## DISTRIBUTED INFORMATION SYSTEMS





The Pig Experience:
Building High-Level Data
flows on top of Map-Reduce

**VLD**B paper

Source: Javeria Iqbal, Martin Theobald

Alan F. Gates, Olga Natkovich, Shubham Chopra, Pradeep Kamath, Shravan M. Narayanamurthy, Christopher Olston, Benjamin Reed, Santhosh Srinivasan, Utkarsh Srivastava,

Building a High-Level Dataflow System on top of Map-Reduce: The Pig Experience, VLDB 2009.

 Map-Reduce and the need for Pig Latin

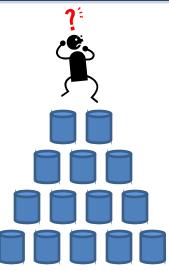
Pig Latin

Compilation into Map-Reduce

Optimization

Future Work





- Internet companies swimming in data
  - TBs/day for Yahoo! Or Google!
  - PBs/day for FaceBook!
- Data analysis is "inner loop" of product innovation



#### Data Warehousing ...?

Scale

- High level declarative approach
- Little control over execution method

Price

Prohibitively expensive at web scale

Up to \$200K/TB

SQL

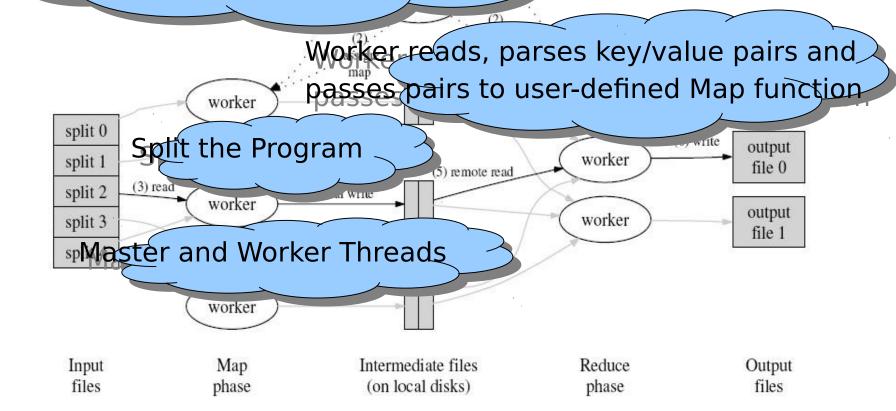
Often not scalable enough

- Map: Performs filtering
- Reduce : Performs the aggregation
- These are two high level declarative primitives to enable parallel processing
- BUT no complex Database Operations e.g. Joins



# **Execution Overview of**<a href="#">Map-Reduce</a>

ered pairs are written to local disk partitions, tion of buffered pairs are sent to reduce workers

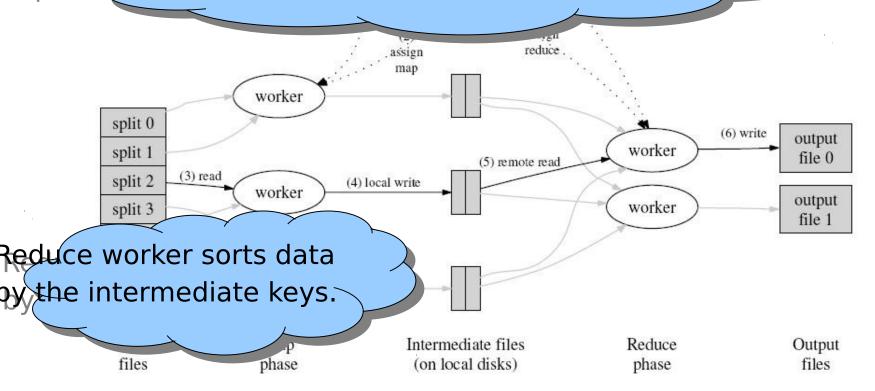




### **Execution Overview of** Map-Reduce

User

nique keys, values are passed to user's Reduce function. utput is appended to the output file for this reduce partition.





### The Map-Reduce Appeal

Scale

Scalable due to simpler design

Explicit programming model

Price

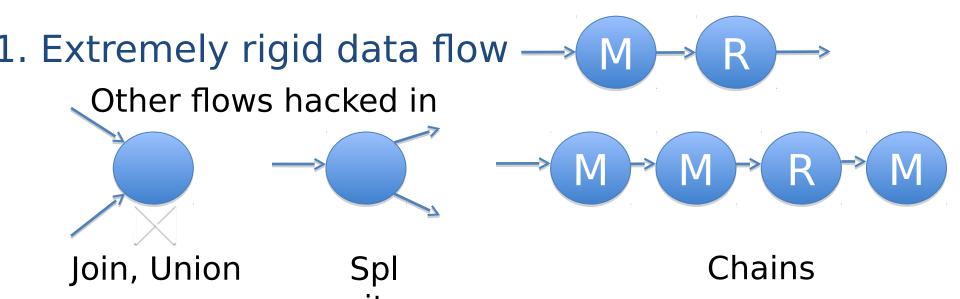
• Only parallelizable Runs on cheap commodity hardware

Less Administration

SQL

Procedural Control- a processing "pipe"





