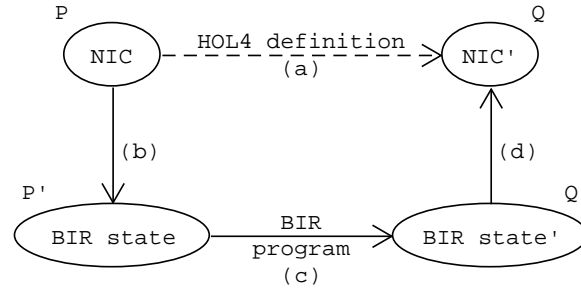


Figure 7.1: Visual structure of the proof. References like (a) to the arrows of this Figure are used throughout the proof to refer to a particular step.

Remark 1. $Q_{NIC} \text{ nic } \text{nic}'$ is defined on both initial and final states, in order to be able to reason about the initial state in the postcondition. This allows us to write $\text{nic}'.x = \text{nic}.x + 1w$ instead of $\text{nic}'.x = 1w$ for example.

Equation 7.3 uses a relation $exec_prog$ that we shall define now. As we want to make a proof on an undefined HOL4 definition (a), we must establish an equivalence between the HOL4 definition (a) and the BIR program (c). In real proofs, this can either be produced by a lifter which generates the BIR program from a given input program and gives a “certificate”, i.e. a theorem stating the equivalence, or be a definition which would then mean that we trust that the BIR program is equivalent to the HOL4 definition. In this proof, we will use a definition. This definition shall state that $exec_prog \text{ nic } \text{bir_prog } \text{nic}'$ (a) is equivalent to executing the BIR program from a state bir_state to a state $\text{bir_state}'$ (c), where nic is somehow equivalent to bir_state (b) and nic' somehow equivalent to $\text{bir_state}'$ (d). To express an equivalence between HOL4 states (“NIC”) and BIR states, we introduce a relation R . The relation $R \text{ nic } \text{bir_state}$ is defined as a simple mapping between the BIR state and the NIC state, as shown in Listing 7.3. Then, we define the relation $exec_prog$ as shown in Equation 7.4¹.

¹Definition 7.4 has been annotated to visualize how it is connected to the structure of the proof on Figure 7.1. In addition, the BIR_exec relation is used as a shorthand for $\text{bir_exec_to_labels}$ in order to simplify the proof.

Listing 7.3: Definition of the relation R

```

val R_def = Define `
  R (nic: nic_state) (bir_state: bir_state_t) <=>
    (bir_env_lookup "nic_dead" bir_state.bst_environ
     = SOME (BType_Bool,
             SOME (BVal_Imm (Imm1 nic.dead))))
  /\ (bir_env_lookup "nic_x" bir_state.bst_environ
      = SOME (BType_Imm Bit32,
              SOME (BVal_Imm (Imm32 nic.x)))) `

```

$$\begin{aligned}
& \vdash \forall nic\ nic'. \text{exec_prog } nic\ bir_prog\ nic' & (a) \\
& \stackrel{\text{def}}{=} \forall bir_state\ bir_state'. & \\
& \quad (R\ nic\ bir_state & (b) \\
& \quad \wedge bir_state' = BIR_exec\ prog\ bir_state) & (c) \\
& \quad \implies R\ nic'\ bir_state' & (d)
\end{aligned} \tag{7.4}$$

Proof of Equation 7.3. In order to begin the proof, as the goal 7.3 is defined over nic states, we need a theorem about the injectivity of the relation R , stating that for all nic exists a bir_state such that $R\ nic\ bir_state$ (b). Additionally, the BIR_exec relation will also need some properties on bir_state , that we shall add in this injectivity theorem now.

Theorem 7.0.1. *Injectivity theorem of R*

$$\begin{aligned}
& \vdash \forall nic. \exists bir_state. \\
& \quad R\ nic\ bir_state \\
& \quad \wedge bir_state.bst_pc.bpc_index = 0 \\
& \quad \wedge bir_state.bst_pc.bpc_label = entry_label \\
& \quad \wedge bir_state.bst_status = BST_Running
\end{aligned}$$

Proof. After rewriting the relation R , we prove theorem 7.0.1 by exhibiting a satisfying bir_state . \square

In possession of a bir_state in relation with a nic , we now need to lift the precondition $P_{NIC}nic$ on this bir_state . First, we need to introduce equivalent pre- —and post-—conditions on the BIR states, then we shall prove that the precondition lifts.

