Introduction to SQL

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Basic Query Structure

A typical SQL query has the form:

select
$$A_1$$
, A_2 , ..., A_n
from r_1 , r_2 , ..., r_m
where P

			section	-	1		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name building budget	ID name dept_name salary	id id secid semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



- 1. Find the names of all instructors
- 2. Find the department names of all instructors, and remove duplicates
- 3. Find the department names of all instructors, not removing duplicates
- 4. Find all attributes of instrctor

show the entire instructor table

- 5. Find a relation that is the same as the *instructor* relation, except that the value of the attribute *salary* is divided by 12
- 6. Find all instructors in Comp. Sci. dept with salary > 70000
- 7. Find the names of all instructors who have taught some course and the course_id
- 8. Find the names of all instructors in the Comp. Sci. department who have taught some course and the course_id
- 9. Find the names of all instructors who have a higher salary than some instructor in 'Comp. Sci'.

			section	_			
		teaches	<u>course_id</u>	takes			I
department	instructor	<u>ID</u>	<u>sec_id</u>	<u>ID</u>	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	semester year building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



- Find the names of all instructors whose name includes the substring "in".
- String Operations
 - The operator like uses patterns that are described using two special characters:
 - percent (%). The % character matches any substring.
 - underscore (_). The _ character matches any character
- Find the names of all instructors whose name has 4 characters.
- Find the names of all instructors whose name has at leaset 4 characters.

			section		1		
		teaches	<u>course_id</u>	takes			Ī
department	instructor	<u>ID</u>	sec_id	<u>ID</u> . ,	student	course	
dept name building budget	<u>ID</u> name dept_name salary	id secid semester year	semester year building room_number time_slot_id	course id sec id semester year grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



- 1. List in alphabetic order the names of all instructors
- 2. List in descending alphabetic order the names of all instructors
- 3. List in order of the combination of the names and salary of all instructors
- 4. Find the names of all instructors with salary between \$90,000 and \$100,000
- 5. Find courses that ran in Fall 2017 or in Spring 2018
- 6. Find courses that ran in Fall 2017 and in Spring 2018
- 7. Find courses that ran in Fall 2017 but not in Spring 2018
- 8. Find courses that ran in Fall 2017 or in Spring 2018, retain all duplications
- 9. Find all instructors whose salary is null
- 10. Find all instructors whose salary is not null
- 11. Null under and, or, with true/false

			section		1		
		teaches	course_id	takes			1
department	instructor	<u>ID</u>	sec_id	<u>ID</u> . ,	student	course	
dept name building budget	ID name dept_name salary	id . <u>sec_id</u> . <u>semester</u> . <u>year</u>	semester year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



 These functions operate on the multiset of values of a column of a relation, and return a value

avg: average valuemin: minimum valuemax: maximum valuesum: sum of values

count: number of values

select $A_1, A_2, ..., A_n$ Aggregation function over values over multiple rows

from $r_1, r_2, ..., r_m$

where P

group by columns **having** condition



New clauses

			section		I		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u>	<u>ID</u>	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	semester year building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



 These functions operate on the multiset of values of a column of a relation, and return a value

avg: average valuemin: minimum valuemax: maximum valuesum: sum of values

count: number of values

- 1. Find the highest salary of any instructor
- 2. Find the average salary of instructors in the Computer Science department
- 3. Find the lowest salary of an instructor who have taught a course
- 4. Find the total number of instructors who teach a course in the Spring 2018 semester
- 5. Find the number of tuples in the course relation

department instructor ID sec id semester vear ID student course id sec id				Section	_	I		
department instructor ID sec id semester ID student course id			teaches	course id	takes			
ID course id	department	instructor		<u>sec_id</u>		student		
$\frac{acpt name}{name}$ $\frac{acpt}{name}$ \frac{acpt}	dept name			<u>year</u>	sec_id	<u>ID</u>		advisor
building dept_name semester room_number year dept_name dept_name semester dept_name salary year time_slot_id grade tot_cred rodules i_id		dept_name	<u>semester</u>	room_number	<u>year</u>	dept_name	dept_name	