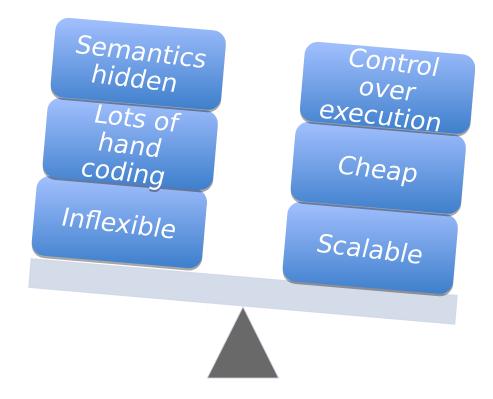
Common operation<sup>it</sup> must be coded by hand Join, filter, projection, aggregates, sorting, distinct

Semantics hidden inside map-reduce functions Difficult to maintain, extend, and optimize

No combined processing of multiple Datasets Joins and other data processing operations

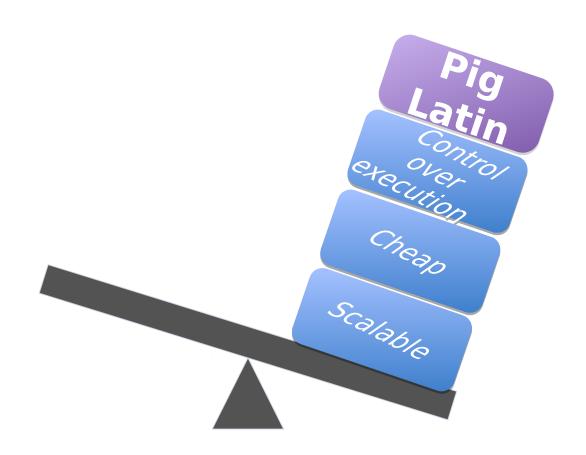


Need a high-level, general data flow language





Need a high-level, general data flow language



 Map-Reduce and the need for Pig Latin

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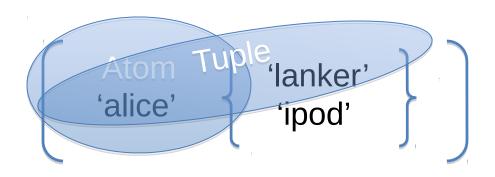
### **Pig Latin: Data Types**

- Rich and Simple Data Model Simple Types: int, long, double, chararray, bytearray **Complex Types:**  Atom: String or Number e.g. ('apple') • Tuple: Collection of fields e.g. (áppe', 'mango') Bag: Collection of tuples ('apple', 'mango') (ápple', ('red', 'yellow'))
- Map: Key, Value Pair



### **Example: Data Model**

- Atom: contains Single atomic value
  Tuple: sequence of fields
  Bag: collection of tuple with possible duplicates



## Pig Latin: Input/Output Data

#### **Input:**

```
queries = LOAD `query log.txt'
USING myLoad()
AS (userId, queryString, timestamp);
Output:
STORE query revenues INTO
  'myoutput'
USING myStore();
```

# Pig Latin: General Syntax

- Discarding Unwanted Data: FILTER
- Comparison operators such as ==, eq, !=, neq
- Logical connectors AND, OR, NOT

## Y.

## Pig Latin: Expression able

$$t = \left( \text{'alice'}, \left\{ \begin{array}{c} (\text{'lakers'}, 1) \\ (\text{'iPod'}, 2) \end{array} \right\}, \left[ \text{'age'} \rightarrow 20 \right] \right)$$

Let fields of tuple t be called f1, f2, f3

Expression Type	Example	Value for t
Constant	'bob'	Independent of t
Field by position	\$0	'alice'
Field by name	f3	$\left[ \text{`age'} \rightarrow 20 \right]$
Projection	f2.\$0	<pre>{ ('lakers') }   ('iPod') }</pre>
Map Lookup	f3#'age'	20
Function Evaluation	SUM(f2.\$1)	1 + 2 = 3
Conditional Expression	f3#'age'>18? 'adult':'minor'	'adult'
Flattening	FLATTEN(f2)	'lakers', 1 'iPod', 2

## Pig Latin: FOREACH with Flatten

expanded\_queries = FOREACH queries GENERATE userId, expandQuery(queryString);

expanded\_queries = FOREACH queries GENERATE userId,

queries:
(userId, queryString, timestamp)

(alice, lakers, 1)
(bob, iPod, 3)

FOREACH queries GENERATE expandQuery(queryString)

(without flattening)

(alice, lakers rumors)
(alice, lakers rumors)
(alice, lakers rumors)
(bob, iPod nano)
(bob, iPod nano)
(bob, iPod shuffle)



