

STAT 309: MATHEMATICAL COMPUTATIONS I
FALL 2015
LECTURE 15

1. MULTIPLE RIGHT-HAND SIDES AND INVERSE

- let $A \in \mathbb{R}^{m \times n}$ and $\mathbf{b}_1, \dots, \mathbf{b}_p \in \mathbb{R}^m$
- suppose we need to solve p linear systems with the same coefficient matrix but different right-hand sides

$$A\mathbf{x}_1 = \mathbf{b}_1, \quad A\mathbf{x}_2 = \mathbf{b}_2, \quad \dots, \quad A\mathbf{x}_p = \mathbf{b}_p \quad (1.1)$$

- this is equivalent to solving the matrix equation

$$AX = B$$

where $X = [\mathbf{x}_1, \dots, \mathbf{x}_p] \in \mathbb{R}^{n \times p}$ and $B = [\mathbf{b}_1, \dots, \mathbf{b}_p] \in \mathbb{R}^{m \times p}$

- for example, this is what we do when we need to compute the inverse of an $n \times n$ nonsingular matrix A :

$$AX = I,$$

which is equivalent to the systems of equations

$$A\mathbf{x}_j = \mathbf{e}_j, \quad j = 1, \dots, n$$

- since only the right-hand side is different in each of these systems, we need only compute the LU factorization of A once
- more generally, this is how we should compute $A^{-1}B$ for matrices A and B , we should solve (1.1) instead of finding the explicit inverse A^{-1} and then multiplying it to B (exercise: what if you need AB^{-1} ?)
- we didn't say too much about why it's a bad idea to compute the explicit inverse of a matrix, for more information about this topic, see Chapter 14 in: N. J. Higham, *Accuracy and Stability of Numerical Algorithms*, 2nd Ed, SIAM, 2002

2. BLOCK FACTORIZATIONS AND SCHUR COMPLEMENT

- a surprisingly simple and powerful idea that appeared implicitly several times in our earlier discussions is that of *block elimination* and *block factorization*
- all it involves is to consider a matrix $A \in \mathbb{R}^{n \times n}$ as a 2×2 block matrix

$$A = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}$$

where $A_{11} \in \mathbb{R}^{p \times p}$, $A_{22} \in \mathbb{R}^{q \times q}$, $A_{12} \in \mathbb{R}^{p \times q}$, $A_{21} \in \mathbb{R}^{q \times p}$ for some p and q where $p + q = n$

- this works for rectangular matrices too but we keep our discussion to square matrices for simplicity
- many of the stuff that we discussed can be carried over to block matrices
- for example, if A_{11} is nonsingular, we could define an $n \times n$ block elimination matrix

$$M_1 = I - U_1 V_1^T$$

where $U_1, V_1 \in \mathbb{R}^{n \times p}$ are

$$U_1 = \begin{bmatrix} 0 \\ A_{21}A_{11}^{-1} \end{bmatrix}, \quad V_1 = \begin{bmatrix} I_p \\ 0 \end{bmatrix}$$

- in other words

$$M_1 = \begin{bmatrix} I_p & 0 \\ 0 & I_q \end{bmatrix} - \begin{bmatrix} 0 \\ A_{21}A_{11}^{-1} \end{bmatrix} \begin{bmatrix} I_p & 0 \end{bmatrix} = \begin{bmatrix} I_p & 0 \\ -A_{21}A_{11}^{-1} & I_q \end{bmatrix}$$

- applying this to A gives

$$M_1 A = \begin{bmatrix} I_p & 0 \\ -A_{21}A_{11}^{-1} & I_q \end{bmatrix} \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} A_{11} & A_{12} \\ 0 & S \end{bmatrix}$$

where

$$S = A_{22} - A_{21}A_{11}^{-1}A_{12}$$

is called the *Schur complement* of A_{11} in A

- we can easily verify that

$$L_1 := M_1^{-1} = \begin{bmatrix} I_p & 0 \\ A_{21}A_{11}^{-1} & I_q \end{bmatrix}$$

