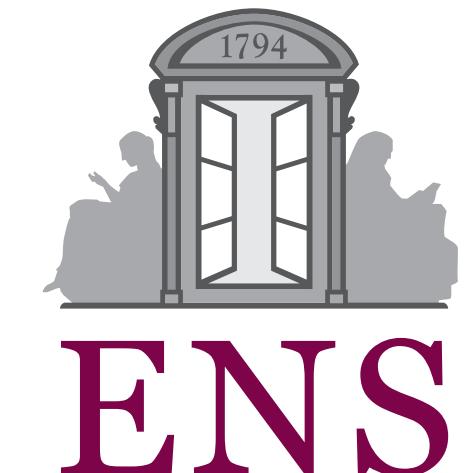


# SSE and SSD: Page-Efficient Searchable Symmetric Encryption

- Angèle Bossuat Quarkslab
- Raphael Bost DGA
- Pierre-Alain Fouque Université de Rennes
- Brice Minaud ENS – CNRS – PSL University – INRIA
- Michael Reichle ENS – CNRS – PSL University – INRIA



# Motivation

Outsource confidential information



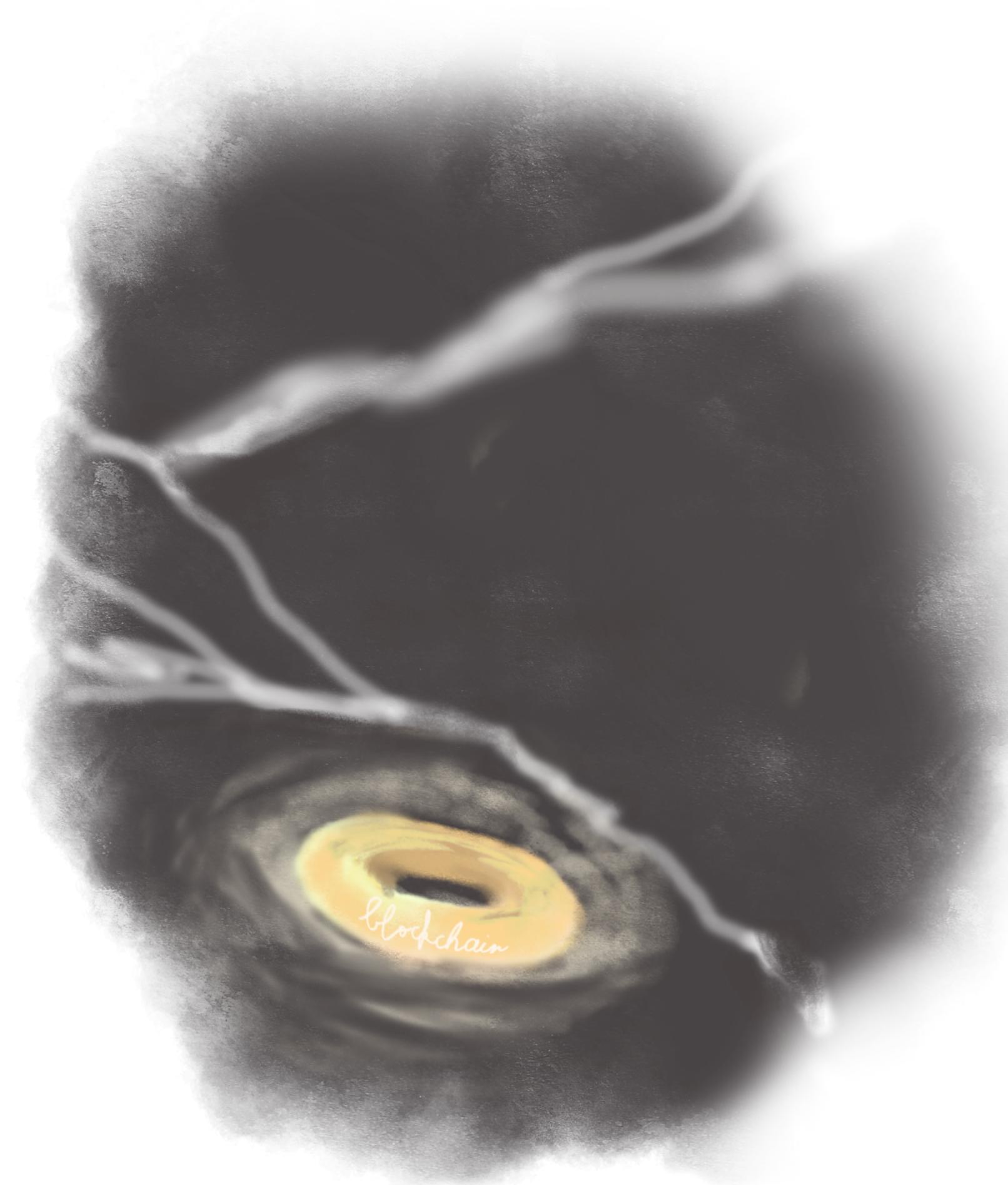
# Motivation

Outsource confidential information



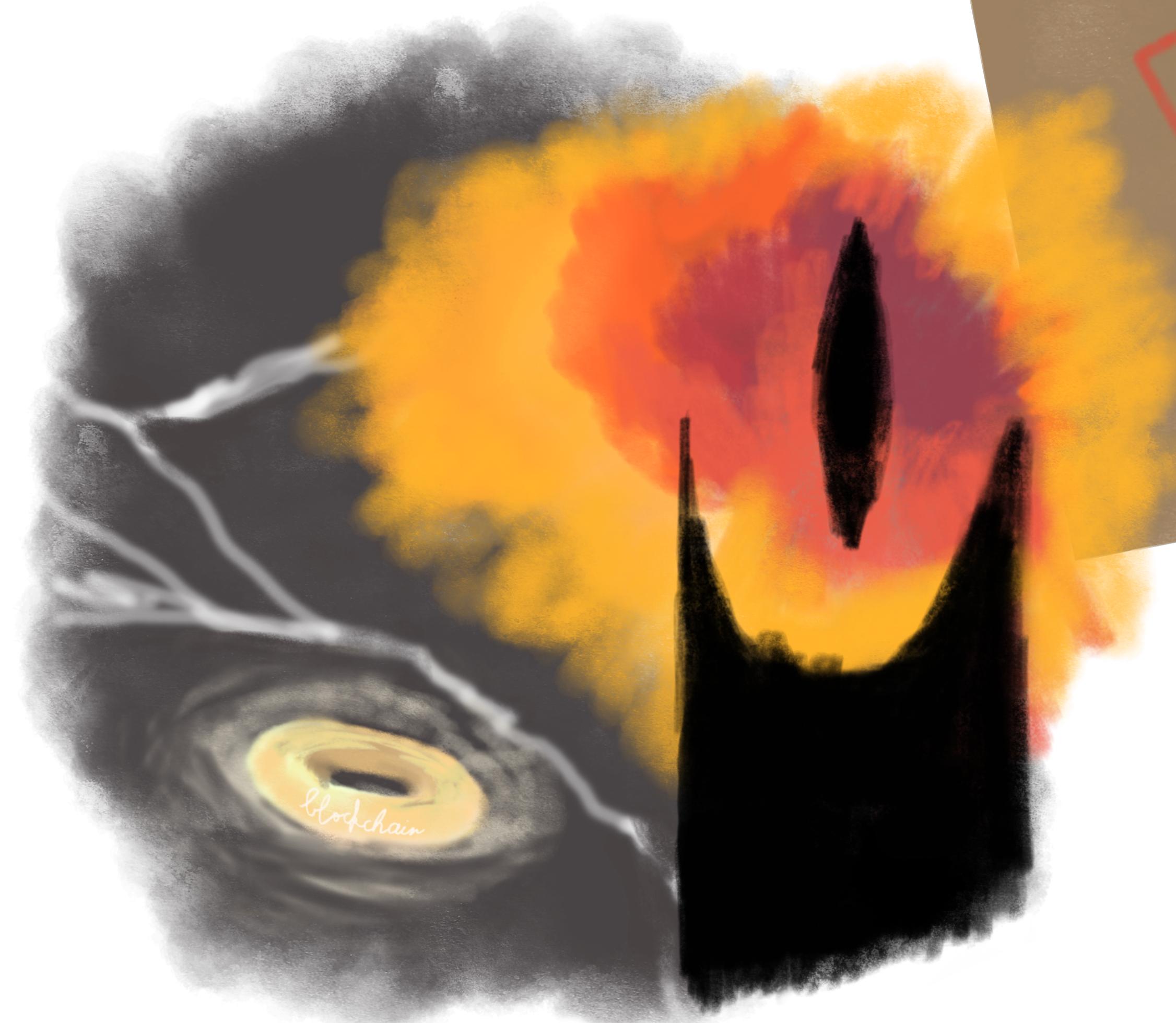
# Motivation

Outsource confidential information



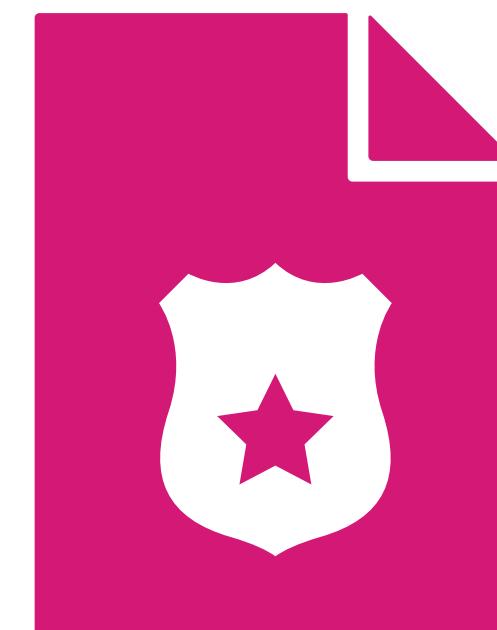
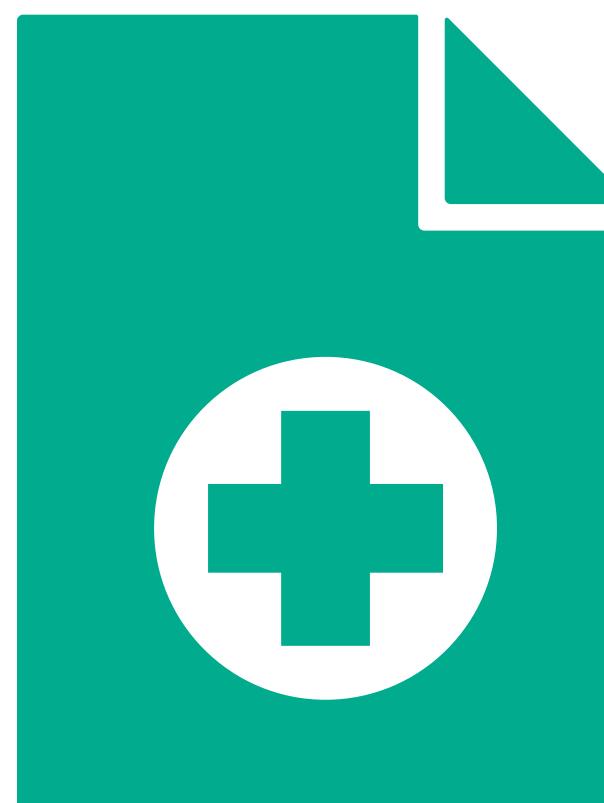
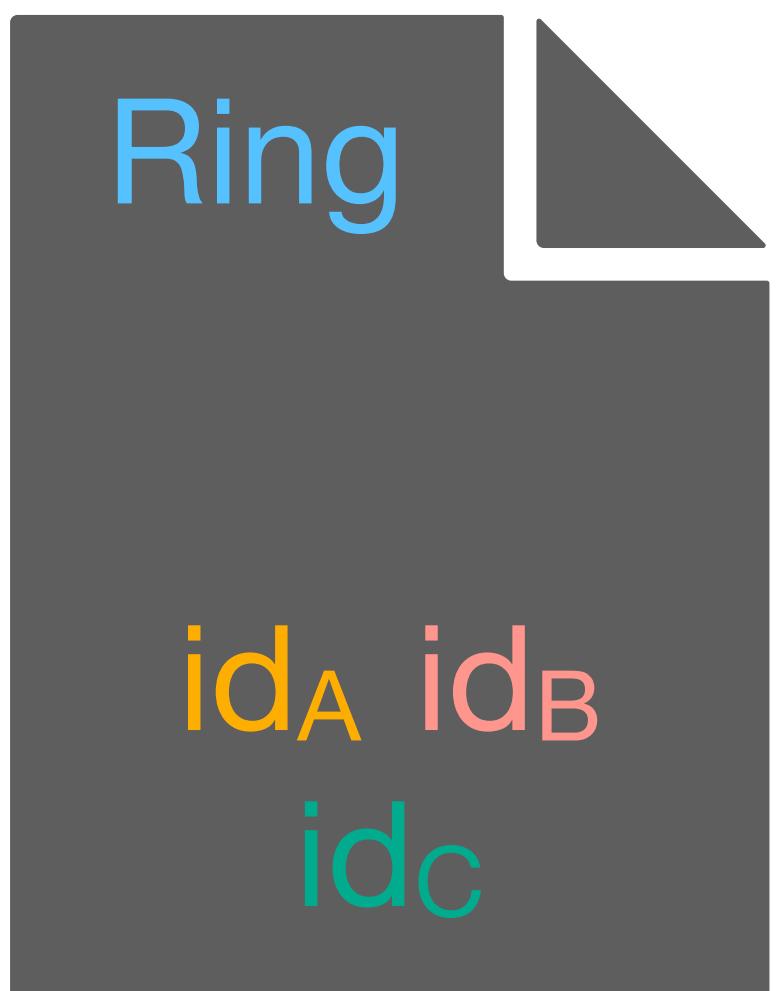
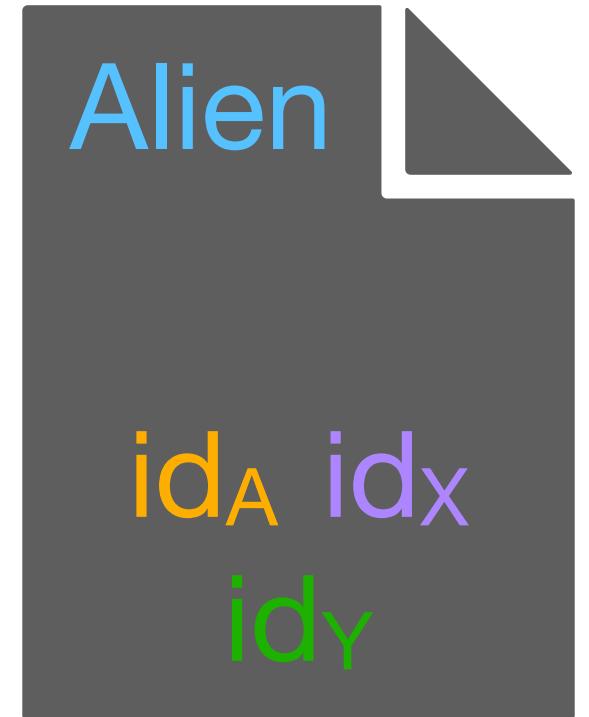
# Motivation

Outsource confidential information



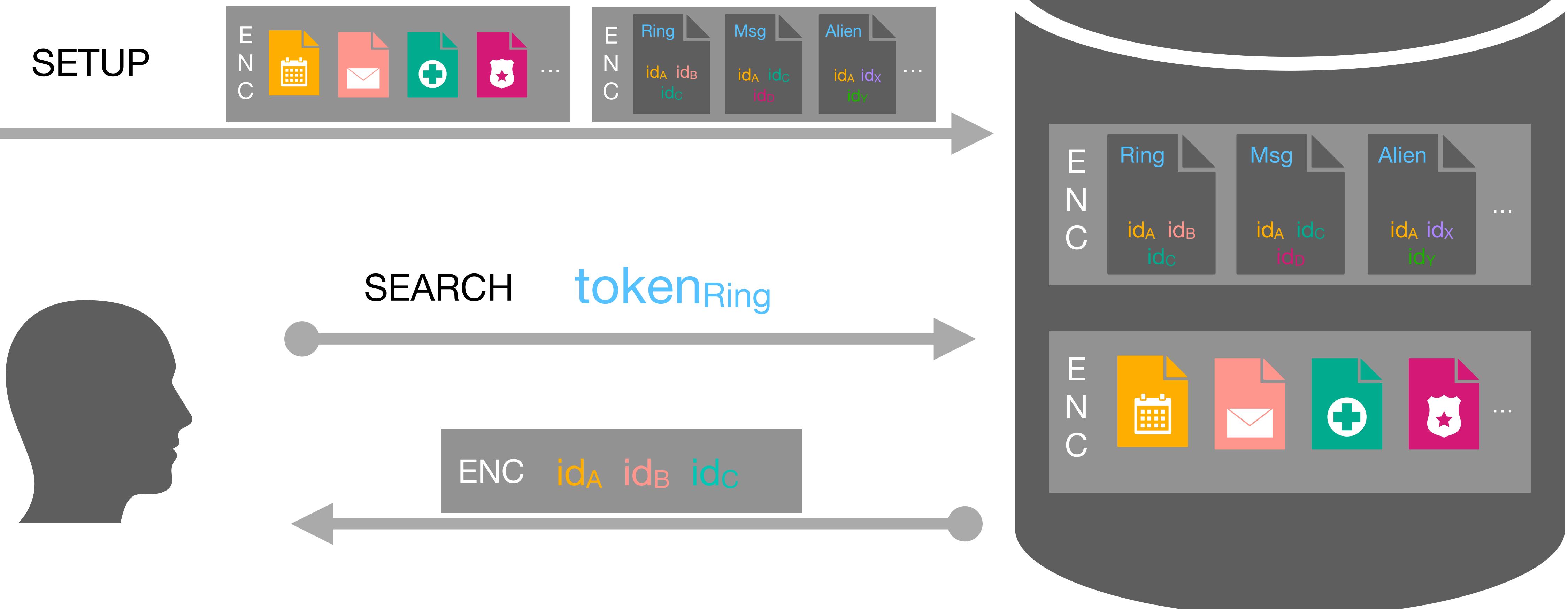
# Motivation

Outsource confidential information



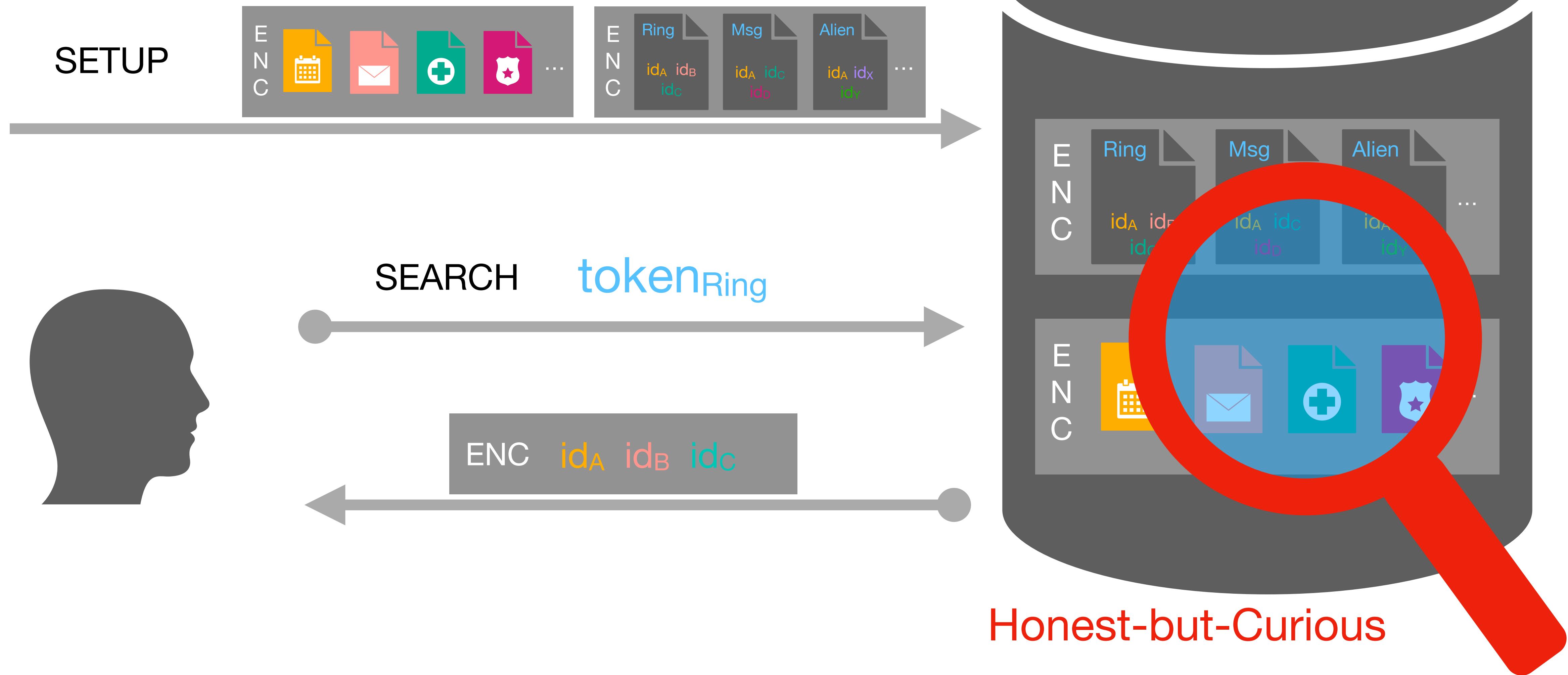
# Motivation

Outsource confidential information



# Motivation

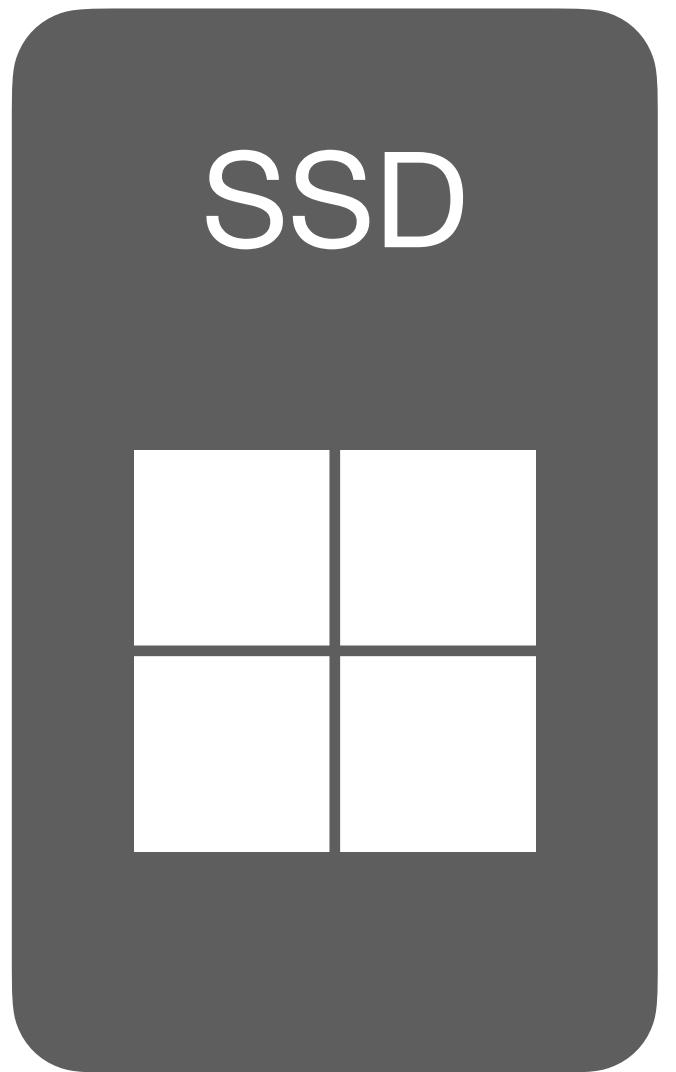
Outsource confidential information



# HDDs vs SSDs



HDD



SSD

## Locality:

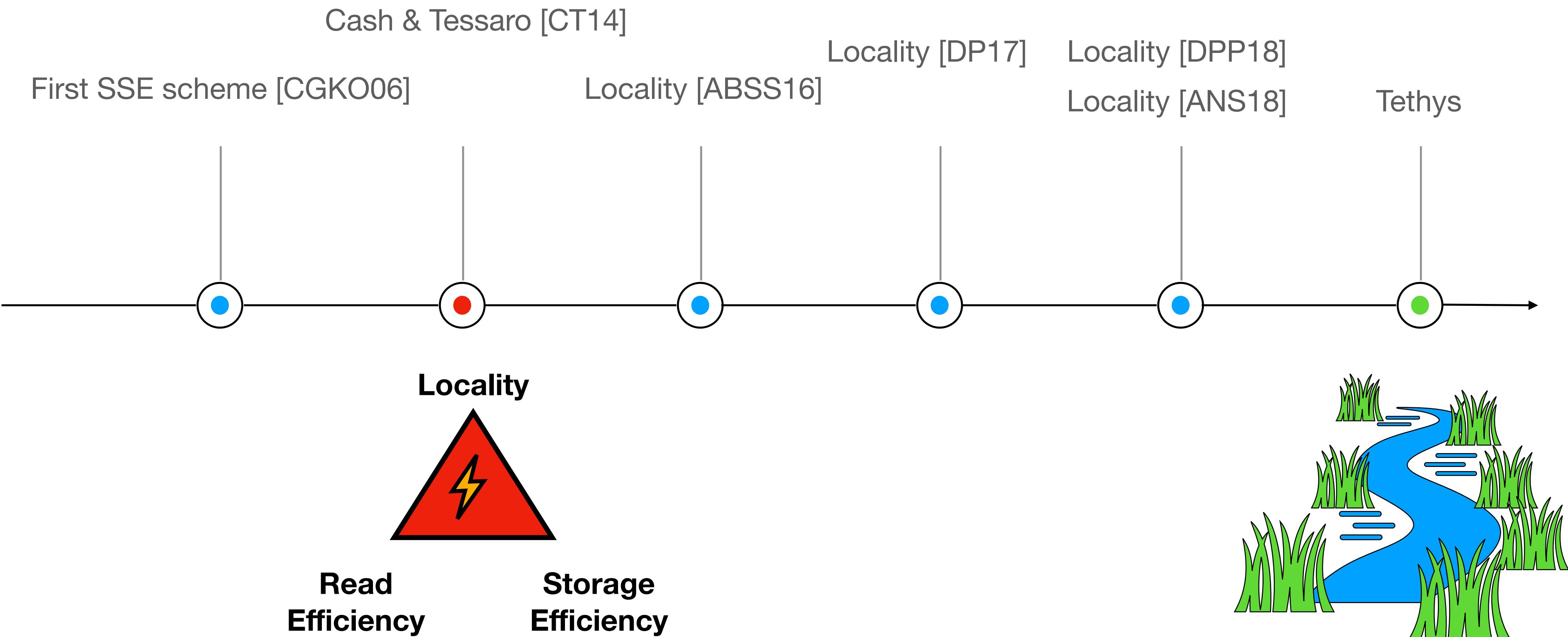
Number of Read  
(non-adjacent)  
Memory Locations