Aggregate – Group By - Having

- 1. Find the average salary of instructors in each department
- 2. Find the names and average salaries of all departments whose average salary is greater than 42000
- 3. Find the names and average salaries of all departments over instructors whose salary is greater than 70000
- 4. Find the names and average salaries of all departments whose average salary is greater than 70000
- 5. Find the average salaries of instructors who have taught a course

			section	_	1		
		teaches	<u>course_id</u>	takes			I
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	<i>advisor</i> <u>s_id</u> i_id



Nested Subqueries

- A subquery is a select-from-where expression that is nested within another query.
- The nesting can be done in the following SQL query

select
$$A_1, A_2, ..., A_n$$
 from $r_1, r_2, ..., r_m$ **where** P

as follows:

- From clause: r_i can be replaced by any valid subquery
- Where clause: *P* can be replaced with an expression of the form:

B is an attribute and operation> to be defined later.

Select clause:

 A_i can be replaced be a subquery that generates a single value.

			section		1		
		teaches	course_id	takes			•
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name building budget	<u>ID</u> name dept_name salary	id . <u>sec_id</u> . <u>semester</u> . <u>year</u>	year building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Subquery in Where Clause

- 1. Name all instructors whose name is either "Mozart" or Einstein"
- 2. Name all instructors whose name is neither "Mozart" nor Einstein"
- 3. Find the total number of (distinct) students who have taken course sections taught by the instructor with *ID* 10101 or 12121 or 15151
- 4. Find names of instructors with salary greater than that of some (at least one) instructor in the Computer Science department.
- 5. Find all instructors earning the highest salary (there may be more than one with the same salary).
- 6. Find the names of all instructors whose salary is greater than the salary of all instructors in the Biology department.

			section	-	1		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	course id sec id semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id

Definition of "some" Clause

■ F <comp> some $r \Leftrightarrow \exists t \in r \text{ such that (F <comp> } t)$ Where <comp> can be: <, ≤, >, =, ≠

Definition of "all" Clause

• F <comp> all $r \Leftrightarrow \forall t \in r \text{ (F <comp> } t)$

$$(5 < \mathbf{all} \quad \begin{array}{c} 0 \\ 5 \\ \hline 6 \\ \end{array}) = \text{false}$$

$$(5 < \mathbf{all} \quad \begin{array}{c} 6 \\ 10 \\ \end{array}) = \text{true}$$

$$(5 = \mathbf{all} \quad \begin{array}{c} 4 \\ 5 \\ \end{array}) = \text{false}$$

$$(5 \neq \mathbf{all} \quad \begin{array}{c} 4 \\ \hline 6 \\ \end{array}) = \text{true (since } 5 \neq 4 \text{ and } 5 \neq 6)$$

$$(\neq \mathbf{all}) \equiv \mathbf{not in}$$
However, $(= \mathbf{all}) \neq \mathbf{in}$



Subquery in Where Clause

- The exists construct returns the value true if the argument subquery is nonempty.
 - Find all courses taught in both the Fall 2017 semester and in the Spring 2018 semester
- The unique construct tests whether a subquery has any duplicate tuples in its result.
- The unique construct evaluates to "true" if a given subquery contains no duplicates.
 - Find all courses that were offered at most once in 2017

			section	-	1		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	course id sec id semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Query Quest

- 1. Find names of instructors with salary greater than that of some (at least one) instructor in the Computer Science department.
 - Use self-join, some, exists, aggregation
- 2. Find all instructors earning the highest salary (there may be more than one with the same salary).
 - Use self-join, all, exists, aggregation

			section	_	1		
		teaches	<u>course_id</u>	takes			I
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	<i>advisor</i> <u>s_id</u> i_id



Subquery in From and Select Clauses

- 1. Find the average instructors' salaries of those departments where the average salary is greater than \$42,000"
- 2. Find all departments with the maximum budget (with clause)
- 3. Find all departments where the total salary is greater than the average of the total salary at all departments
- 4. List all departments along with the number of instructors in each department

			section	_	I		
		teaches	<u>course_id</u>	takes			ı
department in	istructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u> course id	student	course	
building do	<u>D</u> name lept_name alary	course_id sec_id semester year	year building room_number time_slot_id	sec id semester year grade	<u>ID</u> name dept_name tot_cred	course_id title dept_name credits	<i>advisor</i> <u>s_id</u> i_id