### Pig Latin: COGROUP

#### Getting Related Data Together: COGROUP

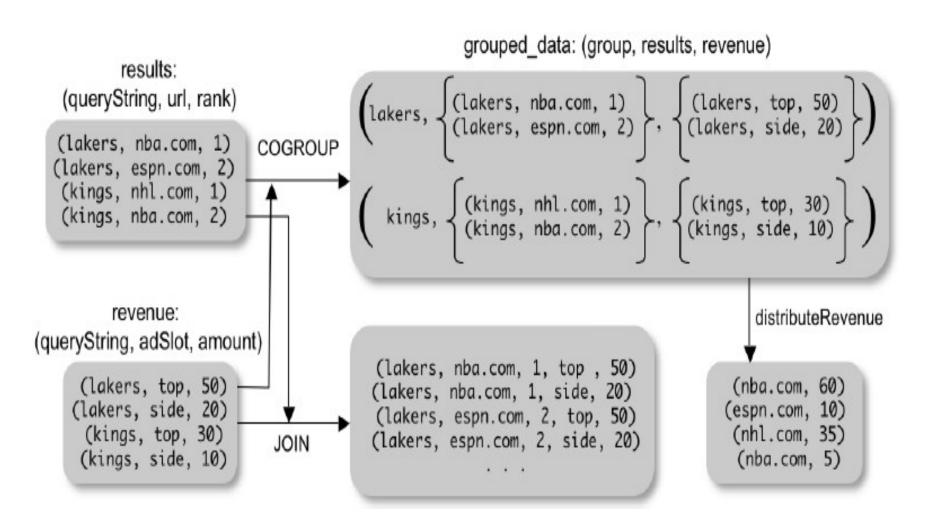
Suppose we have two data sets

result: (queryString, url, position)

revenue: (queryString, adSlot, amount)

```
grouped_data = COGROUP result BY
queryString, revenue BY
queryString;
```

## Pig Latin: COGROUP vs.



## Pig Latin: Map-Reduce

Map-Reduce in Pig Latin

```
map_result = FOREACH input GENERATE
FLATTEN(map(*));
```

```
key_group = GROUP map_result BY $0;
```

```
output = FOREACH key_group GENERATE
reduce(*);
```

# Pig Latin: Other Commands

- UNION: Returns the union of two or more bags
- CROSS: Returns the cross product
- ORDER: Orders a bag by the specified field(s)
- DISTINCT: Eliminates duplicate tuple in a bag

# Pig Latin: Nested Operations

```
grouped revenue = GROUP revenue BY
 queryString;
query revenues = FOREACH
 grouped revenue {
top slot = FILTER revenue BY
adSlot eq `top';
GENERATE queryString,
SUM(top slot.amount),
SUM(revenue.amount);
```



## Pig Pen: Screen Shot

Operators		
LOAD GROUP COGROUP FILTER FOREACH ORDER		
= LOAD USING Default ✓ AS ( ) Generate Query		
visits = LOAD 'visits.txt' AS (user, url, time);	visits:	(Amy, cnn.com, 8am) (Amy, frogs.com, 9am) (Fred, snails.com, 11am)
pages = LOAD 'pages.txt' AS (url, pagerank);	pages:	(cnn.com, 0.8) (frogs.com, 0.8) (snails.com, 0.3)
v_p = JOIN visits BY url, pages BY url;	v_p:	(Amy, cnn.com, 8am, cnn.com, 0.8) (Amy, frogs.com, 9am, frogs.com, 0.8) (Fred, snails.com, 11am, snails.com, 0.3)
users = GROUP v_p BY user;	users:	(Amy, { (Amy, cnn.com, 8am, cnn.com, 0.8), (Amy, frogs.com, 9am, frogs.com, 0.8) }) (Fred, { (Fred, snails.com, 11am, snails.com, 0.3) })
useravg = FOREACH users GENERATE group, AVG(v_p.pagerank) AS avgpr;	useravg:	(Amy, 0.8) (Fred, 0.3)
answer = FILTER useravg BY avgpr > '0.5';	answer:	(Amy, 0.8)

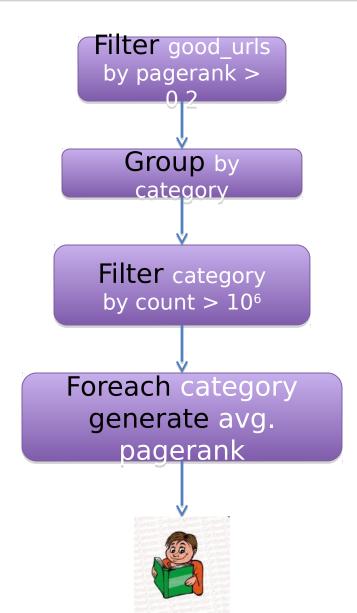


### Pig Latin: Example 1

Suppose we have a table urls: (url, category, pagerank) Simple SQL query that finds, For each sufficiently large category, the average pagerank of highpagerank urls in that category SELECT category, Avg(pagetank) FROM urls WHERE pagerank > 0.2 GROUP BY category HAVING COUNT(\*) > 106



#### **Data Flow**





#### **Equivalent Pig Latin**

- good\_urls = FILTER urls BY pagerank > 0.2;
- groups = GROUP good\_urls BY category;
- big\_groups = FILTER groups BY COUNT(good\_urls) > 10<sup>6</sup>;
- output = FOREACH big\_groups GENERATE category, AVG(good\_urls.pagerank);