- the analogue of LU factorization of A as a 2×2 block matrix

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} L_{11} & 0 \\ L_{21} & L_{22} \end{bmatrix} \begin{bmatrix} U_{11} & U_{12} \\ 0 & U_{22} \end{bmatrix}$$

is called a *block LU factorization*

- note that L_{11} and L_{22} can be any matrices, not necessarily lower triangular, ditto for U_{11} and U_{22}
- multiplying out the RHS, we see that

$$A_{11} = L_{11}U_{11}$$

- it is also easy to see that

$$L_{22}U_{22} = S = A_{22} - A_{21}A_{11}^{-1}A_{12}$$

- we omitted permutation matrices but they can be easily incorporated: for example, if

$$A_{11} = \Pi_1^T L_1 U_1 \Pi_2^T, \quad S = \Pi_3^T L_2 U_2$$

then we have

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} \Pi_1^T & 0 \\ 0 & \Pi_3^T \end{bmatrix} \begin{bmatrix} L_1 & 0 \\ \Pi_3 A_{21} \Pi_2 U_1^{-1} & L_2 \end{bmatrix} \begin{bmatrix} U_1 & L_1^{-1} \Pi_1 A_{12} \\ 0 & U_2 \end{bmatrix} \begin{bmatrix} \Pi_2^T & 0 \\ 0 & I \end{bmatrix}$$

- what we discuss here also apply to LDU , LDL^T , and Cholesky factorizations
- for example if A is symmetric positive definite, then its Cholesky factorization written in 2×2 block form

$$\begin{bmatrix} A_{11} & A_{21}^T \\ A_{21} & A_{22} \end{bmatrix} = A = R^T R = \begin{bmatrix} R_{11}^T & 0 \\ R_{12}^T & R_{22}^T \end{bmatrix} \begin{bmatrix} R_{11} & R_{12} \\ 0 & R_{22} \end{bmatrix} = \begin{bmatrix} R_{11}^T R_{11} & R_{11}^T R_{12} \\ R_{12}^T R_{11} & R_{12}^T R_{12} + R_{22}^T R_{22} \end{bmatrix}$$

is called *block Cholesky factorization*

- again R_{11} and R_{22} need not be upper triangular
- note that since A is symmetric positive definite, so is A_{11} (why?)
- multiplying out the RHS, we see that

$$A_{11} = R_{11}^T R_{11}$$

- it is also easy to see that

$$R_{22}^T R_{22} = A_{22} - A_{21}A_{11}^{-1}A_{21}^T$$

3. MORE ON THE SCHUR COMPLEMENT

- Schur complement is a very useful notion
- in the following we will assume that A is partitioned as in the previous section with A_{11} nonsingular
- the first useful property is that A is nonsingular if and only if S is nonsingular
- a second very useful application is in *solving linear equations by block elimination*, i.e., solving $A\mathbf{x} = \mathbf{b}$ by partitioning it into

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} \begin{bmatrix} \mathbf{x}_1 \\ \mathbf{x}_2 \end{bmatrix} = \begin{bmatrix} \mathbf{b}_1 \\ \mathbf{b}_2 \end{bmatrix} \quad (3.1)$$

where $\mathbf{b}_1 \in \mathbb{R}^p$, $\mathbf{b}_2 \in \mathbb{R}^q$

- plugging the first equation

$$\mathbf{x}_1 = A_{11}^{-1}(\mathbf{b}_1 - A_{12}\mathbf{x}_2) \quad (3.2)$$

into the second equation yields

$$(A_{22} - A_{21}A_{11}^{-1}A_{12})\mathbf{x}_2 = \mathbf{b}_2 - A_{21}A_{11}^{-1}\mathbf{b}_1 \quad (3.3)$$

