

Presented by Michael Ford















- Non SQL? Non relational? Not only SQL?
- DBMS without relational model
- Key Features:
  - Simplicity of data model
  - Control of availability
  - Horizontal scaling
  - Flexibility

## NoSQL – CAP Theorem

- Impossible for a distributed system to accomplish all three of:
  - Consistency
  - Availability
  - Partition Tolerance
- NoSQL systems often sacrifice consistency by choosing eventual consistency over ACID Transactions











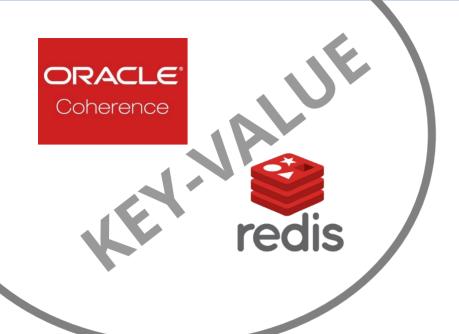








Key	Value
"India"	{"B-25, Sector-58, Noida, India – 201301"
"Romania"	{"IMPS Moara Business Center, Buftea No. 1, Cluj-Napoca, 400606",City Business Center, Coriolan Brediceanu No. 10, Building B, Timisoara, 300011"}
"US"	{"3975 Fair Ridge Drive. Suite 200 South, Fairfax, VA 22033"}









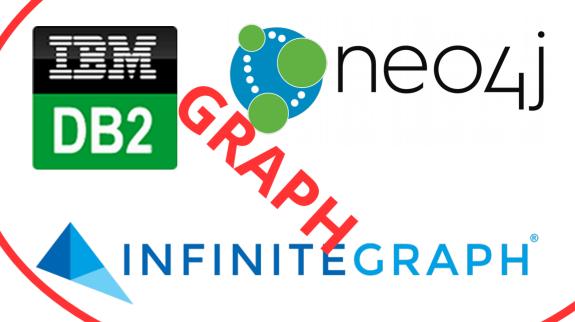












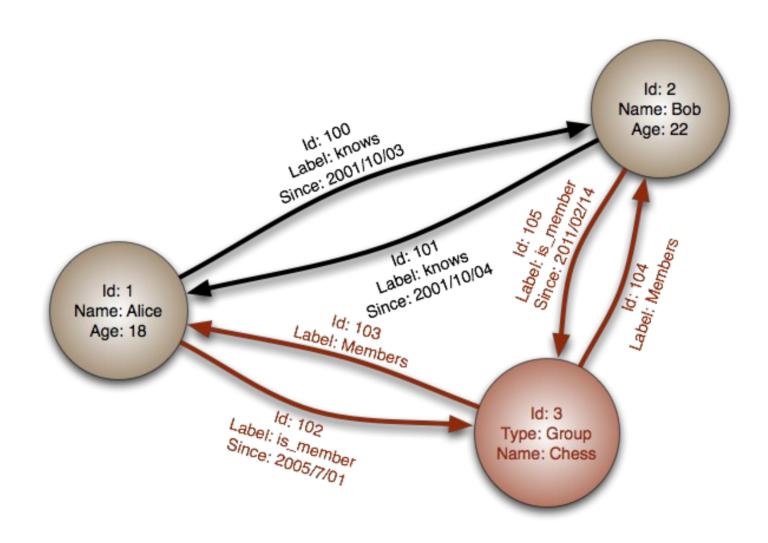
































# BigTable

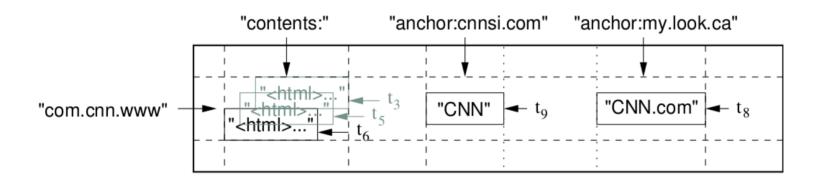


### Overview

- Multi-dimensional, sparse, sorted map with tabular structure
- Dynamic control of schema:
  - Data layout
  - Data format
- Schema choice determines data locality
- Therefore user choices significantly affect speed