With Clause

- The with clause provides a way of defining a temporary relation whose definition is available only to the query in which the with clause occurs.
- Find all departments with the maximum budget

with max_budget (value) as
 (select max(budget)
 from department)
select department.name
from department, max_budget
where department.budget = max_budget.value;

			section	_	I		
		teaches	<u>course_id</u>	takes			ı
department in	istructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u> course id	student	course	
building do	<u>D</u> name lept_name alary	course_id sec_id semester year	year building room_number time_slot_id	sec id semester year grade	<u>ID</u> name dept_name tot_cred	course_id title dept_name credits	<i>advisor</i> <u>s_id</u> i_id

Modification of the Database

- Deletion of tuples from a given relation.
- Insertion of new tuples into a given relation
- Updating of values in some tuples in a given relation

delete from a relation where condition

Deletion



Delete all instructors

delete from instructor

- Delete all instructors from the Finance department delete from instructor where dept_name= 'Finance';
- Delete all tuples in the instructor relation for those instructors associated with a department located in the Watson building.

			section		1		
		teaches	course_id	takes			,
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept_name building budget	ID name dept_name salary	id . <u>sec_id</u> . <u>semester</u> . <u>year</u>	vear building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Deletion (Cont.)

 Delete all instructors whose salary is less than the average salary of instructors

- Problem: as we delete tuples from instructor, the average salary changes
- Solution used in SQL:
 - 1. First, compute avg (salary) and find all tuples to delete
 - Next, delete all tuples found above (without recomputing avg or retesting the tuples)

				section	-	1		
			teaches	course_id	takes			ı
depai	rtment	instructor	<u>ID</u>	sec_id	<u>ID</u>	student	course	
dept build budg	_	<u>ID</u> name dept_name salary	id secid semester year	semester year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	<i>advisor</i> <u>s_id</u> i_id

Insertion



Add a new tuple to course

insert into course
 values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);

or equivalently

insert into course (course_id, title, dept_name, credits)
 values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);

Add a new tuple to student with tot_creds set to null

insert into student
 values ('3003', 'Green', 'Finance', null);

			section		1		
		teaches	course_id	takes			•
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name building budget	<u>ID</u> name dept_name salary	id . <u>sec_id</u> . <u>semester</u> . <u>year</u>	year building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id

Insertion (Cont.)



 Make each student in the Music department who has earned more than 144 credit hours an instructor in the Music department with a salary of \$18,000.

```
insert into instructor
    select ID, name, dept_name, 18000
from student
    where dept_name = 'Music' and total_cred > 144;
```

 The select from where statement is evaluated fully before any of its results are inserted into the relation.

Otherwise queries like

insert into table1 select * from table1

would cause problem

			section	•	I		
		teaches	course_id	takes			1
department	instructor	<u>ID</u>	sec_id	<u>ID</u>	student	course	
dept name building budget	ID name dept_name salary	id <u>secid</u> <u>semester</u> <u>year</u>	semester year building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id

Updates



- 1. Give a 5% salary raise to all instructors
- 2. Give a 5% salary raise to those instructors who earn less than 70000
- 3. Give a 5% salary raise to instructors whose salary is less than average

			section		1		
		teaches	<u>course_id</u>	takes			ı
department	instructor	<u>ID</u>	sec_id	<u>ID</u>	student	course	
dept name building budget	I <u>D</u> name dept_name salary	id . <u>sec_id</u> . <u>semester</u> . <u>year</u>	semester year building room_number time_slot_id	course id sec id semester year grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id

Updates (Cont.)