## Page Efficiency:

Number of Read  
Pages per Query

# SSE Research Progress

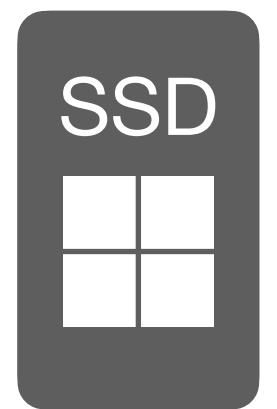
## selected articles



# Goals

## Efficient SSE

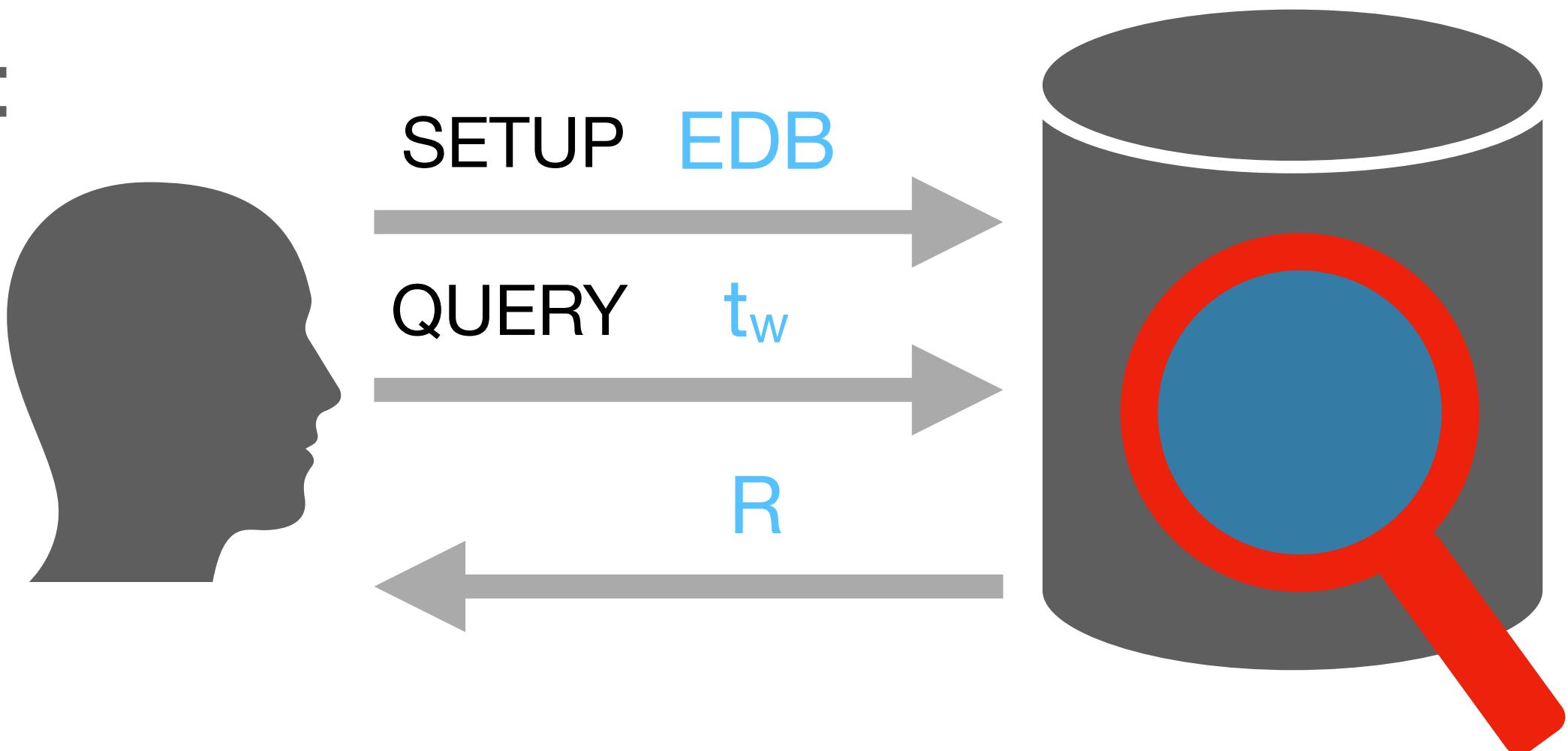
- **Page Efficiency:**
  - Minimize Page Accesses per Query
- **Storage Efficiency:**
  - Minimize Server Storage
- **Security:**
  - Guarantees on Leaked Information



# Security

## Static SSE

- Game:

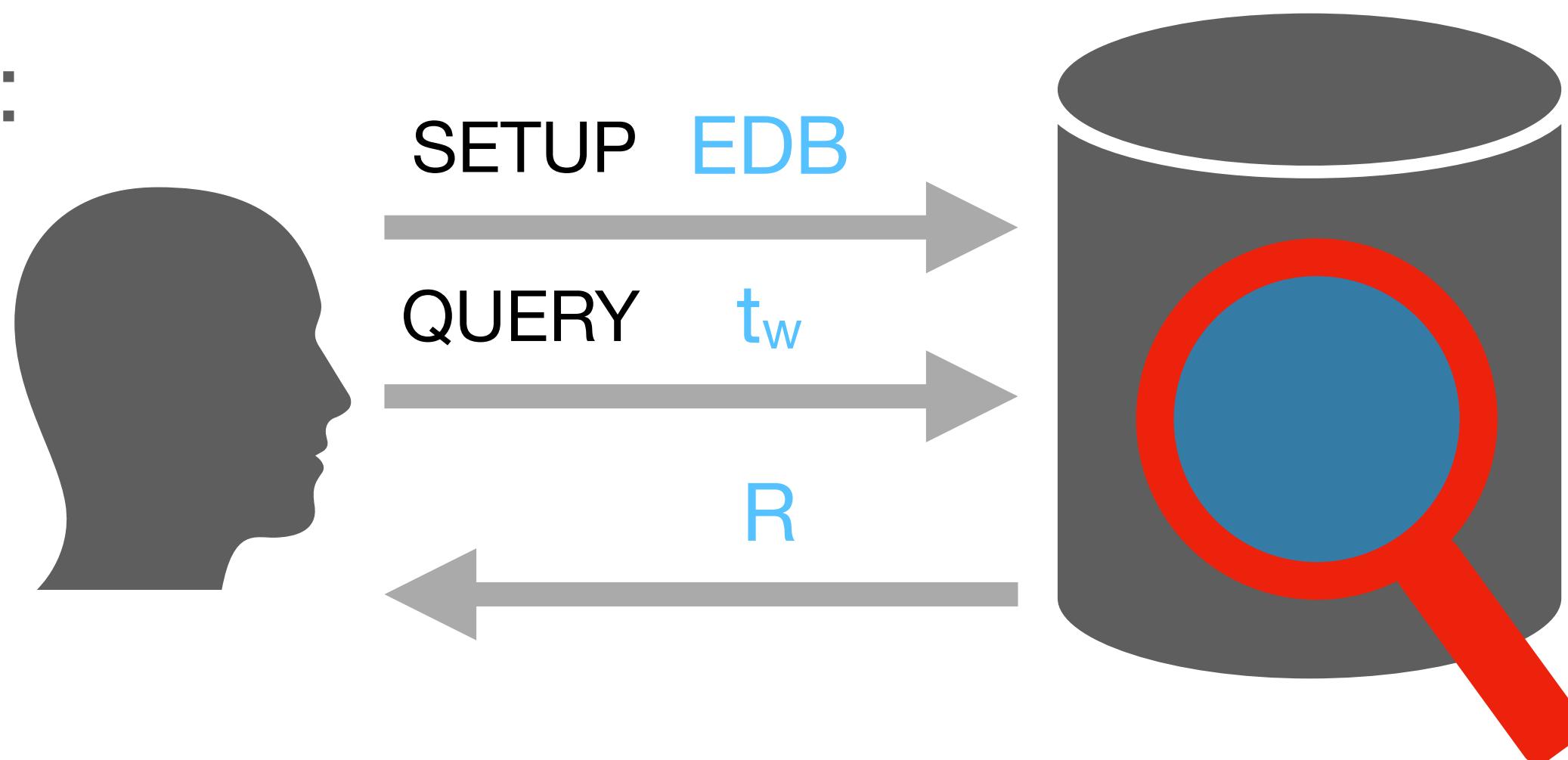


- $\text{EDB} = \text{SSE}.\text{Setup}(\text{DB})$
- $\text{t}_w = \text{SSE}.\text{Search}(\text{DB}, \text{w})$

# Security

## Static SSE

- Game:



- $\text{EDB} = \text{SSE}.\text{Setup}(\mathcal{L}_{\text{setup}}(\text{DB}))$
- $\text{t}_w = \text{SSE}.\text{Search}(\mathcal{L}_{\text{query}}(\text{DB}, w))$

- Leakage Function  $\mathcal{L} = (\mathcal{L}_{\text{setup}}, \mathcal{L}_{\text{query}})$ :

- $\mathcal{L}_{\text{setup}}(\text{DB})$

- $\text{size(DB)}$

- $\mathcal{L}_{\text{query}}(\text{DB}, w)$

- $\text{size(ids}_w)$ , search pattern

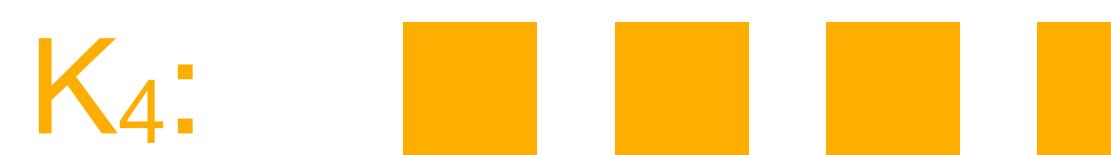
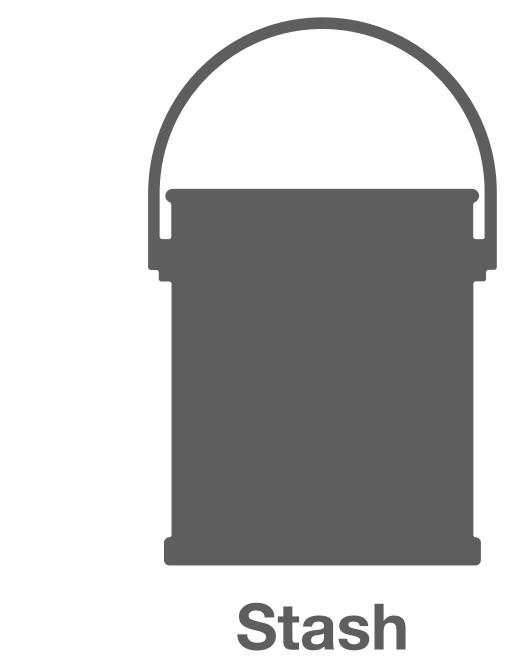
# Contribution

## This Work

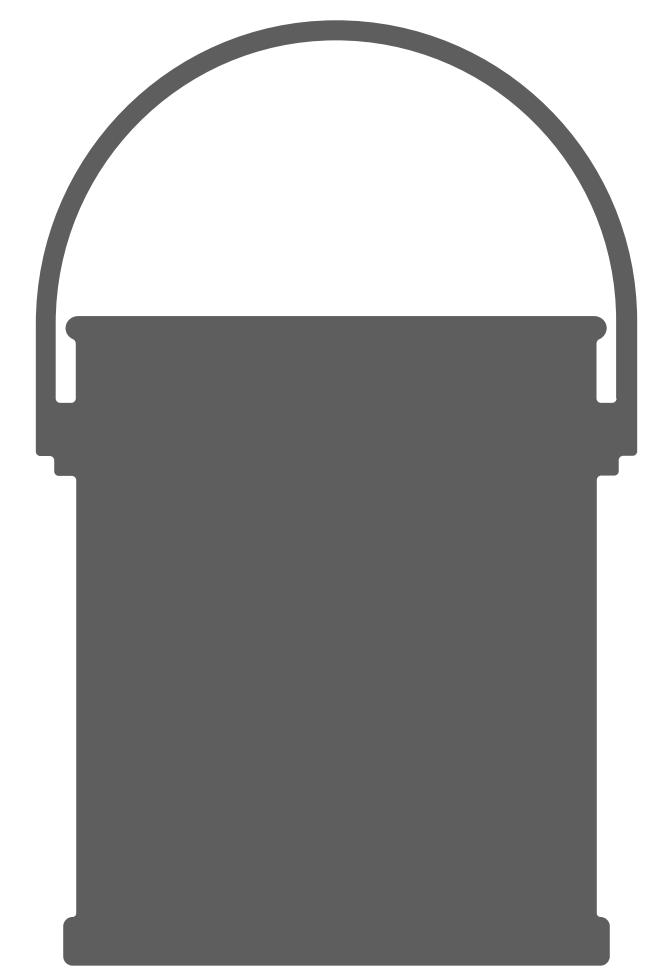
- Page Efficiency:
  - Definition
- Data Independent Packing (DIP):
  - Framework: SSE from DIP
- Tethys:
  - Efficient DIP scheme
  - Based on Cuckoo-Hashing for Weighted Items
  - SSE: optimal page and storage efficiency

# DIP

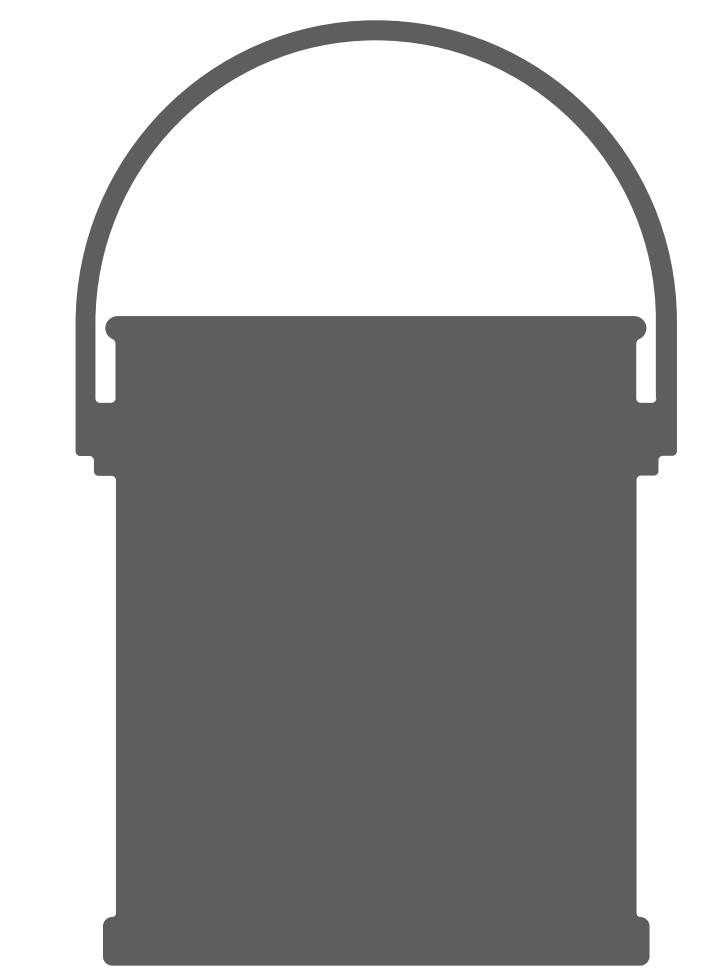
## Data Independent Packing



$n$



...

