**ECE564 ASIC & FPGA Design**

**Final Project**

**Name:** Soumil Krishnanand Heble

**UnityID:** sheble

**StudentID:** 200207079

**Logic Area (um2):** 15852.5362

**Memory:** N/A

**1/(delay.area) (ns-1.um-2)**

7.58E-8

Delay (ns to run provided example).

**Clock period:** 8 ns

**# Clock Cycles:** 104 (Message: Hello)

**1/(delay.area) (TA)**

**Delay (TA provided example. TA to complete)**

**Abstract:** Designed a simplified SHA256 hashing module using synthesizable Verilog RTL and achieved a cell area of 15852.5362 um2 using 45 nm Nandgate OpenCell library. The designed module reads in the message (ASCII) to be hashed, H and K values from an SRAM and writes the result to an output SRAM. The module starts its computation when a go signal is asserted and after writing the result to the output SRAM a finish signal is asserted indicating the hash value has been computed and written to memory. The module is also supplied with the number of characters in the message before the go signal is asserted. The simplification in this implementation of the SHA256 comes from the fact that the message length is restricted to 55 characters leading to the creation of only one 512-bit block. On an average the design takes 121 cycles @ 125 MHz and all the inputs and outputs are registered leading to glitch free IO.

**Simplified SHA256 Hashing Algorithm**

Soumil Krishnanand Heble

**Introduction**

**Hardware being designed:**

* A hardware implementation of the SHA256 algorithm that takes in a message of maximum 55 characters (ASCII) from an SRAM and writes the computed has to an output SRAM. The message being limited to 55 characters results in the creation of only one 512-bit message block.
* The hardware module reads the H and K values from SRAMs and starts the hash computation after a go signal is asserted. The module also takes in the number of characters in the message SRAM.
* Upon completion of the hash value write to the output SRAM a finish signal is asserted.

**Summary of Key Innovations:**

* Since the message is limited to 55 characters, only 447 flip-flops are required to create the initial padded block as per SHA2 specifications.
  + Assuming a message of 55 characters
    - 55 ASCII characters = 440 bits
    - Append 0x80 = 1 bit (Only requires the MSB rest can be logic 0)
    - Zero Padding = 0 bit (logic 0 – no flip flops required)
    - Message Length = 6 bits (Length = 440, of the 16 bits used for length only requires 6 bits can be either 1 or 0 rest can be logic 0)
    - Total Flip-Flops = 447
* Since the W values are required for the message digest calculation sequentially the values are calculated and supplied on the fly to the hash computation module. This saves the need to use 64 32-bit wide registers. This requirement also means that as soon as the 512 bit padded message block is ready the hash computation can be started and can run parallel to the W value calculation.
* The W computation is pipelined and the usage of parallel prefix adders from the Designware library allows the clock period to be squeezed down to 6 ns.
* The sequence of operations required for H computation is parallelized and uses parallel prefix adders from the Designware library.
* The above design choices result in a cell area of 15852.5362 um2.

**Summary of Results Achieved:**

* Area
  + 15852.5362 um2
* Number of Cycles
  + Message\_1 = 100, Length = 1
  + Message\_5 = 104, Length = 5
  + Message\_27 = 126, Length = 27
  + Message\_55 = 154, Length = 55
  + \*\*\* Average = 121
* Clock Period
  + 8 ns
* Performance (Area \* # Cycles \* Clock Period, units = um2 ns)
  + Message\_1 = 12682028.96
  + Message\_5 = 13189310.12
  + Message\_27 = 15979356.49
  + Message\_55 = 19530324.59
  + \*\*\* Average = 15345255.04

**Structure of the Rest of this Report:**

**Micro-Architecture**

**Hardware “Algorithmic” Approach Used:**

While the synthesizable Verilog RTL was being written it was developed as three separate modules namely gen\_padded, gen\_w and gen\_h but for the final submission the code from all the three modules was combined into a single synthesizable Verilog RTL file named MyDesign.v. This was done so that the interfaces as specified in the project description can be maintained.