## Indexing

- Rows
  - Sorted alphabetically, keeps data together if using reverse naming
- Column-families
  - Need to provide column-family:column to retrieve data
  - Columns in family are compressed together so should be same type
  - Designed to have few families but many columns
- Timestamps
  - Default assignment real-time (microseconds)
  - Can be client assigned
  - Garbage collection can be set to automatically keep only n most recent



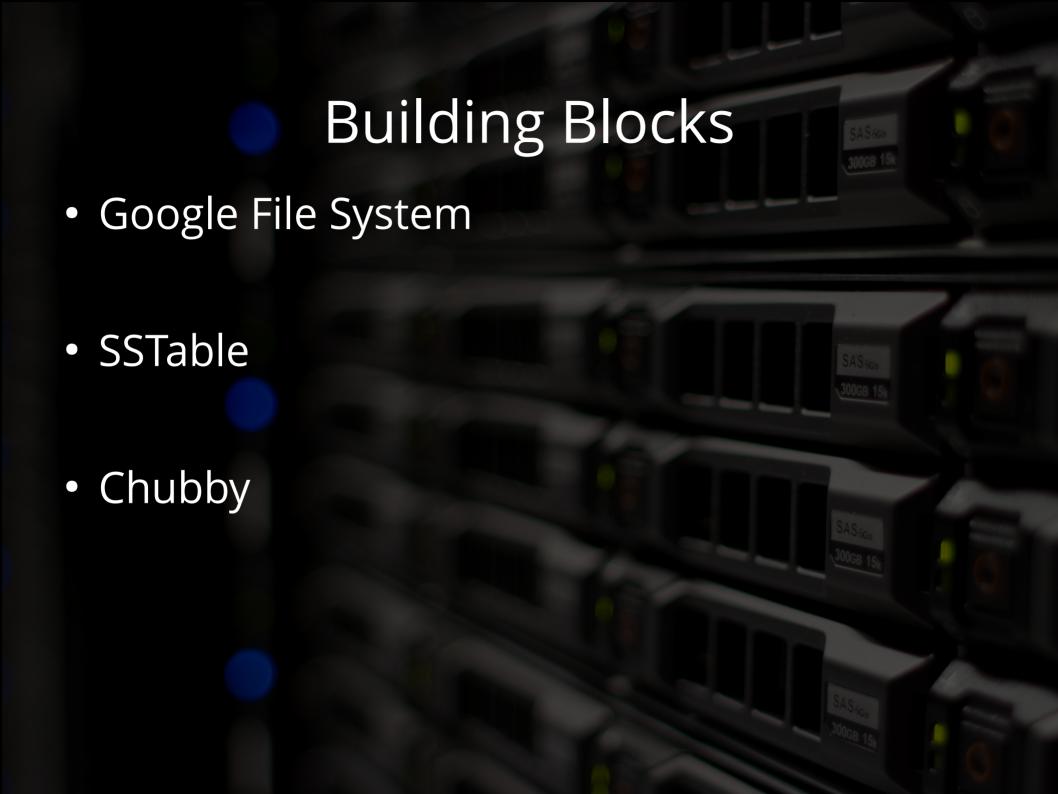
# Partitioning: Tablets

Range of rows

Unit of data distribution

• 100-200 MB each

Automatic splitting



## Google File System

- Distributed data storage system
- Files divided into fixed-size chunks
- Chunks replicated 3 times across chunk servers
- Chunk location and metadata managed with master server

# Sorted String Table - SSTable

Key-value map datatype

Immutable

Sorted sequentially for easy lookup

## Chubby

- Distributed lock service
- Determines lock status with Paxos consensus algorithm
- 5 chubby replicas
  - 1 of which is master
- Used for:
  - Ensure there is only 1 active master
  - Store location of root tablet
  - Discover and delete tablet servers
  - Store column family information
  - Store access control lists