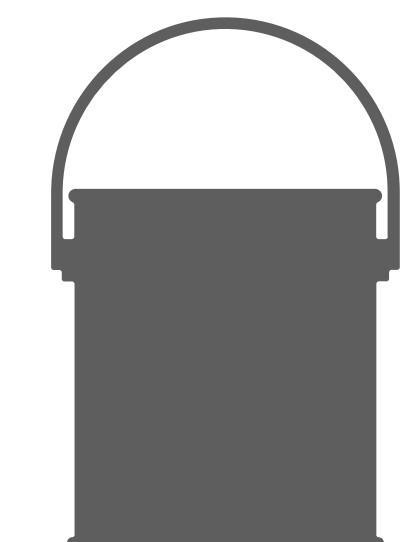


$m$

$p$

# DIP

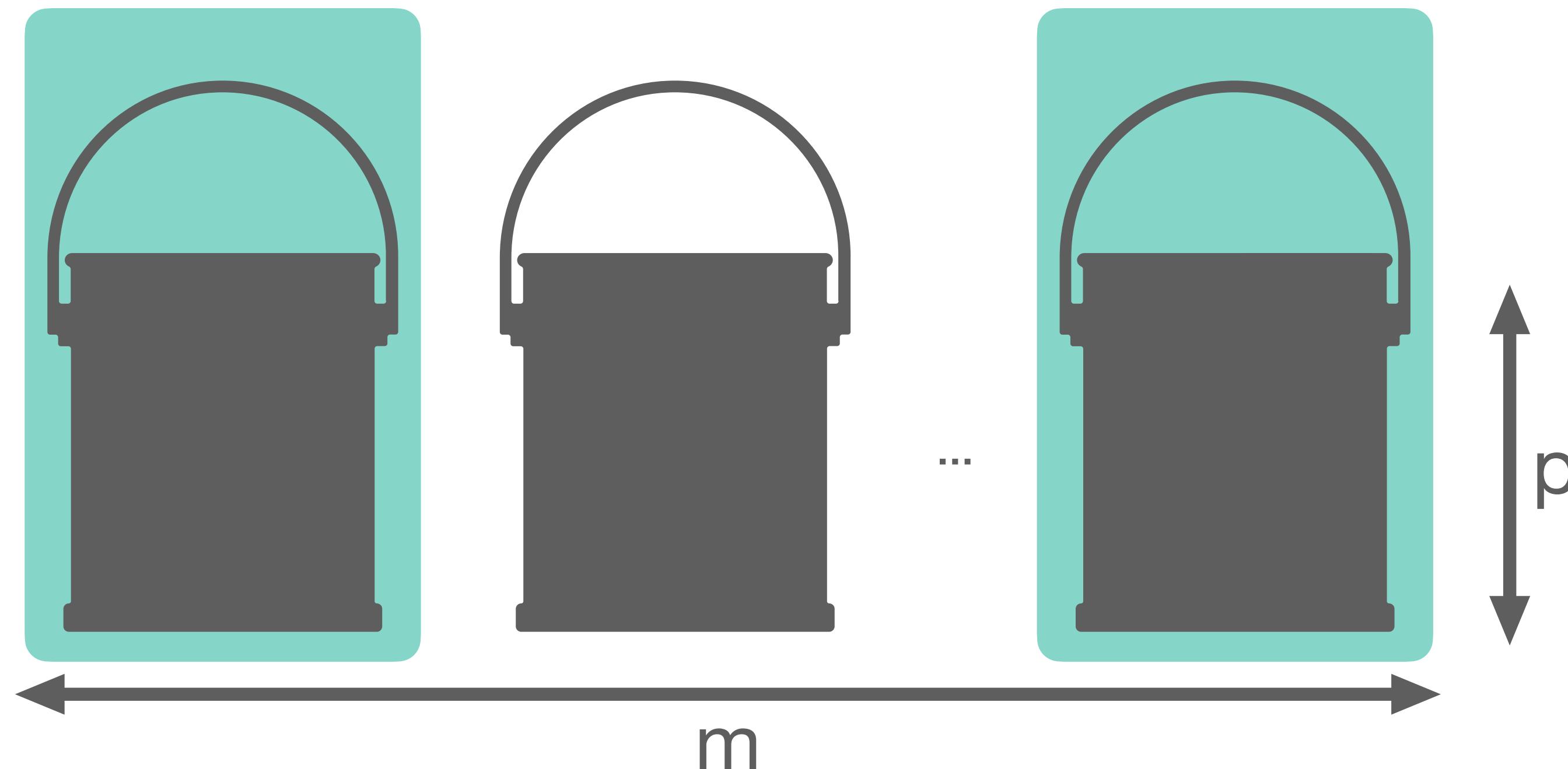
## Data Independent Packing



Stash

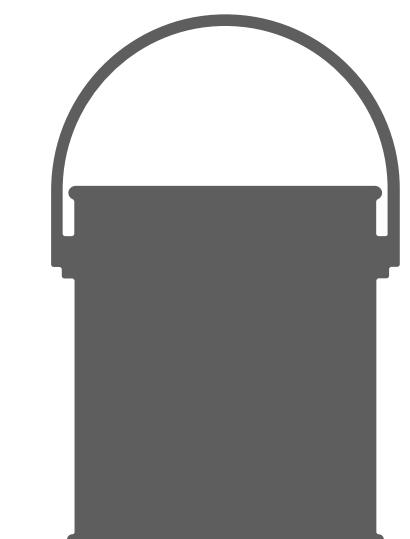


$n$

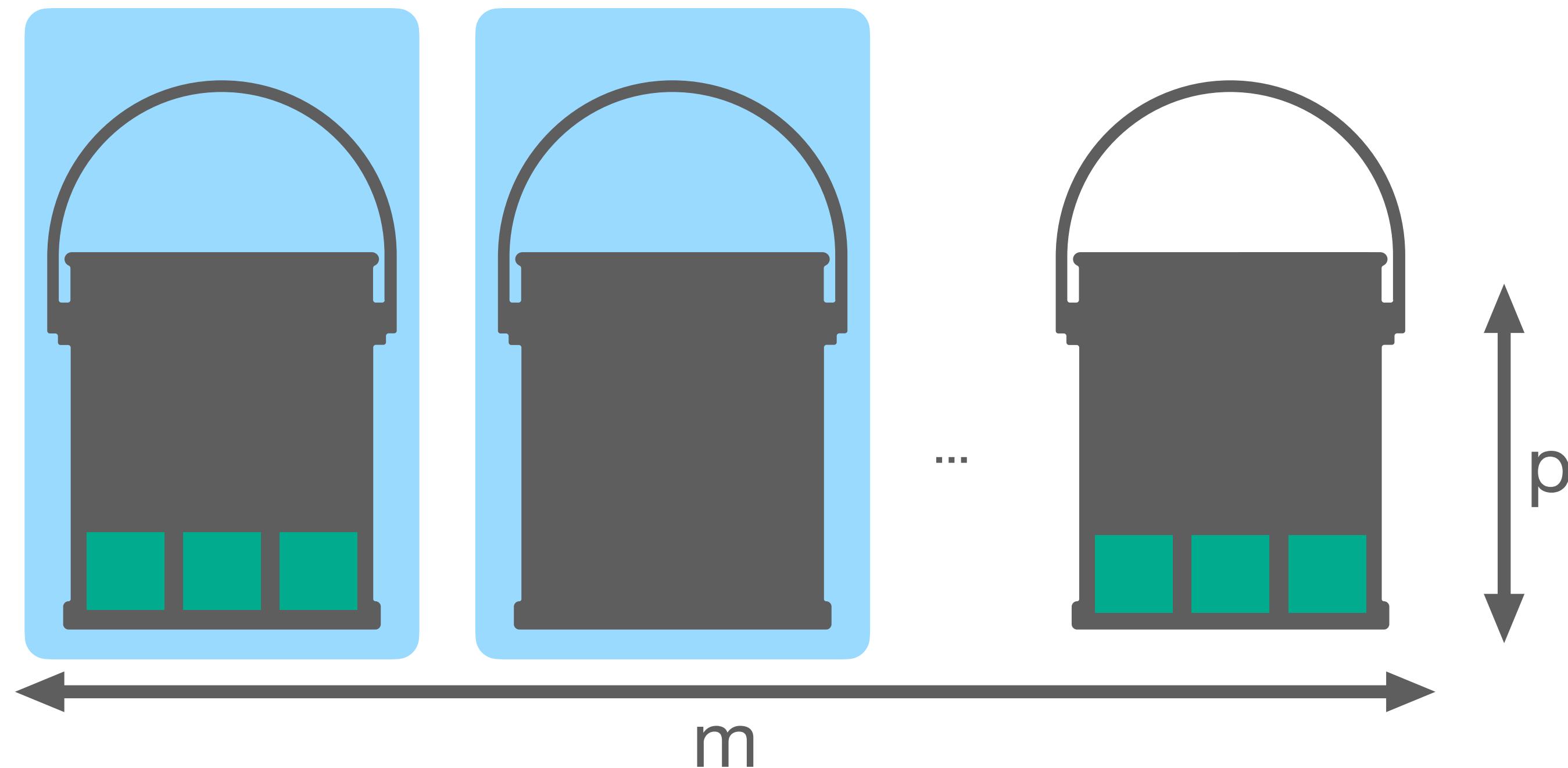


# DIP

## Data Independent Packing

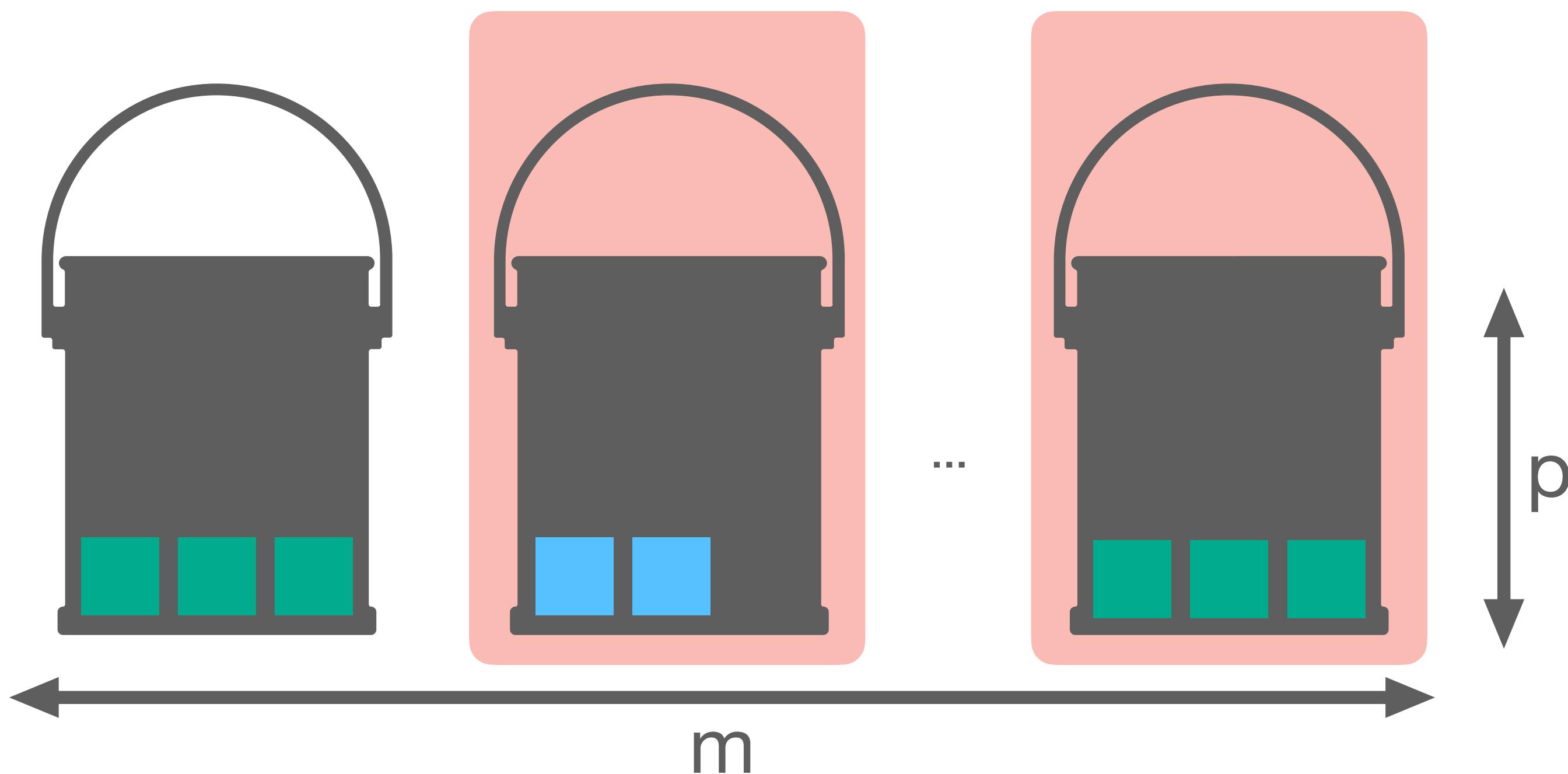
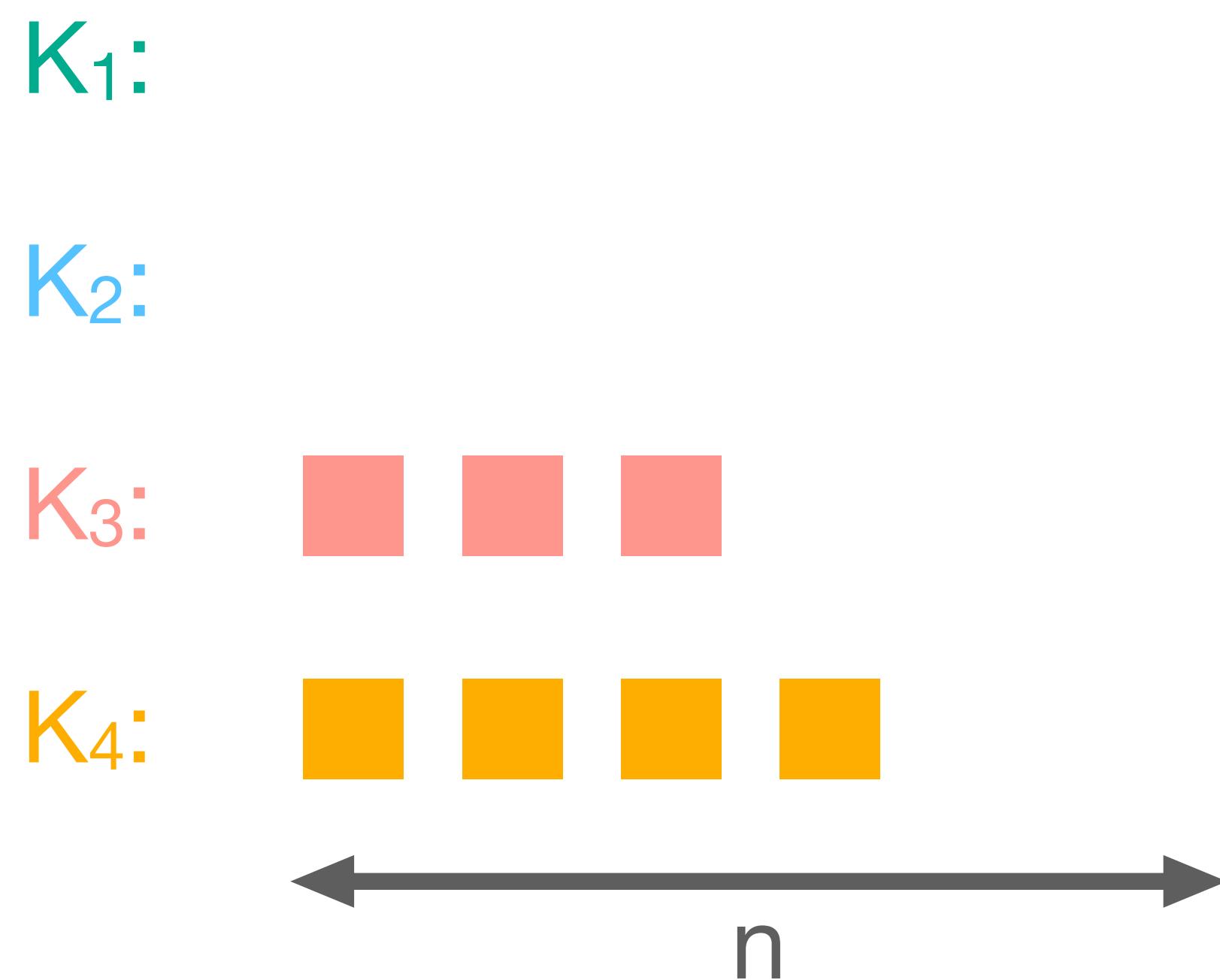
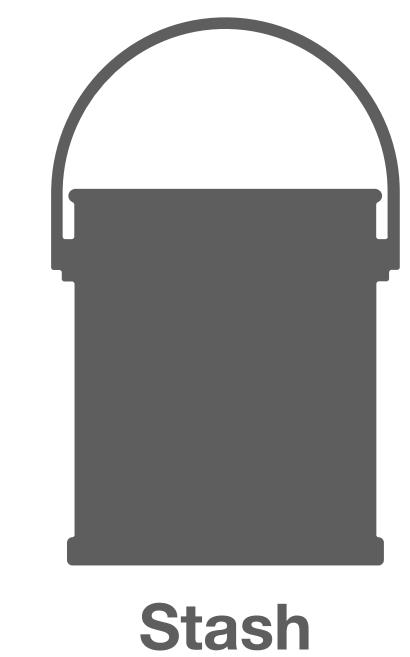


Stash



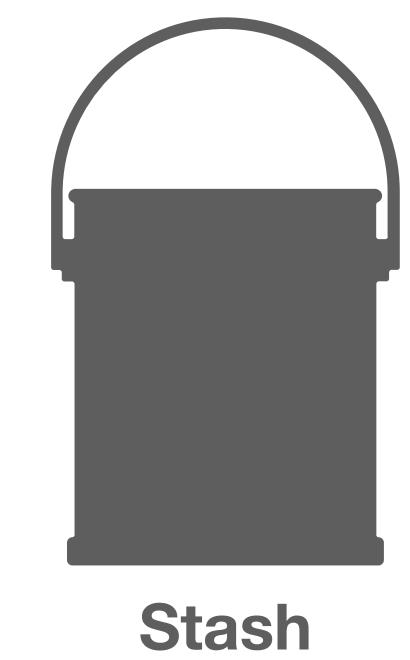
# DIP

## Data Independent Packing



# DIP

## Data Independent Packing



$K_1:$

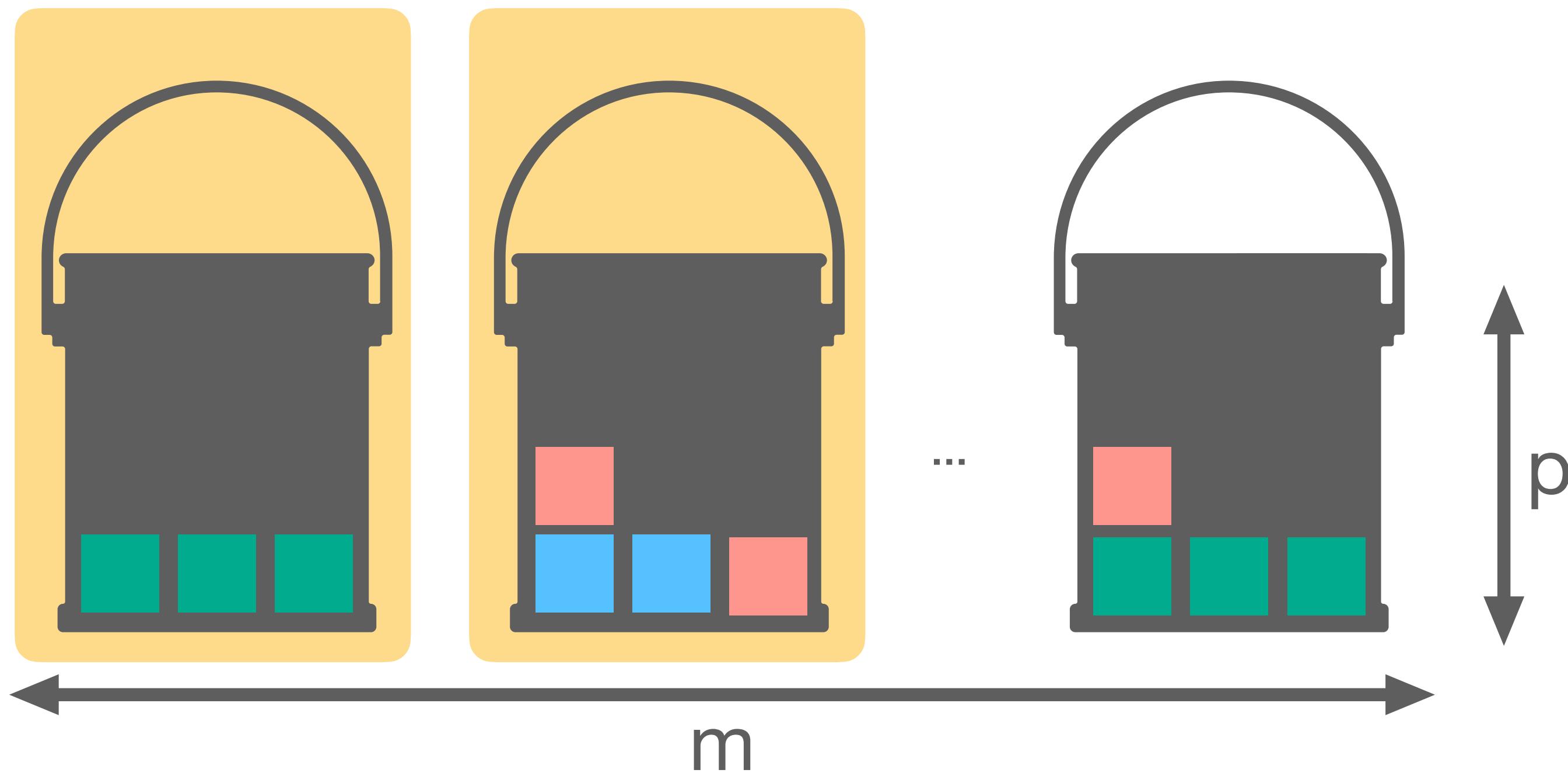
$K_2:$

$K_3:$

$K_4:$



$n$



# DIP

## Definition

**Size( $n$ ):**

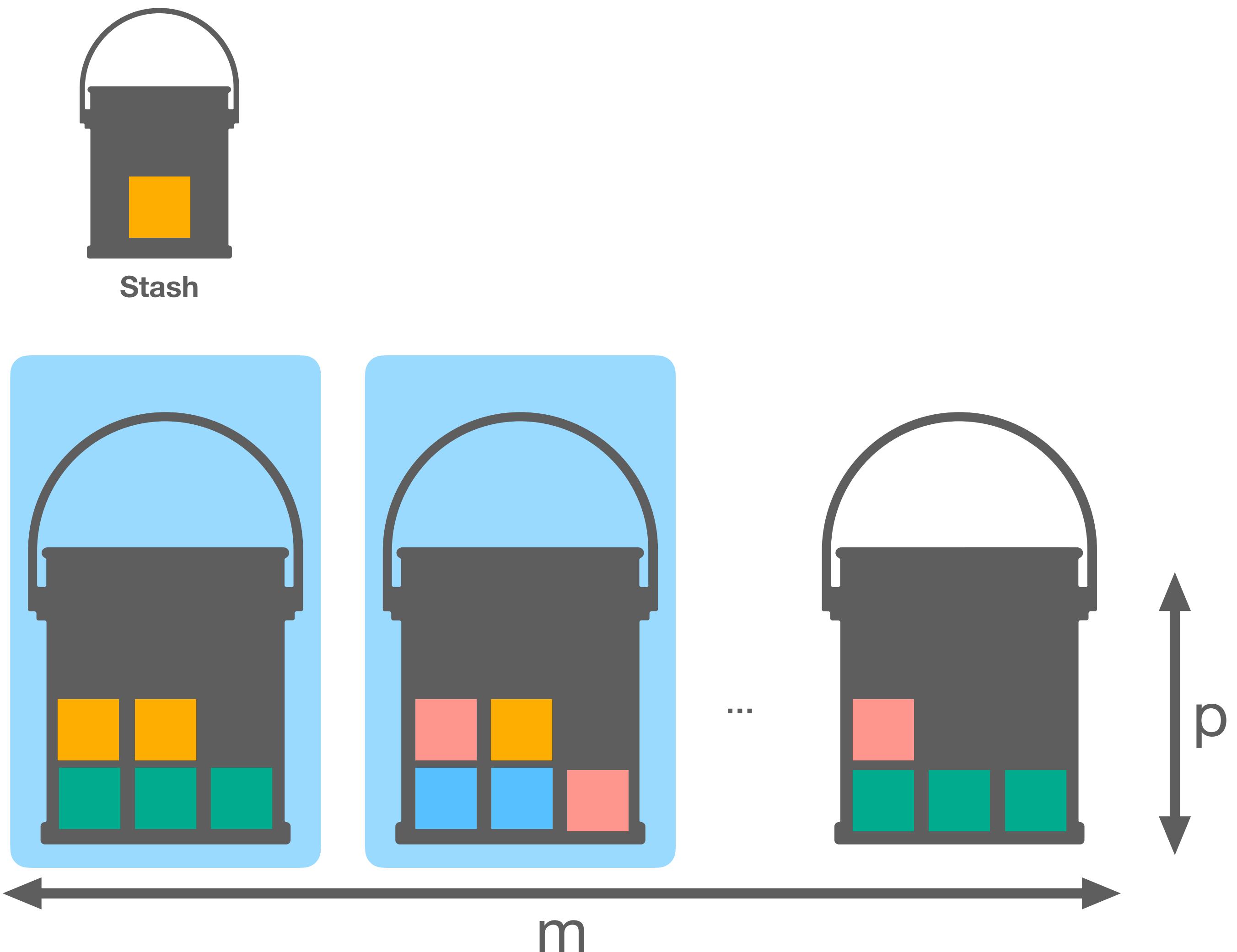
- Return  $m$ , the Number of Buckets

**Build( $M$ ):**

- Returns Buckets filled with Items from Multi-Map  $M$
- Allows for Stash  $S$

**Lookup( $m, K_i, \ell_i$ ):**

- Return Bucket Indices



# DIP

## Efficiency

### Lookup Efficiency:

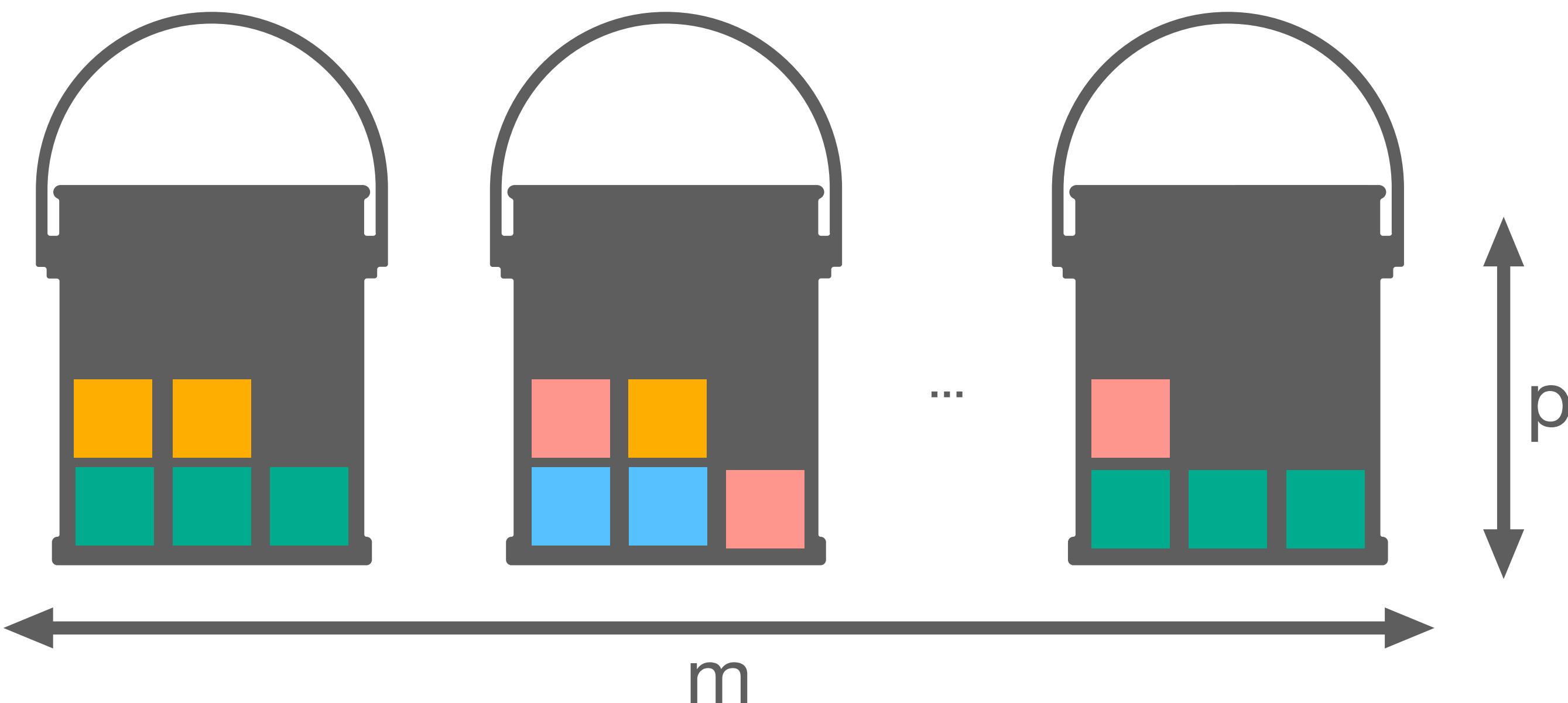
- number of buckets per keyword

### Storage Efficiency:

- $(m \cdot p) / n$

### Stash Size:

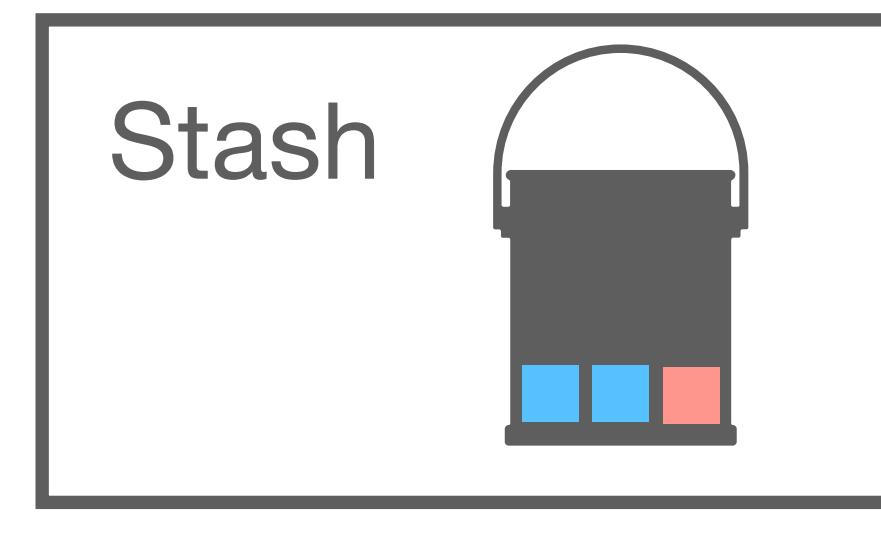
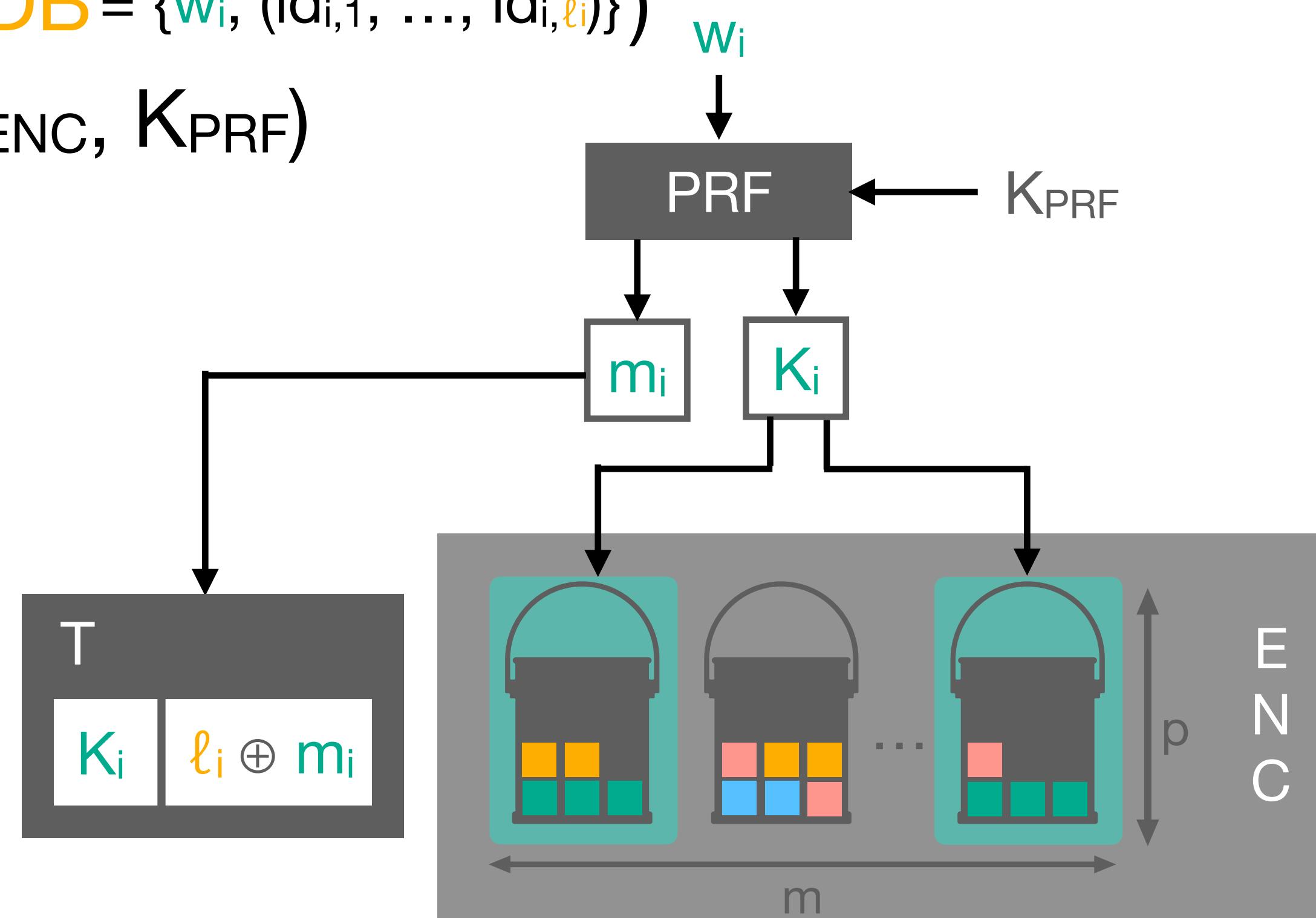
- Size of Stash  $S$