- Increase salaries of instructors whose salary is over \$100,000 by 3%, and all others by a 5%
 - update instructor set salary = salary * 1.03 where salary > 100000;

update instructor
set salary = salary * 1.05

where *salary* <= 100000;

- The order is important
- Can be done better using the case statement

end

update instructor
set salary = case

when salary <= 100000 then salary * 1.05 else salary * 1.03

			section	_	I		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u>	<u>ID</u>	student	course	
dept name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	semester year building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Updates with Scalar Subqueries

Recompute and update tot_creds value for all students

- Sets tot_creds to null for students who have not taken any course
 - Instead of sum(credits), use:

end

case
 when sum(credits) is not null then sum(credits)
 else 0

department	instructor	teaches <u>ID</u>	<u>course_id</u> <u>sec_id</u> semester	takes <u>ID</u> <u>course id</u>	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	year building room_number time_slot_id	sec id semester year grade	ID name dept_name tot cred	course_id title dept_name credits	advisor s_id i id

section



Join conditions

- List the names of students instructors along with the titles of courses that they have taken
 - Natural Join with Using Clause
 - Join condition with on condition

			section		1		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name building budget	I <u>D</u> name dept_name salary	course_id sec_id semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Join types

- Three forms of outer join:

course_id	title	dept_name	credits	prereq_id
	State and the state of the stat	Biology	19	BIO-101
CS-190 CS-315	Game Design Robotics	Comp. Sci. Comp. Sci.		CS-101 null

course_id	title	dept_name	credits	prereq_id
BIO-301	Genetics	Biology	- 88	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	null	null	null	CS-101

course_id	title	dept_name	credits	prereq_id
BIO-301 CS-190 CS-315 CS-347	Game Design	Biology Comp. Sci. Comp. Sci. null	15	BIO-101 CS-101 <i>null</i> CS-101

course_id	title	dept_name	credits
BIO-301	Genetics	Biology	4
CS-190	Game Design		4
CS-315	Robotics	Comp. Sci.	3

course_id	prereg_id
BIO-301	BIO-101
CS-190	CS-101
CS-347	CS-101

			Section		Ī		
		teaches	course_id	takes			
department	instructor	<u>ID</u>	<u>sec_id</u>	<u>ID</u>	student	course	
dept_name building budget	I <u>D</u> name dept_name salary	course_id sec_id semester year	semester year building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id

section



Join

Join types
inner join
left outer join
right outer join
full outer join

Join co	onditions
natura	ıl
on < p	redicate>
using	$(A_1, A_2,, A_n)$
using	(A_1, A_2, \ldots, A_n)

course_id	title	dept_name	credits
BIO-301	Genetics	Biology	4
CS-190	Game Design	Comp. Sci.	4
CS-315	Robotics	Comp. Sci.	3

course_id	prereg_id
BIO-301	BIO-101
CS-190	CS-101
CS-347	CS-101

			section	-	1		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u> course id	student	course	
dept name	<u>ID</u> name	<u>course_id</u> sec_id	<u>year</u> building	<u>sec_id</u>	<u>ID</u>	<u>course_id</u> title	advisor
building	dept_name	<u>semester</u>	room_number	<u>semeste</u> r <u>year</u>	name dept_name	dept_name	<u>s_id</u>
budget	salary	<u>year</u>	time_slot_id	grade	tot_cred	credits	i_id



Joined Relations – Examples

course natural right outer join prereq

course_id	title	dept_name	credits	prereq_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	null	null	null	CS-101

course full outer join prereq using (course_id)

course_id	title	dept_name	credits	prereq_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null
CS-347	null	null	null	CS-101

			section		1		
		teaches	<u>course_id</u>	takes			•
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
<u>dept_name</u>	<u>ID</u> name	 course_id sec_id	<u>year</u> building	<u>course i</u> d <u>sec id</u> semester	<u>ID</u> name	<u>course_id</u> title	advisor
building budget	dept_name salary	<u>semester</u> <u>year</u>	room_number time_slot_id	<u>year</u> grade	dept_name tot_cred	dept_name credits	. <u>s_id</u> i_id



Joined Relations – Examples

course inner join prereq on course.course_id = prereq.course_id

course_id	title	dept_name	credits	prereq_id	course_id
	Genetics Game Design	Biology Comp. Sci.	333	BIO-101 CS-101	BIO-301 CS-190

- What is the difference between the above, and a natural join?
- course left outer join prereq on course.course_id = prereq.course_id

course_id	title	dept_name	credits	prereq_id	course_id
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190
CS-315	Robotics	Comp. Sci.	3	null	null

			section	-	1		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name building budget	ID name dept_name salary	id . <u>sec_id</u> . <u>semester</u> . <u>year</u>	vear building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