## General Structure

### **Tablet servers**

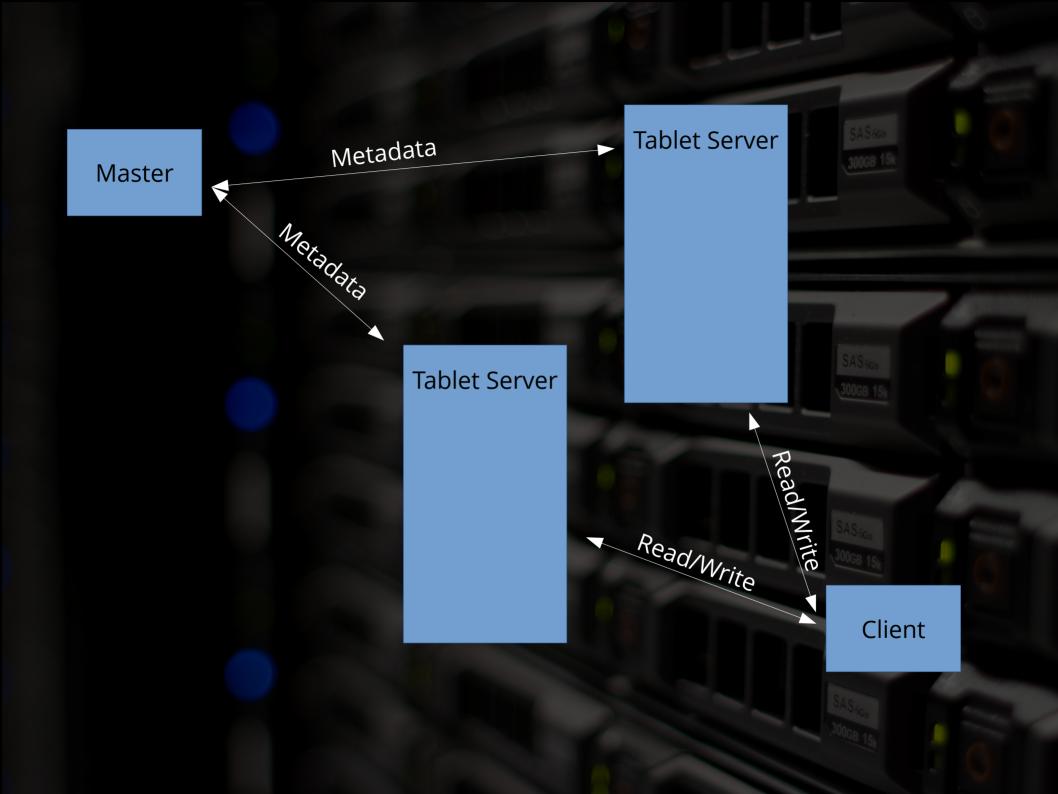
- Manages set of tablets (10-1000 tablets)
- Manages read/write directly with client
- Splits tablets when too big