The three modules if chained can perform the same operation and result in the same area, performance and their IO’s are registered as well. The code fragments from the respective modules start with the comment “/\*>>>>> <module\_name> \*\*\*\*\*\*/” with the module names being either of gen\_padded, gen\_w or gen\_h. Since the IOs are registered the delay in the signal/data/address arrival is handled by using state machines and delay using register chains in all the three modules.

The algorithmic approach explained is split in three as per the above modules:

**gen\_padded:** Creates the padded message required by SHA2 algorithm by reading in data from the message SRAM.

* **State Machine:**
  + **State P0 (IDLE):** Wait in this state until the go signal is asserted.
  + **State P1 (LATCH\_MESSAGE\_LEN):** Reads in the message length provided and clear the required bits in the 512 bit padded message register then move to the next state.
  + **State P2 (READ\_SRAM):** Present address to SRAM and write message length to the required position in the 512 bit padded message register. Move to the next state once the SRAM request address equals the message length.
  + **State P3 (DELAY):** Delay state so that the data from SRAM is written (This state is required due to registering IOs). Move to the next state unconditionally.
  + **State P4 (WRITE\_0x80):** Append 0x80 after the message is read and move to the next state unconditionally.
  + **State P5 (FINISH\_WAIT):** Assert a pad register ready signal to the gen\_w module and wait for the next go and finish to be high then start the cycle again from State P1.
* The address presented to the SRAM and the enable signal is delayed internally by using chained registers and used to index into the pad register. The delayed enable is used as a write enable for the pad register. This saves from using additional counter and a larger FSM for handling the delay of registered IO.

**gen\_w:** Reads in the 512 bit padded message from gen\_padded module and computes and presents the W values upon request from gen\_h module (hash computation module). The gen\_h module has an SRAM like interface to the gen\_w module to request for W values. It has to present the gen\_w module with an address and enable signal to read out the values.

* **State Machine:**
  + **State W0 (IDLE):** Wait in this state until the pad message ready signal is asserted by the gen\_padded module.
  + **State W1 (READ\_IN\_PAD):** Read in the padded message in to the 16 32-bit register for W value computation then move to next state.
  + **State W2 (SERVICE\_REQ):** Wait for W value request from gen\_h module and service the value if the requested address matched the current value to be serviced count. Move to State W3 if the current value to be serviced counter overflows.
  + **State W3 (WAIT\_!PAD\_RDY):** Wait for the pad ready signal to be de-asserted then move to State W0.
* The W value calculation is pipelined and happens on the fly. First 16 W values are ready as soon as the padded message is read in to the 16 32-bit registers. The W values to be read for the next value to be computed (initially the 17th value) is picked from the 1st, 2nd, 9th and 15th locations from the 16 32-bit registers. The 16 32-bit registers act as a FIFO queue upon receiving a read request from the gen\_h module the value in the top or 1st register is placed on the data line and the data in the following registers are moved up one position.
* The W computation is pipelined and the results are registered into the base of the 16 32-bit registers. The pipelining is discussed in detail in the next section.