- this allows us to solve $A\mathbf{x} = \mathbf{b}$ as follows
 - form $A_{11}^{-1}A_{12}$ and $A_{11}^{-1}\mathbf{b}_1$ by solving a system with multiple right hand sides
 - form $S = A_{22} - A_{21}A_{11}^{-1}A_{12}$ and $\tilde{\mathbf{b}} = \mathbf{b}_2 - A_{21}A_{11}^{-1}\mathbf{b}_1$
 - solve $S\mathbf{x}_2 = \tilde{\mathbf{b}}$ for \mathbf{x}_2
 - solve $A_{11}\mathbf{x}_1 = \mathbf{b}_1 - A_{12}\mathbf{x}_2$ for \mathbf{x}_1
- this would be very useful if A_{11} is an ‘easy to invert’ matrix, e.g., A_{11} is diagonal, banded, orthogonal, Toeplitz, sparse, etc
- such situations where the ‘top left corner’ of a matrix A has special structure arise more often than you think, especially in
 - numerical optimization (KKT matrix — A_{11} corresponds to the Hessian, the other blocks correspond to the constraints)
 - numerical PDE (discretized version of differential operator with boundary conditions — A_{11} corresponds to the operator, the other blocks to the boundary conditions)
- another way to view the above method is via the factorization

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} A_{11} & 0 \\ A_{21} & S \end{bmatrix} \begin{bmatrix} I & A_{11}^{-1}A_{12} \\ 0 & I \end{bmatrix} \quad (3.4)$$

- so solving $A\mathbf{x} = \mathbf{b}$ can be broken up into two steps

$$\begin{cases} \begin{bmatrix} A_{11} & 0 \\ A_{21} & S \end{bmatrix} \begin{bmatrix} \mathbf{y}_1 \\ \mathbf{y}_2 \end{bmatrix} = \begin{bmatrix} \mathbf{b}_1 \\ \mathbf{b}_2 \end{bmatrix} \\ \begin{bmatrix} I & A_{11}^{-1}A_{12} \\ 0 & I \end{bmatrix} \begin{bmatrix} \mathbf{x}_1 \\ \mathbf{x}_2 \end{bmatrix} = \begin{bmatrix} \mathbf{y}_1 \\ \mathbf{y}_2 \end{bmatrix} \end{cases}$$

- or equivalently

$$\begin{cases} A_{11}\mathbf{y}_1 = \mathbf{b}_1 \\ S\mathbf{y}_2 = \mathbf{b}_2 - A_{21}\mathbf{y}_1 \\ \mathbf{x}_2 = \mathbf{y}_2 \\ \mathbf{x}_1 = \mathbf{y}_1 - A_{11}^{-1}A_{12}\mathbf{y}_2 \end{cases}$$

- a third application is to use (3.4) to evaluate determinant

$$\det(A) = \det(A_{11}) \det(S)$$

- while a fourth is in inverting block matrices

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}^{-1} = \begin{bmatrix} A_{11}^{-1} + A_{11}^{-1}A_{12}S^{-1}A_{21}A_{11}^{-1} & -A_{11}^{-1}A_{12}S^{-1} \\ -S^{-1}A_{21}A_{11}^{-1} & S^{-1} \end{bmatrix}$$

- the trick to derive this expression is to consider (3.1) and try to express

$$\begin{bmatrix} \mathbf{x}_1 \\ \mathbf{x}_2 \end{bmatrix} = \begin{bmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{bmatrix} \begin{bmatrix} \mathbf{b}_1 \\ \mathbf{b}_2 \end{bmatrix}$$

and in which case $B = A^{-1}$

- we already have (3.3) which expresses \mathbf{x}_2 in terms of \mathbf{b}_1 and \mathbf{b}_2
- we need something similar for \mathbf{x}_1 and so we plug (3.3) back into (3.2) which gives us

$$\mathbf{x}_1 = (A_{11}^{-1} + A_{11}^{-1}A_{12}S^{-1}A_{21}A_{11}^{-1})\mathbf{b}_1 - A_{11}^{-1}A_{12}S^{-1}\mathbf{b}_2 \quad (3.5)$$

- now we just write (3.3) and (3.5) in block form

$$\begin{bmatrix} \mathbf{x}_1 \\ \mathbf{x}_2 \end{bmatrix} = \begin{bmatrix} A_{11}^{-1} + A_{11}^{-1}A_{12}S^{-1}A_{21}A_{11}^{-1} & -A_{11}^{-1}A_{12}S^{-1} \\ -S^{-1}A_{21}A_{11}^{-1} & S^{-1} \end{bmatrix} \begin{bmatrix} \mathbf{b}_1 \\ \mathbf{b}_2 \end{bmatrix}$$

which yields the required formula

4. RANK-1 UPDATING

- suppose that we have solved the problem $A\mathbf{x} = \mathbf{b}$ and we wish to solve the perturbed problem