Listing 7.4: Equivalent pre- and postconditions on BIR states

```

val BIR_P_exp_def = Define `BIR_P_exp = ^ (bandl [

```

```

    beq ((bden o bvarimm1) "nic_dead", bfalse),
    beq ((bden o bvarimm32) "nic_x", bconst32 0)
  ]) `
val BIR_Q_exp_def = Define `BIR_Q_exp = ^(bandl [
    beq ((bden o bvarimm1) "nic_dead", bfalse),
    beq ((bden o bvarimm32) "nic_x", bconst32 1)
  ]) `
val BIR_P_def = Define `BIR_P bstate =
    bir_eval_bool_exp BIR_P_exp bstate.bst_envIRON `
val BIR_Q_def = Define `BIR_Q bstate =
    bir_eval_bool_exp BIR_Q_exp bstate.bst_envIRON `

```

Limitation Q_{BIR} is a function of the end state only. Hence, in order to reason about the initial state in the we need in general to introduce ghost variables postcondition. In this proof, since the contract that we are proving is simple, using the actual value of $nic.x$ is enough. However, this may pose a problem if we want to generalize the proof.

Notation $P_{BIR} bir_state$ and $Q_{BIR} bir_state$ are defined using $bir_eval_bool_exp$, which evaluates respectively the expressions P_{BIR}^{exp} and Q_{BIR}^{exp} in a given BIR state. In order to simplify the proof, let's define a new operator $\stackrel{eval}{=}$ that is used to evaluate given variables, e.g. $bir_state.x \stackrel{eval}{=} 0w$.

Theorem 7.0.2. *Lifting of $P_{NIC} nic$ to bir_state*

$$\vdash \forall bir_state (\exists nic. R nic bir_state \wedge P_{NIC} nic) \implies P_{BIR} bir_state$$

Proof. Let's do this proof in a backward way. By discharging the antecedent of the implication and using the existencial elimination inference rule, we get as assumptions $P_{NIC} nic$ and $R nic bir_state$. From this, we can deduce that $bir_state.x \stackrel{eval}{=} 1w$ and $bir_state.dead \stackrel{eval}{=} \perp$. Then, we can substitute those values in the goal, which proves it. \square

Assuming that we have a Hoare Triple theorem between initial and final BIR states, we can use Definition 7.4 in order to establish that $R nic' bir_state'$. Then, in order to prove $Q_{NIC} nic$ (d), we have to transfer the postcondition back from bir_state to nic .

Theorem 7.0.3. *Lowering Q_{BIR} bir_state to nic .*

$$\begin{aligned} \vdash \forall bir_state'. Q_{BIR} bir_state' \implies \\ (\forall nic\ nic' bir_state. P_{BIR} bir_state \\ \wedge R\ nic\ bir_state \wedge R\ nic' bir_state' \\ \implies Q_{NIC} nic\ nic') \end{aligned}$$

We introduce bir_state and P_{BIR} in this theorem for the reason explained in Remark 1, i.e. reason about both the initial and final state in the postcondition.

Proof. This proof has been done in HOL4. The reasoning is quite similar to the proof of Theorem 7.0.2, as the backward proof mainly involves rewriting and simplification. We will omit this proof here and redirect the reader to the HOL4 proof available in our source repository [28]. \square

We will now prove that the Hoare Triple holds on the BIR program.

Theorem 7.0.4. $\{P_{BIR}^{exp}\} bir_prog \{Q_{BIR}^{exp}\}$

Proof. To prove this Hoare Triple, we used the proof-producing procedure implemented in **HolBA** in order to generate the weakest precondition. The automatically derived weakest precondition is shown in Figure 7.2. Section 5.4 already discussed how to perform this proof: we have to prove Equation 5.12 with **wp** being the expression in Figure 7.2 and **pre** being P_{BIR}^{exp} . Because we want to use a **SMT** solver, we need to turn the goal of the backward proof into a *wordsTheory* expression. *combinTheory* isn't needed in this case since BIR memories are not used. Equation 5.12 uses *bir_eval_exp* which evaluates an expression in the given BIR state. Therefore, to translate the goal into a *wordsTheory* expression, we need to use BIR's semantic. The semantic needs well-typedness and initialization of the variables. At the time of writing, HolBA offers no support for automatic rewriting with the semantic definitions, so multiple lemmas about initialization, well-typedness and type equality must be manually proved for every variable. Those are not shown here because they consist of simple rewriting and simplification.