# DIP $\Rightarrow$ SSE Framework

**Setup**( $\text{DB} = \{w_i, (id_{i,1}, \dots, id_{i,\ell_i})\}$ )

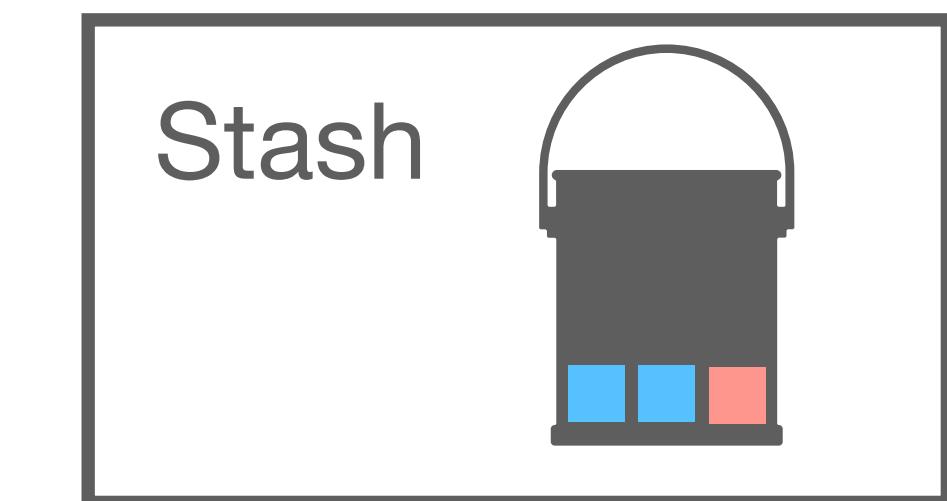
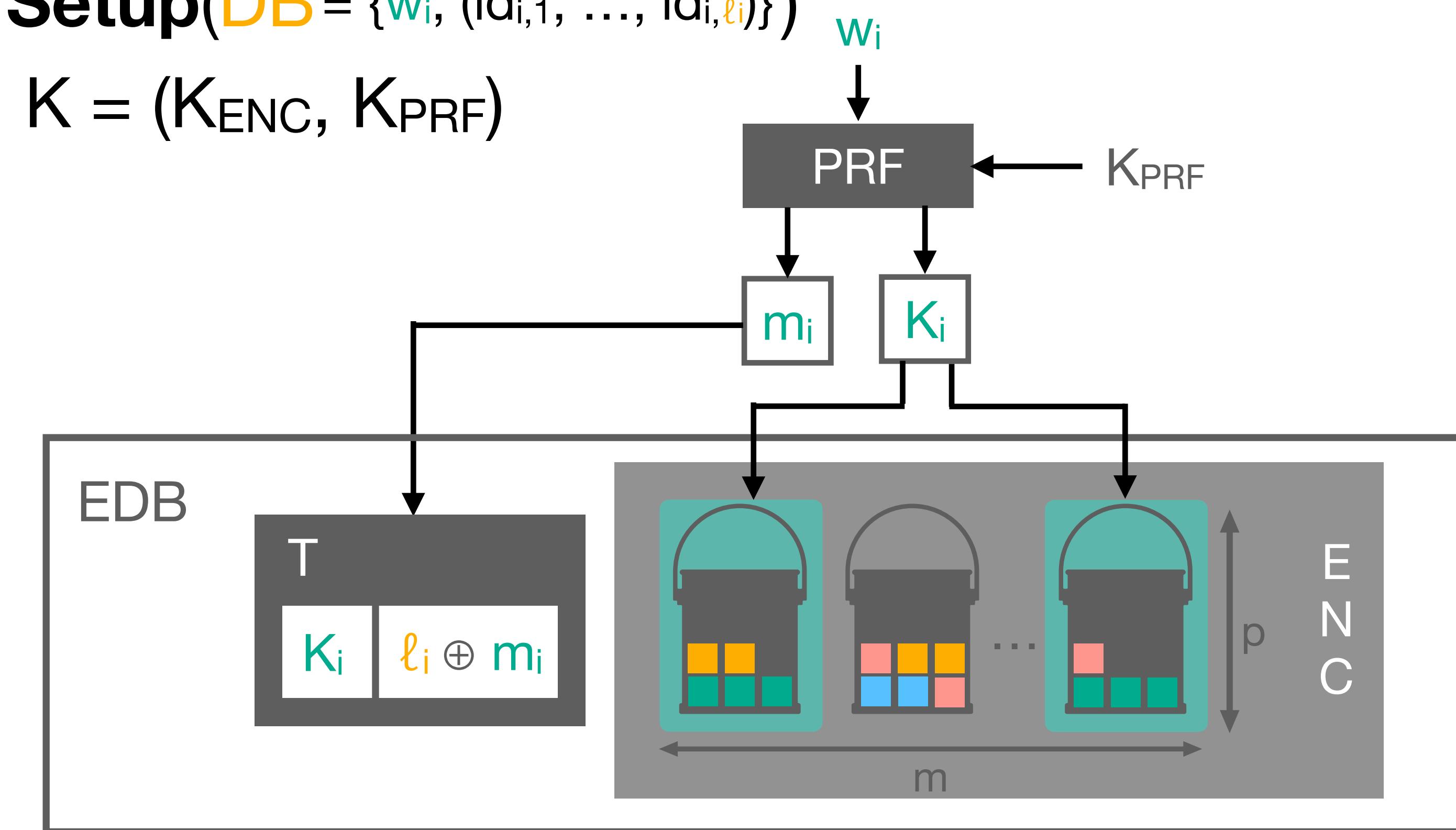
$$K = (K_{\text{ENC}}, K_{\text{PRF}})$$



# DIP $\Rightarrow$ SSE Framework

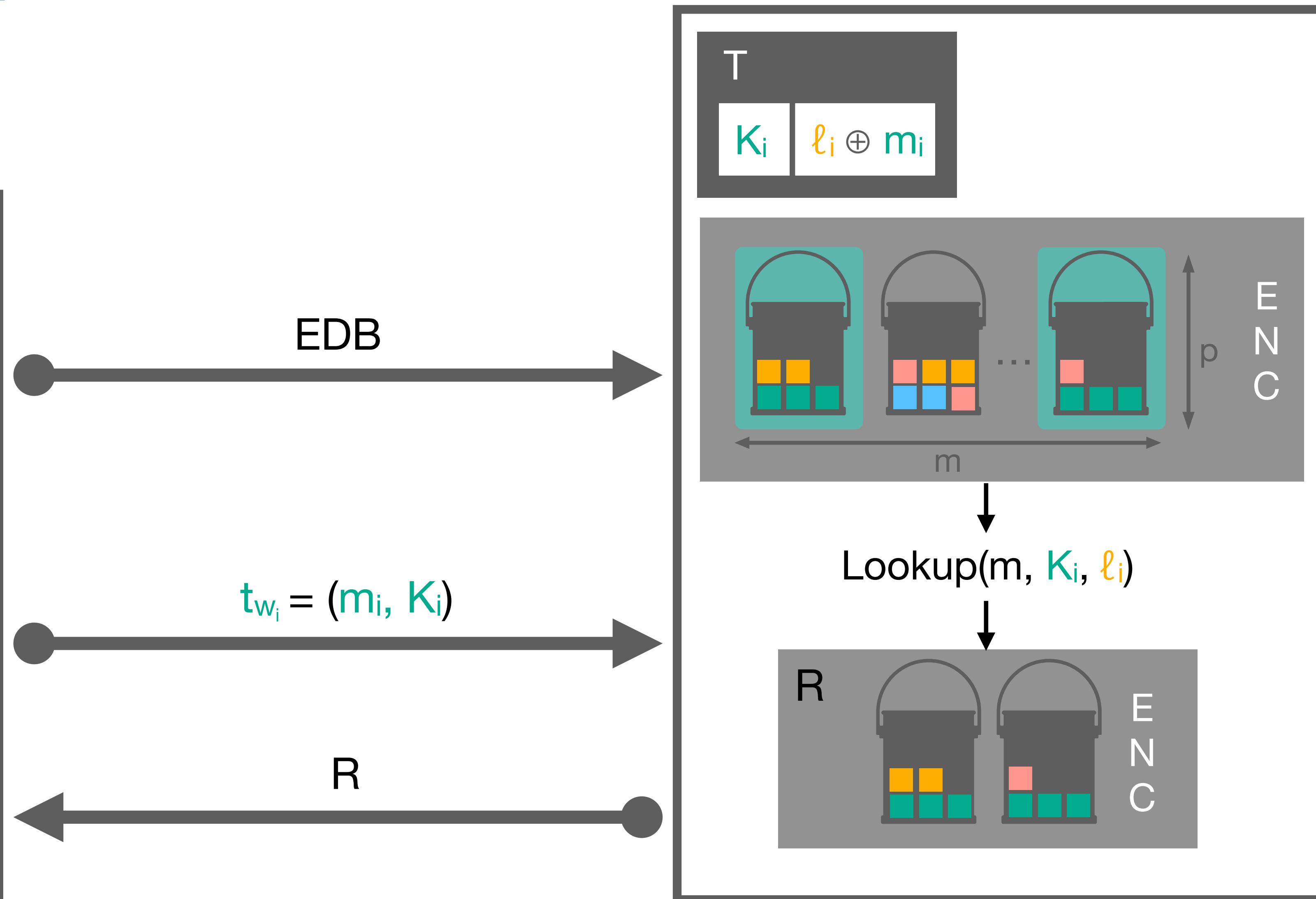
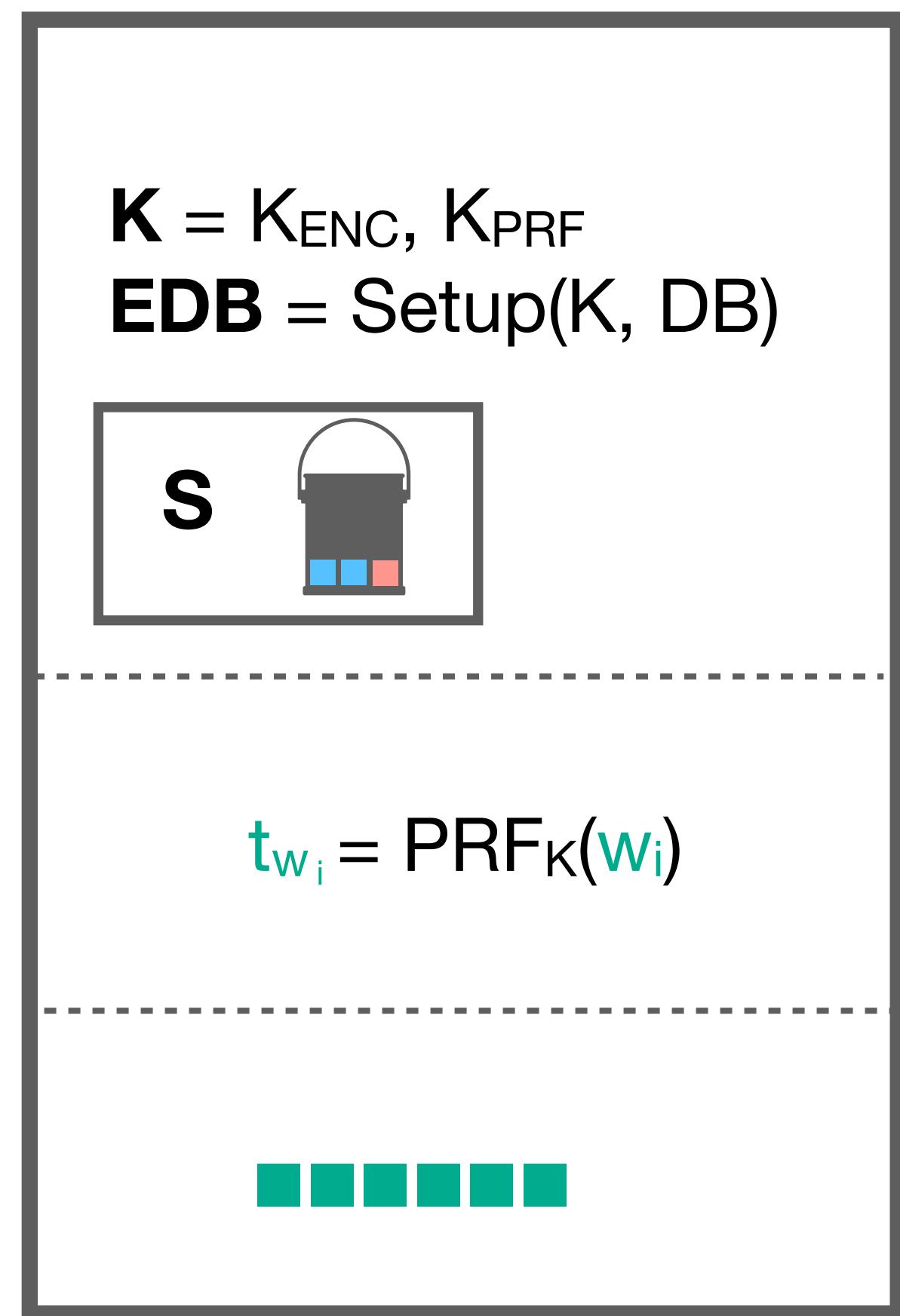
**Setup**( $\text{DB} = \{w_i, (id_{i,1}, \dots, id_{i,\ell_i})\}$ )

$$K = (K_{\text{ENC}}, K_{\text{PRF}})$$



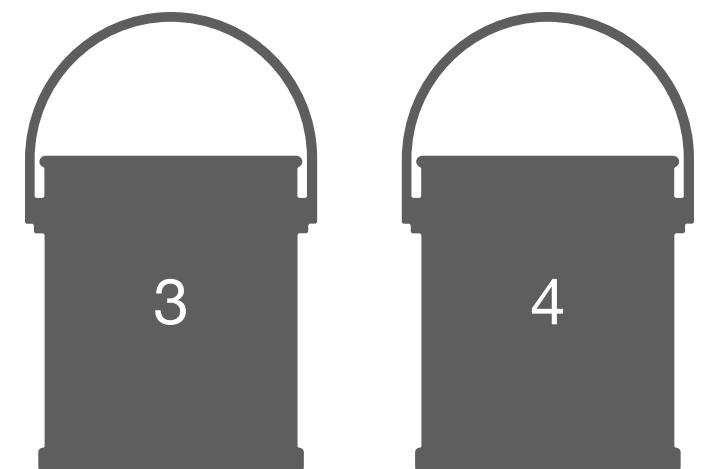
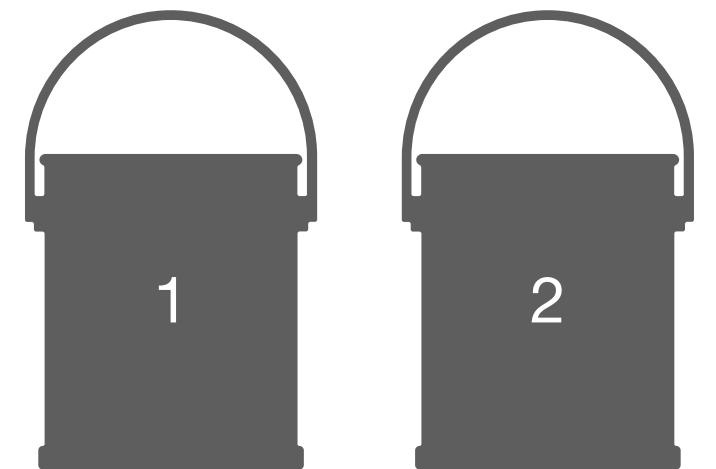
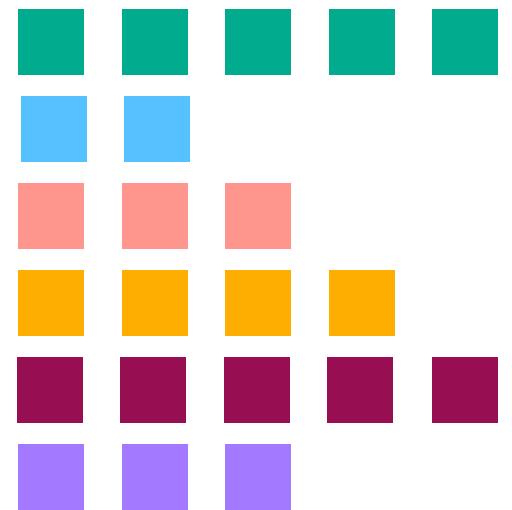
# DIP $\Rightarrow$ SSE

## Framework



# Tethys

## Efficient DIP



1

2

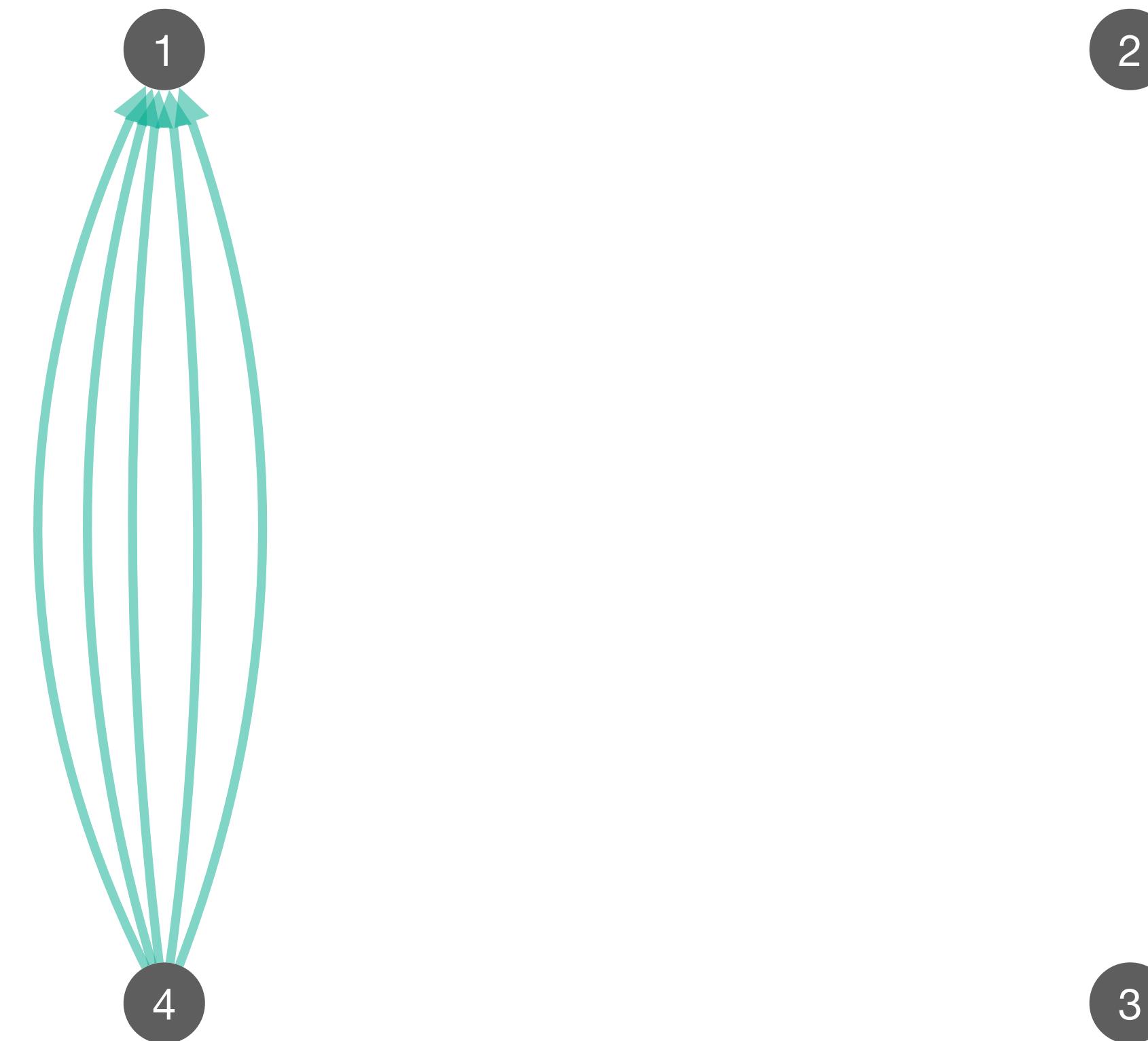
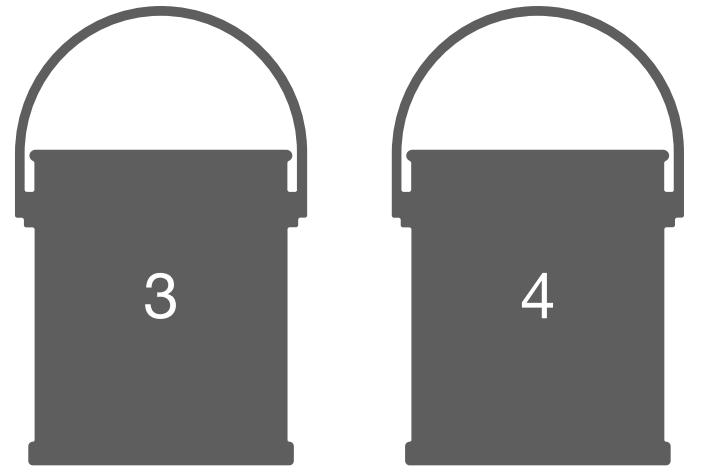
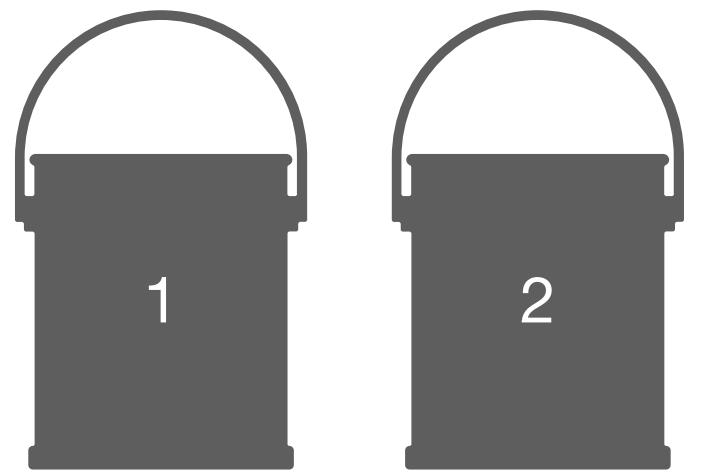
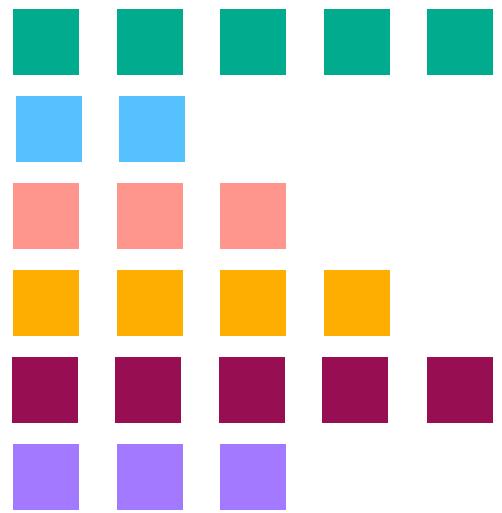
4

