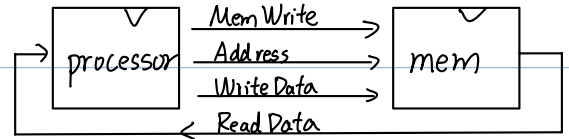
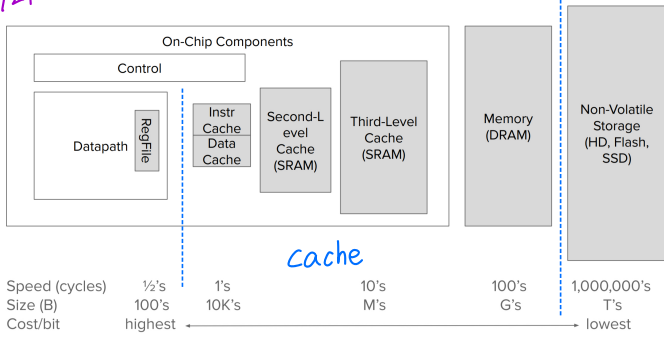


Mem physically separate from processor. Slow: L

To keep up: hierarchy (speed vs size)

1/21



Temporal locality if data recently used → likely used in future

Spatial ~ nearby data used in future

Performance Hit - data found in that level of mem.

Miss - X → next level

$$AMAT_i = \text{AccessTime}_i + \text{MissRate}_i \cdot AMAT_{i+1} \quad (\text{recursive hierarchy})$$

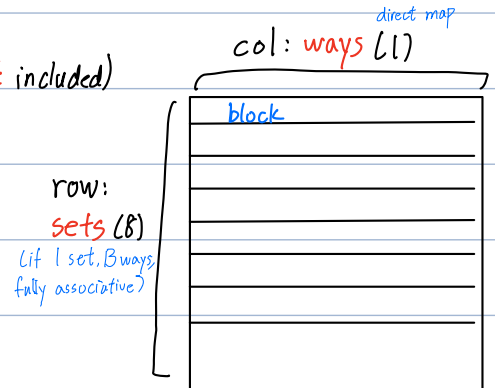
E. 2000 S/L, 1250 in cache, 750 not, AccTime_{cache} = 1 cycle; AccTime_{mem} = 100 cycle
 $1 + \frac{750}{2000} \cdot 100 = 38.5 \text{ cycles}$

or neighboring

Cache: recent data (too old → replaced)

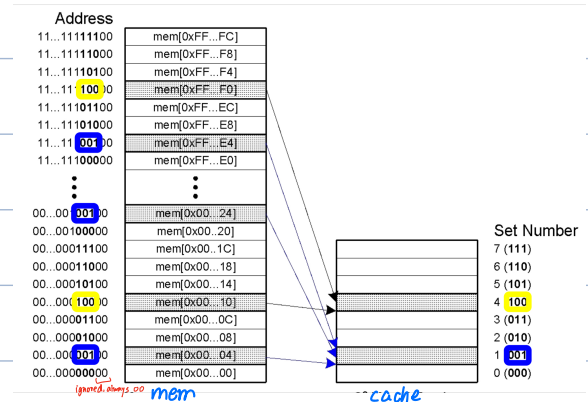
addressed w/ simple address hash

Def. Capacity (C): # bytes = Bb	8 words (B)	(meta not included)
# blocks B = SN	8	
Block size b (byte/block)	1 word/block	
# sets S	8	
Deg. of associativity N	1	



addr $\xrightarrow{\text{map hash}}$ one set, index set, check all ways

Hash: pull a few bits from addr.



V: valid bit, 1 ✓ ; Tag: remainder of addr.

Hit = V and (Tags =)

...00000100

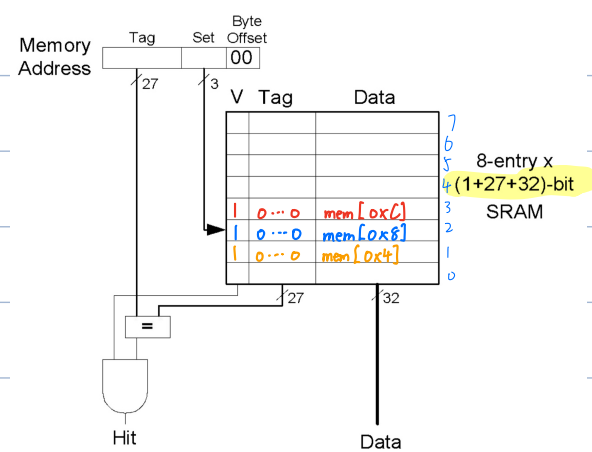
E. load 0x4, 0xC, 0x8 5 times

```

addi $t0, $0, 5
loop: beq $t0, $0, done
      lw $t1, 0x4($0)
      lw $t2, 0xC($0)
      lw $t3, 0x8($0)
      addi $t0, $t0, -1
      j loop
done:

```

set index



First time miss. Once in cache → saved → hit

3.1 Xs

3.15 - 1) ✓s → 20% miss rate

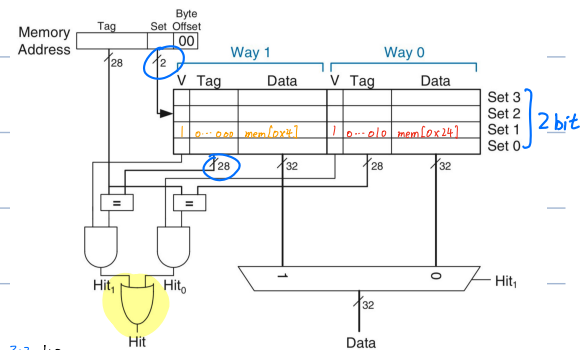
If addr. are 0x4, 0x24 → overwrite same hash → 100% miss : (

2-way associative cache (8x1 → 4x2) idx: 3 → 2

E. load 0x4, 0x24 5 times

more read

Fully-asso. → min. conflict miss, slow, \$\$\$



Spatial block size ↑

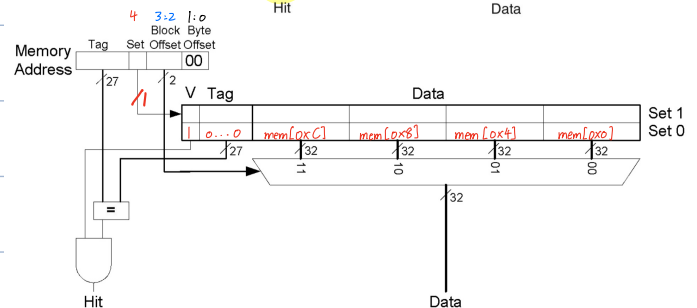
E. b = 4 words → $B = \frac{C}{b} = 2$ blocks

log₂ 4

3b → 1 set + 2 block offset

...00000100

E. 0x4, 0x8, 0xC



only 0x4 L miss → loads 0x0 ~ 0xC → all hit! → miss rate = 1/15

Miss 1. Compulsory: 1st time access → 0

2. Capacity: too much data to cache ← C

3. Conflict: map to same location ← C, N

but as C ↑, AMAT ↓