# Example 2: Data Analysis Task

#### I the top 10 most visited pages in each categ

#### **Visits**

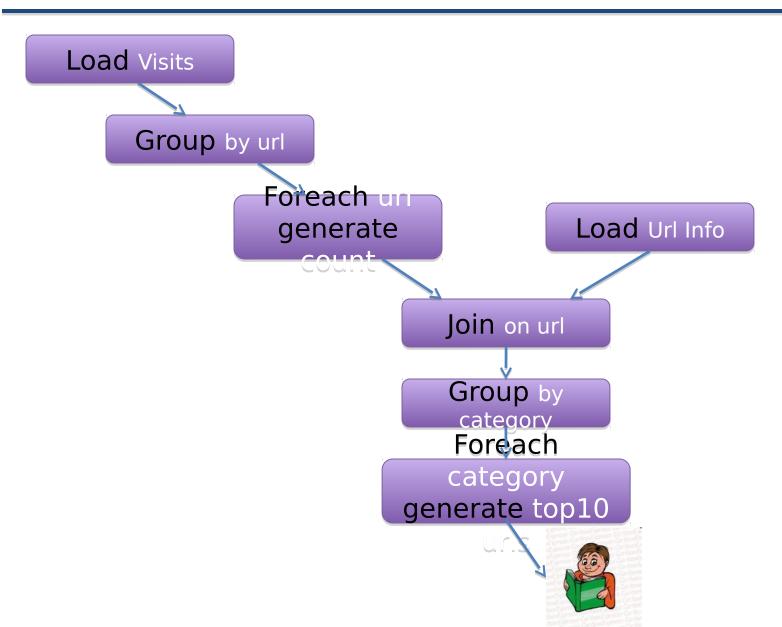
User	Url	Time
Amy	cnn.com	8:00
Amy	bbc.com	10:00
Amy	flickr.com	10:05
Fred	cnn.com	12:00

#### Url Info

Url	Catego ry	PageRa nk
cnn.com	News	0.9
bbc.com	News	0.8
flickr.co m	Photos	0.7
espn.co m	Sports	0.9



#### **Data Flow**





## **Equivalent Pig Latin**

```
visits
              = load '/data/visits' as (user, url,
 time);
gVisits
              = group visits by url;
visitCounts = foreach gVisits generate url,
  count(visits);
             = load '/data/urlInfo' as (url,
urlInfo
  category, pRank);
visitCounts = join visitCounts by url, urlInfo by
  url;
```

gCategories = group visitCounts by category; topUrls = foreach gCategories generate

# Quick Start and Interoperability

```
= load '/data/visits' as (user, url,
visits
 time);
gVisits
             = group visits by url;
visitCounts = foreach gVisits generate url,
  count(urlVisits);
             = load '/data/urlInfo' as (url,
urlInfo
  catego
            Operates directly over
                     files
visitCour
                                      l, urlInfo by
  url;
gCategories = group visitCounts by category;
topUrls = foreach gCategories generate
```

# Quick Start and Interoperability

```
= load '/data/visits' as (user, url,
visits
 time);
gVisits
             = group visits by url;
visitCounts = foreach gVisits generate url,
  count(urlVisits);
             - load '/data/urllnfo' as (url,
urlInfo
              Schemas optional;
 catego
               Can be assigned
                  dynamically
                                       urlinfo by
visitCoun
 url;
gCategories = group visitCounts by category;
topUrls = foreach gCategories generate
```

## Y!

### User-Code as a First-Class Citizen

```
visits
time
gVisits
visitCo
cou
```

User-defined functions (UDFs) can be used in every construct

- Load, Store
- Group, Filter, Foreach

s (user, url,

erate url,

```
urlInfo = load '/data/urlInfo' as (url, category, pRank);
```

```
visitCounts = join visitCounts by url, urlInfo by url;
```

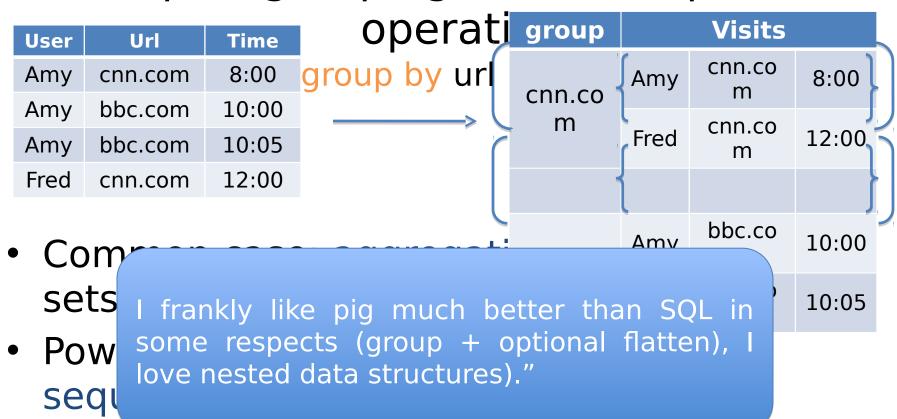
```
gCategories = group visitCounts by category;
topUrls = foreach gCategories generate
```

- Pig Latin has a fully nested data model with:
  - Atomic values, tuples, bags (lists), and maps yahoo, email news

Avoids expensive joins



#### Decouples grouping as an independent



• Efficient Implementation (see paper) Ted Dunning

Chief Scientist, Veoh



#### results

quer y	url	ran k
Laker s	nba.com	1
Laker s	espn.com	2
Kings	nhl.com	1
Kings	nba.com	2