3

$p = 5$

# Tethys

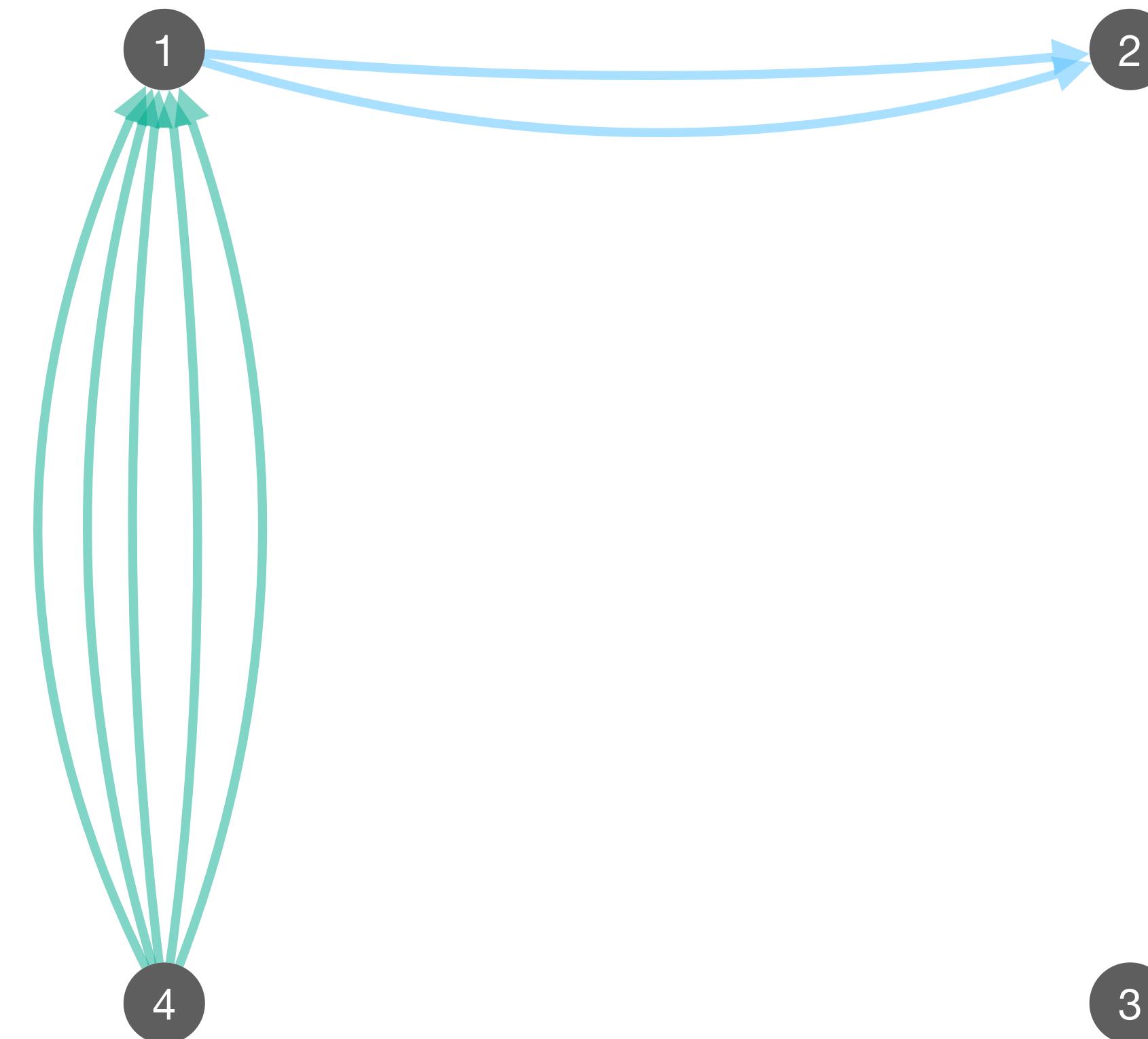
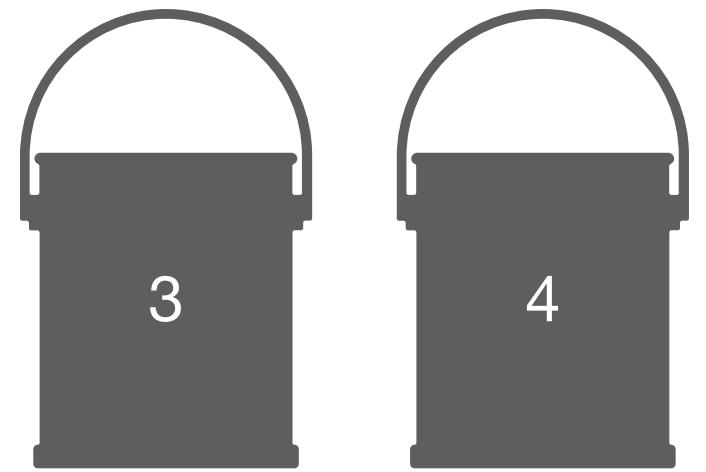
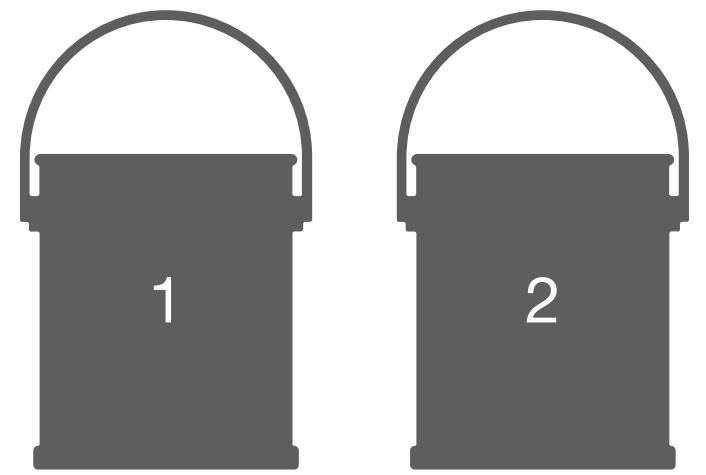
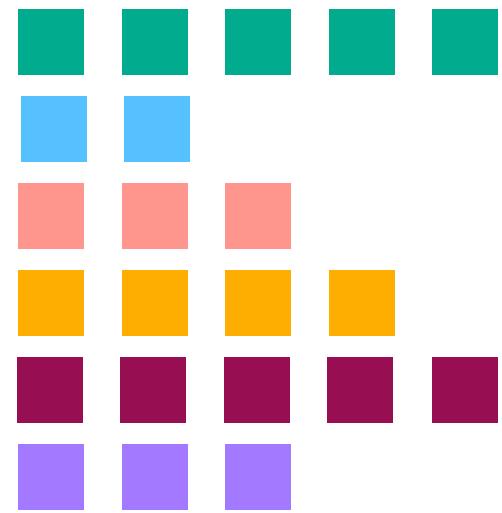
## Efficient DIP



$p = 5$

# Tethys

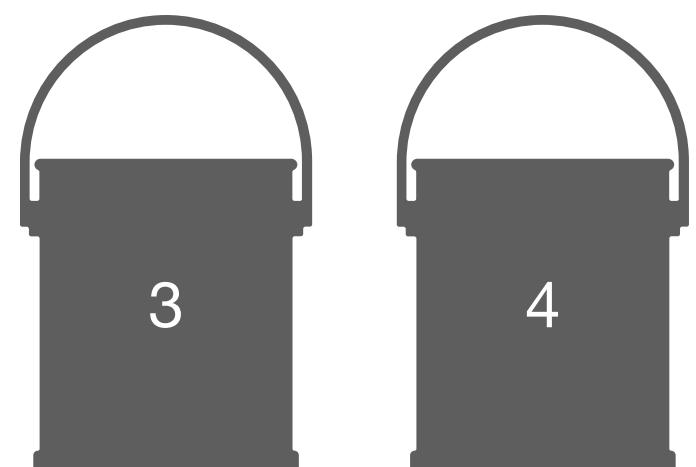
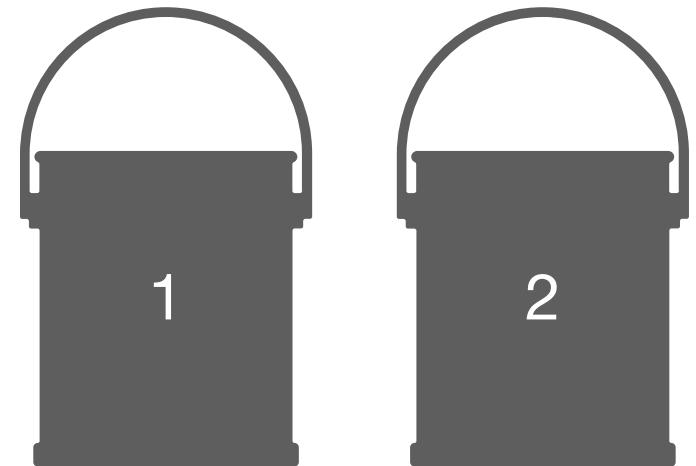
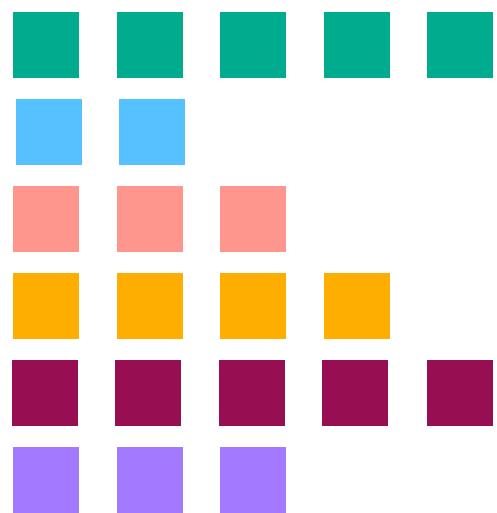
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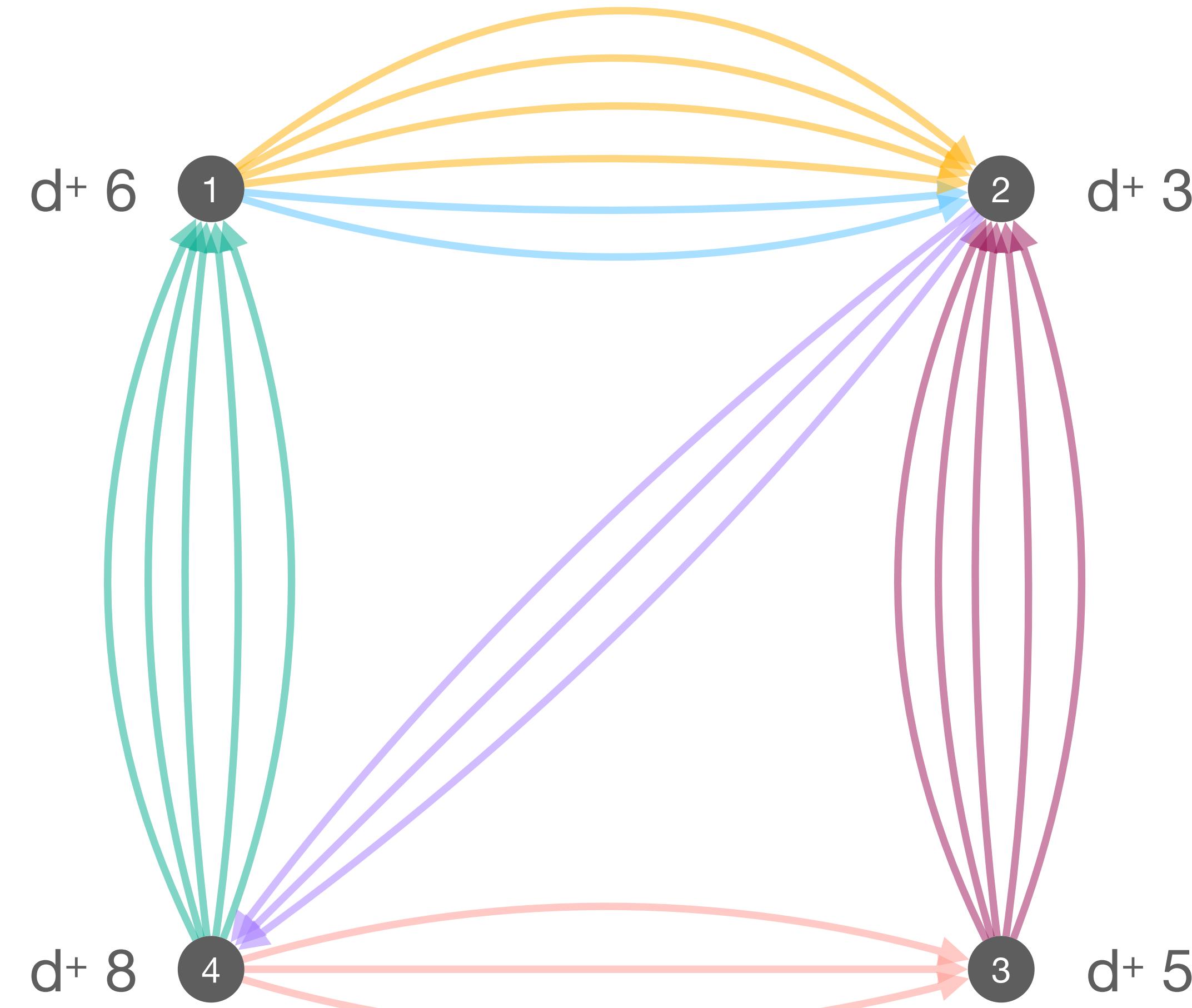
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# Tethys

## Efficient DIP

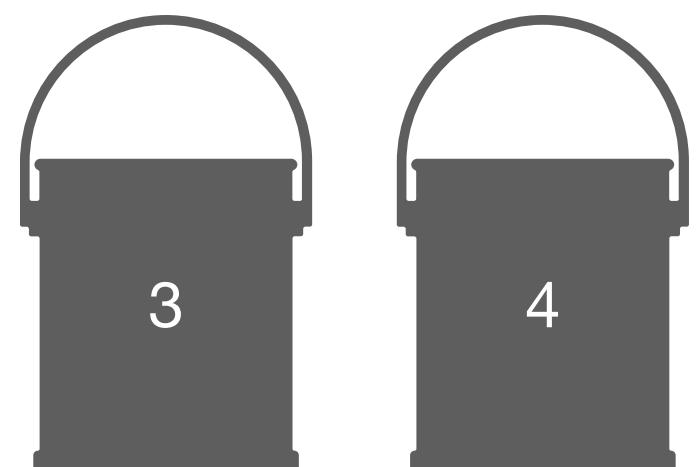
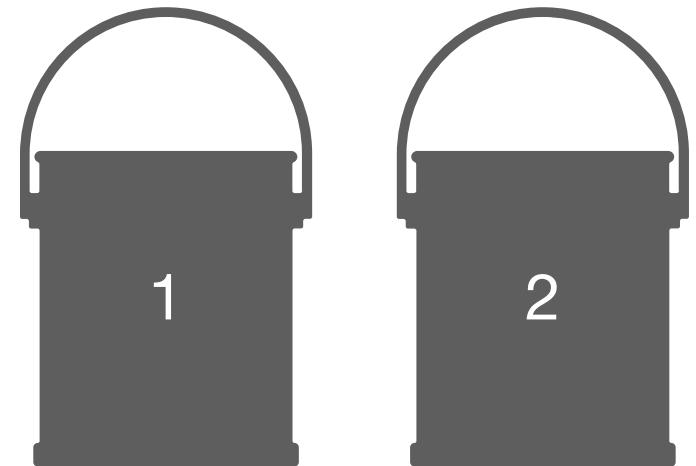
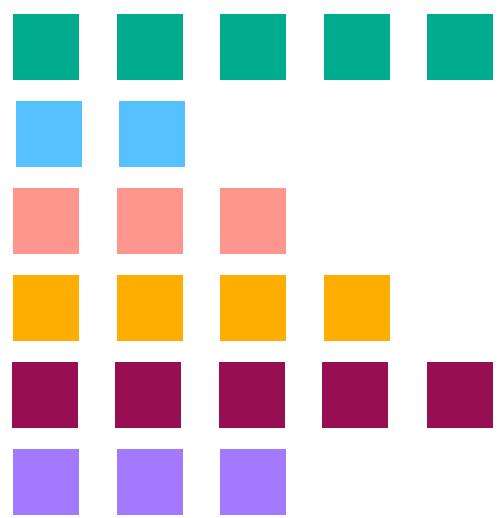


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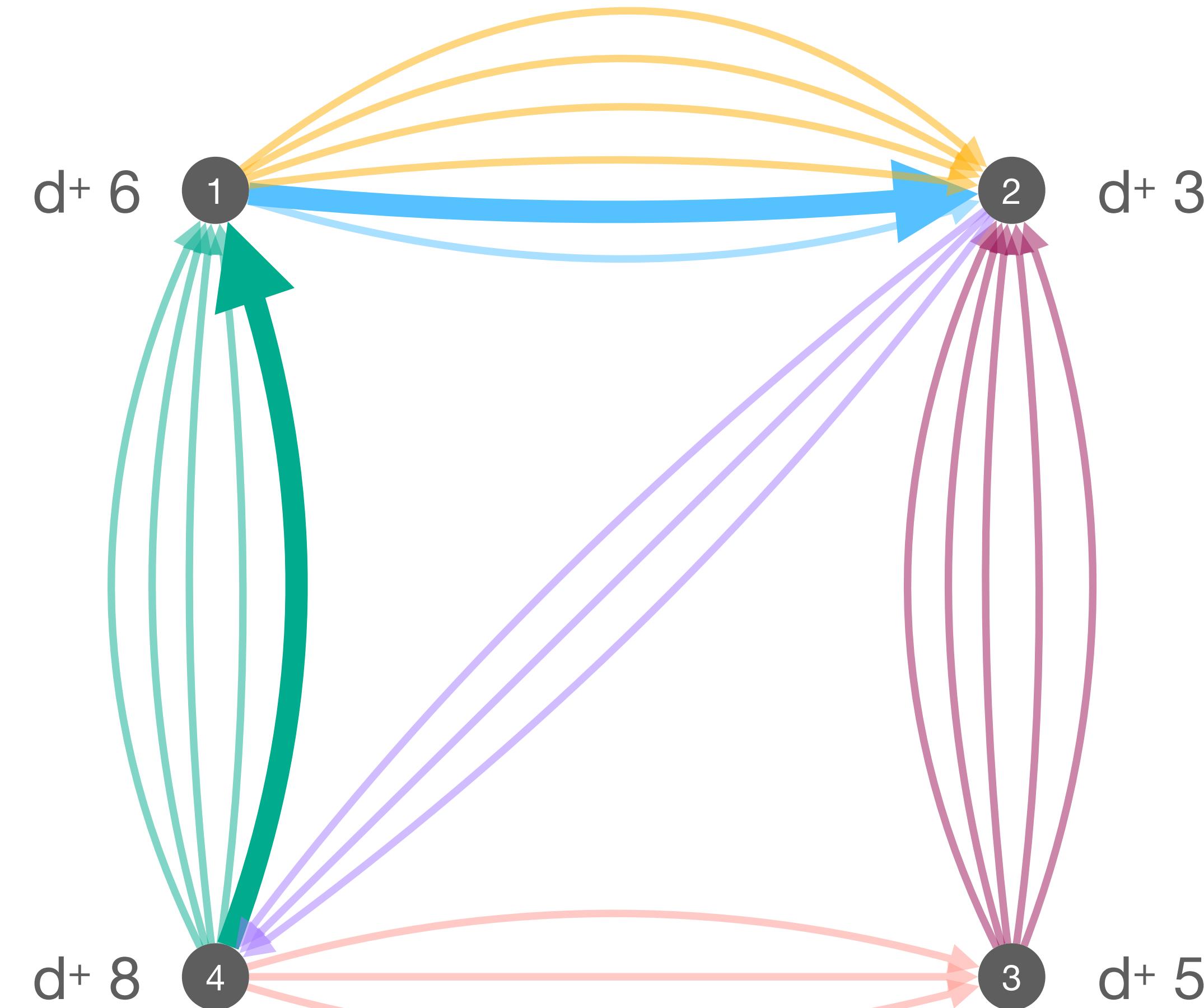


# Tethys

## Efficient DIP

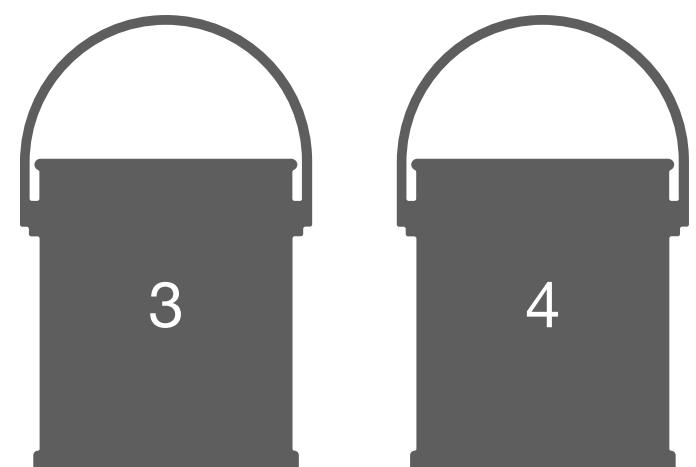
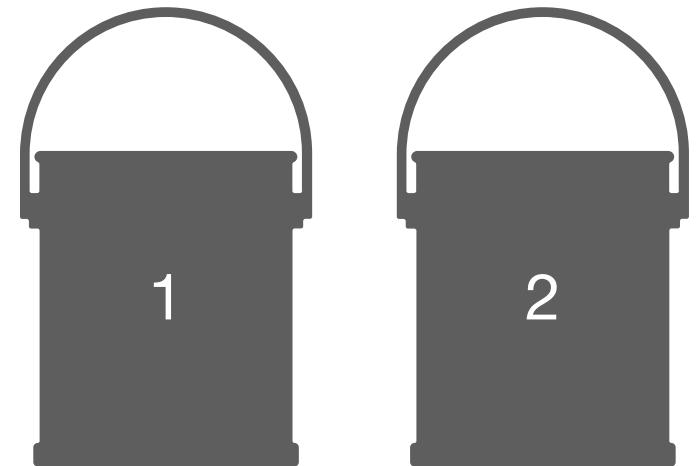
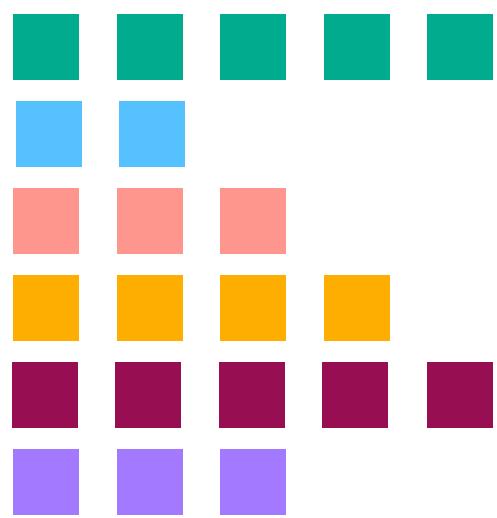


