

2.3 1) nop addi x0, x0, 0

2) ret jalrr x0, 0(x1)

3) call offset { $\text{unipc } x_1, \text{offset}[31:12] + \text{offset}[11];$
jalr x1, offset[11:0](x1).

4) mv rd, rs addi rd, rs, 0.

5) rdycke rd crrs rd, cycleh, x0.

6) sextw rd, rs addiw rd, rs, 0.

2.7 1) slti t3, t2, 0 2) add t0, t1, t2

slt t4, t0, t1 bltu t0, t1, overflow.

3) ARM: ① 进位与借位 (无符号数): 标志寄存器的 C 位表示进位或非借位标志, 进位 C=1

② 溢出: 当操作数与运算结果为二进制补码表示的带符号数, 用 CPSR 的 V=1 表示.

2.8 1) $2^L - 1$ (L 为操作指令的位数) $i \times i - 1, i$.

2) W: 无效操作 DE: 除零异常 OF: 溢出异常 UF: 借位溢出 NX: 不精确

* 不知道, 没查到...

3) ARM: 将除数进行可除性判断后再进行运算.

2.12 1) S 2) M 3) M 4) S 5) V

2.13 add t3, x0, x0 #Assign t3=i, t0=100 2.14

add t0, x0, 100

blt a0, a1, if # a0 < a1 & a < b.

sub a2, a1, a0

Loop: bge t3, t0, exit

j end

sll t5, t3, 2

If: add a2, a0, a1

add t0, t0, t5

j end

add t1, t1, t5

end: ret.

lw t1, 0(t0)

mul t0, t1, t2

addi t3, t3, 1

j rec Loop

exit: ret

2.15. add t0, t0, x0
 addi t1, x0, 3
 add t2, x0, x0.
 sll t3, t2, 2.
 add t0, t0, x0.
 add t0, t0, t3
 add t0, t1, x0.
 addx mul t3, t3, t1.
 add t0, t0, t3
 add t0, t1, x0.

2.17: $a_0=0, a_1=1, a_2=30$
 loop: $a_0 = a_2$ 直接结束循环.
 $a_1 \times 2$ 覆盖 \Rightarrow
 a_0+1 覆盖.

2.16. swap: lw t0, 0(t0)
 add t2, x0, t0.
 lw t1, 0(t1)
 add t0, x0, t1.
 sw t0, 0(t0)
 add t1, t2, x0.
 sw t1, 0(t1)
 ret.

int $a_0=0, a_1=1, a_2=30$;
 for ($a_0=0, a_0 < a_2, a_0++$) {
 $a_1 = a_1 \times 2$;
 }
 printf("%d", a_1)
 \Rightarrow 计算 $a_1 = 2^{30} = ?$