Multiparty Computation made Practical Using the Virtual Ideal Functionality Framework

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Outline

Background

Multiparty Computation Related Projects

Design Goals

Automatic Parallel Execution Program Counters

Benchmark Results

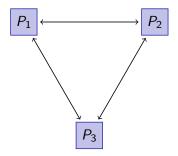
Multiplications
Comparisons

Possible Improvements

Program Counters Memory Overheads Debugging

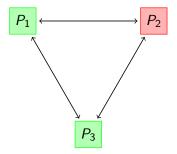
Conclusion

Multiparty Computation

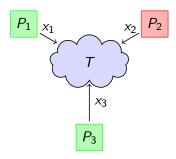


- n players
- ▶ Wish to jointly compute *f*
- ▶ Player P_i has input x_i
- Players learn $y = f(x_1, x_2, \dots, x_n)$

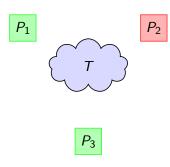
Multiparty Computation



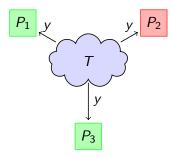
- n players
- ▶ Wish to jointly compute *f*
- Player P_i has input x_i
- Players learn $y = f(x_1, x_2, \dots, x_n)$
- ► Up to t players are corrupt
- ► Must keep inputs private
- ► Players only learn *y*



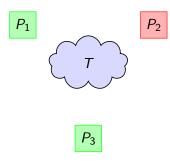
► Send inputs to trusted party over secure channels



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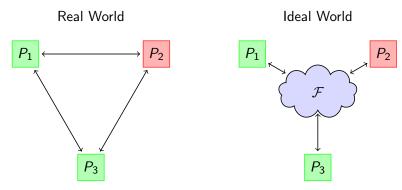


- Send inputs to trusted party over secure channels
- T computes $y = f(x_1, x_2, \dots, x_n)$
- Distributes y to all players

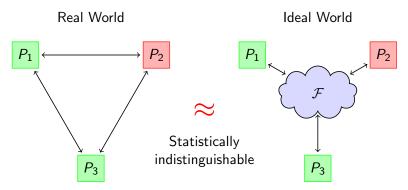


- Send inputs to trusted party over secure channels
- T computes $y = f(x_1, x_2, \dots, x_n)$
- Distributes y to all players
- Obviously secure!
- ▶ But who should play *T*?

Security Model: UC Framework



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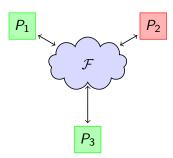
Functionality provided by VIFF

Supported operations:

- ► Input
- Output
- Addition
- Multiplication
- Comparison

All over arbitrary finite fields.

Ideal World



Related Projects

- FairPlay project (Haifa, Israel)
 - ▶ Yao-garbled circuits for 2 players
 - Java implementation
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- SIMAP project (Aarhus, Denmark)
 - General multiparty computations
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 - Some work done on a domain specific language

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▶ Easy to use for the programmer:

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x = a * b

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z = e * f
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Automatically run things in parallel:

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x = a * b

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z = e * f
```

Extensible with new operations:

```
\begin{aligned} & \textbf{def} \ \mathsf{max}(\mathsf{a}, \ \mathsf{b}): \\ & \mathsf{c} = \mathsf{a} > \mathsf{b} \\ & \textbf{return} \ \mathsf{c} * \mathsf{a} + (1 - \mathsf{c}) * \mathsf{b} \end{aligned}
```

Parallel Execution

- Networks have significant latency
- ▶ Want to run many operations in parallel
- Including primitive and compound operations

Example: Hamming Distance

```
\begin{array}{l} \textbf{def} \ \mathsf{xor}(\mathsf{a}, \ \mathsf{b}): \\ \textbf{assert} \ \mathsf{a}. \mathsf{field} \ \textbf{is} \ \mathsf{b}. \mathsf{field} \\ \textbf{if} \ \mathsf{a}. \mathsf{field} \ \textbf{is} \ \mathsf{GF256}: \\ \textbf{return} \ \mathsf{a} + \mathsf{b} \\ \textbf{else}: \\ \textbf{return} \ \mathsf{a} + \mathsf{b} - 2 * \mathsf{a} * \mathsf{b} \end{array}
```

- Straight-forward exclusive-or
- ▶ Fast for $GF(2^8)$ elements
- ▶ Slower for \mathbb{Z}_p elements
- ► (Already part of VIFF)

Example: Hamming Distance

