

SYSTEM PROGRAMMING

TOPIC 6: PROCESS RELATIONSHIPS

Seongjin Lee

Updated by: Muhammad Faris Fathoni

Updated: 2018/11/08

06-process_rel

insight@gnu.ac.kr

<http://open.gnu.ac.kr>

Systems Research Lab.

Gyeongsang National University



Table of contents

1. Terminal Logins
2. Process Groups and Sessions
3. Controlling Terminal and Job Controls



Overview

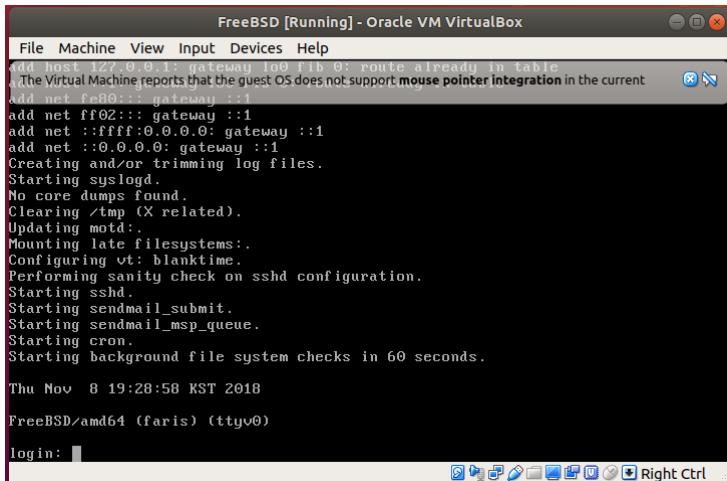
- Terminal Logins
- Sessions
- Controlling Terminal
- Job Control



TERMINAL LOGINS



BSD Terminal



```
FreeBSD [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
add host 127.0.0.1: gateway 100 fib 0: route already in table
The Virtual Machine reports that the guest OS does not support mouse pointer integration in the current
add net fe80::: gateway ::1
add net ff02::: gateway ::1
add net ::ffff:0.0.0.0: gateway ::1
add net ::0.0.0.0: gateway ::1
Creating and/or trimming log files.
Starting syslogd.
No core dumps found.
Clearing /tmp (X related).
Updating motd:.
Mounting late filesystems:.
Configuring vt: blanktime.
Performing sanity check on sshd configuration.
Starting sshd.
Starting sendmail_submit.
Starting sendmail_msp_queue.
Starting cron.
Starting background file system checks in 60 seconds.

Thu Nov  8 19:28:58 KST 2018

FreeBSD/amd64 (faris) (ttyv0)

login: 
```



Terminal Logins

In early UNIX systems

- Users logged in using dumb terminals that were connected to the host with hard-wired connections



Terminal Logins

In early UNIX systems

- Users logged in using dumb terminals that were connected to the host with hard-wired connections
- The terminals were either local or remote



Terminal Logins

In early UNIX systems

- Users logged in using dumb terminals that were connected to the host with hard-wired connections
- The terminals were either local or remote
- Users login through a terminal device driver in the kernel



Terminal Logins

In early UNIX systems

- Users logged in using dumb terminals that were connected to the host with hard-wired connections
- The terminals were either local or remote
- Users login through a terminal device driver in the kernel
- A host had a fixed number of terminal devices



Terminal Logins

In early UNIX systems

- Users logged in using dumb terminals that were connected to the host with hard-wired connections
- The terminals were either local or remote
- Users login through a terminal device driver in the kernel
- A host had a fixed number of terminal devices

type who in the shell

James console Oct 12 15:36

James ttys000 Oct 13 22:02

James ttys001 Oct 13 22:08

```
root@faris:~ # who
faris_sample _ ttys000 Nov 8 19:36
```



BSD Terminal Logins

Mac OS X and Linux login procedure follows essentially the same steps as the BSD

The system administrator creates `/etc/ttys`, `ttys(5)`, that has one line per terminal device

Each line specifies the name of the device and other parameters that are passed to the `getty(8)` program



BSD Terminal Logins cont'd

1. the kernel creates process ID 1, the init process

process ID 1



BSD Terminal Logins cont'd

1. the kernel creates process ID 1, the `init` process
2. the `init` process reads the file `/etc/ttys`
3. creates empty environment