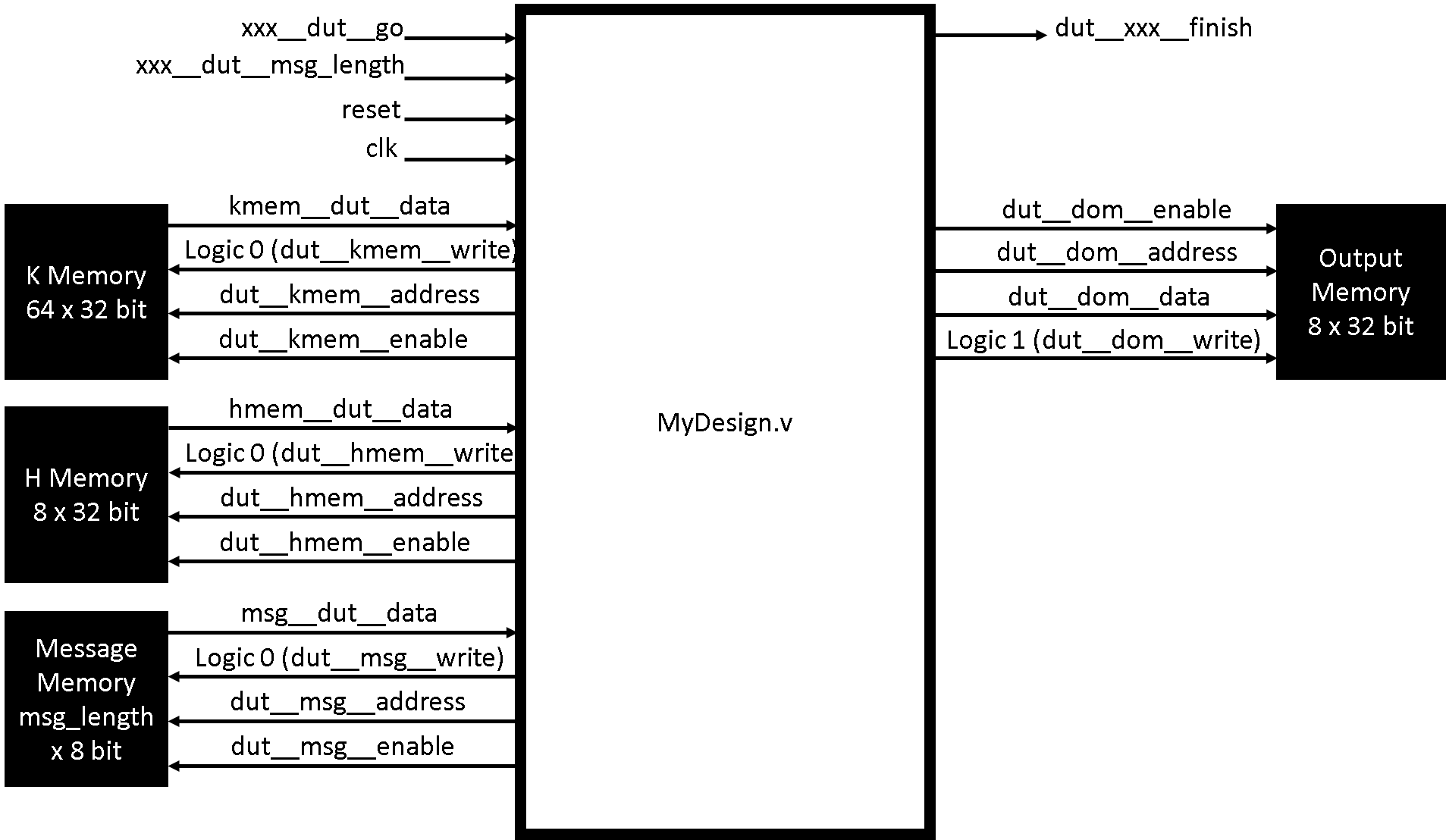
**gen\_h:** Reads a-h values from H SRAM, W from gen\_w module and constants from the K SRAM and computes the hash of the message and writes it to the output SRAM.

* **State Machine:**
  + **State M0 (IDLE):** Wait in this state until the go signal is asserted.
  + **State M1 (READ\_AH):** Copy the a-h values from H SRAM into the 8 32-bit registers for has computation. When the address counter feeding address to the SRAM overflows move to the next state.
  + **State M2 (WAIT\_FOR\_W):** Wait for the W module to read in the padded message and get ready to service the address request. Move to the next state once W module is ready.
  + **State M3 (READ\_WK\_0):** Start reading the W values from the W module and the K values from the K SRAM.Move to the next state unconditionally.
  + **State M4 (WAIT\_0):** Wait for the request to propagate and W and K data to arrive. Move to the next state unconditionally.
  + **State M5 (WAIT\_1):** Wait for the request to propagate and W and K data to arrive. Move to the next state unconditionally.
  + **State M6 (WAIT\_2):** Wait for the request to propagate and W and K data to arrive. Move to the next state unconditionally.
  + **State M7 (H\_COMPUTE):** Perform the H operations. Increment the current iteration counter. Move on to State M8 if the current iteration counter overflows else go to State M7.
  + **State M8 (ADD\_AH):** Read in the a-h values from H SRAM and add them to the a-h hash computation registers.Move to the next state once the SRAM request address counter overflows.
  + **State M9 (WAIT\_ADD\_0):** Wait for the pending additions to complete. Move to the next state unconditionally.
  + **State M10 (WAIT\_ADD\_1):** Wait for the pending additions to complete. Move to the next state unconditionally.
  + **State M11 (WRITE\_OP):** Write the computed hash values to the output SRAM. Move to the next state once the output SRAM write address overflows.
  + **State M12 (WAIT):** Assert finish signal. Move to State M1 if go signal is asserted again.
* The H SRAM enable signals and address are passed to a chained register delay chain and used as write enable and index address to the a-h registers.