$$(A + \mathbf{u}\mathbf{v}^T)\mathbf{y} = \mathbf{b}$$

- such a perturbation is called a *rank-one update* of A , since the matrix $\mathbf{u}\mathbf{v}^T$ has rank 1 (unless \mathbf{u} or \mathbf{v} is zero)
- as an example, we might find that there was an error in the element a_{11} and we update it with the value \bar{a}_{11}
- we can accomplish this update by setting

$$\bar{A} = A + (\bar{a}_{11} - a_{11})\mathbf{e}_1\mathbf{e}_1^T, \quad \mathbf{e}_1 = \begin{bmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

- for a general rank-one update, we can use the *Sherman–Morrison formula*, which we will derive here
- multiplying through the equation $(A + \mathbf{u}\mathbf{v}^T)\mathbf{y} = \mathbf{b}$ by A^{-1} yields

$$(I + A^{-1}\mathbf{u}\mathbf{v}^T)\mathbf{y} = A^{-1}\mathbf{b} = \mathbf{x}$$

- we therefore need to find $(I + \mathbf{w}\mathbf{v}^T)^{-1}$ where $\mathbf{w} = A^{-1}\mathbf{u}$
- we assume that $(I + \mathbf{w}\mathbf{v}^T)^{-1}$ is a matrix of the form $(I + \sigma\mathbf{w}\mathbf{v}^T)$ where σ is some constant
- from the relationship

$$(I + \mathbf{w}\mathbf{v}^T)(I + \sigma\mathbf{w}\mathbf{v}^T) = I$$

we obtain

$$\sigma\mathbf{w}\mathbf{v}^T + \mathbf{w}\mathbf{v}^T + \sigma\mathbf{w}\mathbf{v}^T\mathbf{w}\mathbf{v}^T = 0$$

- however, the quantity $\mathbf{v}^T\mathbf{w}$ is a scalar, so this simplifies to

$$(\sigma + 1 + \sigma\mathbf{v}^T\mathbf{w})\mathbf{w}\mathbf{v}^T = 0$$

which yields

$$\sigma = -\frac{1}{1 + \mathbf{v}^T\mathbf{w}}$$

- it follows that the solution \mathbf{y} to the perturbed problem is given by

$$\mathbf{y} = (I + \sigma \mathbf{w} \mathbf{v}^T) \mathbf{x} = \mathbf{x} + \sigma (\mathbf{v}^T \mathbf{x}) \mathbf{w}$$

and the perturbed inverse is given by

$$\begin{aligned} (A + \mathbf{u} \mathbf{v}^T)^{-1} &= (I + A^{-1} \mathbf{u} \mathbf{v}^T)^{-1} A^{-1} \\ &= \left(I - \frac{1}{1 + \mathbf{v}^T \mathbf{w}} \mathbf{w} \mathbf{v}^T \right) A^{-1} \\ &= A^{-1} - \frac{1}{1 + \mathbf{v}^T A^{-1} \mathbf{u}} A^{-1} \mathbf{u} \mathbf{v}^T A^{-1} \end{aligned} \quad (4.1)$$

which is the Sherman–Morrison formula

- an efficient algorithm for solving the perturbed problem $(A + \mathbf{u} \mathbf{v}^T) \mathbf{y} = \mathbf{b}$ can therefore proceed as follows:
 - solve $A \mathbf{x} = \mathbf{b}$
 - solve $A \mathbf{w} = \mathbf{u}$
 - compute $\sigma = -1/(1 + \mathbf{v}^T \mathbf{w})$
 - compute $\mathbf{y} = \mathbf{x} + \sigma (\mathbf{v}^T \mathbf{x}) \mathbf{w}$
- note that we already have the solution to $A \mathbf{x} = \mathbf{b}$ but we have to solve another system $A \mathbf{w} = \mathbf{u}$
- so how is this better than simply solving $(A + \mathbf{u} \mathbf{v}^T) \mathbf{y} = \mathbf{b}$?
- the answer is that if we have LU factorization of A , then solving $A \mathbf{w} = \mathbf{u}$ requires two back solves, which takes $O(n^2)$ operations whereas solving $(A + \mathbf{u} \mathbf{v}^T) \mathbf{y} = \mathbf{b}$ from scratch would require $O(n^3)$ operations
- note that this also works if we have the QR or any other factorizations of A that facilitate solving linear equations involving A
- an alternative approach is to note that