#### revenue

quer y	adSlot	amou nt
Laker s	top	50
Laker s	side	20
Kings	top	30
Kings	side	10

	group	results		revenue			
	Lakers	Laker s	nba.co m	1 }	Lakers	top	50
		Laker s	espn.co m	2	Lakers	side	20
	•			}			,
		Kings	nhl.co . m	1	Kings	top	30
OSS	-pkioosu	Ct Of t Kings	the 2 b	ags '	WOUld ( Kings	give na	atura 10

- Explicit Data Flow Language unlike SQL
- Low Level Procedural Language unlike Map-Reduce
- Quick Start & Interoperability
- Mode (Interactive Mode, Batch, Embedded)
- User Defined Functions
- Nested Data Model

 Map-Reduce and the need for Pig Latin

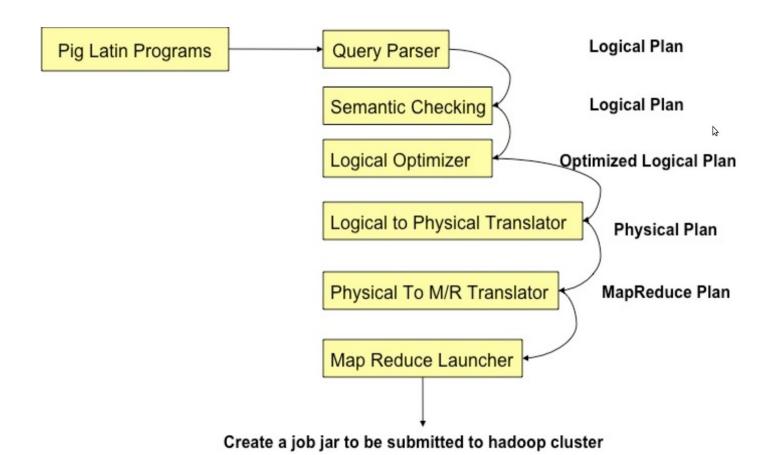
Pig Latin

Compilation into Map-Reduce

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Future Work

## Compilation

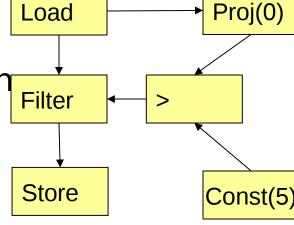




- Type checking with schema
- References verifying
- Logic plan generating
  - One-to-one fashion
  - Independent of execution platform
  - Limited optimization

- Consists of DAG of Logical Operators as nodes and Data Flow represented as edges
  - Logical Operators contain list of i/p's o/p's and schema
- Logical operators
  - Aid in post parse stage checking (type checking)
  - Optimization
  - Translation to Physical Plan

```
a = load 'myfile';
b = filter a by $0 > 5;
store b into 'myfilteredfile';
```



- Layer to map Logical Plan to multiple back-ends, one such being M/R (Map Reduce)
  - Chance for code re-use if multiple back-ends share same operator
- Consists of operators which pig will run on the backend
- Currently most of the physical plan is placed as operators in the map reduce plan
- Logical to Physical Translation
  - 1:1 correspondence for most Logical operators
  - Except Cross, Distinct, Group, Co-group and Order

## Togic Plan

A=LOAD 'file1' AS (x, y, z);

B=LOAD 'file2' AS (t, u, v);

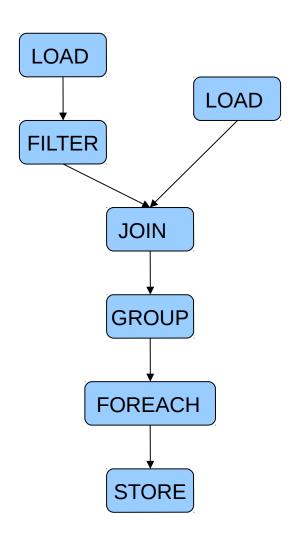
C=FILTER A by y > 0;

D=JOIN C BY x, B BY u;

E=GROUP D BY z;

F=FOREACH E GENERATE group, COUNT(D);

STORE F INTO 'output';



# Physical Plan

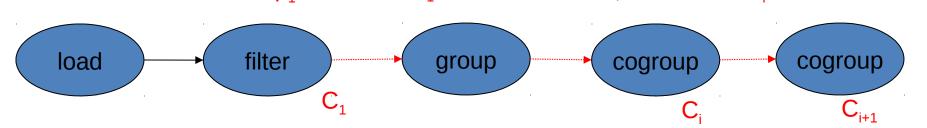
- 1:1 correspondence with most logical operators
- Except for:
  - DISTINCT
  - (CO)GROUP
  - JOIN
  - ORDER