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course full outer join prereq using (course_id)

course_id	title	dept_name	credits	prereq_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null
CS-347	null	null	null	CS-101

			section		1		
		teaches	<u>course_id</u>	takes			•
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
<u>dept_name</u>	<u>ID</u> name	 course_id sec_id	<u>year</u> building	<u>course i</u> d <u>sec id</u> semester	<u>ID</u> name	<u>course_id</u> title	advisor
building budget	dept_name salary	<u>semester</u> <u>year</u>	room_number time_slot_id	<u>year</u> grade	dept_name tot_cred	dept_name credits	. <u>s_id</u> i_id



View

- A view provides a mechanism to hide certain data from the view of certain users.
 - Consider a person who needs to know an instructors name and department, but not the salary. This person should see a relation described, in SQL, by

select *ID*, *name*, *dept_name* **from** *instructor*

- create view v as < query expression >
- 1. A view of instructors without their salary
- 2. Find all instructors in the Biology department
- 3. Create a view of department salary totals

			section		1		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	course id sec id semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Define Views using other Views

- create view physics_fall_2017 as select course.course_id, sec_id, building, room_number from course, section where course.course_id = section.course_id and course.dept_name = 'Physics' and section.semester = 'Fall' and section.year = '2017';
- create view physics_fall_2017_watson as select course_id, room_number from physics_fall_2017 where building= 'Watson';

			section	_	1		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	id secid semester year	<u>semester</u> <u>year</u> building room_number time_slot_id	<u>course</u> id <u>sec id</u> <u>semeste</u> r <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



View Expansion

Expand the view :

```
create view physics_fall_2017_watson as
  select course_id, room_number
from physics_fall_2017
where building= 'Watson'
```

To: create view physics_fall_2017_watson as select course_id, room_number from (select course.course_id, building, room_number from course, section where course.course_id = section.course_id and course.dept_name = 'Physics' and section.semester = 'Fall' and section.year = '2017') where building= 'Watson';

			Section	_	I		
		teaches	course id	takes			
department	instructor	<u>ID</u>	<u>sec_id</u>	<u>ID</u>	student	course	
dept name building budget	ID name dept_name salary	course_id sec_id semester year	semester year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id

section



Views

- Materialized Views
- Update Views
 - Add a new tuple to faculty view which we defined earlier insert into faculty values ('30765', 'Green', 'Music');

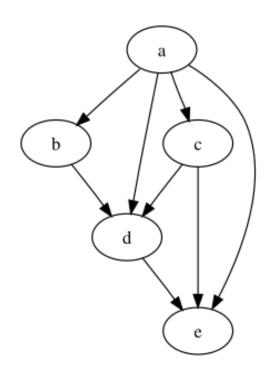
			section		1		
		teaches	<u>course_id</u>	takes			•
department	instructor	<u>ID</u>	sec_id	<u>ID</u>	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	id . <u>sec_id</u> . <u>semester</u> . <u>year</u>	semester year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id

Some Updates Cannot be Translated Uniquely

- create view instructor_info as
 select ID, name, building
 from instructor, department
 where instructor.dept_name = department.dept_name;
- insert into instructor_info values ('69987', 'White', 'Taylor');
- Issues
 - Which department, if multiple departments in Taylor?
 - What if no department is in Taylor?
- Most SQL implementations allow updates only on simple views
 - The from clause has only one database relation.
 - The select clause contains only attribute names of the relation, and does not have any expressions, aggregates, or distinct specification.
 - Any attribute not listed in the select clause can be set to null
 - The query does not have a group by or having clause.



Dependency Graph of Views should be a DAG



			section	_	1		
		teaches	course_id	takes			ı
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u> course id	student	course	
dept name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	year building room_number time_slot_id	sec id semester year grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Query Quest

- 1. Display a list of all instructors, showing their ID, name, and the number of sections that they have taught. Make sure to show the number of sections as 0 for instructors who have not taught any section. Your query should use an outer join, and should not use scalar subqueries.
- 2. Outer join expressions can be computed in SQL without using the SQL outer join operation. To illustrate this fact, show how to rewrite each of the following SQL queries without using the outer join expression.
 - 1. select * from student natural left outer join takes
 - 2. (hard) select * from student natural full outer join takes

		section	-	1		
	teaches	<u>course_id</u>	takes			ı
department instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u> course id	student	course	
dept_nameIDbuildingdept_namebudgetsalary	course_id sec_id semester year	year building room_number time_slot_id	sec id semester year grade	ID name dept_name tot_cred	course_id title dept_name credits	<i>advisor</i> <u>s_id</u> i_id