```
def xor(a, b):
    assert a.field is b.field
    if a.field is GF256:
        return a + b
    else:
        return a + b - 2 * a * b
```

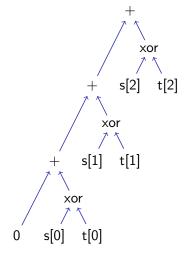
```
\label{eq:def-hamming} \begin{split} \text{def hamming}(s, \ t): \\ & \text{distance} = 0 \\ & \text{for i in range}(\text{len}(s)): \\ & \text{distance} \ += \text{xor}(s[i], \ t[i]) \\ & \text{return distance} \end{split}
```

- Straight-forward exclusive-or
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- Hamming distance
- xor calls should run in parallel

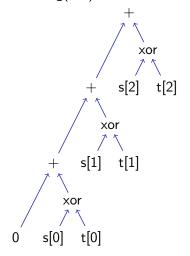
Hamming Distance Execution Tree

hamming(s, t) translates to:

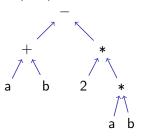


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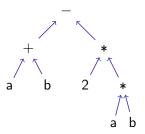
xor(a, b) translates to:



VIFF Execution Strategy

- Create execution tree as we go along
- Destroy execution tree from bottom up
 - Each node waits on nodes below
 - Bottom nodes wait on network traffic
- Composable: just plug new operations into tree!

Program Counters



- No parsing execution tree never fully constructed
- ▶ No fixed evaluation order
- ▶ But we must identify results

Program Counters Implementation

- ► First attempt:
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 - Each player increments a global counter
 - Fails because of asynchronous execution
- Working solution:
 - Manually "weave" a program counter through program
 - ▶ Tedious, easy to forget to increment program counter
- Current solution:
 - Runtime methods wrapped by a decorator
 - Calculated automatically based on call stack

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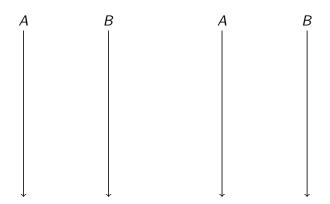
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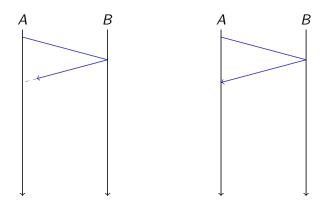
Conclusion

- ► Tested on 3 machines: USA, Norway, and Denmark
- Results for VIFF 0.4 (VIFF 0.6 is similar or better)
- Tested multiplications and comparisons

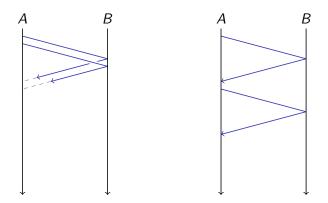
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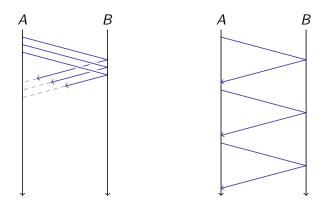
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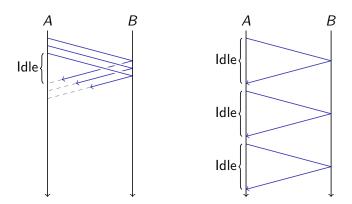
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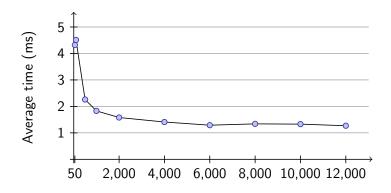


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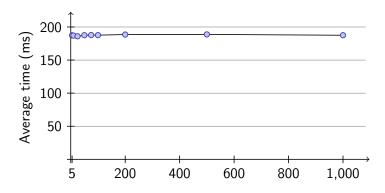
Parallel Multiplications

Multiplying random 65-bit numbers:



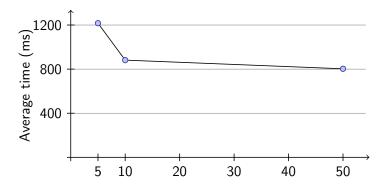
Serial Multiplications

Multiplying random 65-bit numbers:



Parallel Comparisons

Comparing random 32-bit numbers, 65-bit field modulus:



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Program Counter Overheads

- ▶ They work, but it's a bit magic...
- Exactly when must the program counter be updated?
- Excessive wrapping slows down method calls
- Program counter size depends on stack depth

Memory Overheads

- Python objects have a large memory overhead:
 - Reference count (4 bytes)
 - Object attribute dictionary (144 bytes)
 - ► Allocation overhead (>8 bytes)

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- Python objects have a large memory overhead:
 - ▶ Reference count (4 bytes)
 - Object attribute dictionary (144 bytes)
 - ► Allocation overhead (>8 bytes)
- ▶ 100,000 field elements with 65-bit prime:
 - ▶ Optimal: ≈800 KB
 - VIFF: ≈40 MB (expanded 50 times)
 - More memory needed for execution tree

Debugging

- ▶ Something went wrong! What now?
- ▶ Debugging asynchronous programs can be hard

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- ► Something went wrong! What now?
- Debugging asynchronous programs can be hard
- Need better
 - Logging infrastructure
 - Handling of exceptions

Type Safety

- Python is a strongly typed language
 - ▶ "12" * 3 == "121212" but 12 * 3 == 36
 - ▶ "1" + 12 raises TypeError
- ▶ Types are only checked at runtime

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- Python is a strongly typed language
 - ▶ "12" * 3 == "121212" but 12 * 3 == 36
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- ► Types are only checked at runtime
- Unit tests help here
- Better input validation

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- Automatic parallel execution
- ► Free Software: LGPL
- ▶ Please see http://viff.dk/
 - Source code
 - Documentation
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Thank you for listening!