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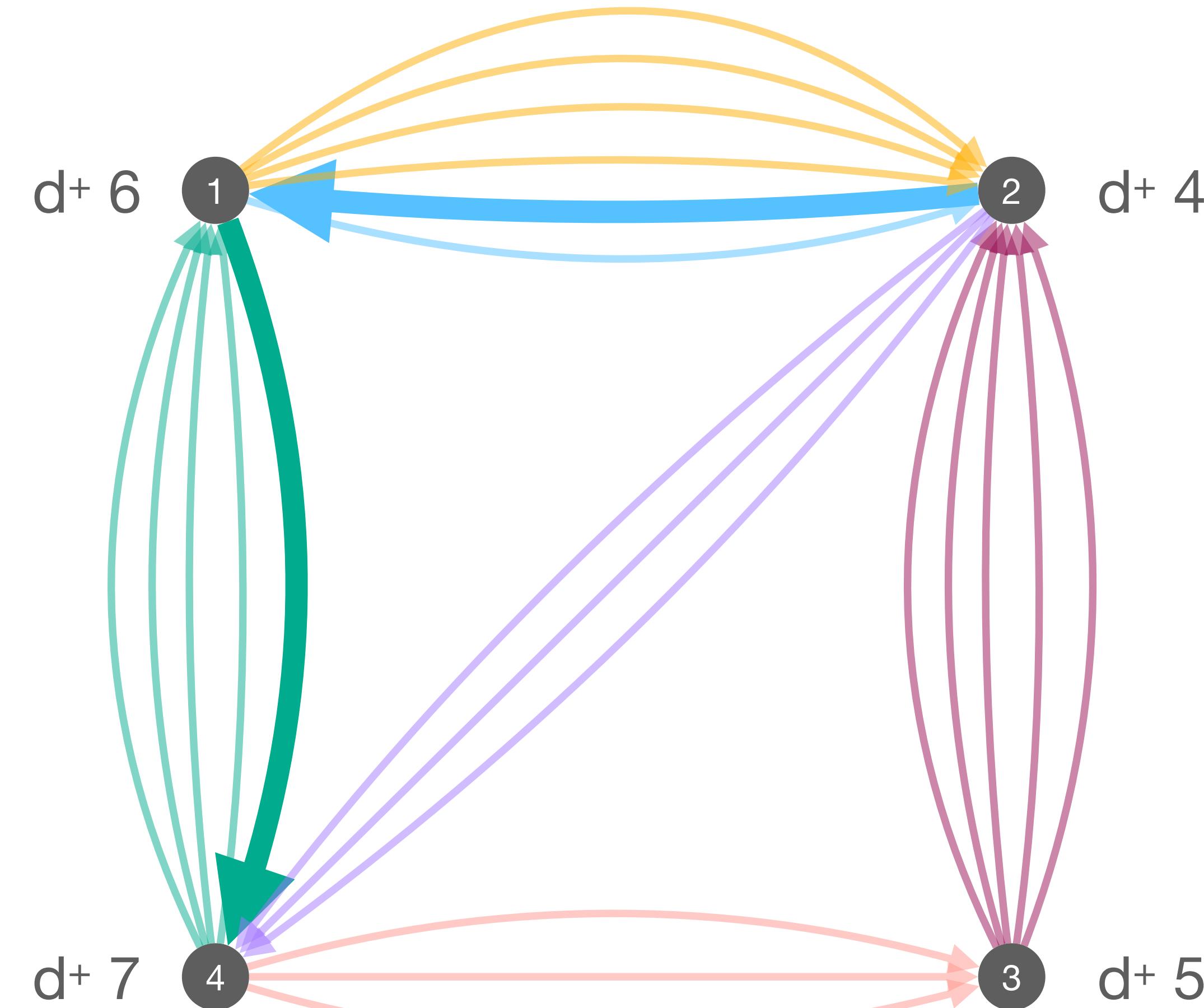


# Tethys

## Efficient DIP

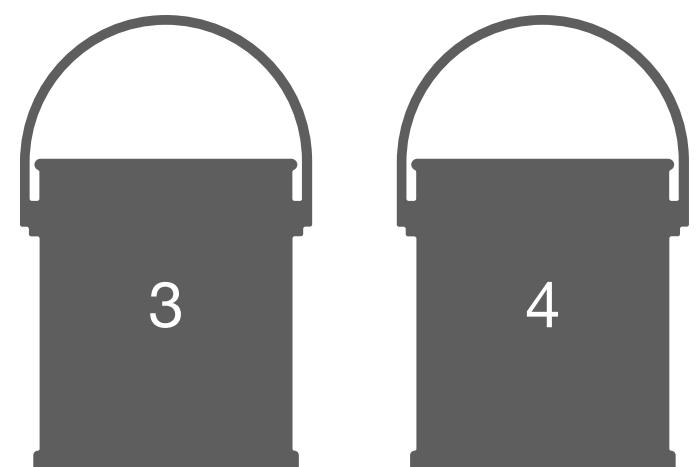
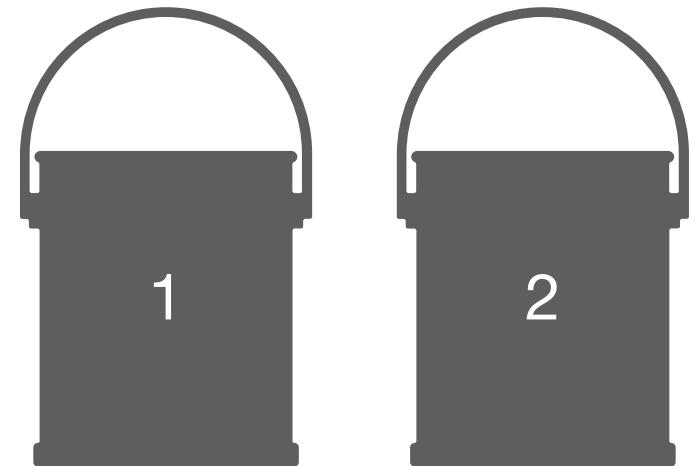
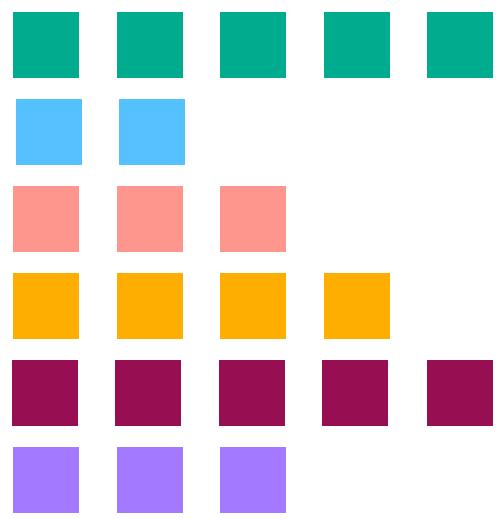


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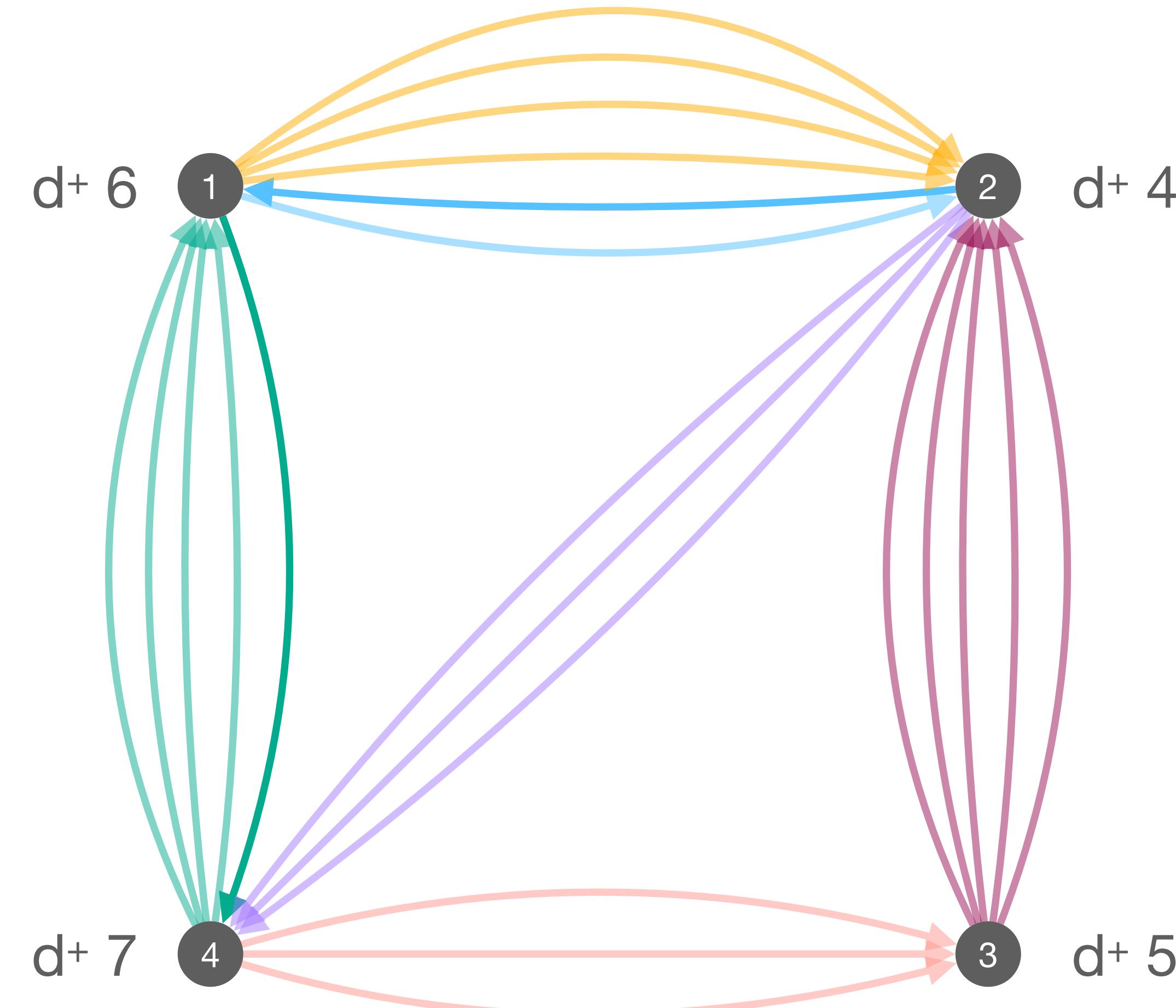


# Tethys

## Efficient DIP

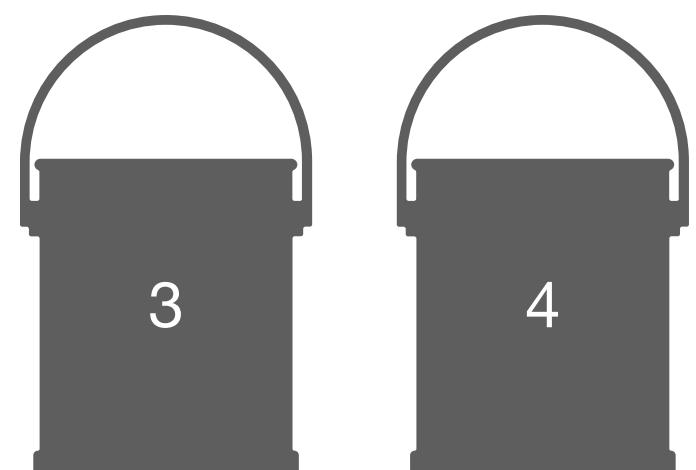
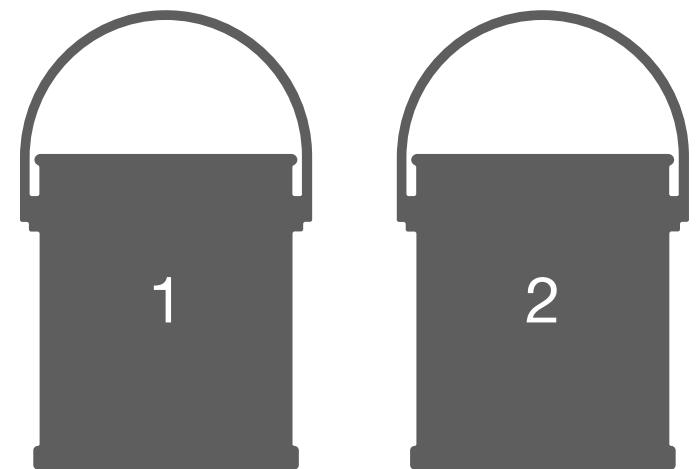
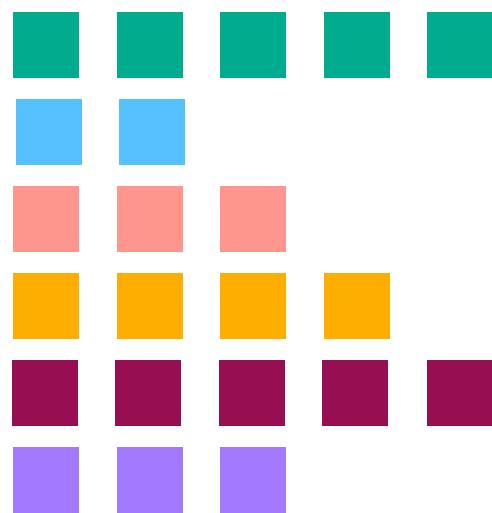


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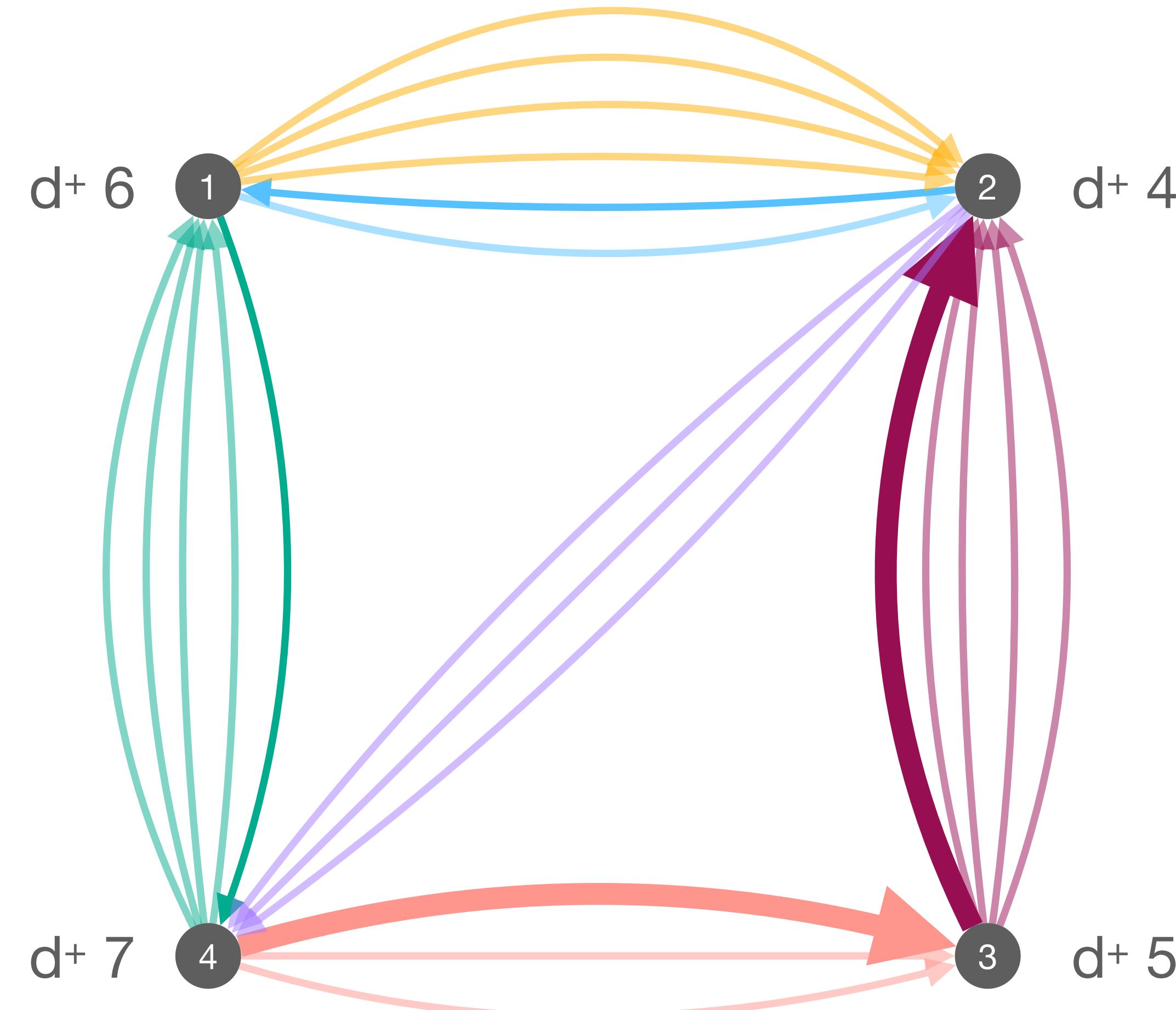


# Tethys

## Efficient DIP

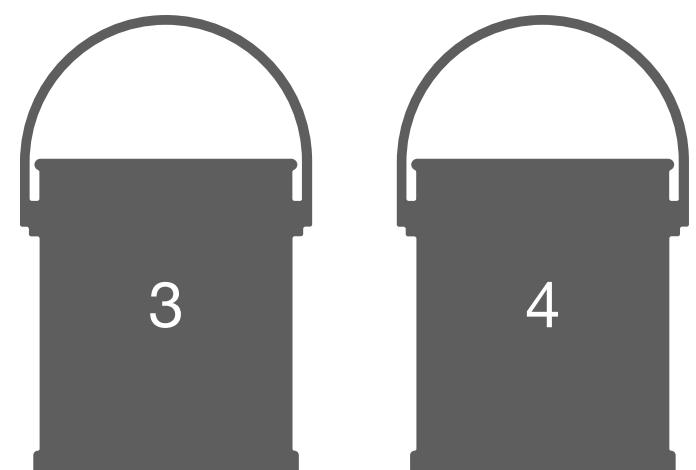
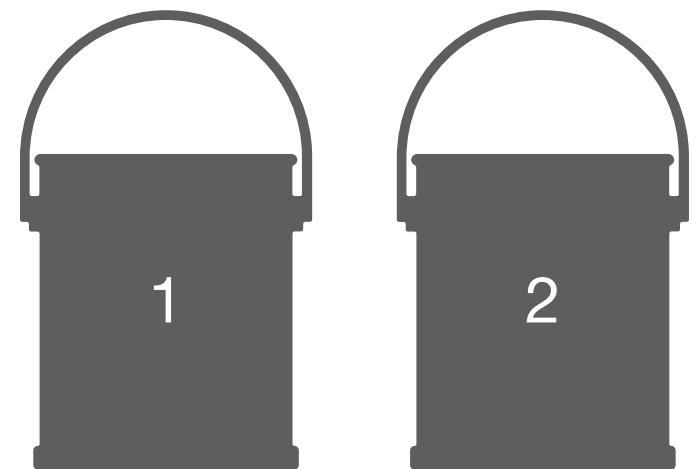
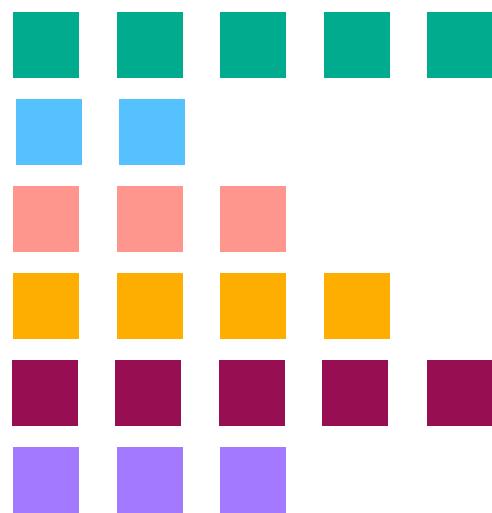


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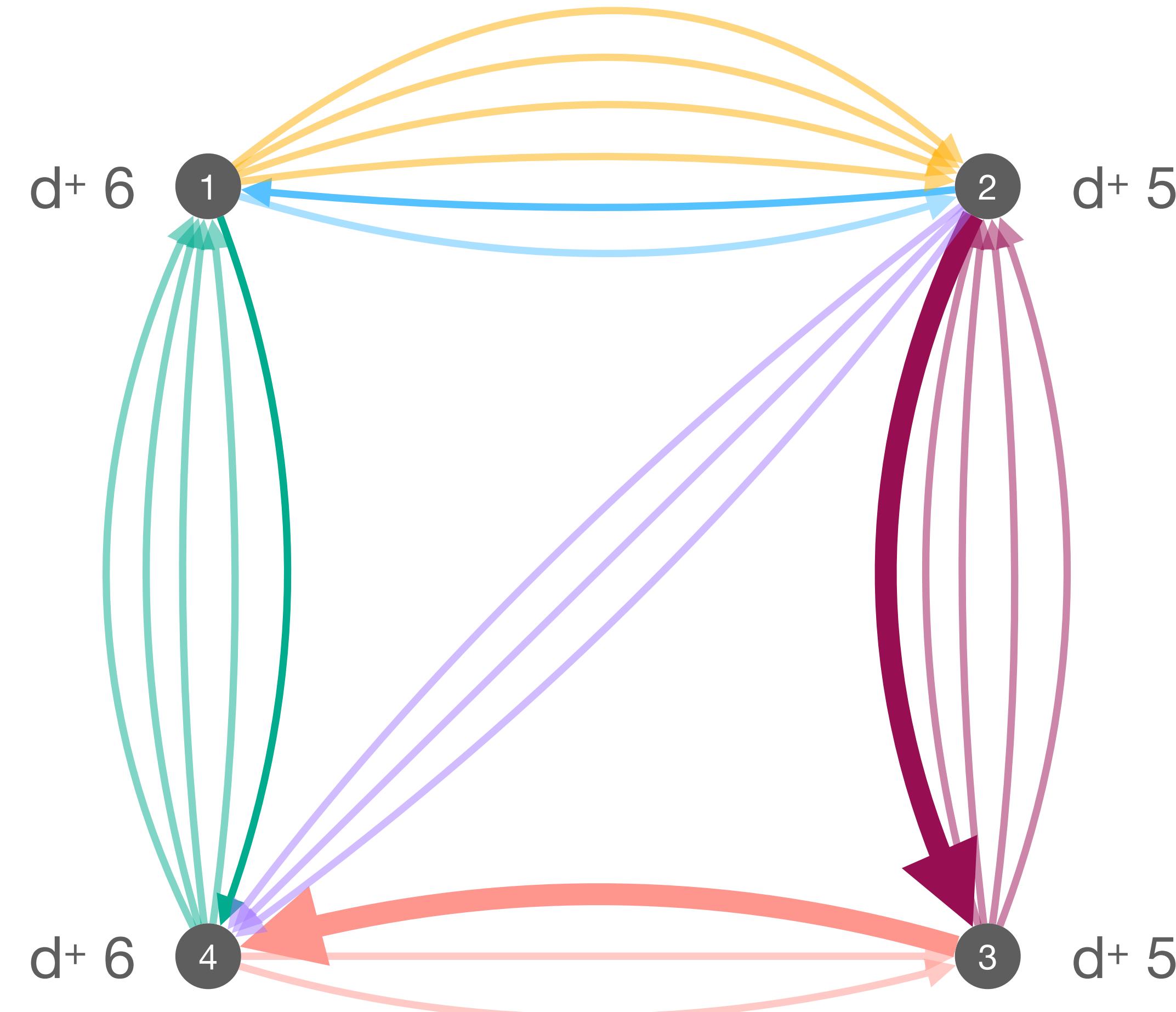