Integrity Constraints

- A checking account must have a balance greater than \$10,000.00
- A salary of a bank employee must be at least \$4.00 an hour
- A customer must have a (non-null) phone number

			section		1		
		teaches	course_id	takes			•
department	instructor	<u>ID</u>	sec_id	<u>ID</u>	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	id <u>sec_id</u> <u>semester</u> <u>year</u>	semester year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id

n I



Constraints on a Single Relation

- not null
 - name varchar(20) not null budget numeric(12,2) not null
- primary key
- unique $(A_1, A_2, ..., A_m)$
 - Candidate keys are permitted to be null (in contrast to primary keys).
- check (P), where P is a predicate (read DDL.sql)

create table department

(dept_name varchar(20), building varchar(15),
budget numeric(12,2) check (budget > 0),
primary key (dept_name));

			section	_	1		
		teaches	<u>course_id</u>	takes			l
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
<u>dept_name</u>	<u>ID</u> name	<u>course_id</u> <u>sec_id</u>	<u>year</u> building	<u>course i</u> d <u>sec id</u> semester	<u>ID</u> name	<u>course_id</u> title	advisor
building budget	dept_name salary	<u>semester</u> <u>year</u>	room_number time_slot_id	<u>year</u> grade	dept_name tot_cred	dept_name credits	s <u>id</u> i_id

The check clause



- The check (P) clause specifies a predicate P that must be satisfied by every tuple in a relation.
- Example: ensure that semester is one of fall, winter, spring or summer create table section

```
(course_id varchar (8),
sec_id varchar (8),
semester varchar (6),
year numeric (4,0),
building varchar (15),
room_number varchar (7),
time slot id varchar (4),
primary key (course_id, sec_id, semester, year),
check (semester in ('Fall', 'Winter', 'Spring', 'Summer')))
Complex check conditions in section (not supported by sqlite)
check (time_slot_id in (select time_slot_id from time_slot))
```

			section				
		teaches	<u>course_id</u>	takes			ı
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name	<u>ID</u> name	 course_id sec_id	<u>year</u> building	<u>course_i</u> d <u>sec_id</u> semester	<u>ID</u> name	<u>course_id</u> title	advisor
building budget	dept_name salary	<u>semester</u> <u>year</u>	room_number time_slot_id	<u>year</u> grade	dept_name tot_cred	dept_name credits	. <u>s_id</u> i_id



Referential Integrity

- foreign key (dept_name) references department
 - By default, a foreign key references the primary-key attributes of the referenced table.
 - SQL allows a list of attributes of the referenced relation to be specified explicitly.

foreign key (dept_name) **references** department (dept_name)

The specified list of attributes must be declared as a superkey of the referenced relation, using either a primary key constraint or a unique constraint.

			section	_	1		
		teaches	<u>course_id</u>	takes			I
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	<i>advisor</i> <u>s_id</u> i_id



Actions in Referential Integrity

- When a referential-integrity constraint is violated, the normal procedure is to reject the action that caused the violation.
- An alternative, in case of delete or update is to cascade

- Instead of cascade we can use :
 - set null,
 - set default

			section	_	I		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	sec_id	<u>ID</u> . ,	student	course	
dept name building budget	ID name dept_name salary	course_id sec_id semester year	semester year building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Integrity Constraint Violation

- How to insert a tuple without causing constraint violation?
 - Insert father and mother of a person before inserting person
 - OR, set father and mother to null initially, update after inserting all persons (not possible if father and mother attributes declared to be **not null**)
 - OR defer constraint checking

			section		1		
		teaches	course_id	takes			1
department	instructor	<u>ID</u>	sec_id	<u>ID</u>	student	course	
dept name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	semester year building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Built-in Data Types in SQL

date: Example: date '2005-7-27'

time: Example: time '09:00:30' time '09:00:30.75'

