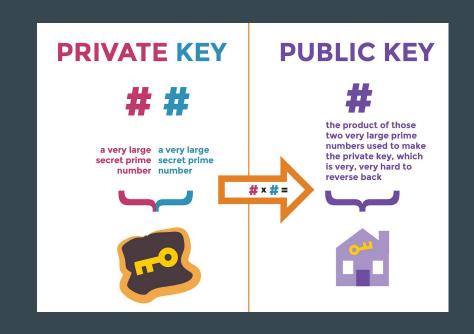
Widening the Sieve

•••

Richard Muniu, Mickey Haregot, Kevin Zheng

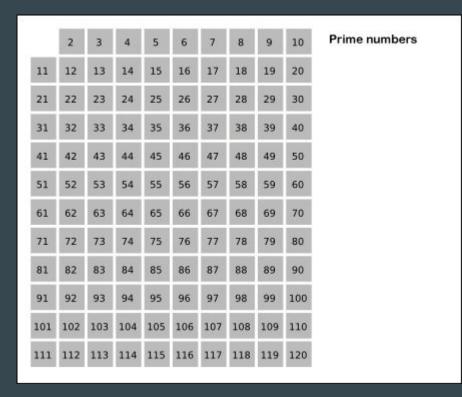
Finding Prime Numbers

- Prime numbers hard to find efficiently
- This makes them useful in security protocols
 - E.g. public/private key encryption
 - Rivest-Shamir-Aldeman (RSA) algorithm
- 2000 year-old algorithms still surprisingly effective today



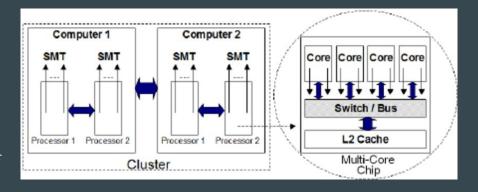
The Sieve of Eratosthenes

- 1. Have an array for each integer between 0 and some max N
- Iterate through the array,
 whenever you find an
 "unmarked" element, mark all of
 its multiples.
- 3. All unmarked elements by end of the run are prime numbers
- 4. Features: unbalanced load, not embarrassingly parallel, low computation per step



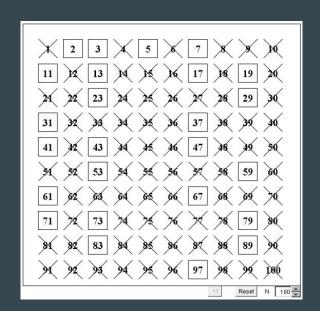
System Traits Desired

- Decided to use single machine
 - No communication
 - Used single node on cluster (Strelka)
- Reasons why
 - Sieve algorithm relatively simple
 - Using multiple nodes result in large communication costs
 - Results required to be aggregated at end of runs
 - May become large part of total runtime



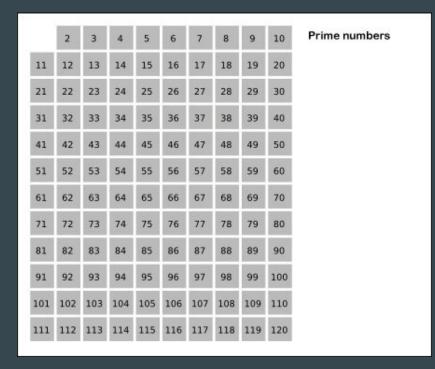
Methods of Implementation

- Sequential
 - Baseline test
- OpenMP
 - Parallelized by base: each thread investigates a "unmarked" number separately
- CUDA
 - Massively parallel version of OpenMP implementation.
 - Takes advantage of SIMD architecture
- Balanced Pthreads
 - Each thread is responsible for a different section of the array and the base process broadcasts "bases" for the other threads to investigate



OpenMP

- Uses OpenMP's loop primitives
- Marks multiples in parallel
- Load balancing problem: heavy work assigned to one thread
- Resolved using dynamic scheduling of iteration chunks





CUDA

- Similar to OpenMP implementation
- Massively parallel
- Every possible "base" has its multiples simultaneously marked off
- Takes advantage of SIMT architecture
- Multiple kernel calls to combat max-thread limitations

```
if (threadIdx.x < 4) {
    A;
    B;
} else {
    X;
    Y;
}</pre>
```