# Tethys

## Efficient DIP

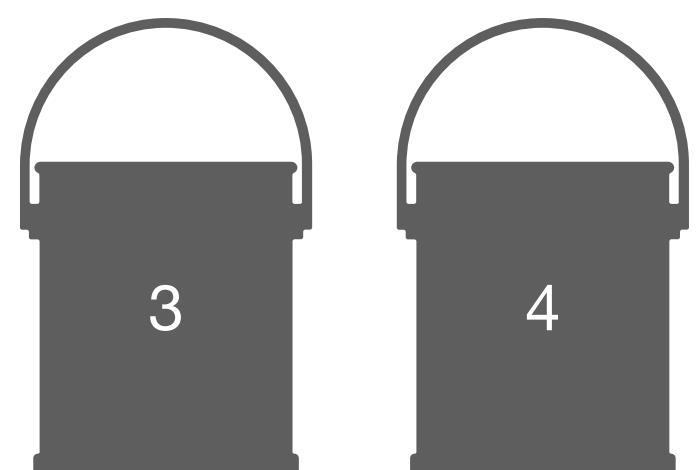
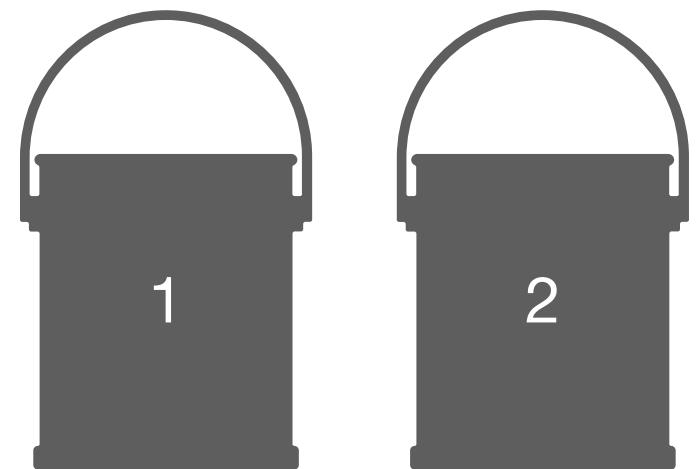
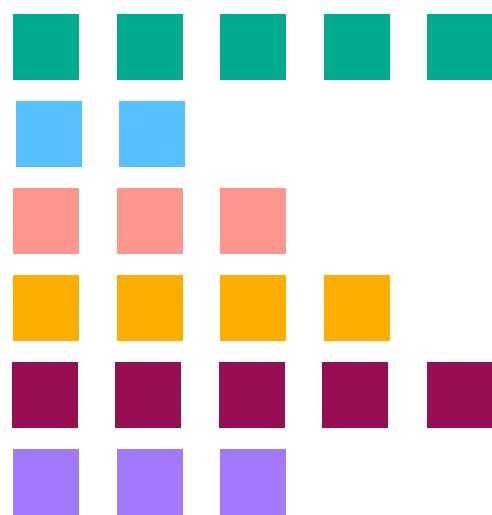


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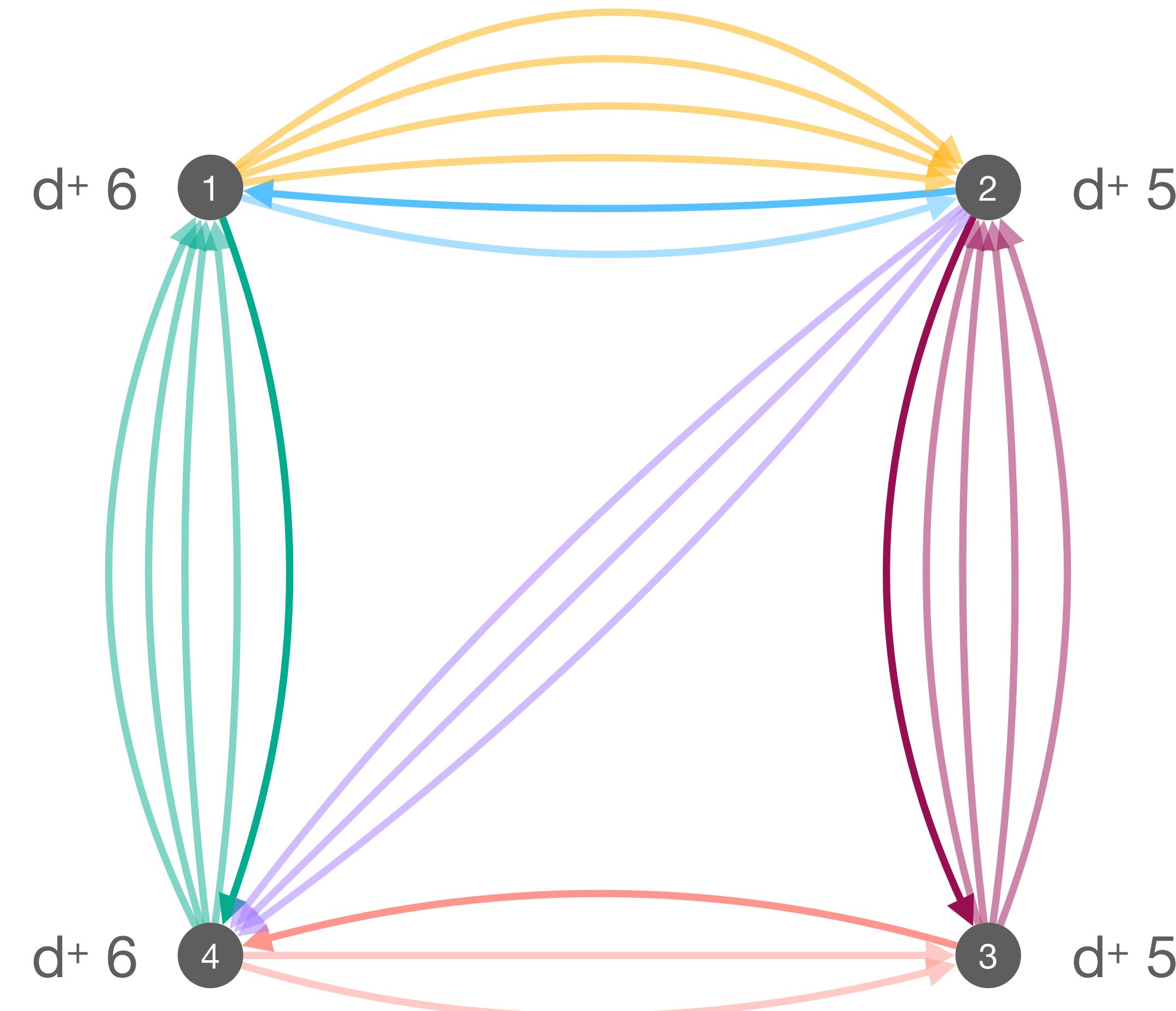


# Tethys

## Efficient DIP

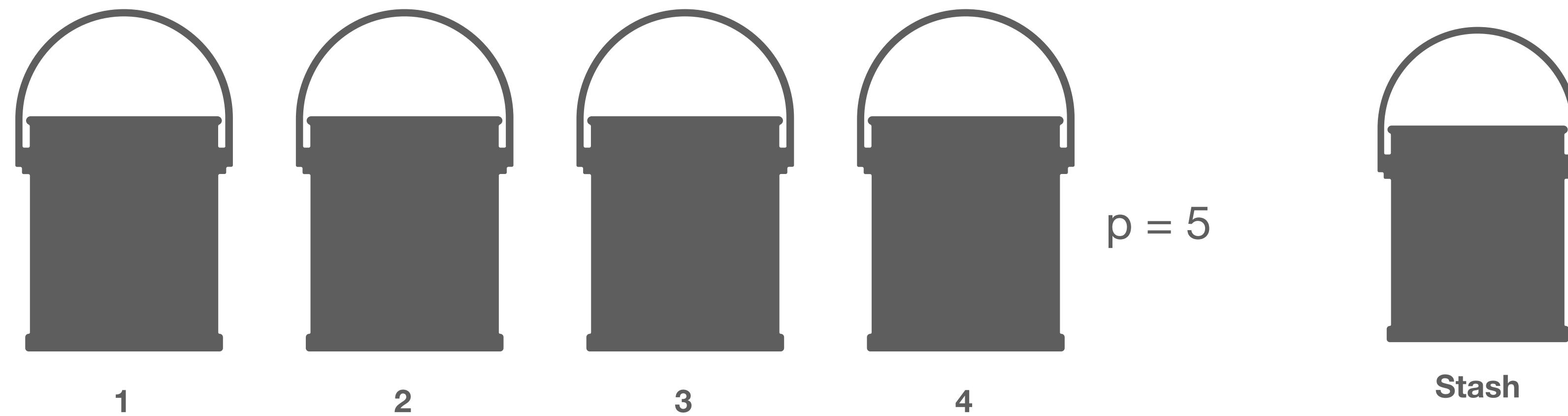
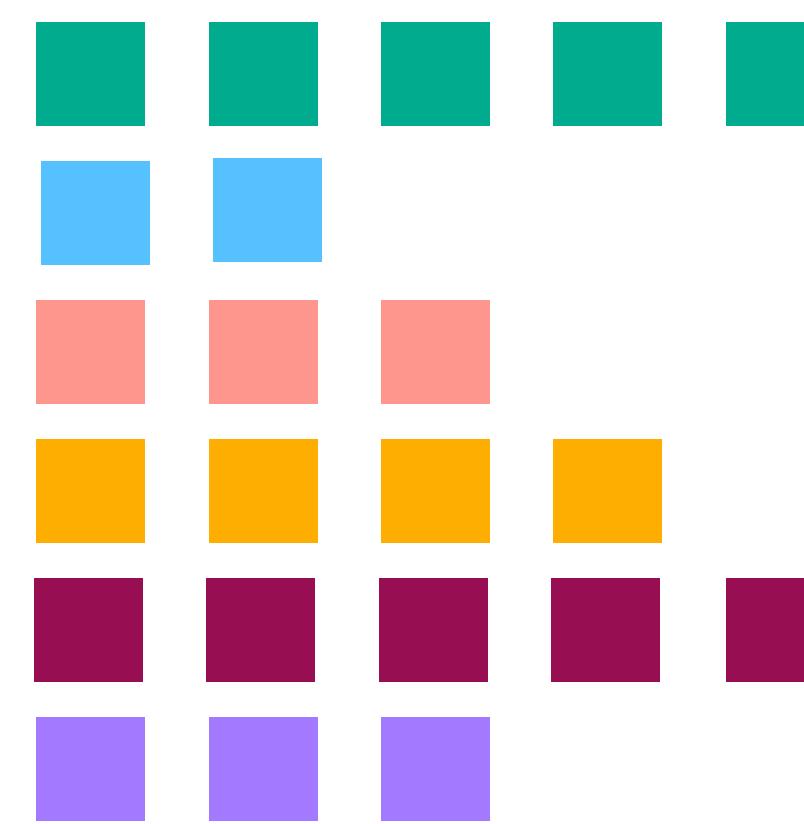
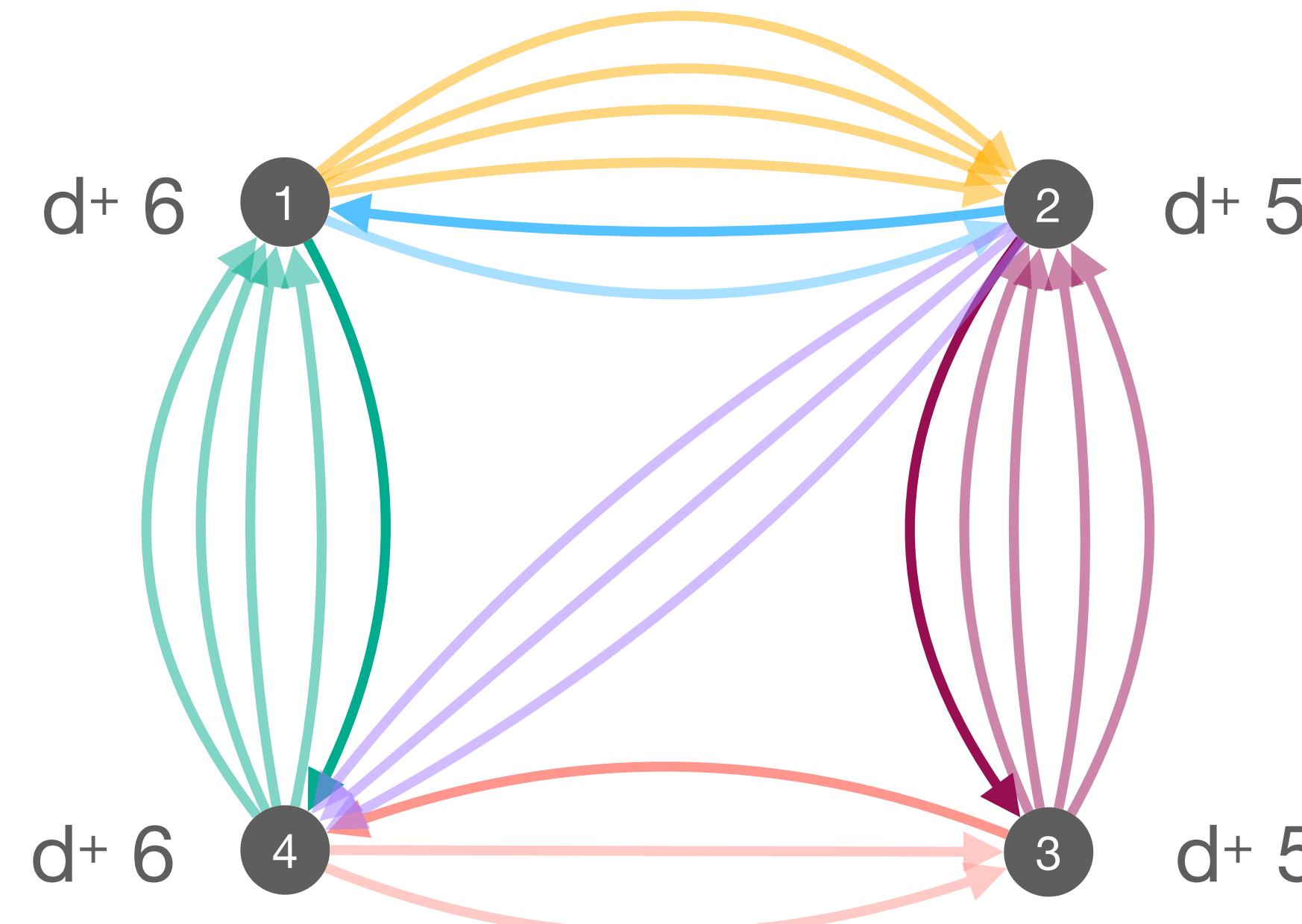


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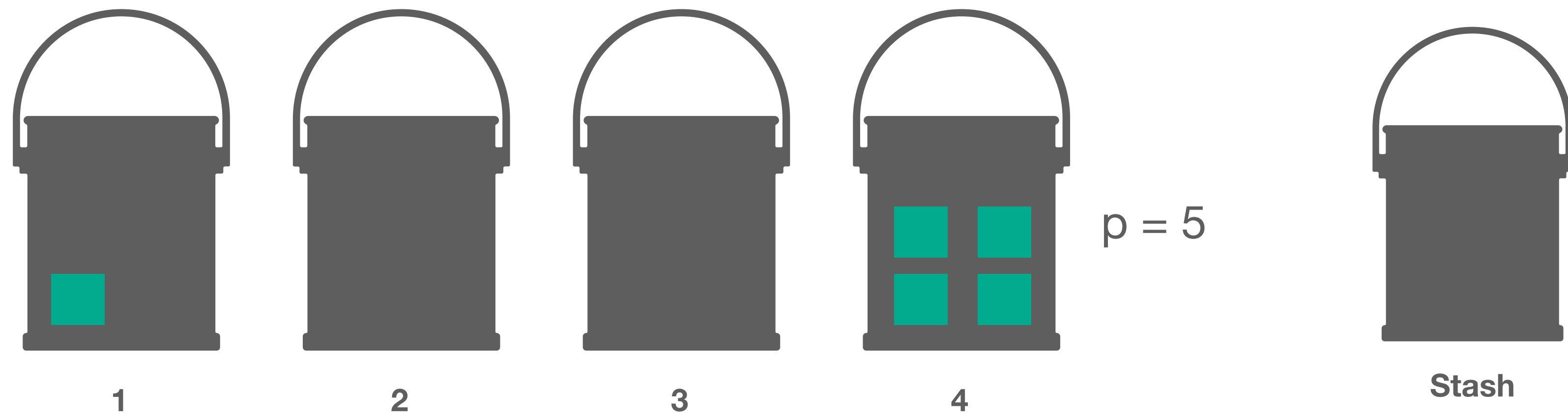
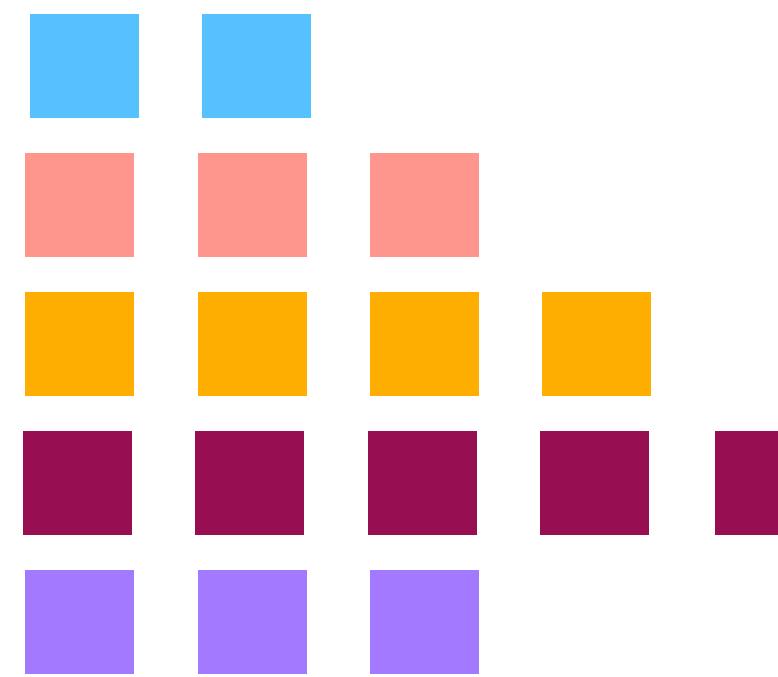
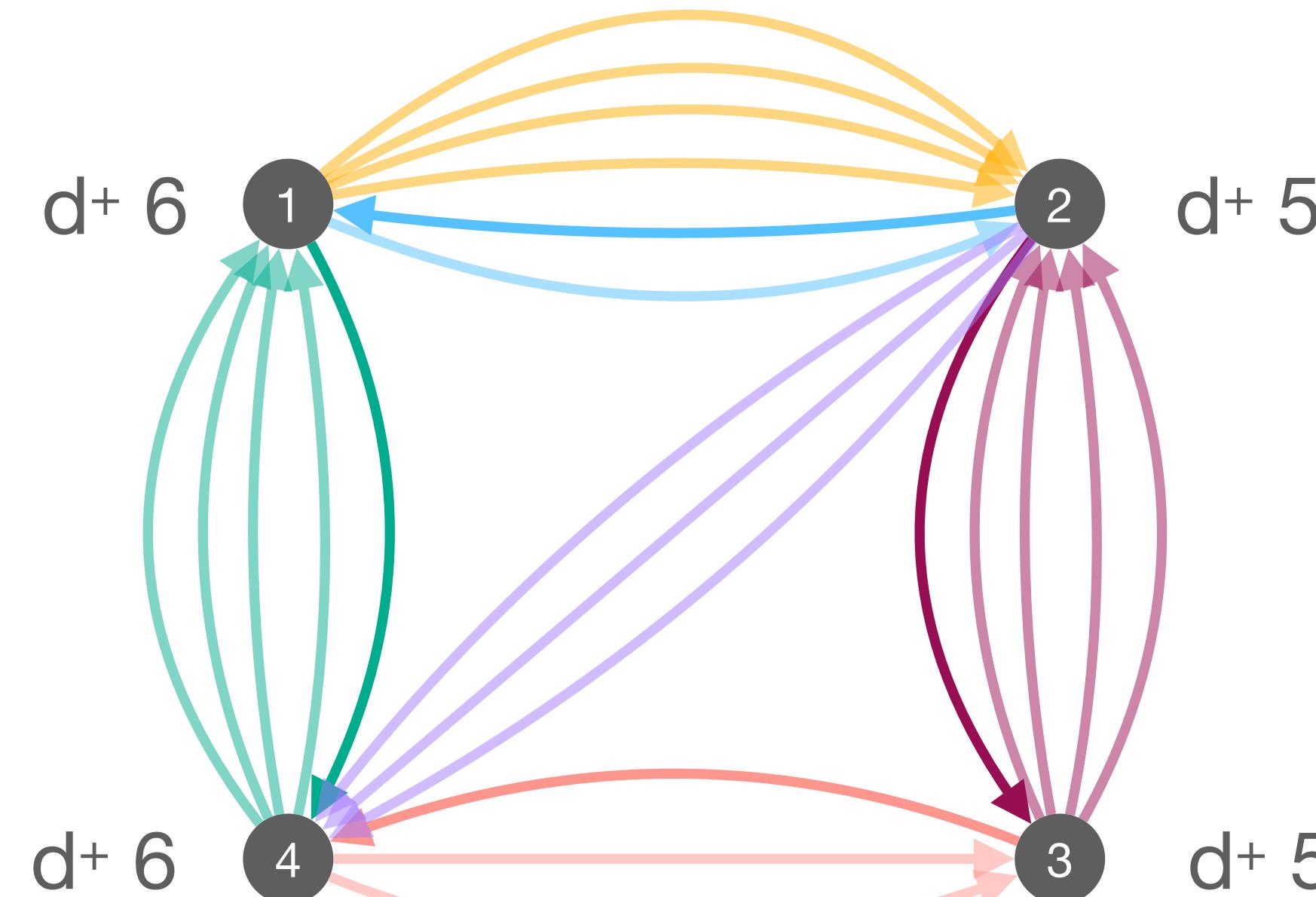
# Tethys

