

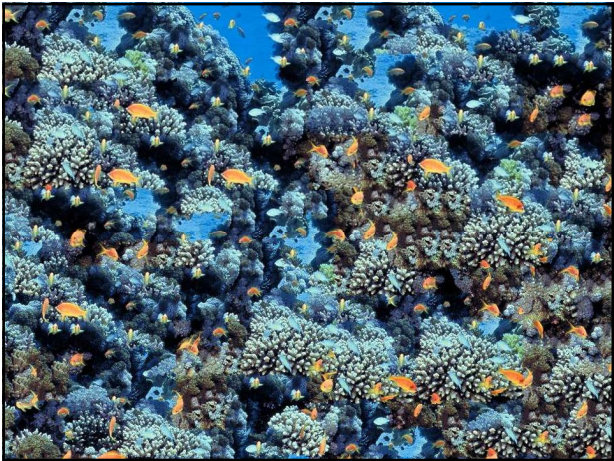
COMP20003
Algorithms and Data Structures
Introduction to Graphs

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Semester 2

Application: Image synthesis

A. A. Efros & W. T. Freeman. [Image quilting for texture synthesis](#). SIGGRAPH 2002

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The key to solving many algorithmic problems is to think of them in terms of graphs. Graph theory provides a language for talking about the properties of relationships, and it is amazing how often messy applied problems have a simple description and solution in terms of classical graph properties.

Designing truly novel graph algorithms is a very difficult task. The key to using graph algorithms effectively in applications lies in correctly modeling your problem so you can take advantage of existing algorithms. Becoming familiar with many different algorithmic graph *problems* is more important than understanding the details of particular graph algorithms, particularly since Part II of this book will point you to an implementation as soon as you know the name of your problem.

Here we present basic data structures and traversal operations for graphs, which will enable you to cobble together solutions for basic graph problems. Chapter 6 will present more advanced graph algorithms that find minimum spanning trees, shortest paths, and network flows, but we stress the primary importance of correctly modeling your problem. Time spent browsing through the catalog now will leave you better informed of your options when a real job arises.

S. S. Skiena. The Algorithm Design Manual.
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Graph definition

Graph:

- a **representation** of a **set of objects**
- and the **relationships** between them

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Graph definition

Graph $G = \{V, E\}$ Set of:

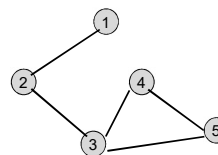
- **Vertices** V : can contain **information**
- **Edges** E (links between vertices): can have **direction** and/or **weight**

Compared to **trees** and linked **lists**:

- vertices = nodes
- edges = links

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Undirected graph

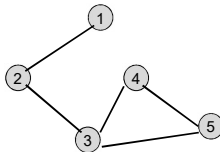


Edges have **no direction** specified

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Connected
Undirected graph

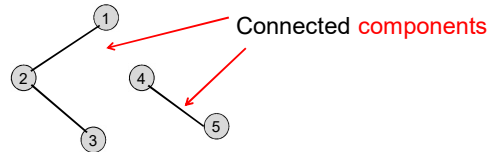


Every pair of **vertices** is **connected**
(possibly indirectly)

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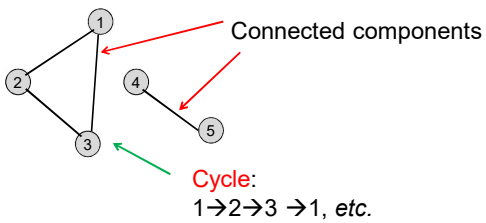
Unconnected
Undirected graph



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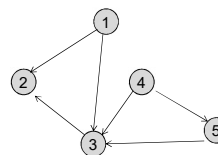
Unconnected
Undirected graph with cycle



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Directed graph

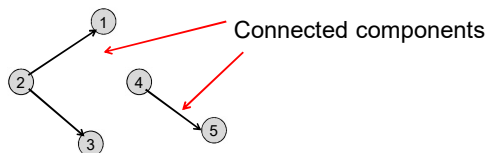


- Edge **direction** is specified
- Links are **not symmetrical**

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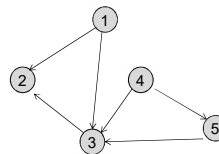
**Acyclic, unconnected
directed graph**



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Directed graph



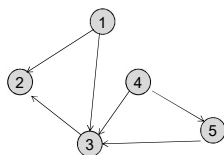
Reachability:

Can you get from Vertex 1 to Vertex 5?

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**Weakly connected
Directed graph**

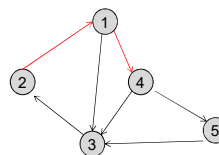


Replace all directed edges with undirected edges,
to obtain a connected (undirected) graph

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**Strongly connected
Directed graph**

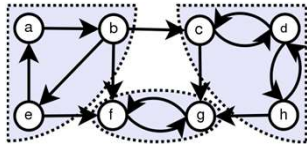


If every vertex is reachable
from every other vertex

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Strongly connected Components in a Digraph

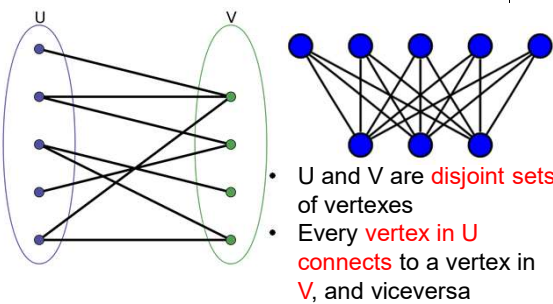


If **every** vertex is **reachable**
from every other vertex
within the **same** component

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Bipartite Graph

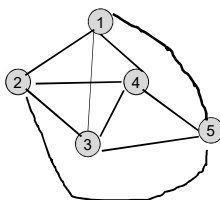


- U and V are **disjoint sets** of vertexes
- Every **vertex in U** **connects** to a vertex in **V**, and viceversa

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Complete graph



$V(V-1)/2$ **edges**

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Trees and Graphs

Tree, undirected graph that is:

- Connected
- Acyclic
- *n.b.* Any two vertices are **connected** by **exactly one simple path**
- *n.b.* All vertices are connected

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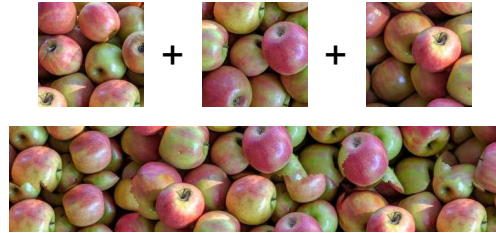
Demos...

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Application: image synthesis

- How does it use graphs?

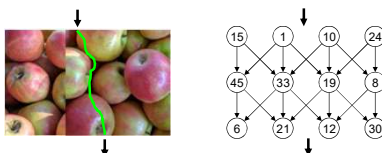


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Application: image synthesis

- To find the best splice, represent overlapping images as a graph:
 - Edge weight = difference in colour
 - Problem = find minimum path



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