### bgical to Physical Plan for Group operator

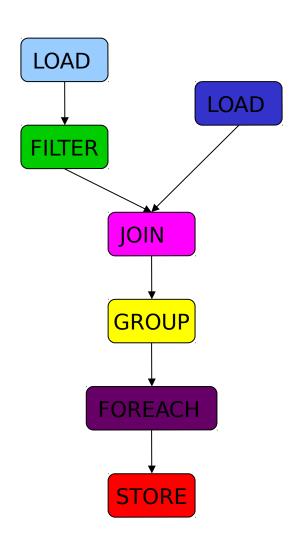
Logical operator for co-group/group is converted to 3 Physical opertors, {(2,Y)2}} - Local Rearrange (LA)1,R)1, (1,B)2}} - Global Rearrange (GR)<sup>2, (2,G)1</sup>} {Key,{ListofValues}} GR  $\{1,(1,R)\}^{1}$ Package (PKG)  $\{2,(2,G)\}^1$  $\{1,(1,B)\}^2$ • Example: {Key,(Value)}(table no)  ${2,(2,Y)}^2$ LR Co-group A by Acol1, B by Bcol1 Tuples Acol1 - (2,G)Bcol1-

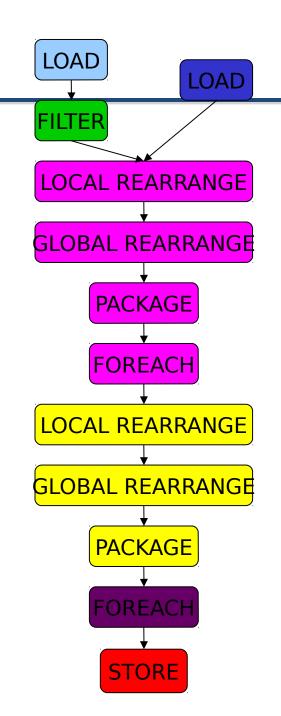
### Map Reduce Plan

- Physical to Map Reduce (M/R) Plan conversion happens through the MRCompiler
  - Converts a physical plan into a DAG of M/R operators
- Boundaries for M/R include cogroup/group, distinct, cross, order by, limit (in some cases)
  - Push all subsequent operators between cogroup to next cogroup into reduce
  - order by is implemented as 2 M/R jobs
- JobControlCompiler then uses the M/R plan to construct a JobControl object map, reduce,

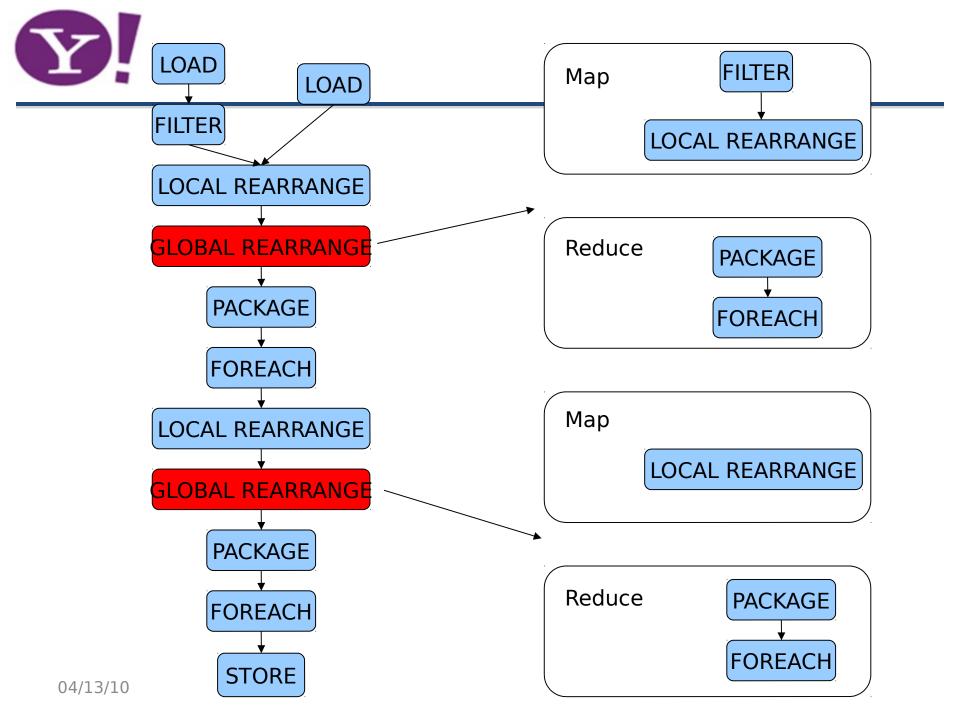








- Determine MapReduce boundaries
  - GLOBAL REARRANGE
- Some operations are done by MapReduce framework
- Coalesce other operators into Map & Reduce stages
- Generate job jar file

