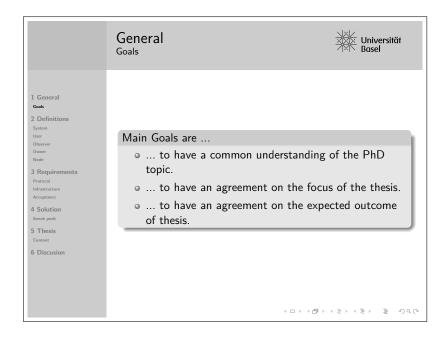


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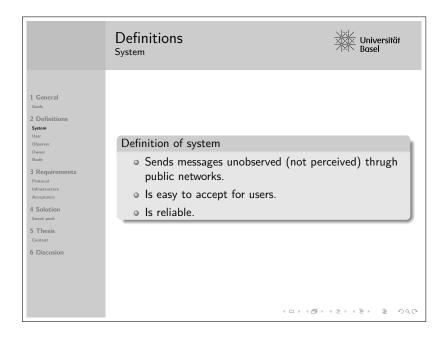
1 General

1.1 Goals

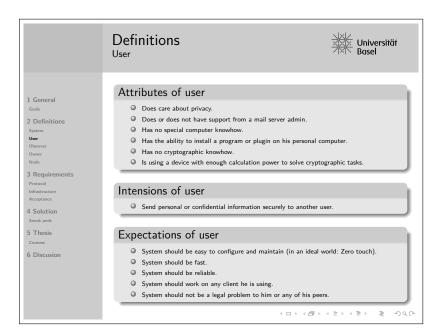


2 Definititons

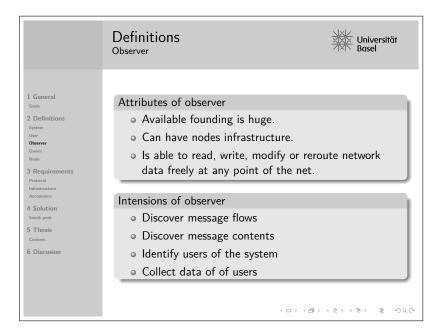
2.1 System



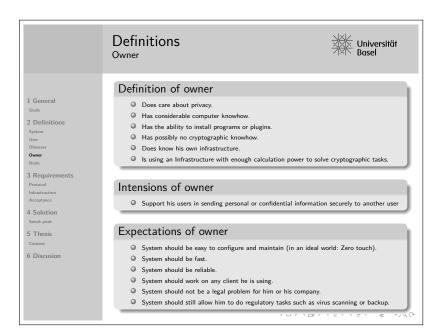
2.2 User



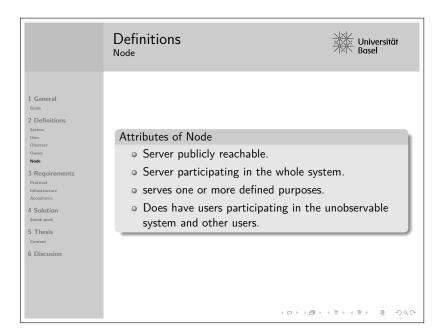
2.3 Observer



2.4 Owner

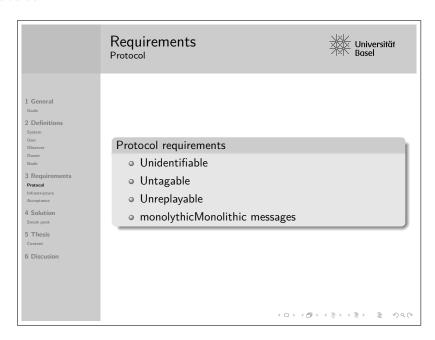


2.5 Node



3 Requirements

3.1 Protocol



• Unidentifiable

If a message or a participating node is identifiable then it is easy for an observer to block some or all parts of the system. This makes the system unreliable and may force users to use specific nodes (such as nodes which are under the control of the observer) and therefore compromise the overall security. Only a service that is able to hide its messages in legitimate network traffic is not subject to selective blocking.

• Untagable

If messages going through the system are tagable by any of the participants (nodes) then an observer might tag messages and then follow them while they are propagating the network. If information is appended to a message it must be cloaked with the same reliability as the original message itself.