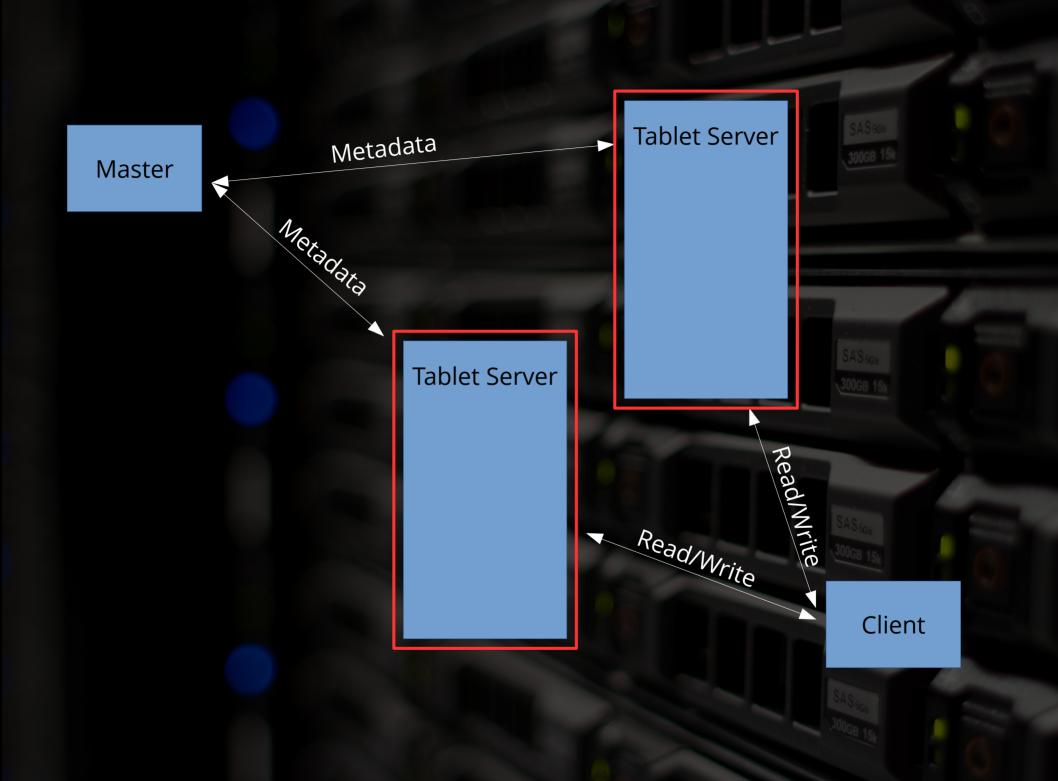
#### Master server

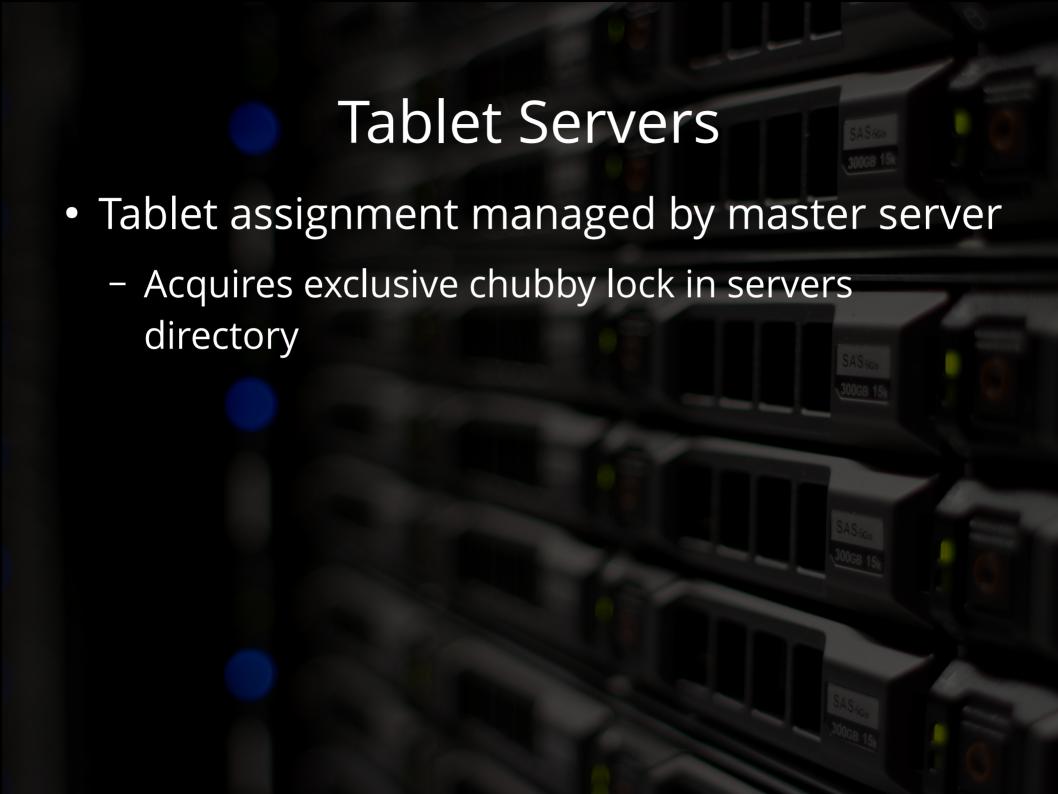
- Assigns tablets to tablet servers
- Managing tablet servers and load
- Garbage collection