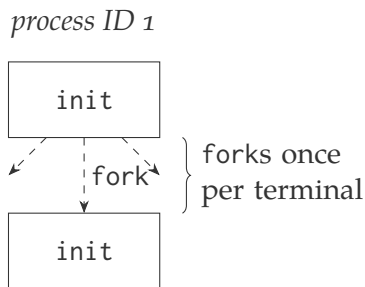
process ID 1

`init`



BSD Terminal Logins cont'd

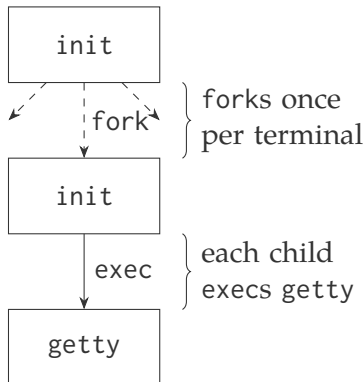
1. the kernel creates process ID 1, the init process
2. the init process reads the file `/etc/ttys`
3. creates empty environment
4. forks for every terminal device



BSD Terminal Logins cont'd

1. the kernel creates process ID 1, the init process
2. the init process reads the file /etc/ttys
3. creates empty environment
4. forks for every terminal device
5. followed by exec of the program getty
6. opens terminal device (fd 0, 1, 2)
7. reads user name
8. initial environment set

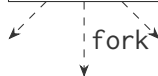
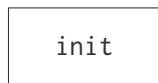
process ID 1



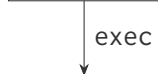
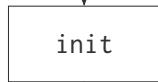
BSD Terminal Logins cont'd

1. the kernel creates process ID 1, the init process
2. the init process reads the file /etc/ttys
3. creates empty environment
4. forks for every terminal device
5. followed by exec of the program getty
6. opens terminal device (fd 0, 1, 2)
7. reads user name
8. initial environment set
9. followed by exec of the program login

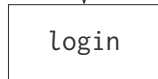
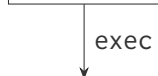
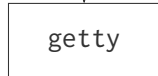
process ID 1



} forks once per terminal



} each child execs getty



BSD Terminal Logins

init(8)

PID 1

PPID 0

EUID 0

reads /etc/ttys



BSD Terminal Logins

init(8)	PID 1	PPID 0	EUID 0
reads /etc/ttys			
getty(8)	PID N	PPID 0	EUID 0



BSD Terminal Logins

init(8)	PID 1	PPID 0	EUID 0
reads /etc/ttys			
getty(8)	PID N	PPID 0	EUID 0
opens terminal			
prints "login:"			
read username			



BSD Terminal Logins

init(8)	PID 1	PPID 0	EUID 0
reads /etc/ttys			
getty(8)	PID N	PPID 0	EUID 0
opens terminal			
prints "login:"			
read username			
login(1)	PID N	PPID 0	EUID 0
getpass(3), encrypt, compare with getpwnam(3)			
register login in system database			
read/display various files			
initgroups(3)/setgid(2), initialize environment			
chdir(2) to home directory			
chown(2) terminal device			
setuid(2) to user's uid, exec(3) shell			



BSD Terminal Logins

init(8)	PID 1	PPID 0	EUID 0
reads /etc/ttys			
getty(8)	PID N	PPID 0	EUID 0
opens terminal			
prints "login:"			
read username			
login(1)	PID N	PPID 0	EUID 0
getpass(3), encrypt, compare with getpwnam(3)			
register login in system database			
read/display various files			
initgroups(3)/setgid(2), initialize environment			
chdir(2) to home directory			
chown(2) terminal device			
setuid(2) to user's uid, exec(3) shell			
\$SHELL	PID N	PPID 0	EUID U

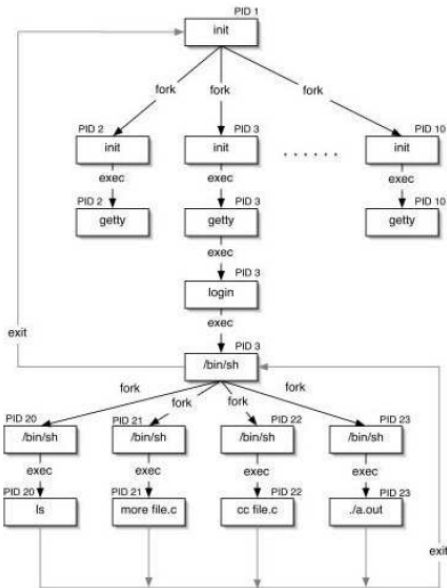


BSD Terminal Logins

init(8)	PID 1	PPID 0	EUID 0
reads /etc/ttys			
getty(8)	PID N	PPID 0	EUID 0
opens terminal			
prints "login:"			
read username			
login(1)	PID N	PPID 0	EUID 0
getpass(3), encrypt, compare with getpwnam(3)			
register login in system database			
read/display various files			
initgroups(3)/setgid(2), initialize environment			
chdir(2) to home directory			
chown(2) terminal device			
setuid(2) to user's uid, exec(3) shell			
\$SHELL	PID N	PPID 0	EUID U
ls(1)	PID M	PPID N	EUID U