Then, the proof consist of consecutively rewriting following the definition of *bir_eval_exp* and the definition it uses until the goal only contains

$$bir_env_read\ (BVar\ "nic_x"\ (BType_Imm\ Bit32))\ bir_state.bst_environ$$

and similarly for *nic.dead*. Let's call those expressions *x_val* and *dead_val* respectively. For expressions, BIR semantic is defined over immutable and constant values. Therefore, we need to establish an equivalence between the values that we currently have.

```

(BExp_And
  (BExp_Or
    (BExp_Not
      (BExp_LessThan
        (BExp_Const (Imm32 10w))
        (BExp_Plus
          (BExp_Den (BVar "nic_x" (BType_Imm Bit32)))
          (BExp_Const (Imm32 1w))))))
    (BExp_And
      (BExp_Equal (BExp_True) (BExp_False))
      (BExp_Equal
        (BExp_Plus
          (BExp_Den (BVar "nic_x" (BType_Imm Bit32)))
          (BExp_Const (Imm32 1w)))
        (BExp_Const (Imm32 1w))))))
  (BExp_Or
    (BExp_LessThan
      (BExp_Const (Imm32 10w))
      (BExp_Plus
        (BExp_Den (BVar "nic_x" (BType_Imm Bit32)))
        (BExp_Const (Imm32 1w))))
    (BExp_And
      (BExp_Equal
        (BExp_Den (BVar "nic_dead" (BType_Imm Bit1)))
        (BExp_False))
      (BExp_Equal
        (BExp_Plus
          (BExp_Den (BVar "nic_x" (BType_Imm Bit32)))
          (BExp_Const (Imm32 1w)))
        (BExp_Const (Imm32 1w))))))
  )
)

```

Figure 7.2: Autogenerated weakest precondition for proof of Theorem 7.0.4.

Lemma 7.0.5.

$$\exists x_{imm}. x_{val} = BVal_Imm\ x_{imm}$$

Proof. Assuming well-typedness and initialization, this theorem immediately results from the BIR semantic. \square

Lemma 7.0.6.

$$\exists x_{word}. x_{imm} = Imm32\ x_{word}$$

Proof. This lemma is part of the BIR semantic, as one of the six conjuncts of the `bir_imm_t_nchotomy` theorem, which establishes this existence theorem for every BIR immutable types. \square

Now, using Lemma 7.0.5, we are able to substitute all occurrences of x_val into $BVal_Imm(Imm32\ x_word)$, and similarly for $dead_val$. Finally, rewriting the goal using the full set of BIR semantic theorems and some rewriting rules, the goal reduces to an expression free of BIR terms:

$$\begin{aligned}
 (dead_w = 0w) \wedge (x_w = 0w) &\implies \\
 \left(\neg(10w <_+ x_w + 1w) \vee ((1w = 0w) \wedge (x_w + 1w = 1w)) \right) \wedge & \quad (7.0.6.1) \\
 \left((10w <_+ x_w + 1w) \vee ((dead_w = 0w) \wedge (x_w + 1w = 1w)) \right)
 \end{aligned}$$

An **SMT** solver is able to prove this goal. Interestingly, HOL4 simplification procedure are also able to prove it. However, they won't be able to prove it for more complicated ones, or will be less effective than SMT solvers. \square

Finally, using the deduction rule with Theorems 7.0.1, 7.0.2, 7.3, 7.0.4, 7.4 and 7.0.3, in that order, concludes this proof. \blacksquare

Chapter 8

Conclusions

8.1 Discussion

The goal of this paper was to perform experiments about automation of formal verification of devices at the binary level. The NIC model of [5] has served as a central theme throughout the work. Work in this thesis has been divided in practice into three parts, in addition to the learning process due to the very steep learning curve of HOL4:

1. **Implementation of the NIC BIR model:** the formal model of the NIC of [5] has been partially implemented as a BIR program, and BSL has been added to the HolBA platform in order to make this task more convenient.
2. **Non proof-producing automatic contract verification library:** a user-friendly (adjective rarely used to describe HOL4), convenient and automatized library for contract verification have been implemented. This library brings HolBA one step closer to “one-button solutions” for performing software verification. It features a non proof-producing translation of BIR expressions to expressions that can be exported to SMT solvers, and an extension of HOL4’s SMT library has been realized in order to add reasoning about BIR memories.
3. **Formal proof of a BIR program on a HOL4 state:** a novel usage of BIR and method for transferring properties from BIR program to formal HOL4 specification has been carried out. Additionally, while not automated, the work needed in order to implemented automatic HOL4 tactics has been clearly identified. Every automation work should be preceded by a manual implementation of the task, that this proof intends to be.