Balanced pthreads

- Splits up array into segments
- Master-worker model
- Meant to balance workload
- Master discovers "bases" and broadcasts them to workers
- Broadcast uses shared variable and barriers

	2	3	4	5	6	7	8	9	10	
11	12	13	14	15	16	17	18	19	20	
21	22	23	24	25	26	27	28	29	30	
31	32	33	34	35	36	37	38	39	40	
41	42	43	44	45	46	47	48	49	50	
51	52	53	54	55	56	57	58	59	60	
61	62	63	64	65	66	67	68	69	70	
71	72	73	74	75	76	77	78	79	80	
81	82	83	84	85	86	87	88	89	90	
91	92	93	94	95	96	97	98	99	100	

Testing

- Compared sequential vs parallel runtimes on Strelka
- OpenMP/balanced pthreads: tested both seq_sieve and par_sieve
 - o 32-core CPU
 - o hi-mem partition
- CUDA: ran cudaSieve on
 - o 4x NVIDIA 2080 Ti GPUs
 - o gpu-02 partition.
- For both OpenMP and balanced pthread, dynamic allocation in RAM failed for arrays of size 10¹³ and above.
- For CUDA, dynamic allocation in DRAM failed for arrays of size 10¹² and above.
- Used average values from three runs of seq_sieve and par_sieve for each method.



RESULTS AND DISCUSSION

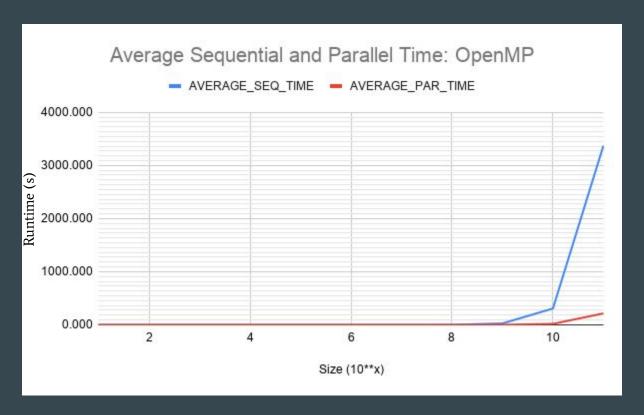
OpenMP

• Nonlinear speedup of up to 16

• No speedup until arrays of size 10^7

	Ι Δ	Α	
Size	Average	Average	
(10^{x})	Sequential	Parallel	Speedup
(10)	Time(s)	Time(s)	
1	0.001	0.017	0.058
2	0.001	0.016	0.063
3	0.001	0.017	0.059
4	0.001	0.016	0.061
5	0.001	0.018	0.057
6	0.008	0.017	0.462
7	0.073	0.027	2.750
8	2.050	0.201	10.182
9	27.495	1.923	14.298
10	308.771	19.372	15.939
11	3371.608	216.177	15.596

OpenMP



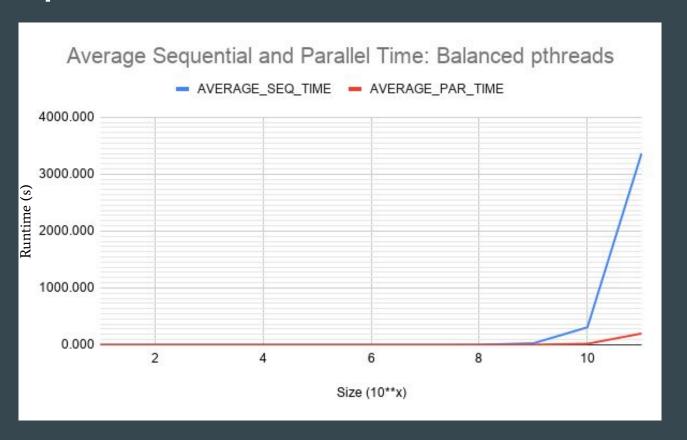
Balanced pthreads

Nonlinear speedup of close to 17

 No consistent speedup until arrays of size 10⁷

Size	Average	Average	Speedup	
	Sequential	Parallel		
(10^{x})	Runtime	Runtime		
1	0.005	0.002	2.500	
2	0.001	0.002	0.429	
3	0.001	0.003	0.375	
4	0.001	0.005	0.200	
5	0.001	0.009	0.154	
6	0.007	0.016	0.458	
7	0.071	0.046	1.547	
8	2.204	0.228	9.653	
9	27.518	2.440	11.278	
10	308.638	18.786	16.429	
11	3364.497	199.944	16.827	