BSD Terminal Login Process



BSD Terminal

```
FreeBSD 11.2-RELEASE (GENERIC) #0 r335510: Fri Jun 22 04:32:14 UTC 2018

Welcome to FreeBSD!

Release Notes, Errata: https://www.FreeBSD.org/releases/
Security Advisories:  https://www.FreeBSD.org/security/
FreeBSD Handbook:     https://www.FreeBSD.org/handbook/
FreeBSD FAQ:          https://www.FreeBSD.org/faq/
Questions List: https://lists.FreeBSD.org/mailman/listinfo/freebsd-questions/
FreeBSD Forums:       https://forums.FreeBSD.org/

Documents installed with the system are in the /usr/local/share/doc/freebsd/
directory, or can be installed later with:  pkg install en-freebsd-doc
For other languages, replace "en" with a language code like de or fr.

Show the version of FreeBSD installed:  freebsd-version ; uname -a
Please include that output and any error messages when posting questions.
Introduction to manual pages:  man man
FreeBSD directory layout:      man hier

Edit /etc/motd to change this login announcement.
If you accidentally end up inside vi, you can quit it by pressing Escape, colon
(:), q (q), bang (!) and pressing return.
farris sample@farris:~ %
```



PROCESS GROUPS AND SESSIONS

Process Groups

Each process belongs to a process group

- it is a collection of one or more processes
- usually associated with the same job
- the group has a unique process group ID
- the process group exist as long as one process is in the group

```
#include <unistd.h>
pid_t getpgrp(void);
pid_t getpgid(pid_t pid);
// Returns: process group ID of calling process
```



Process group cont'd

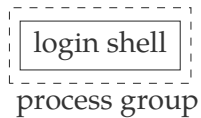
A process joins an existing process group or creates a new process group by calling `setpgid`

```
#include <unistd.h>
int setpgid(pid_t pid, pid_t pgid);
// Returns: 0 if OK, 1 on error
```

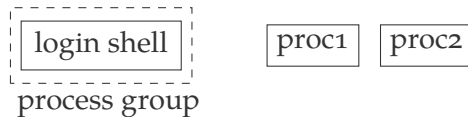
Each process group can have a process group leader

- Process group leader identified by its pgid (if `pgid == pid`)
- Leader can create a new process group, create processes in the group
- if `pid == 0`, caller process ID is used
- if `pgid == 0`, group ID == `pid`
- A process can set the process group ID of itself or its children





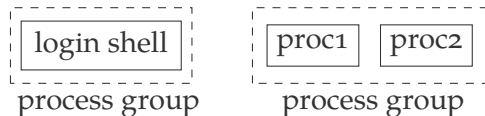




Sessions

The processes in a process group are usually placed there by a shell pipeline

```
proc1 | proc2 &
```

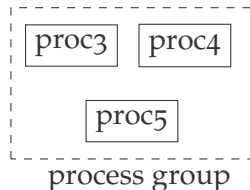
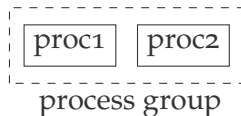
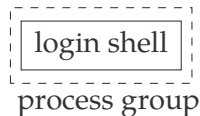


Sessions

The processes in a process group are usually placed there by a shell pipeline

```
proc1 | proc2 &
```

```
proc3 | proc4 | proc5
```

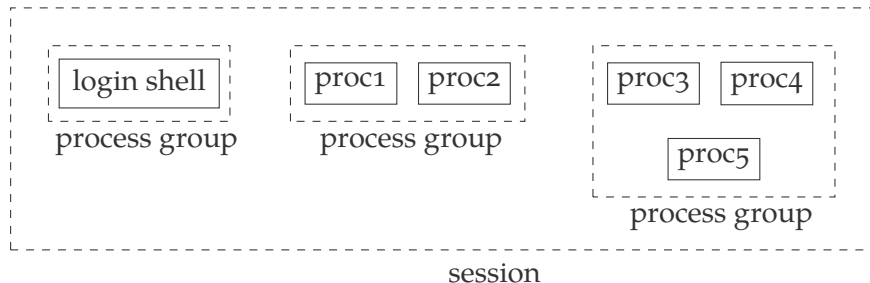


Sessions

The processes in a process group are usually placed there by a shell pipeline

```
proc1 | proc2 &
```

```
proc3 | proc4 | proc5
```



Session is a collection of one or more process groups



Sessions cont'd

A process establishes a new session by calling the `setsid` function

```
#include <unistd.h>
pid_t setsid(void);
// Returns: process group ID if OK, 1 on error
```

If the calling process is not a process group leader, this function creates a new session. Three things happen