### Client

 Can iterate over: column families, columns within a family, rows







Chubby file

## **Tablet Servers**

- Tablet assignment managed by master server
  - Acquires exclusive chubby lock in servers directory
- Tablet location storing
  - Chubby File
  - Root tablet
  - Metadata tablets
  - Tables containing data tablets

# Tablet Serving

### Tablet Storage

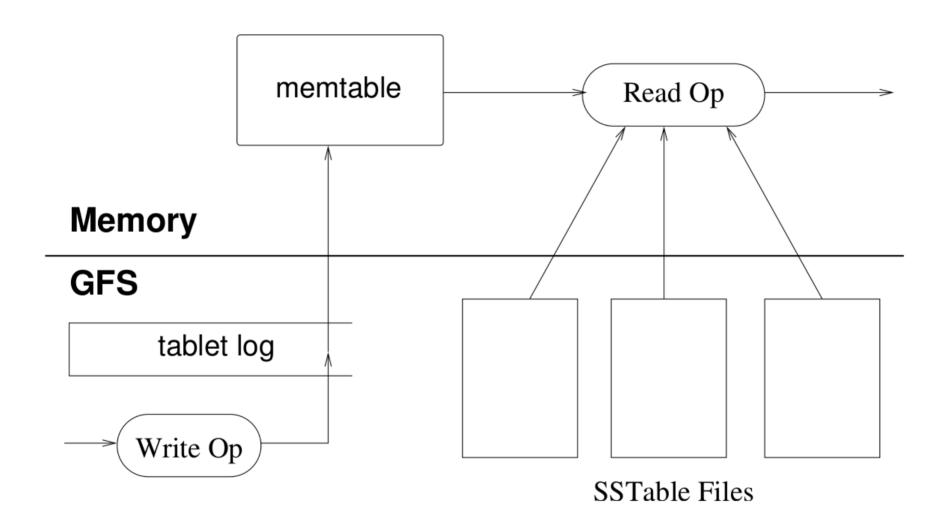
- Recent commits logged in memtable
- Old commits logged in SSTables

#### **Tablet Read**

- SSTable locations from metadata
- Constructs tablet by merging SSTables and memtable using commits