■ **Timestamp** Example: **timestamp** '2005-7-27 09:00:30.75'

- interval: Subtracting a date/time/timestamp value from another gives an interval value
- Large objects (photos, videos, CAD files, etc.) are stored as a large object:
 - blob: binary large object
 - clob/text: character large object
 - When a query returns a large object, a pointer is returned rather than the large object itself.

			section		1		
		teaches	course_id	takes			1
department	instructor	<u>ID</u>	sec_id	<u>ID</u>	student	course	
dept name building budget	<u>ID</u> name dept_name salary	id id sem_ester vear	semester year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Index Creation

- We create an index with the create index command
 create index <name> on <relation-name> (attribute);
- create table student (ID varchar (5), name varchar (20) not null, dept_name varchar (20), tot_cred numeric (3,0) default 0, primary key (ID))
- create index studentID_index on student(ID)
- The query:

select *
from student
where ID = '12345'

Demo: EXPLAIN QUERY PLAN

			section		1		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name building budget	I <u>D</u> name dept_name salary	course_id sec_id semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Transactions

- The transaction must end with one of the following statements:
 - Commit work.
 - Rollback work.
- Atomic transaction
- Isolation from concurrent transactions
- ACID
 - Consistency and Durability.

Demo: transaction

			section	_	I		
		teaches	<u>course_id</u>	takes			ı
department in	istructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u> course id	student	course	
building do	<u>D</u> name lept_name alary	course_id sec_id semester year	year building room_number time_slot_id	sec id semester year grade	<u>ID</u> name dept_name tot_cred	course_id title dept_name credits	<i>advisor</i> <u>s_id</u> i_id



Assertions

- An assertion is a predicate expressing a condition that we wish the database always to satisfy.
- The following constraints, can be expressed using assertions:
 - For each tuple in the student relation, the value of the attribute tot_cred must equal the sum of credits of courses that the student has completed successfully.
 - An instructor cannot teach in two different classrooms in a semester in the same time slot
- An assertion in SQL takes the form:

create assertion <assertion-name> check (<predicate>);

			section	-	1		
		teaches	<u>course_id</u>	takes			ı
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u> course id	student	course	
<u>dept_nam</u> e	<u>ID</u> name	<u>course_id</u> sec_id	<u>year</u> building	<u>sec_id</u>	<u>ID</u> name	<u>course_id</u> title	advisor
building budget	dept_name salary	<u>semester</u> <u>year</u>	room_number time_slot_id	<u>semeste</u> r <u>year</u> grade	dept_name tot_cred	dept_name credits	<u>s_id</u> i_id



Authorization

- We may assign a user several forms of authorizations on parts of the database.
 - Read allows reading, but not modification of data.
 - Insert allows insertion of new data, but not modification of existing data.
 - Update allows modification, but not deletion of data.
 - Delete allows deletion of data.
- Each of these types of authorizations is called a privilege. We may authorize the user all, none, or a combination of these types of privileges on specified parts of a database, such as a relation or a view.

		section	-	1		
	teaches	<u>course_id</u>	takes			ı
department instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u> course id	student	course	
dept_nameIDbuildingdept_namebudgetsalary	course_id sec_id semester year	year building room_number time_slot_id	sec id semester year grade	ID name dept_name tot_cred	course_id title dept_name credits	<i>advisor</i> <u>s_id</u> i_id



Authorization (Cont.)

- Forms of authorization to modify the database schema
 - Index allows creation and deletion of indices.
 - Resources allows creation of new relations.
 - Alteration allows addition or deletion of attributes in a relation.
 - Drop allows deletion of relations.

			section	-	1		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	course id sec id semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Authorization Specification in SQL

- The grant statement is used to confer authorization
 grant <privilege list> on <relation or view > to <user list>
- <user list> is:
 - a user-id
 - public, which allows all valid users the privilege granted
 - A role (more on this later)
- Example:
 - grant select on department to Amit, Satoshi
- Granting a privilege on a view does not imply granting any privileges on the underlying relations.
- The grantor of the privilege must already hold the privilege on the specified item (or be the database administrator).

		teaches	course_id	takes			
department	instructor	<u>ID</u>	sec_id	<u>ID</u>	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	semester year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Privileges in SQL