A process establishes a new session by calling the `setsid` function

```
#include <unistd.h>
pid_t setsid(void);
// Returns: process group ID if OK, 1 on error
```

If the calling process is not a process group leader, this function creates a new session. Three things happen

1. process becomes the *session leader*, and is only process in this new session



A process establishes a new session by calling the `setsid` function

```
#include <unistd.h>
pid_t setsid(void);
// Returns: process group ID if OK, 1 on error
```

If the calling process is not a process group leader, this function creates a new session. Three things happen

1. process becomes the *session leader*, and is only process in this new session
2. the process becomes the process group leader (`pgid == pid`)
 - if the caller is already a process group leader, then returns an error



A process establishes a new session by calling the `setsid` function

```
#include <unistd.h>
pid_t setsid(void);
// Returns: process group ID if OK, 1 on error
```

If the calling process is not a process group leader, this function creates a new session. Three things happen

1. process becomes the *session leader*, and is only process in this new session
2. the process becomes the process group leader (`pgid == pid`)
 - if the caller is already a process group leader, then returns an error
3. No controlling terminal



Sessions cont'd

getsid function returns the process group ID of a process's session leader

```
#include <unistd.h>
pid_t getsid(pid_t pid);
// Returns: session leader's process group ID if OK, 1 on error
```

if `pid == 0`, it returns the pgid of calling process's session leader



CONTROLLING TERMINAL AND JOB CONTROLS

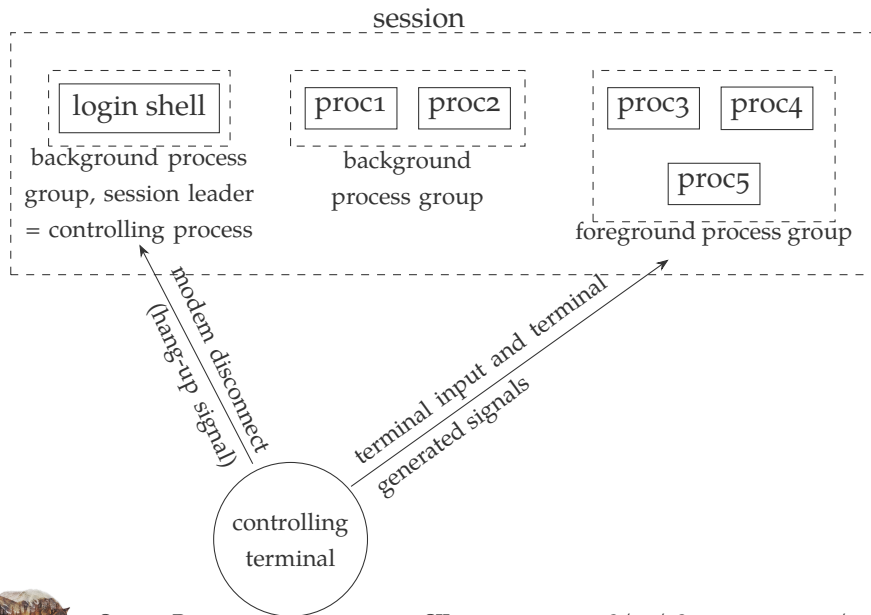


Controlling terminal

- Session can have a single controlling terminal
- session leader that connects to controlling terminal is *controlling process*
- process groups are divided into a single *foreground process group* and one or more *background process groups*
- interrupt signals are sent to foreground process group



Controlling Terminal



Job Control

We can start a job in either the foreground or the background

Examples:

foreground `vi main.c` → starts a job in the foreground

background `make all` → start a job in the background

```
$ make all > Make.out &
```

```
[1] 1475
```

```
$ pr *.c | lpr &
```

```
[2] 1490
```

```
$ just press RETURN
```

```
[2] + Done pr *.c | lpr &
```

```
[1] + Done make all > Make.out &
```



Job Control cont'd

The foreground jobs are affected by some special characters, which generate signals

- Interrupt character (typically Ctrl + C) generates SIGINT
- Quit character (typically Ctrl + backslash) generates SIGQUIT
- Suspend character (typically Ctrl + Z) generates SIGTSTP
- Pause character (typically Ctrl + S) generates SIGSTOP
- Resume character (typically Ctrl + Q) generates SIGCONT



Job Control cont'd

```
1 $ cat > temp.foo &  
   [1] 1681  
2 $  
   [1] + Stopped (SIGTTIN) cat > temp.foo &  
3 $ fg %1  
4 cat > temp.foo  
5 hello, world  
6 ^D  
7 $ cat temp.foo  
   hello, world
```

1. start in background, but it'll read from standard input
2. we press RETURN
3. bring job number 1 into the foreground
4. the shell tells us which job is now in the foreground
5. enter one line
6. type the end-of-file character
7. check that the one line was put into the file