#### **Tablet Write**

- Committed to log
- Contents inserted into memtable



## Compaction

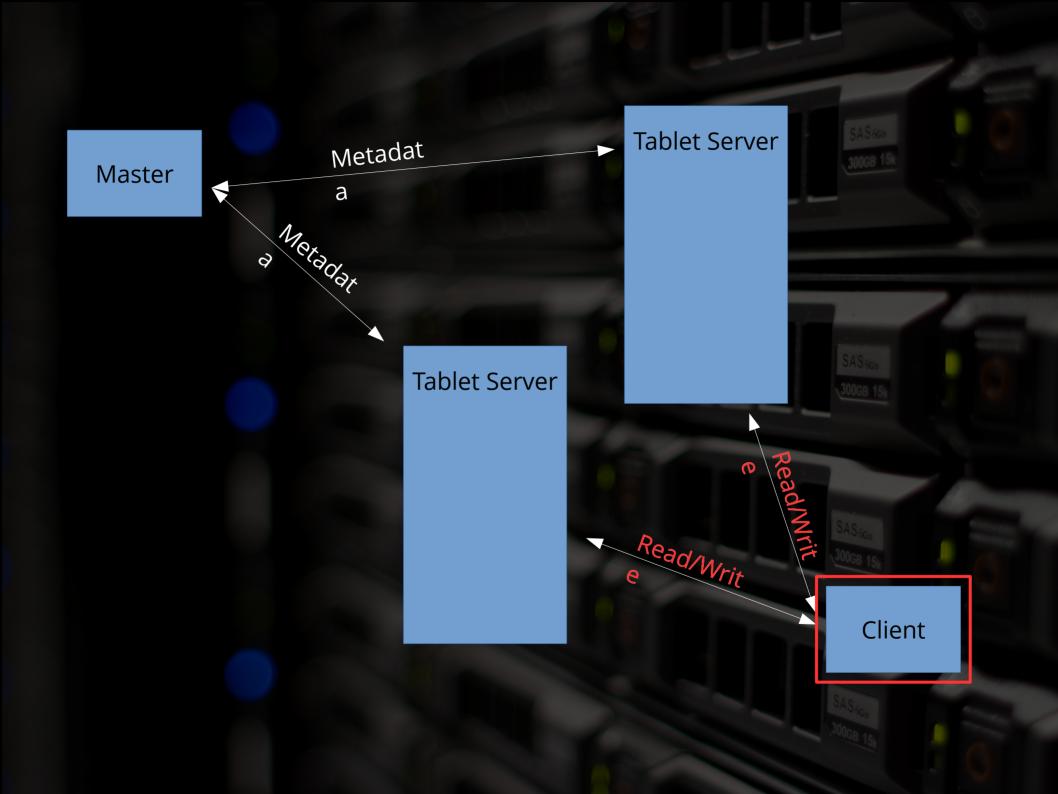
Problem: memtable fills up with commits

### Minor compaction

- Memtable made into new SSTable
- New memtable created

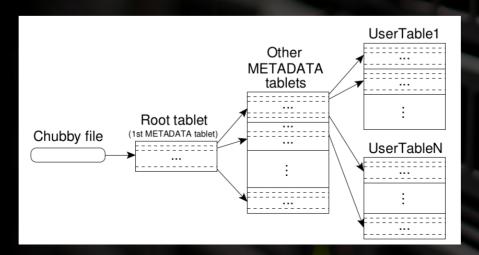
### Major compaction

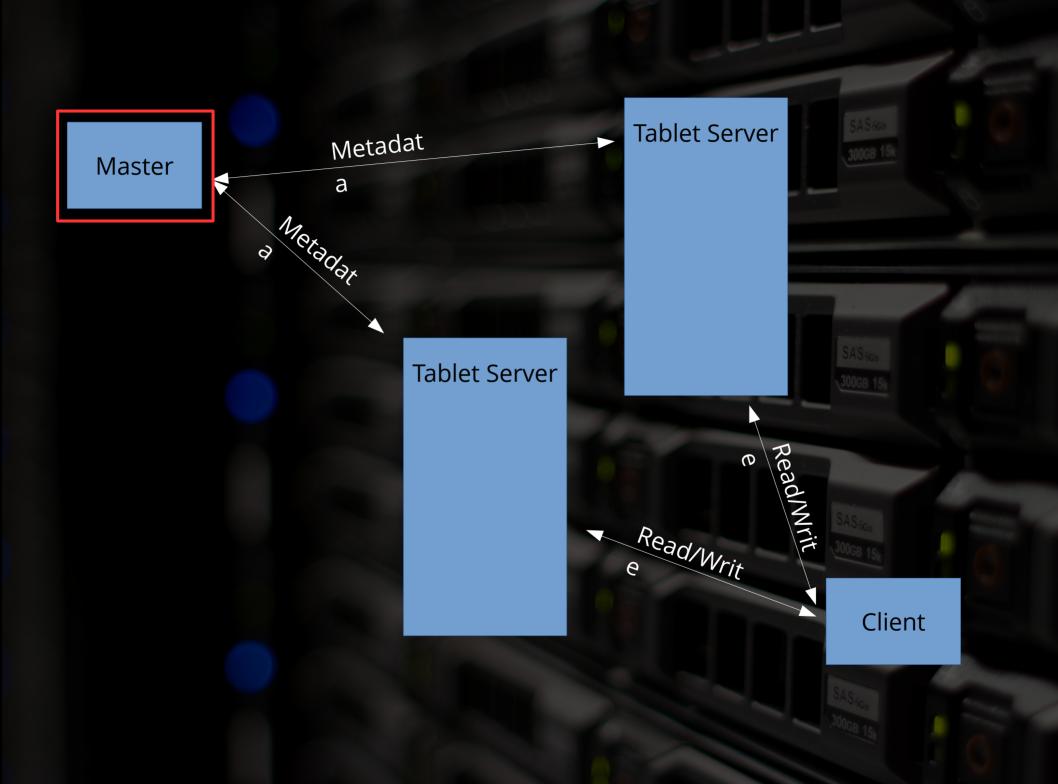
Combines memtable and SSTables with no deletions



