

# Homework 2 - LL Parsing - Solution

## Objectives

1. Demonstrate an ability to convert a grammar into an equivalent LL grammar.
2. Demonstrate an ability to produce an LR automata from a grammar.

## LL Problems

(5 pts. each) For each of these grammars, convert it to an equivalent LL grammar or else assert that it is already LL.

Problems

Solutions

```
----- 1 -----
S   ::= x E
      | z
E   ::= E a E
      | E b E
      | c

S   ::= x E
      | z
E   ::= c E'
E'  ::= a E E'
      | b E E'
      | eps
```

```
----- 2 -----
S   ::= z y
      | y E
E   ::= S a
      | x a

Already LL
```

```
----- 3 -----
S   ::= x E
      | x E y
      | z
E   ::= a E
      | b

S   ::= x E S'
      | z
S'  ::= y
      | eps
E   ::= a E
      | b
```

# LR Problem

Consider the following grammar:

```
1  S    ::= ( S S )
2      | * V
3  V    ::= * V
4      | p
```

## Part 1

(5 pts.) First, calculate the first and follow sets for the non-terminals of the grammar.

$\text{First}(S) = \{ (, *, \} \}$

$\text{First}(V) = \{ *, p \}$

$\text{Follow}(S) = \{ (, *, ), \$ \}$

$\text{Follow}(V) = \{ (, *, ), \$ \}$

## Part 2

(30 pts.) Now, describe the LR automata for the grammar by listing the item sets (states) and filling out the action and goto tables.

### States

0.  $S ::= . ( S S )$   
    $S ::= . * V$

1.  $S ::= ( . S S )$   
    $S ::= . ( S S )$   
    $S ::= . * V$

2.  $S ::= * . V$   
    $V ::= . * V$   
    $V ::= . p$

3.  $S ::= ( S . S )$   
    $S ::= . ( S S )$   
    $S ::= . * V$

4.  $S ::= * V .$

5.  $V ::= * . V$   
    $V ::= . * V$   
    $V ::= . p$

6.  $V ::= p .$

7.  $S ::= ( S S . )$

8.  $V ::= * V .$

9.  $S ::= ( S S ) .$

## Action

	p	*	(	)	\$
0		S	S		
1		S	S		
2	S	S			
3		S	S		
4		R2	R2	R2	R2
5	S	S			
6		R4	R4	R4	R4
7				S	
8		R3	R3	R3	R4
9		R1	R1	R1	R1

## Goto

	p	*	(	)	\$	S	V
0		2	1				
1		2	1			3	
2	6	5					4
3		2	1			7	
4							
5	6	5					8
6							
7				9			
8							
9							