**High Level Architecture Drawing and Description of Data Flow**

**Interface Specification**

**Top Level Module Interfaces**

****

**Top Level Module Interface Table and Description**

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| **Direction** | **Type** | **Width (Bus)** | **Name** | **Description** |
| Input | Wire | 1 bit | clk | Clock Signal |
| Input | Wire | 1 bit | reset | Synchronous Reset Signal |
| Input | Wire | 1 bit | xxx\_\_dut\_\_go | Go Pulse To Start SHA256 Computation |
| Input | Wire | 6 bit | xxx\_\_dut\_\_msg\_length | Number of Characters in Message SRAM |
| Output | Reg | 1 bit | dut\_\_xxx\_\_finish | Hash Written to Output SRAM Signal |
| Output | Reg | 1 bit | dut\_\_msg\_\_enable | Enable SRAM |
| Output | Reg | 1 bit | dut\_\_msg\_\_write | SRAM !R/W Select (Always Logic 0) |
| Output | Reg | 6 bit | dut\_\_msg\_\_address | SRAM Access Address |
| Input | Wire | 8 bit | msg\_\_dut\_\_data | Data from SRAM Read Port |
| Output | Reg | 1 bit | dut\_\_kmem\_\_enable | Enable SRAM |
| Output | Reg | 1 bit | dut\_\_kmem\_\_write | SRAM !R/W Select (Always Logic 0) |
| Output | Reg | 6 bit | dut\_\_kmem\_\_address | SRAM Access Address |
| Input | Wire | 32 bit | kmem\_\_dut\_\_data | Data from SRAM Read Port |
| Output | Reg | 1 bit | dut\_\_hmem\_\_enable | Enable SRAM |
| Output | Reg | 1 bit | dut\_\_hmem\_\_write | SRAM !R/W Select (Always Logic 0) |
| Output | Reg | 3 bit | dut\_\_hmem\_\_address | SRAM Access Address |
| Input | Wire | 32 bit | hmem\_\_dut\_\_data | Data from SRAM Read Port |
| Output | Reg | 1 bit | dut\_\_dom\_\_enable | Enable SRAM |
| Output | Reg | 1 bit | dut\_\_dom\_\_write | SRAM !R/W Select (Always Logic 1) |
| Output | Reg | 3 bit | dut\_\_dom\_\_address | SRAM Access Address |
| Output | Reg | 32 bit | dut\_\_dom\_\_data | Data to SRAM Write Port |

**Verification**

The synthesizable Verilog RTL and the synthesized netlist were verified rigorously for maximum possible operating conditions using the provided sample test bench and a custom test bench. The obtained hash values were validated by using the results from a high level language implementation of the module in python for the same input messages.

The provided sample test bench applies a go signal and waits for the finish line to be asserted then repeats the process again one more time. Meanwhile it snoops the output memory write line when the output memory enable line is asserted and writes the data to a result text file for each time the go signal is asserted. Using this test script the hash of four different messages of were computed.

The four different messages tested are:

* Message\_1 = a
* Message\_5 = Hello
* Message\_27 = abcdefghijklmnopqrstuvwxyza
* Message\_55 = abcdefghijklmnopqrstuvwxyzabcdefghijklmnopqrstuvwxyzabc

The second custom test bench tested the synthesizable Verilog RTL for two test cases.