## Client Features

- Caches tablet locations
- If location inaccurate, moves up tablet server hierarchy
- Atomic single-row transactions



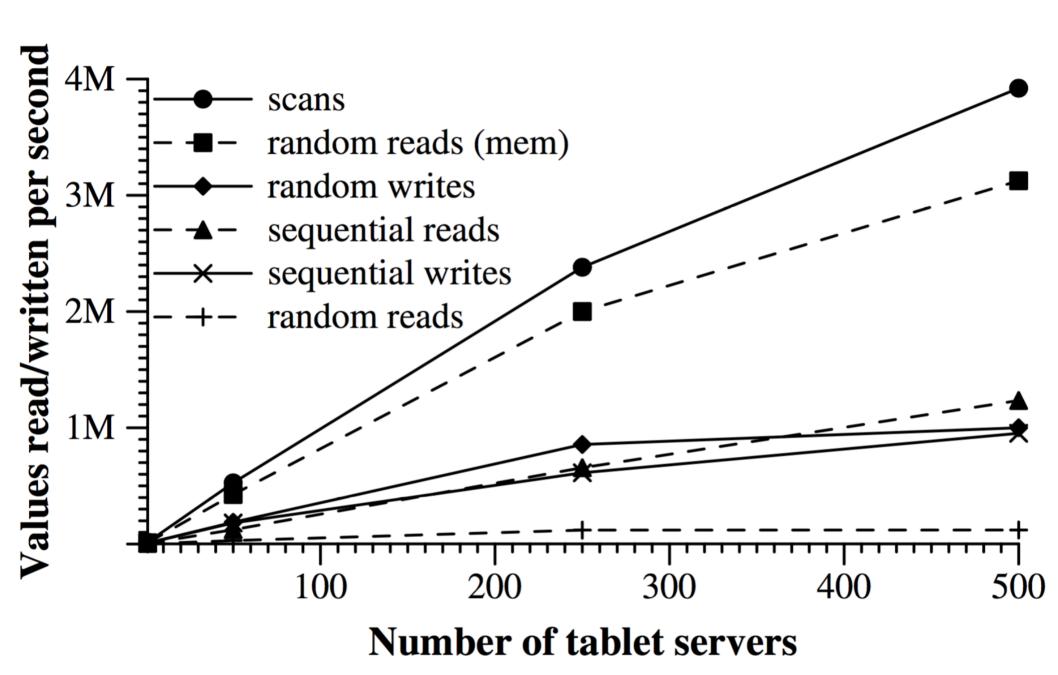




# Experiments and Examples

### Benchmarks

- Sequential write
- Sequential read
- Scan
- Random reads
- Random write





Apache Cassandra, Hbase

Google Cloud

Google Spanner