## Efficient DIP



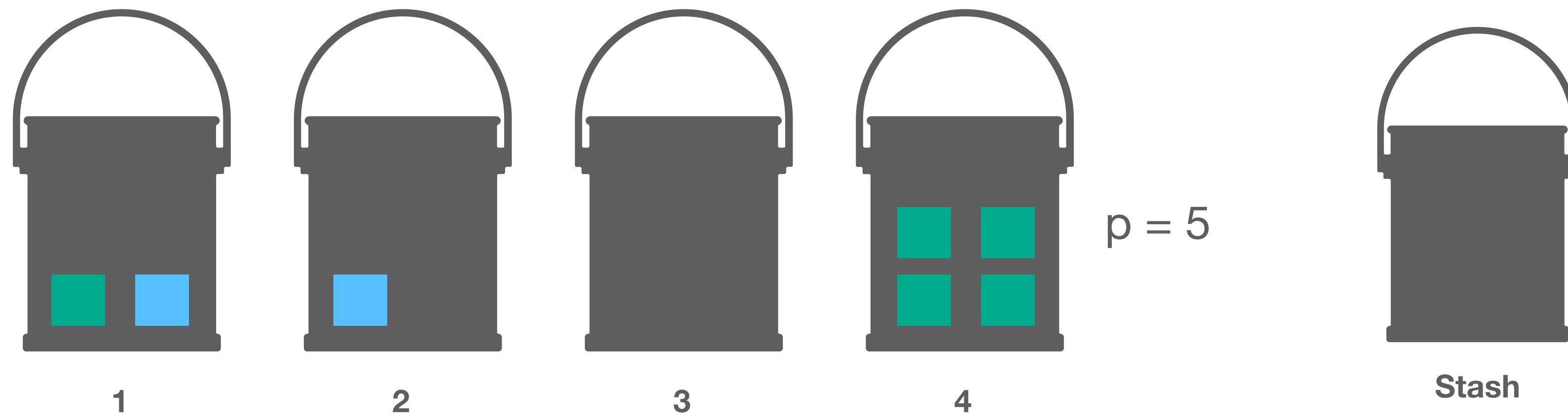
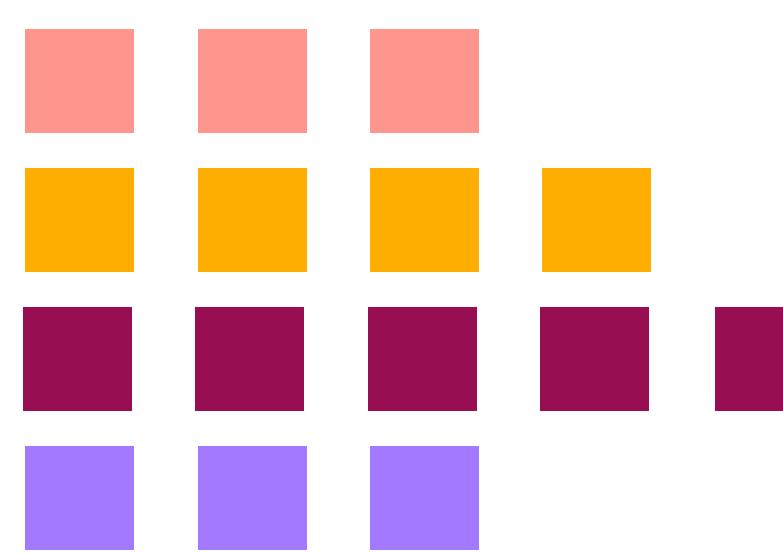
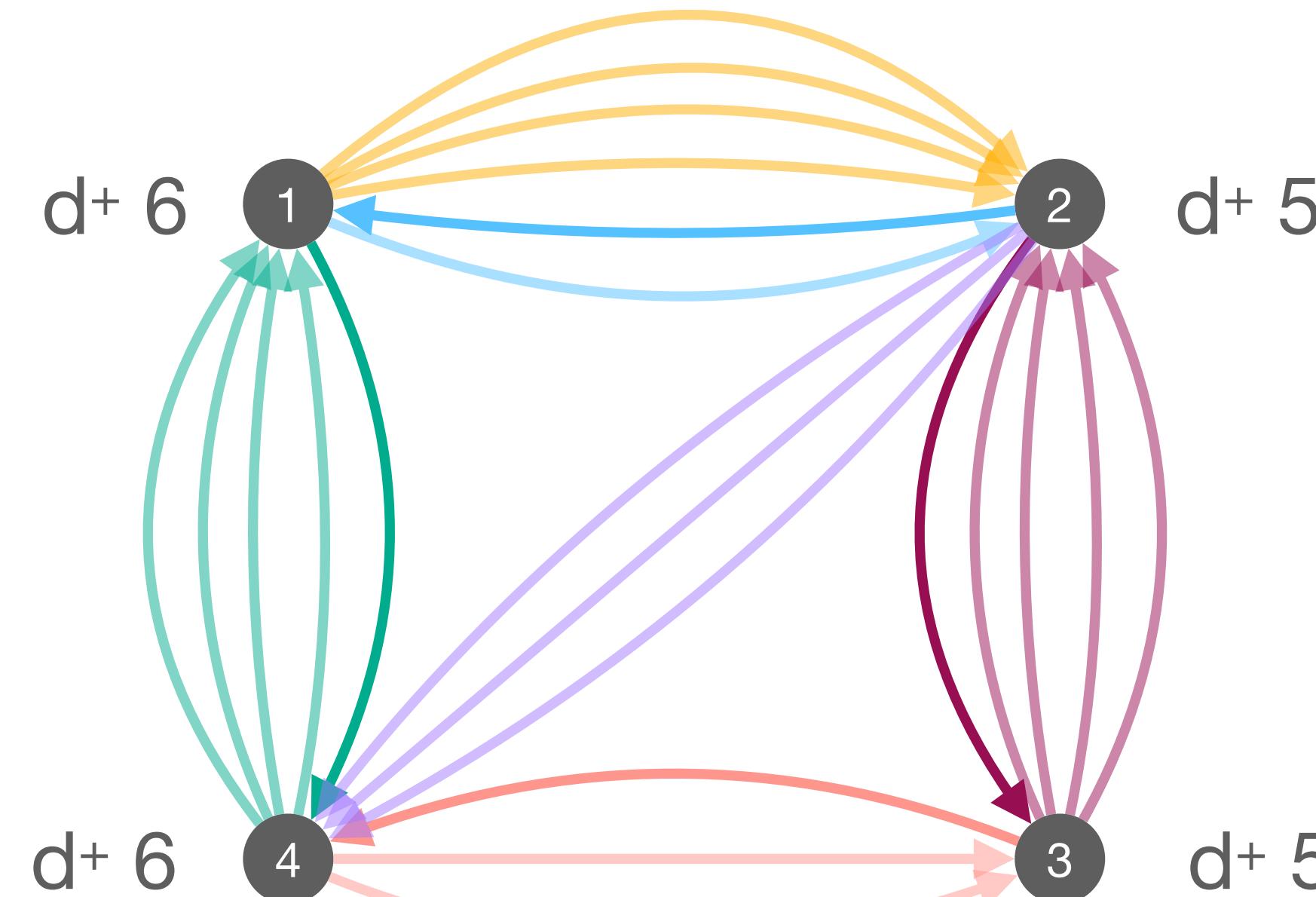
# Tethys

## Efficient DIP



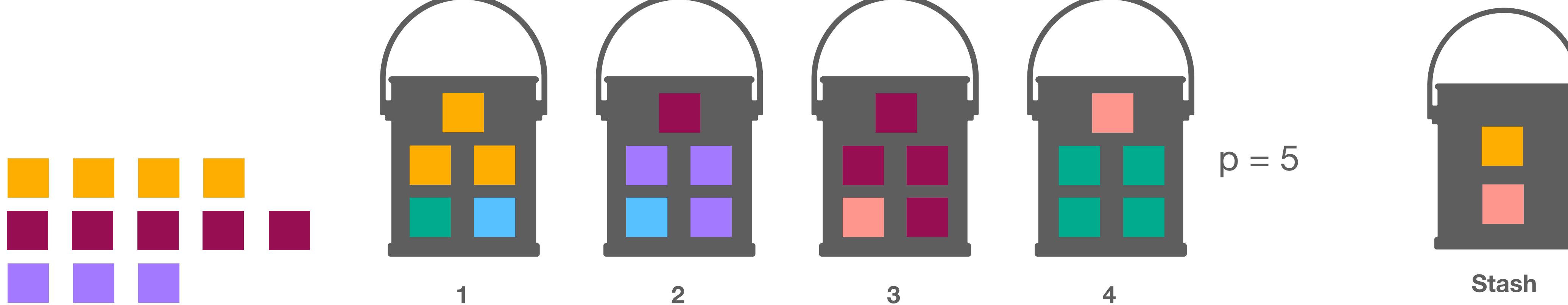
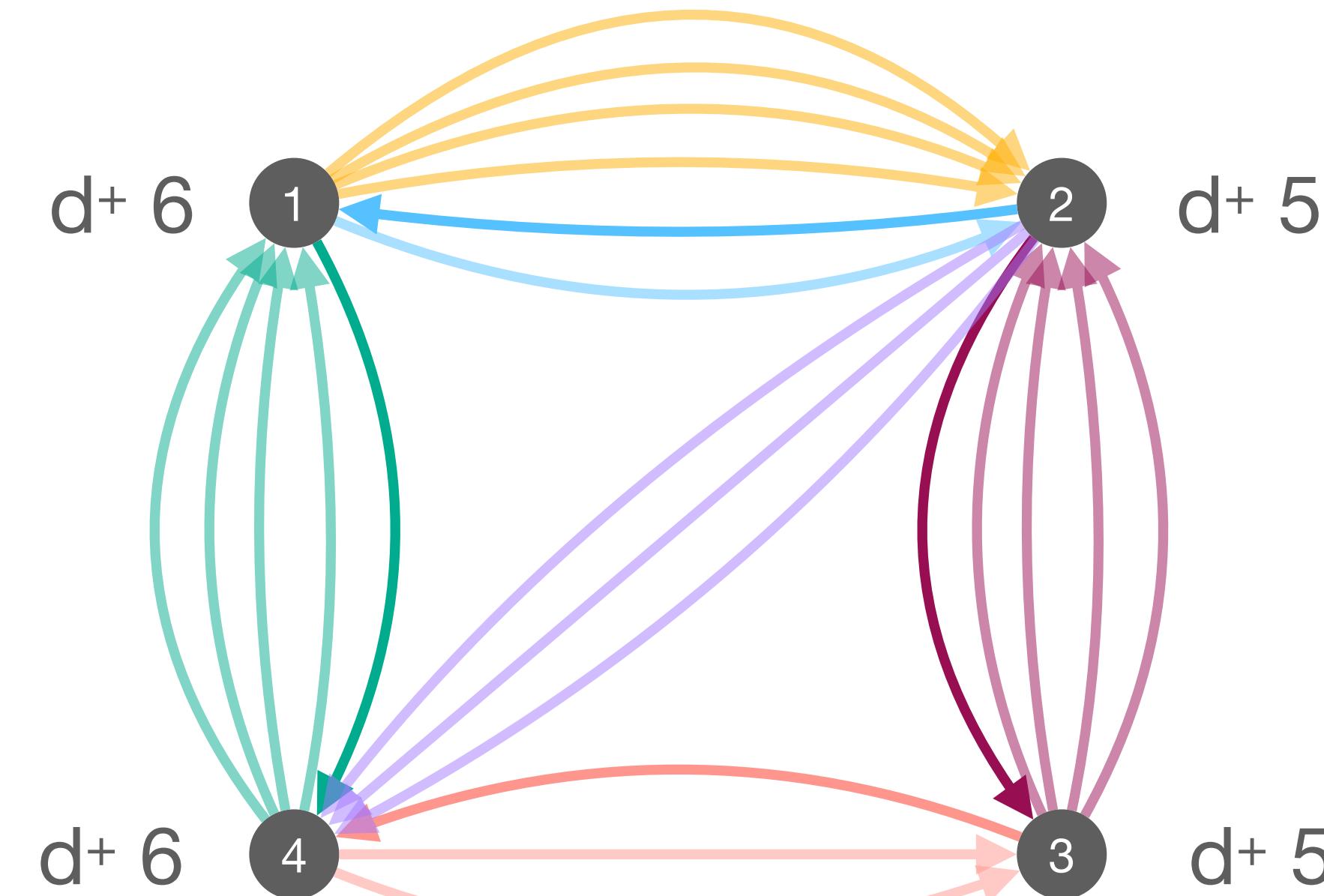
# Tethys

## Efficient DIP



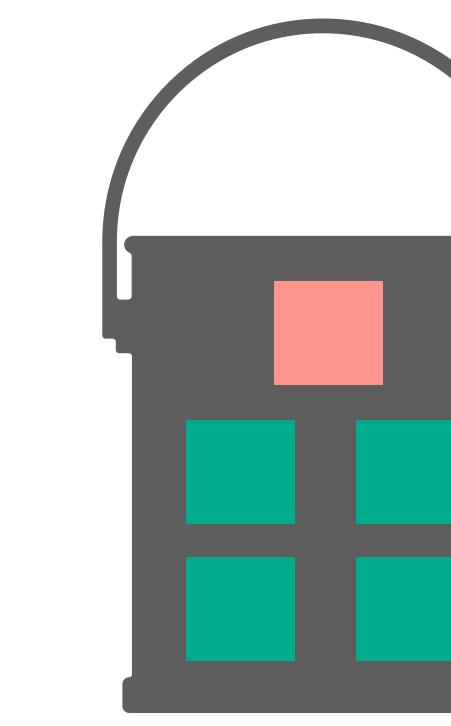
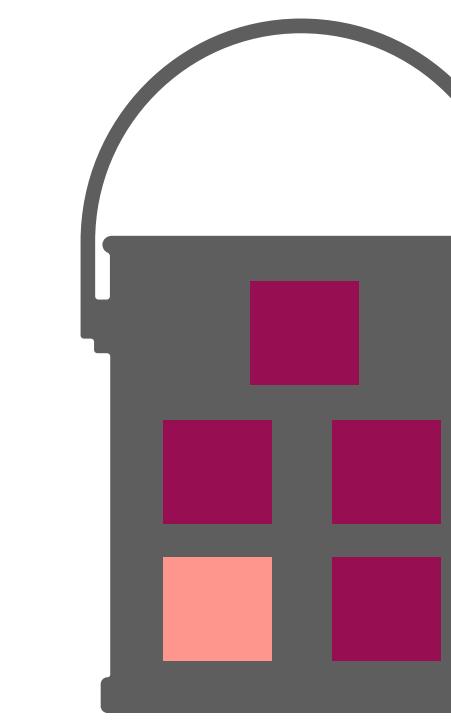
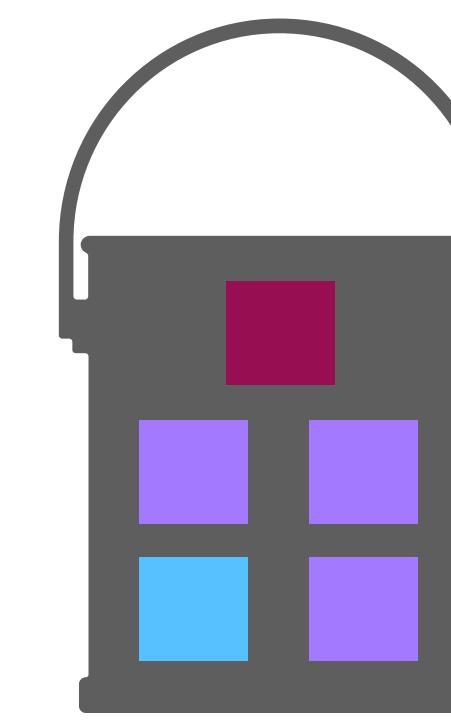
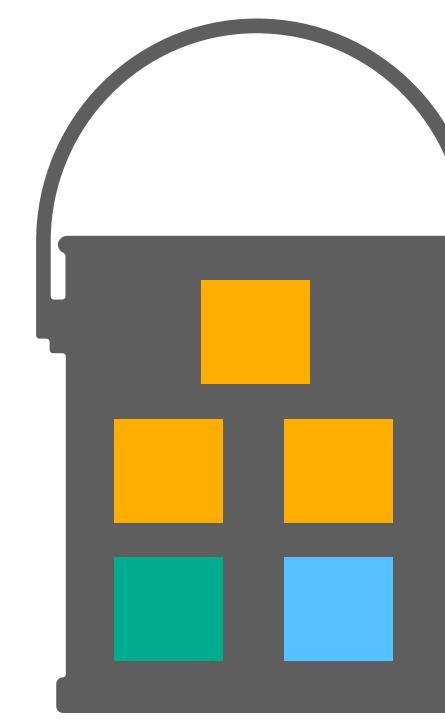
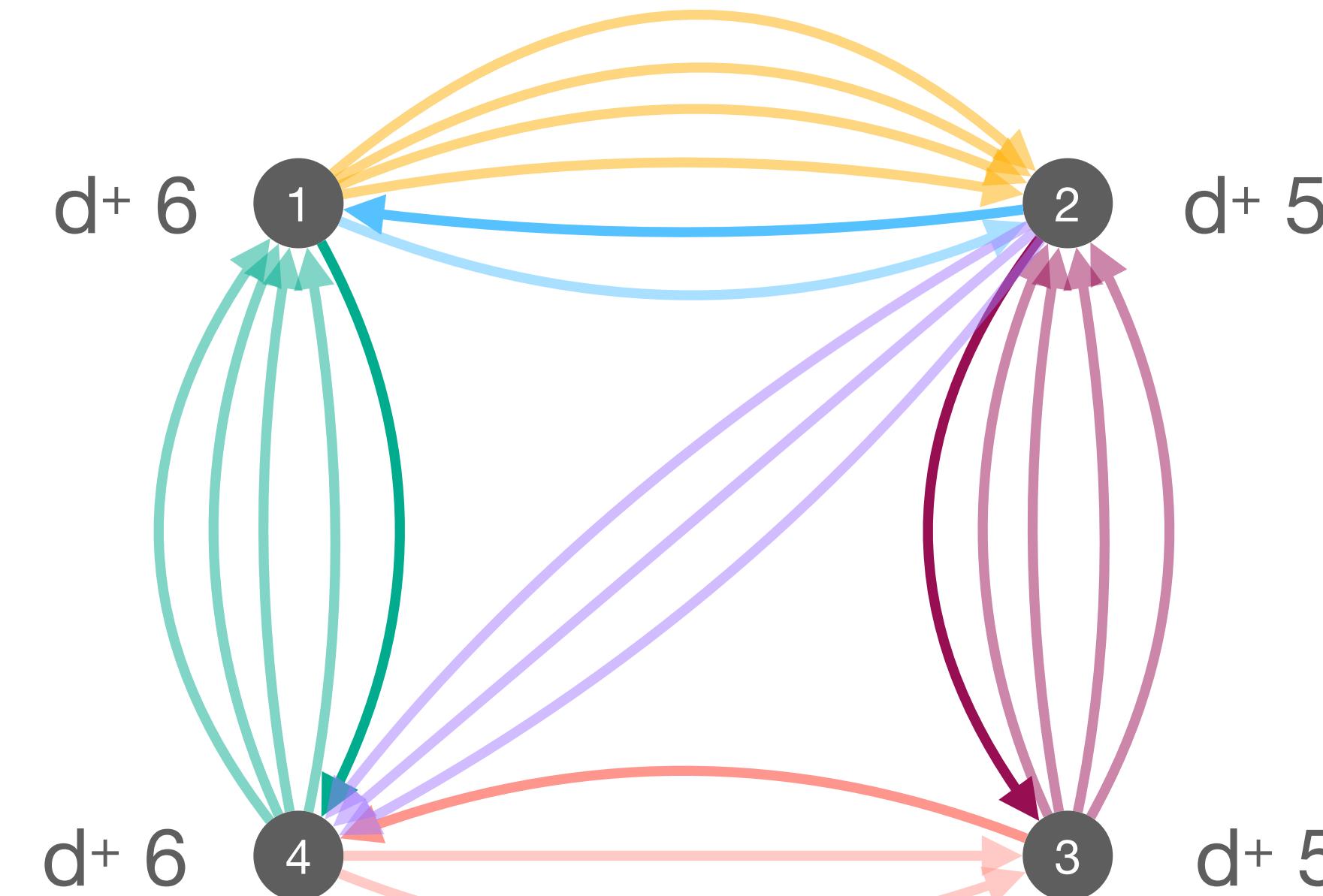
# Tethys

## Efficient DIP

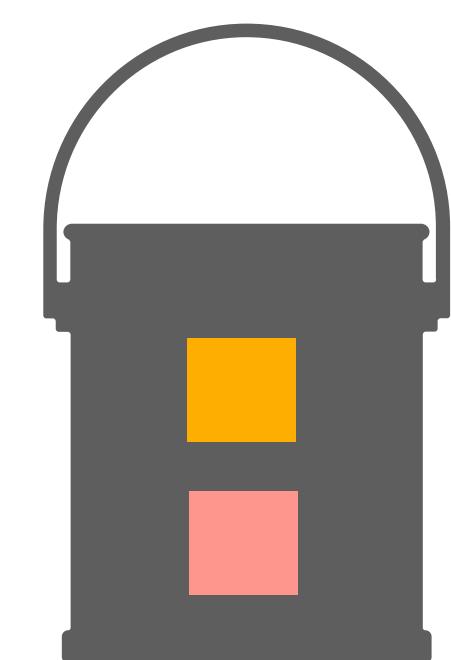


# Tethys

## Efficient DIP



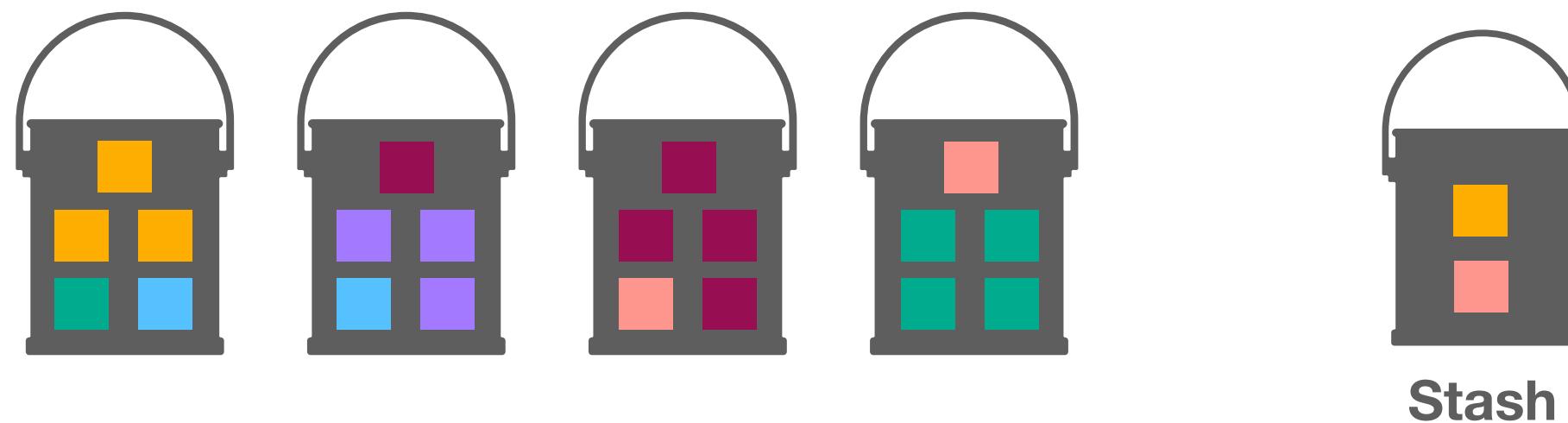
$p = 5$



Stash

# Tethys

## Efficient DIP



- Main Theorem:
  - There Exists a Valid Assignment w.o.p. such that:
    - $m = (2 + \varepsilon)n/p$
    - **Stash:**  $\omega(\log \lambda)/\log(n)$  pages
  - Direct Generalization of Cuckoo-Hashing in the Static Case

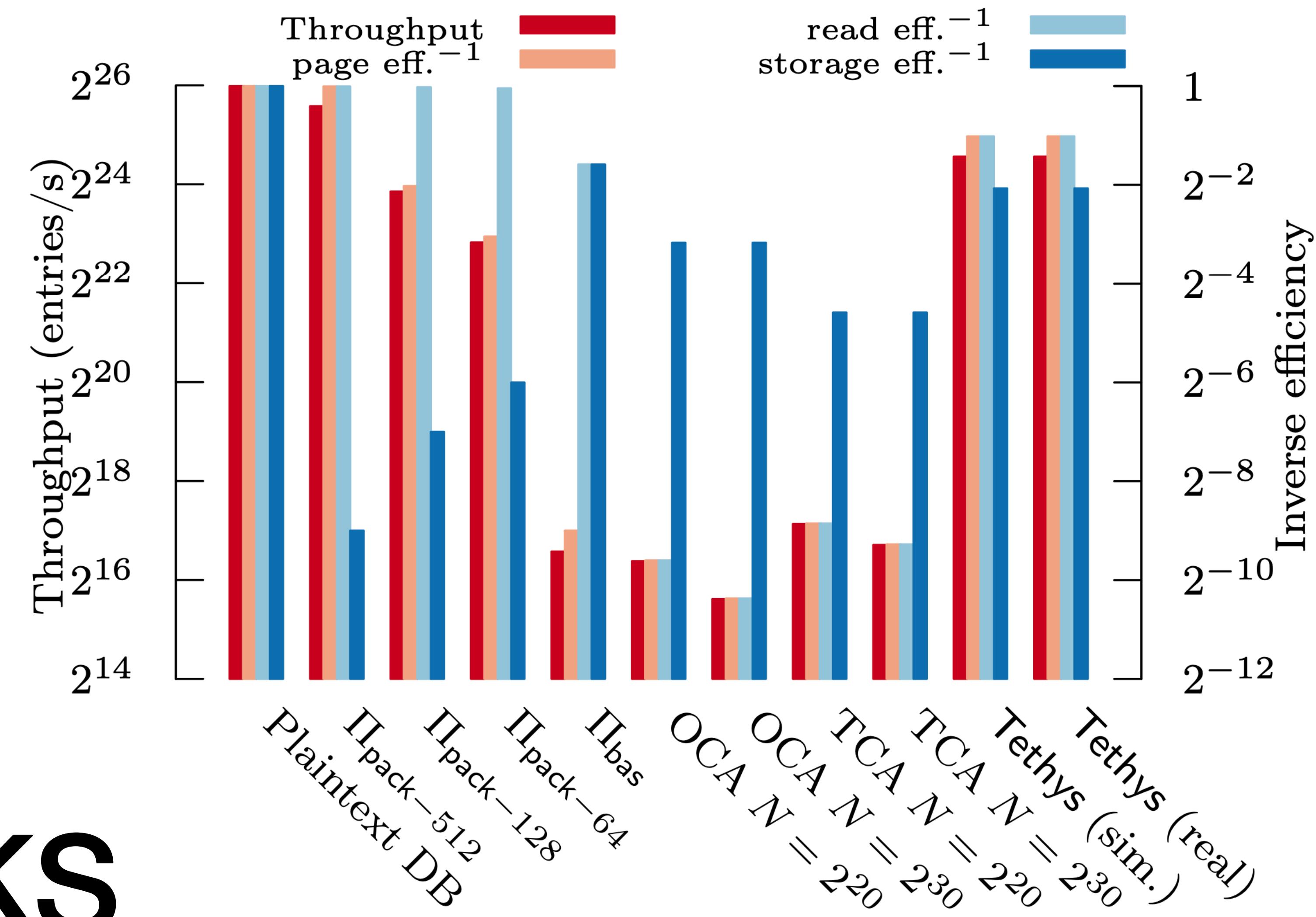
# Conclusion

## Efficiency

- Cuckoo-Hashing with Stash for Items of Variable Size
- **DIP**: Data-Independent Packing
- **Tethys**:  $O(1)$  Storage + Page Efficiency

# Conclusion

## Efficiency



THANKS