Program Slicing:

Program Dependence Graphs

source

Interprocedural Slicing
Using Dependence Graphs,
Horwitz, Reps & Binkley
ACM TOPLAS, Jan., 1990.

available on course web site

http://www.dur.ac.uk/k.b.gallagher/local/Slicing/toplas.pdf pp1-9

The Program Dependence Graph 1

- Multigraph
- Vertices
 - Assignment
 - Control predicate, no side effects
 - Entry,
 - InitialState(v)
 - FinalUse(v) == End
- Control Dependence Edge: v -> w
 - v is entry; w is not in loop or conditional
 - v is control; w is immediately controlled by v

The Program Dependence Graph - 2

- Flow Dependent Edge: v -> w
 - v defines variable x
 - w uses (references) x
 - no intervening defs of x
- Loop Carried:
 - flow dependent, plus
 - edge back to controlling predicate
 - v, w in loop

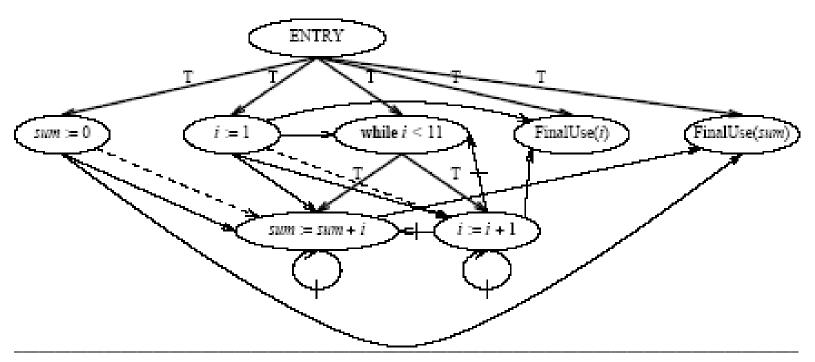
The Program Dependence Graph - 3

- Loop Independent:
 - flow Dependent, plus
 - NO edge back to controlling predicate
- Def-Order Dependence: v -> w
 - v, w define the same variable
 - v, w are in same branch of conditional
 - there is $z: v \rightarrow z$ and $w \rightarrow z$
 - v is "left" of w

Example page 7

```
main
 sum := 0;
  i = 1;
  while (i < 11) do
     sum := sum + I;
     i = I + 1;
   od;
end(sum, i)
```

Example page 7



Slicing as Graph Reachability

Slice on v is the subgraph of all nodes and edges that can reach v (see page 8/9)

```
procedure Mark Vertices Of Slice (G, S) declare

G: a program dependence graph

S: a set of vertices in G

Work List: a set of vertices in G

v, w: vertices in G

begin

Work List := S

while Work List \neq \emptyset do

Select and remove vertex v from Work List

Mark v

for each unmarked vertex w such that edge w \mapsto_f v or edge w \mapsto_e v is in E(G) do

Insert w into Work List

od

od

end
```

Application of the algorithm

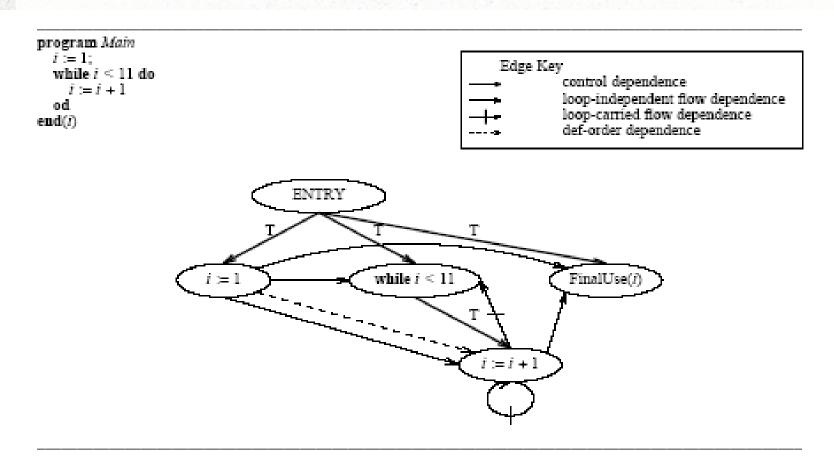


Figure 3. The graph and the corresponding program that result from slicing the program dependence graph from Figure 1 with respect to the final-use vertex for i.