* The first case tests the design by pulsing the go signal while a hash computation is running. This test makes sure that the design does not respond to rogue go pulses while a computation is running.
* The second case asserts the go line indefinitely and lets the computation run twice. This test makes sure that the module is ready for the next computation immediately after it finishes a run.

**Verification Result:** The synthesizable Verilog RTL and synthesized netlist passed all the tests successfully.

Following are some screenshots of the signals of the top module when the sample test bench provided is run with the message “Hello” of length 5:

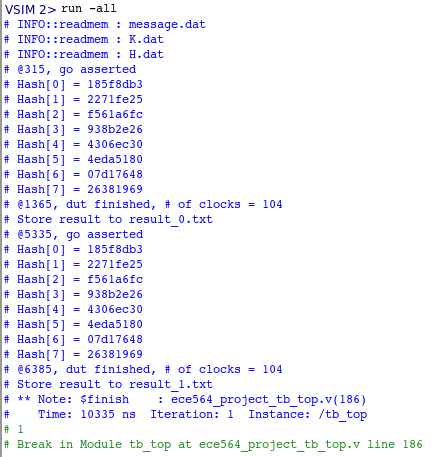


Figure 1 - Modelsim Test bench Run Transcript

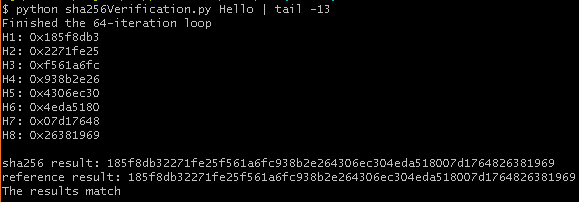


Figure 2 - Hash of the message "Hello" using a High Level Language Implementation of the Module in Python

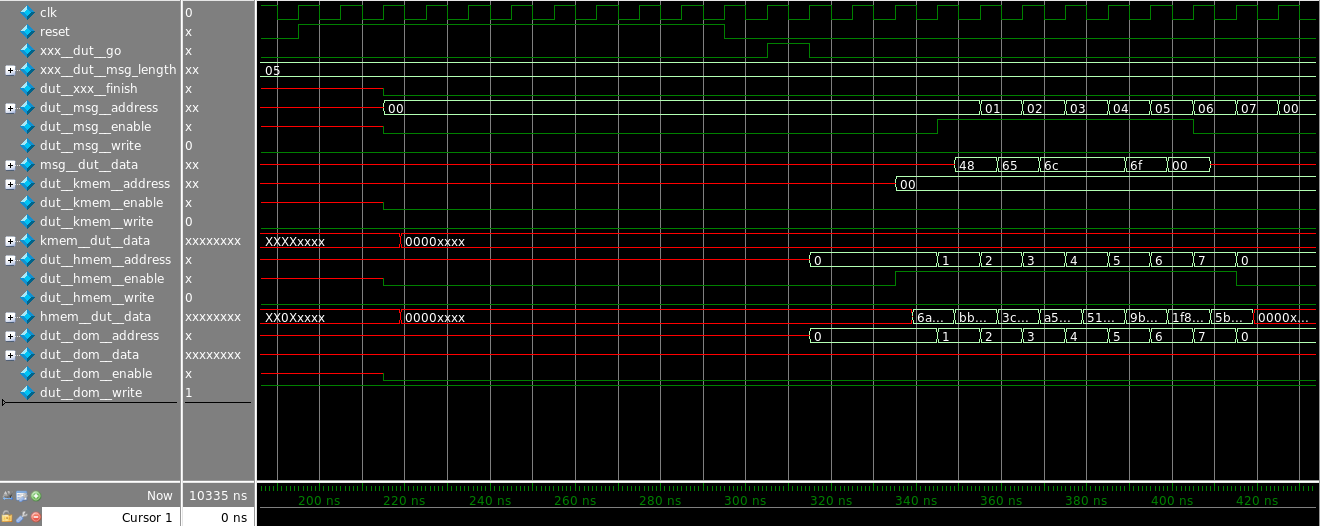
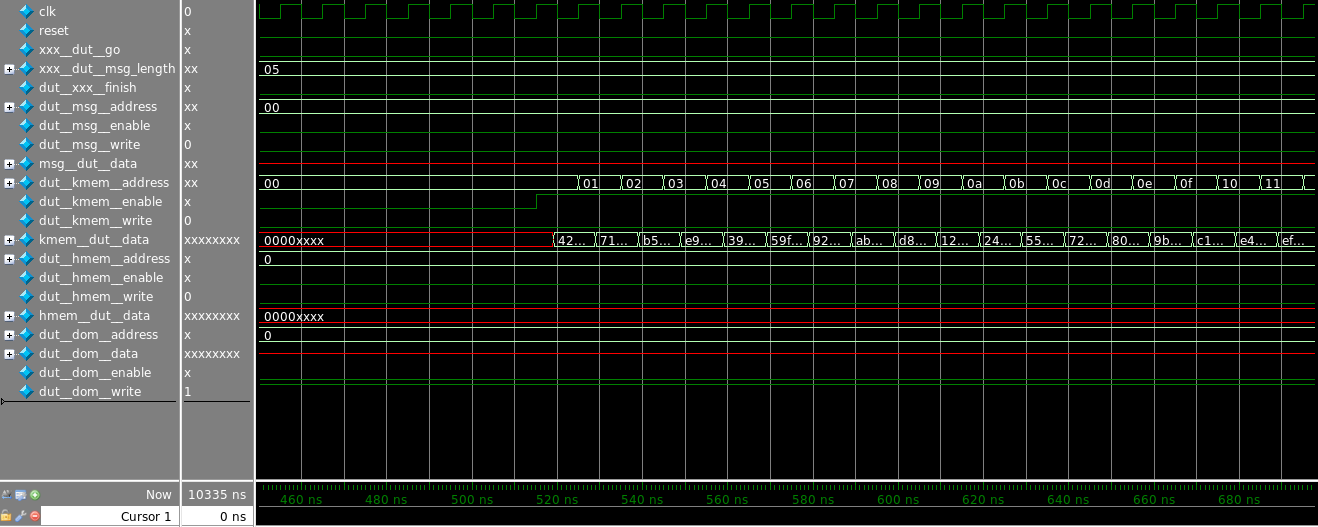