Throughout this work, a strong focus has been placed on ease of use and user-friendliness with BSL, the pretty-printer, the implementation of new logging utilities, guidelines about error handling in SML and with HOL4, and a friendly exception printer. I

strongly believe that user experience of such complex platforms are of primary importance, because it can dramatically reduce time spent in debugging and the need of an exhaustive documentation, and give to users to desire to keep using the platform. If it is decided to put more work in HolBA in order to make it a powerful binary analysis platform, the research team is advised to consider this aspect while expanding HolBA capabilities.

Some failed experiments have also been realized, including the DepGraph tool that revealed to not be adapted to its planned usage or which would have needed to much work for too few benefits (though DepGraph’s design is future-proof and the tool can easily be extended later if new needs or novel ideas appear), and the former two implementation attempts of the NIC model using flowcharts or C which have been found to not be suitable for device models.

8.2 Future work

The topic of software verification, while being as old as Computer Science, still have a lot of work to be done and paths to be explored. This thesis opens some paths that should be considered by someone interested in the topic.

The pretty-printer implemented in this work is only a small experiment only scratching the surface of what can be done in HOL4. Further work in this domain should follow the concrete needs that arise with extended usage of BIR. Current limitations and possible further work of the pretty-printer presented in this document include:

- the generated output isn’t parsable, which is against HOL4 conventions where every printed terms should be parsable by HOL4 quoting system. A solution to this problem would be to implement a theory introducing definitions for the new abbreviations. However, some liberties taken in the pretty-printer cannot be parsed: (a) the flat printing of nested binary expressions of the same state, because they would need multiple arity functions which are not available in HOL4, and (b) **if-then-else** statements in the current representation.
- the pretty-printer doesn’t use infix operators, feature widely used in custom syntax used in some theories, like *wordsTheory*. Infix syntax for binary operators can dramatically reduce verbosity of expressions, and give a more familiar representation that *could* be simpler to work with.
- the choice of colors could be reviewed in order to more closely follow HOL4 convention and usage. For example, BIR types are highlighted in the same color than free variables. While creating a very limited chance of ambiguities, users familiar with HOL4 conventions can be annoyed with this inconsistency.

The non proof-producing contract verification library implemented in this work can be really useful for prototyping, but it cannot be used for trustful verification. The

limitations it suffers could be reasonably fast to overcome, and they should be fixed especially if HolBA were to become a strong binary analysis platform. Limitations include (a) the impossibility to compose contract in order to reason about longer functions, to compose functions or to prove properties on not loop-free programs, and (b) the obvious limitation of not being proof-producing. Work is ongoing at the time of writing about weak-correctness and contract composition, answering to (a). The main blocker of (b) is the absence of a proof-producing way of using SMT solvers for proving contracts (Equation 5.12). The proof of Chapter 7 and its SML implementation identify what automatic procedures should be added to HolBA in order to perform step (b) and progress towards proof-producing automatic verification. Additionally, the signature of `prove_contract` should be rework in order to enable mixing it with proof-producing code, after both (a) and (b) are solved, for example it should work with HOL4 definitions instead of raw terms.

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Appendix A

Tracing and error handling in HOL4

TODO: Is this annex relevant in this thesis?

Reference in Section 6.3, but the need arised in general when debugging HOL4 code for tracing code. LogLib represent an higher-level implementation of HOL4 tracing system. Verbosity levels are specified and output is colored, helping to read long outputs. Also include exception wrapping.

Content:

- Present the problems specific to HOL4 that arise when debugging proofs and libs.
- Introduce tracing in HOL4 in general (levels, HOL4 warnings and error **messages**).
- Present logLib
- Introduce exception wrapping and explain how they help to locate the location of errors and can help interactive use of HOL4 by giving more precise information about what went wrong. Analogous to backtraces.

Appendix B

DepGraph's design

TODO: Is this annex relevant in this thesis?

Content:

- Present the modular architecture: Frontend \rightarrow [middlewares] \rightarrow backend structure.
- This tool can be used for other languages or data via new front-ends.
- Other visualization techniques/rendering can be implemented via backends.
- Middlewares can be any graph algorithm. Signature: `graph -> graph`.
- Time to implement: 4 days