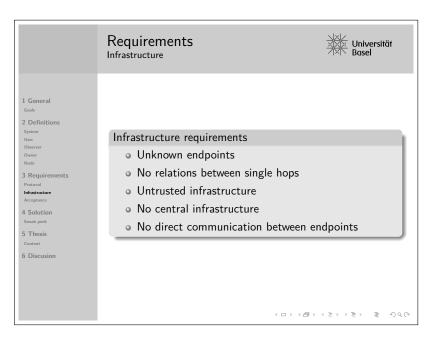
Unreplayable

If an observer can replay any part of the message (send it multiple times), he can identify the traffic generated by those messages by statistical means. This would enable him to identify traffic which is caused by a specific message and thus narrow down the possible final recipients.

• Monolithic messages

Messages should not depend on external content (such as images). If a message is not self-contained then "bugging" is an easy way to identify the message on its way up until they reach the recipient or the recipient itself.

3.2 Infrastructure



• Unknown endpoints

Every endpoint should behave the same as an intermediate routing point. They should receive and send messages so that they are not identifiable as endpoints. Identifiable endpoints simplify analysis.

No relations between single hops

Messages transferred from server to server must be unrelated. Server identifiable to send messages due to received messages are potential targets for analysis.

• Untrusted infrastructure

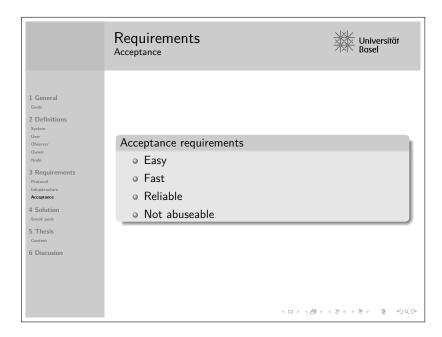
Unlike in a company owned net, in a public network trusting an infrastructure is not sensible. It is very often not clear who owns a server and who else does have access to it. The motivation of an infrastructure owner is often not clear and his intentions may or may not be sincere. So an unobservable system may not build its unobservability based on behaviour of the transporting infrastructure.

• No central infrastructure

Central infrastructure may be attacked or shut down. They are easier to monitor than an unknown number of participants. Furthermore a central infrastructure may be used to compromise security of messages or nodes. It enables an observer to identify nodes by monitoring the traffic of a central infrastructure.

• No direct communication between endpoints
If sender and receiver communicate directly then they are easily identified. So – all communications between endpoints should normally be done via intermediate nodes.

3.3 Acceptance



Easy

A system must be easy to use. The possibilities should be similar to common elaborated systems and the usage should be alike or the same. This offers a steep learning curve to the user.

If ignored, only the users heavily concerned about their privacy would be willing to use the system. All others would ignore it as they are not ready to invest efforts into a system that offers them not sufficient benefits but new limitations.

Fast

In today's world we already adapted to fast moving messages. It is quite common that people talk to each other and send at the same time additional informations by chat or mail. They do expect that this information propagates fast through public networks. For some messages even an almost instant reply of the recipient is expected by the sender. Therefore any system must allow a fast transport of messages from the sender to the recipient.

• Reliable

Messages are expected to arrive at the recipient's device. Today there are numerous common systems such as email, chat, sms and mms offering reliable transfers. Any system not sending reliably will not be used due to the limitations given by an unreliable system.

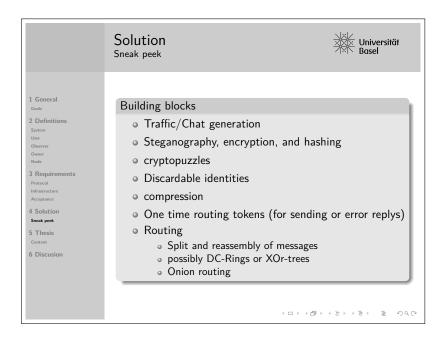
Another part of reliability is the protection. The message protection must be unbreakable (within reasonable bounds). If the system can be attacked easily then it offers "no value" for ädditional effort". For most users this would be a reason to discard such a solution.

• Not abuseable

Any system may be abused. The willingness of using a system if it is to easily abusable is very limited. A user will not be using a system which increases UBM (unsolicited bulk messages) or enables someone to blackmail him easily.

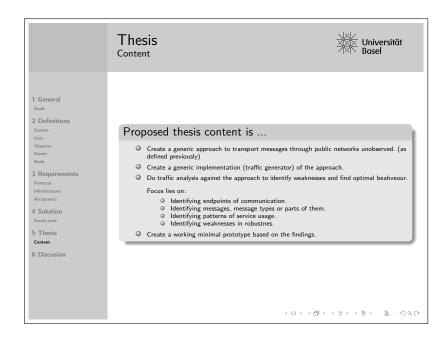
4 Solution

4.1 Sneak peek



5 Thesis

5.1 Content



6 Discussion