Balanced pthreads



CUDA

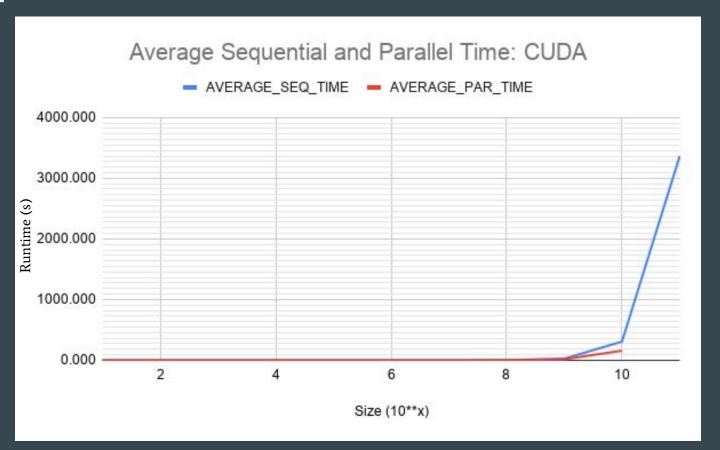
• Nonlinear speedup of up to 2

• No consistent speedup until arrays of size 10⁹, but relatively consistent runtime until that point

• Lower performance than both other implementations

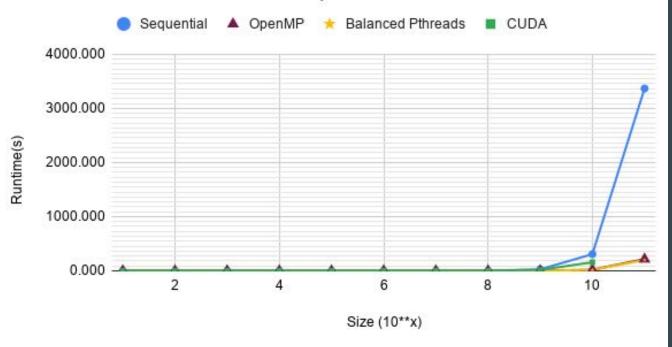
e e			
Size	AVERAGE	AVERAGE	
(10^x)	SEQ	PAR	Speedup
(10)	TIME	TIME	
1	0.005	3.128	0.002
2	0.001	3.039	0.000
3	0.001	3.060	0.000
4	0.001	3.057	0.000
5	0.001	3.076	0.000
6	0.007	3.090	0.002
7	0.071	3.295	0.021
8	2.204	4.536	0.486
9	27.518	17.802	1.546
10	308.638	158.962	1.942
11	=	-	J a a

CUDA

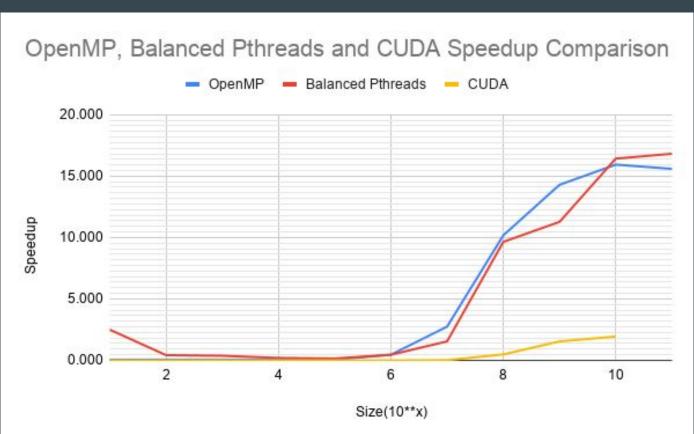


Runtime Trend and Comparison

Sequential, OpenMP, Balanced Pthreads and CUDA Runtime Comparison



Speedup Trend and Comparison



Conclusion and Future Directions

- A lot to gain from parallelization in form of speedup and efficiency
- Benefits are not obvious/present at low task sizes due to paralellization overheads
- CUDA performed slower than other methods
 - o hardware limitations?
 - Potential software optimizations?
- Some future directions to consider:
 - Further optimizations to CUDA implementation
 - Explore distributed memory approach to eliminate max array size limitation

QUESTIONS

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