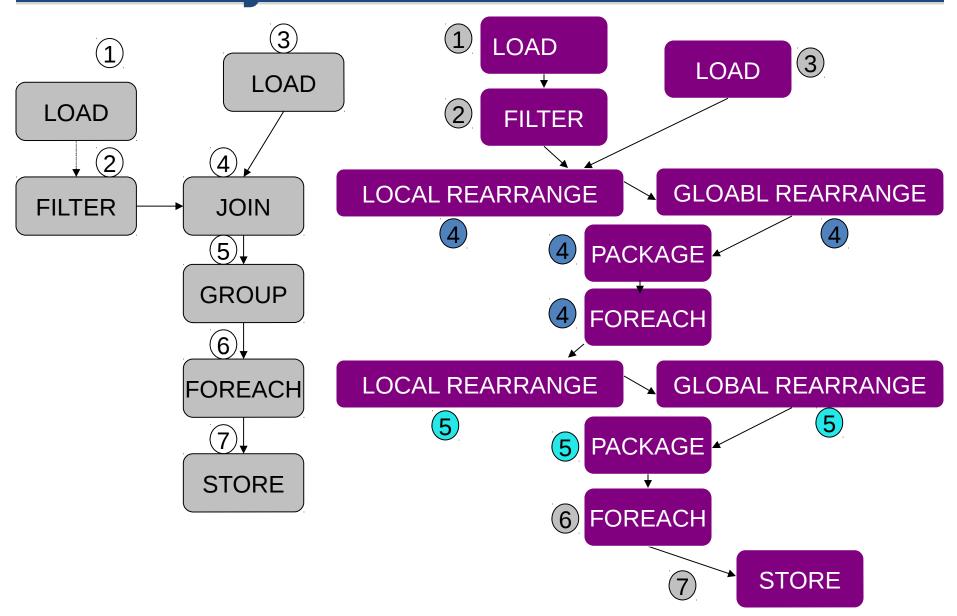


### Pig Latin to Physical Plan

```
A = LOAD 'file1' AS
  (x,y,z);
                                             LOAD
B = LOAD 'file2' AS
                            LOAD
  (t,u,v);
                                         JOIN
                            FILTER
C = FILTER A by y > 0;
                                                    x,y,z,t,u,v
                           X, Y, Z
D = JOIN C by x,B by u;
                                         GROUP
E = GROUP D by z;
                                                   group, count
                                         FOREACH
F = FOREACH E
  generate group,
  COUNT(D);
                                         STORE
STORE F into 'output';
```