- select: allows read access to relation, or the ability to query using the view
 - Example: grant users U_1 , U_2 , and U_3 **select** authorization on the *instructor* relation:

grant select on instructor to U_1 , U_2 , U_3

- insert: the ability to insert tuples
- update: the ability to update using the SQL update statement
- delete: the ability to delete tuples.
- all privileges: used as a short form for all the allowable privileges

			section	_	I		
		teaches	course_id	takes			1
department	instructor	<u>ID</u>	sec_id	<u>ID</u>	student	course	
dept name building budget	<u>ID</u> name dept_name salary	id secid semester year	semester year building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Revoking Authorization in SQL

- The revoke statement is used to revoke authorization.
 revoke <privilege list> on <relation or view> from <user list>
- Example:
 - revoke select on student from U_1 , U_2 , U_3
- <privilege-list> may be all to revoke all privileges the revokee may hold.
- If <revokee-list> includes public, all users lose the privilege except those granted it explicitly.
- If the same privilege was granted twice to the same user by different grantees, the user may retain the privilege after the revocation.
- All privileges that depend on the privilege being revoked are also revoked.

			section	_	I		
		teaches	<u>course_id</u>	takes			ı
department in	nstructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
building	<u>ID</u> name dept_name salary	course_id sec_id semester year	year building room_number time_slot_id	<u>course_id</u> <u>sec_id</u> <u>semester</u> <u>year</u> grade	<u>ID</u> name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Roles

- A role is a way to distinguish among various users as far as what these users can access/update in the database.
- To create a role we use:

create a role <name>

- Example:
 - create role instructor
- Once a role is created we can assign "users" to the role using:
 - grant <role> to <users>

			section	_	1		
		teaches	<u>course_id</u>	takes			I
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	<i>advisor</i> <u>s_id</u> i_id



Roles Example

- create role instructor;
- grant instructor to Amit;
- Privileges can be granted to roles:
 - grant select on takes to instructor,
- Roles can be granted to users, as well as to other roles
 - create role teaching_assistant
 - grant teaching_assistant to instructor,
 - Instructor inherits all privileges of teaching_assistant
- Chain of roles
 - create role dean;
 - grant instructor to dean;

 grant dean to Satoshi; 			section		1			
teaches		course_id	takes			1		
(department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
	dept_name	<u>ID</u> name	 course_id sec_id	<u>year</u> building	<u>course_id</u> <u>sec_id</u> semester	<u>ID</u> name	<u>course_id</u> title	advisor
	building budget	dept_name salary	<u>semester</u> <u>year</u>	room_number time_slot_id	<u>year</u> grade	dept_name tot_cred	dept_name credits	. <u>s_id</u> i_id



Authorization on Views

- create view geo_instructor as
 (select *
 from instructor
 where dept_name = 'Geology');
- grant select on geo_instructor to geo_staff
- Suppose that a geo_staff member issues
 - select * from geo_instructor,
- What if
 - geo_staff does not have permissions on instructor?
 - Creator of view did not have some permissions on instructor?

			section		1		
		teaches	<u>course_id</u>	takes			1
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u>	student	course	
dept name building budget	ID name dept_name salary	course_id sec_id semester year	year building room_number time_slot_id	<u>course id</u> <u>sec id</u> <u>semester</u> <u>year</u> grade	ID name dept_name tot_cred	course_id title dept_name credits	advisor s_id i_id



Other Authorization Features

- references privilege to create foreign key
 - grant reference (dept_name) on department to Mariano;
 - Why is this required?
- transfer of privileges
 - grant select on department to Amit with grant option;
 - revoke select on department from Amit, Satoshi cascade;
 - revoke select on department from Amit, Satoshi restrict;
 - And more!

			section	-	1		
ı		teaches	<u>course_id</u>	takes			ı
department	instructor	<u>ID</u>	<u>sec_id</u> semester	<u>ID</u> course id	student	course	
dept_name building budget	<u>ID</u> name dept_name salary	course_id sec_id semester year	year building room_number time_slot_id	sec id semester year grade	ID name dept_name tot_cred	course_id title dept_name credits	<i>advisor</i> <u>s_id</u> i_id



FIN

Any questions?