Figure 3 - Top Level Module Signal Waveform II

Figure 4 - Top Level Module Signal Waveform I

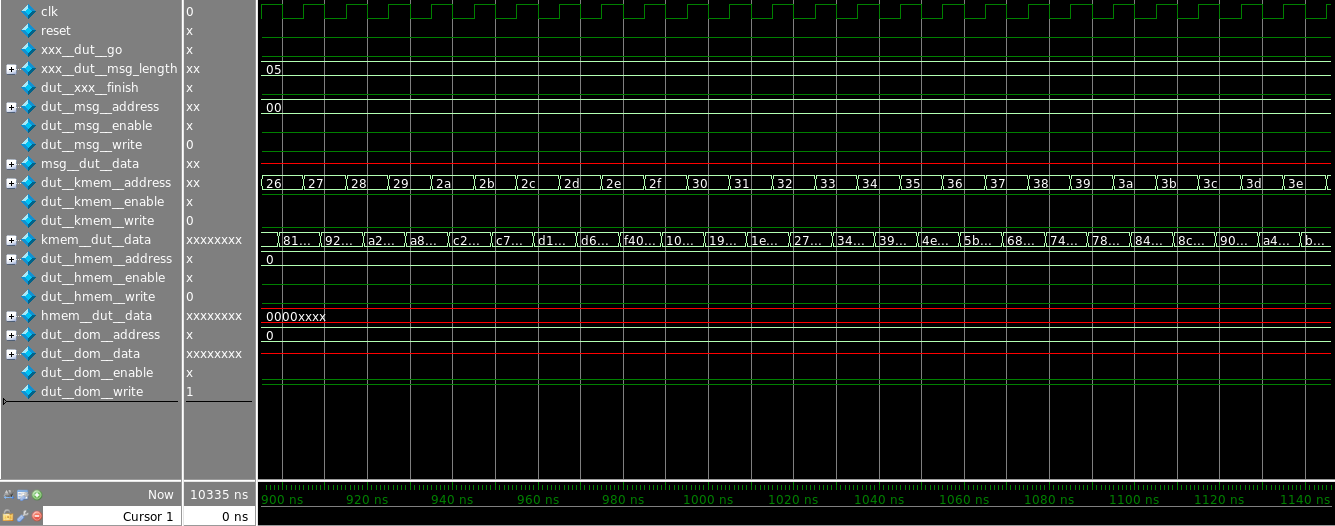


Figure 5 - Top Level Module Signal Waveform IV

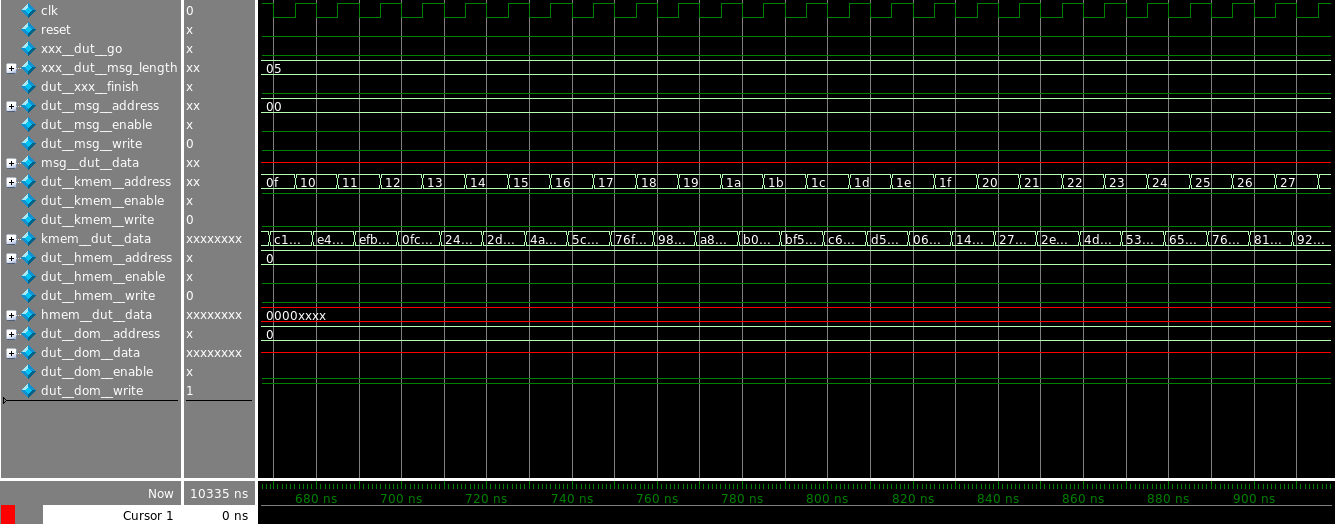


Figure 6 - Top Level Module Signal Waveform III

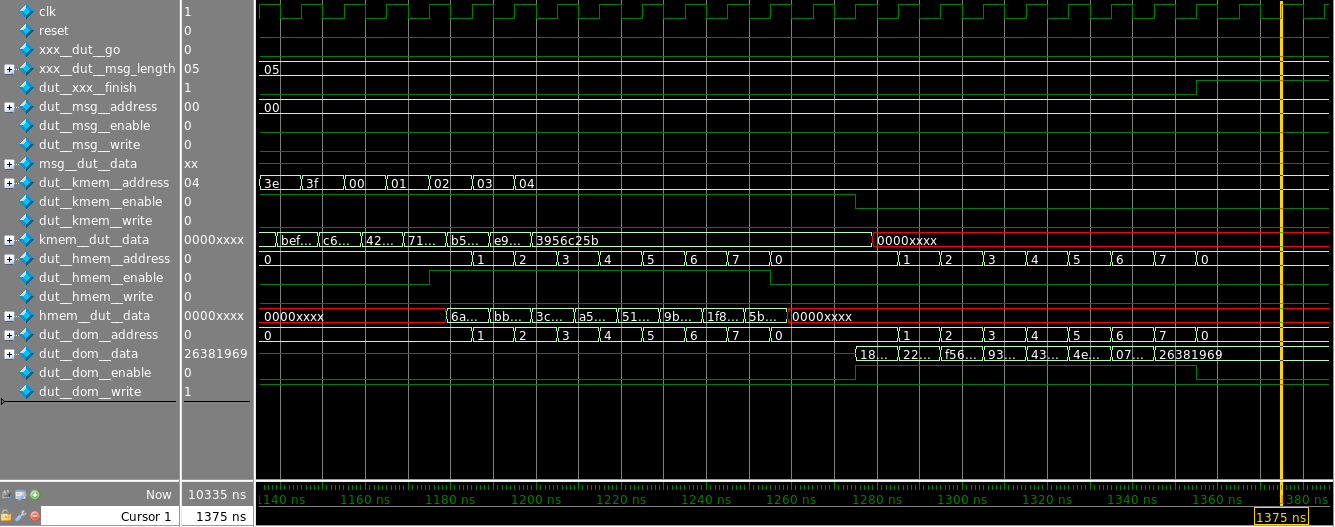


Figure 7 - Top Level Module Signal Waveform V

**Results Achieved**

The synthesizable Verilog RTL was compiled using Synopsys 2015 Design Compiler and used 45nm Nandgate OpenCell Library along with some IP from the Designware library (for blocks such as parallel prefix adders and multiplexers).

The final **area** of MyDesign.v is **15852.5362 um2** (Combinational Logic: um2, Non-Combinational Logic: um2) at a clock period of 8 ns (125 MHz).

* Design Compiler version: Synopsys 2015
* Medium compile effort
* Minimize area while meeting timing requirements
* Clock Skew: 0.05 ns
* The inputs are driven by a DFF: Tcq: 0.2 ns, IP Delay: 0.04 ns
* The outputs drive four DFF: Tsu: 0.25 ns, OP Delay: 0.45 ns

**Setup Violation Check:** 0.0007 ns (Slack **MET**)

**Hold Violation Check:** 0.0188 ns (Slack **MET**)

**Hold fixed Setup Violation Check:** 0.0009 ns (Slack **MET**)

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| **Message Length** | **# Clock Cycles** | **Clock Period (ns)** | **Area (um2)** | **Performance (# Clock Cycles \* Area \* Clock Period, unit um2 ns)** |
| 1 | 100 | 8 | 15852.5362 | 12682028.96 |
| 5 | 104 | 8 | 15852.5362 | 13189310.12 |
| 27 | 126 | 8 | 15852.5362 | 15979356.49 |
| 55 | 154 | 8 | 15852.5362 | 19530324.59 |
| **Average** | **121** | **8** | **15852.5362** | **15345255.04** |

**Conclusion**

The project goal was successfully achieved and the resulting synthesizable Verilog RTL was verified with various messages and the hash values produced were compared with a high level language implementation of the module in python. The RTL was synthesized using Synopsys 2015 Design Compiler and was able to achieve a clock period of 8 ns with a cell area of 15852.5362 um2. The design compiler used parallel prefix adders from the Designware library and 45 nm standard cells from Nandgate OpenCell library. Upon inspection of the synthesized design it was found that the critical path in the design is the hash computation iteration which consists of several chained adders and other logic gates. It is difficult to pipeline the H computation operations since the next computation is dependent on the result of the current iteration.

The secure hashing algorithm (SHA) is used for many applications like authentication, message signing, data integrity check and cryptography. Due to the nature of operations in the algorithm an application specific hardware for SHA computation benefits from the performance boost offered by the intrinsic parallelism of the hardware compared to the execution of the algorithm in a general purpose CPU. Its widely used in many mainstream applications like SSH, TLS etc. hence many modern processors usually have some hardware implementation for SHA operations for e.g. Many of the Intel CPUs have a dedicated cryptoprocessor core for SHA computation and with the Internet of Things gaining popularity among the masses security is becoming paramount hence low power hardware implementations of these cryptographic algorithms are important.