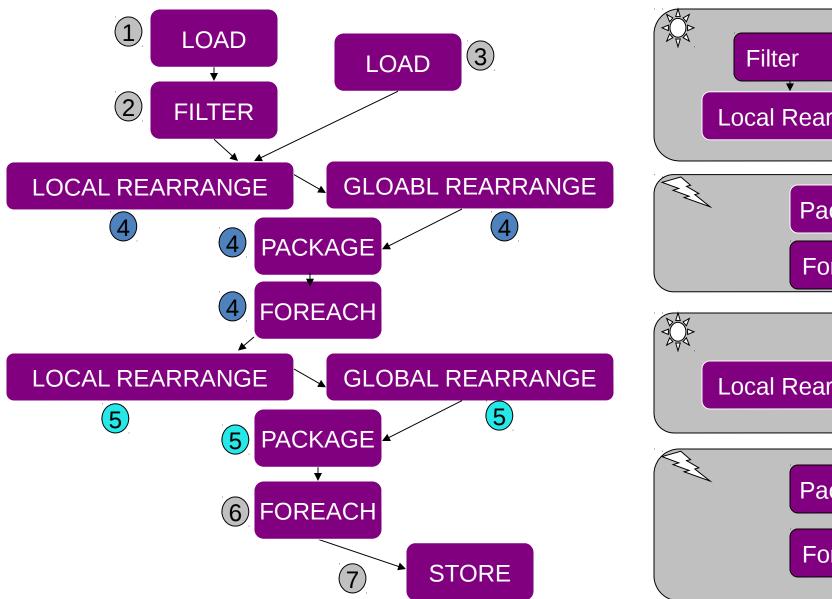


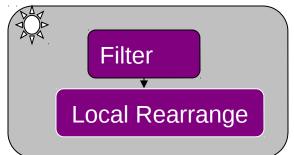
#### Logical Plan to Physical Plan

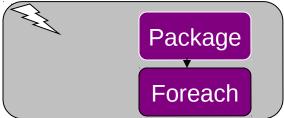


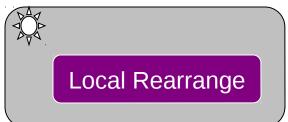


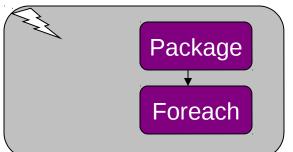
#### Physical Plarto Map-Reduce Plan





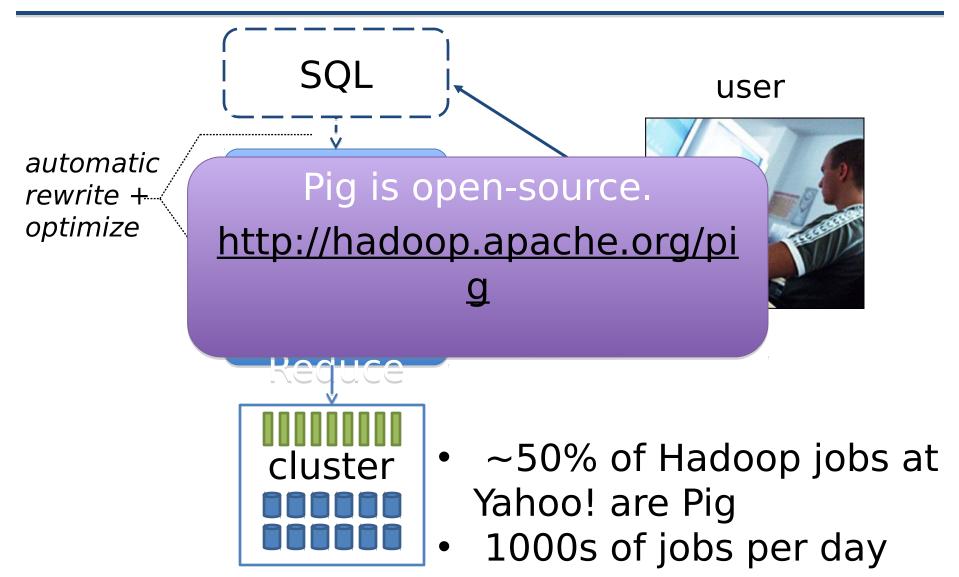






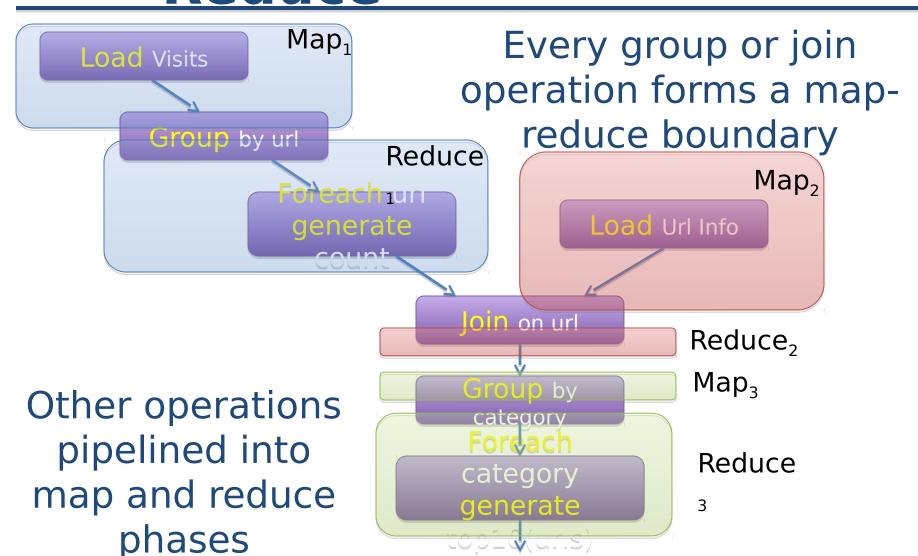


#### **Implementation**



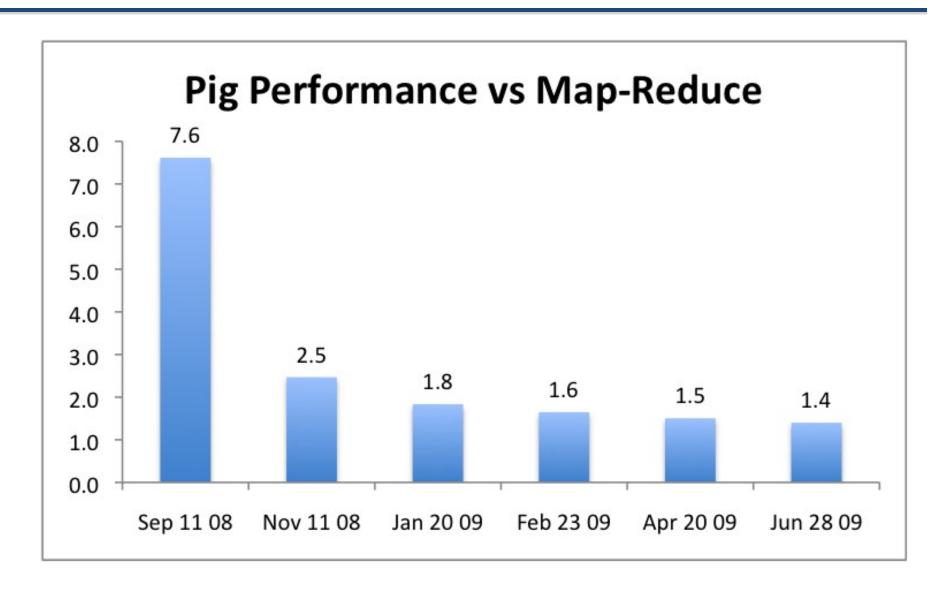


#### Compilation into Map-Reduce

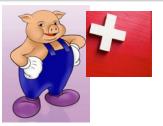




#### **Performance**



## Strong & Weak Points





The step-by-step method of creating a program in Pig is much cleaner and simpler to use than the single block method of SQL. It is easier to keep track of what your variables are, and where you are in the process of analyzing your data.

Scalability

Memory

Jasmine Novak Engineer, Yahoo

With the various interleaved clauses in SQL It is difficult to know what is actually happening sequentially.

on for

#### Processing

Open Source

Non Java Pisie Ciemiewicz Kellence, Yaho

LimitedOptimization



- Big demand for parallel data processing
  - Emerging tools that do not look like SQL DBMS
  - Programmers like dataflow pipes over static files
- Hence the excitement about Man-Pig Latin
   Sweet spot between map-reduce and SQL
- But, Map-Reduce is too low-level and rigid