$$\begin{aligned} (A + \mathbf{u} \mathbf{v}^T)^{-1} &= [A(I + A^{-1} \mathbf{u} \mathbf{v}^T)]^{-1} \\ &= (I + \sigma A^{-1} \mathbf{u} \mathbf{v}^T) A^{-1} \\ &= A^{-1} + \sigma A^{-1} \mathbf{u} \mathbf{v}^T A^{-1} \end{aligned}$$

which yields

$$\begin{aligned} (A + \mathbf{u} \mathbf{v}^T)^{-1} \mathbf{b} &= A^{-1} (I + \sigma \mathbf{u} \mathbf{v}^T A^{-1}) \mathbf{b} \\ &= A^{-1} (\mathbf{b} + \sigma (\mathbf{v}^T A^{-1} \mathbf{b}) \mathbf{u}) \end{aligned}$$

and therefore we can solve $(A + \mathbf{u} \mathbf{v}^T) \mathbf{y} = \mathbf{b}$ by solving a problem of the form $A \mathbf{x} = \mathbf{b}$ where the right-hand side \mathbf{b} is perturbed

5. RANK- r UPDATE

- what we have in the previous section can be generalized by repeated application of the same technique

$$A + \mathbf{u}_1 \mathbf{v}_1^T + \cdots + \mathbf{u}_r \mathbf{v}_r^T = A + UV^T \quad (5.1)$$

where $U = [\mathbf{u}_1, \dots, \mathbf{u}_r]$, $V = [\mathbf{v}_1, \dots, \mathbf{v}_r] \in \mathbb{R}^{n \times r}$

- (5.1) is called a *rank- r update* of A
- this is useful if, for example, r entries of A are modified, requiring us to obtain the solution of $(A + UV^T) \mathbf{x} = \mathbf{b}$ from the original solution $A \mathbf{x} = \mathbf{b}$
- the notion of rank- r update is very much related to that of Schur complement

- if we introduce new variables $\mathbf{y} = C\mathbf{x}$, then

$$(A + BC)\mathbf{x} = \mathbf{b}$$

can be written as

$$\begin{cases} A\mathbf{x} + B\mathbf{y} = \mathbf{b} \\ \mathbf{y} = C\mathbf{x} \end{cases} \quad (5.2)$$

or equivalently

$$\begin{bmatrix} A & B \\ C & -I \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ \mathbf{y} \end{bmatrix} = \begin{bmatrix} \mathbf{b} \\ \mathbf{0} \end{bmatrix}$$

- in other words $A + BC$ is the Schur complement of $-I$ in $\begin{bmatrix} A & B \\ C & -I \end{bmatrix}$
- we now derive a generalization of the Sherman–Morrison formula (4.1) by solving (5.2)
- plug $\mathbf{x} = A^{-1}(\mathbf{b} - B\mathbf{y})$ into $\mathbf{y} = C\mathbf{x}$ to get

$$(I + CA^{-1}B)\mathbf{y} = CA^{-1}\mathbf{b}$$

and plug the expression $\mathbf{y} = (I + CA^{-1}B)^{-1}CA^{-1}\mathbf{b}$ back into $\mathbf{x} = A^{-1}(\mathbf{b} - B\mathbf{y})$ to get

$$\mathbf{x} = [A^{-1} - A^{-1}B(I + CA^{-1}B)^{-1}CA^{-1}]\mathbf{b}$$

- since \mathbf{b} is arbitrary, this must mean that

$$(A + BC)^{-1} = A^{-1} - A^{-1}B(I + CA^{-1}B)^{-1}CA^{-1} \quad (5.3)$$

- this is called the *Sherman–Woodbury–Morrison formula* and is useful for find rank- r updates of solutions to $A\mathbf{x} = \mathbf{b}$
- a word of caution: both (4.1) and (5.3) should not be used for computing explicit inverse (which is a bad idea in the first place) because they are numerically unreliable