• for each object A: identity arrow:  $1_A:A\to A$ • Unit:  $f \circ 1_A = f = 1_B \circ f$ Def. concrete Categories Categories in which Objects are Sets, possibly equipped with some structure, and arrows are  $certain,\ possibly\ structure-perserving,\ functions.$ Def. functor Let  $\mathbf{C}, \mathbf{D}$  be categories, then  $F: \mathbf{C} \to \mathbf{D}$  is a functor with: 1.  $F(f:A \rightarrow B) = F(f):F(A) \rightarrow F(B)$  $2. \ F(g \circ f) = F(g) \circ F(f)$ 3.  $F(1_A) = 1_{F(A)}$ Def. discrete categories categories with only the identity arrows  $\,$ Def. monoid A set M with an associative binary operation Def. isomorphism In any category  $\mathbf{C}$ , an arrof  $f:A\to B$  is an isomorphism, if there is an arrow  $g:B\to A$  in  ${\bf C}$ such that.  $g \circ f = 1_A$  and  $f \circ g = 1_B$ 

Category Theory Anki Study Document

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Def. Category

• for  $f:A\to B, g:B\to C$  with cod(f)=dom(g)

composite of f and  $g \colon g \circ f : A \to C$ 

 $\bullet$  Objects: A, B, C, ...

 $\bullet$  Arrows: f, g, h, ...

• for each  $f:A\to B$ domain: A = dom(f) $\operatorname{codomain:} B = \operatorname{cod}(f)$ 

we write  $f^{-1} = g$ . A is isomorphic to B  $(A \cong B)$  if there exists an isomorphism between them. Def. group A group G is a monoid with an inverse  $g^{-1}$  for every element g. Def. Free Monoid A monoid M is **freely generated** by a subset A of M with: 1. no junk: every element  $m \in M$  can be written as a product of elemens of A  $m = a_1 \cdot_M \dots \cdot_M a_n, \quad a_i \in M$ 2. no noise: No "nontrivial" relations hold in  ${\rm M:}$ if  $a_1...a_n = a'_1...a'_n$  then this is required by the axioms for monoids. Def. Universal Mapping Propperty M(A) given  $i:A \to |M(A)|,$  Monoid N and  $f:A \to |N|$ there is a unique monoid homomorphism  $\,$  $\bar{f}:M(\bar{A})\to N \text{ s.t. } |\bar{f}|\circ i=f$