Analysis and Synthesis Techniques for Hopfield Type Synchronous Discrete Time Neural Networks with Application to Associative Memory

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thinked —In the present paper we establish a qualificative theory for Synchropous discrete time Elopfield-type neural actionsky that objections are accomplished in two phases. First, we address the anothers of the class of neural networks considered. Next, waking use of these results. he develop a violitizate procedure for the class of negrat networks considered horows.

In the prolycle section we estilize techniques from the theory of lurge-weath intergularizing dynamical systems to derive total for the asymptotic stability of an equilibrium of the mount necessary. We place present estimates for the rate at which the trajectures of the network will converge from an initial condition to a final state. In the synthesis, section we made the stability lests from the analysis section as constraints to develop a design algorithm for accomplist memories. The prover algorithm goal enters that each drain dimensity will be stored as an equilibrium and that each deserted themsely salt for assumptioth atts stable.

The applicability of the propert results is demonstrated by means of ewii specific esamples

1. Is rapportunis

 $\P[N][H]$ and [2]. Repaid proving makely for neutral Inclworks that are finds interconnected, in addition, Hopfield was able to show that the trajectories of these networks words not oscillate, but would seek local minimaof a cottom energy function. The main assumption of this analysis is that the interconnection many moist be senmetric with zero diagonal elements. Using the energy function as the basis for a design procedure (see [19]) fight). Tetworks have been utilized to convert analog signals to digital signals, to decompose an maginye signaland to implement associative memories. Several other applications are discussed in [8]

Pio particular application discussed in this paper is the associative memory. The year of the associative memory is responsible reflect particular to a little for a little by a misuch a

manner that the network will recognize an apput pattern (i.e., $P = U^{(0)}(\Delta F)$ that is mast one on the stored vectors tup, VT) as the stored vector. Hopfield showed (no) a neonal network could simulate an associative atennory if each Pricould be stored as a local minimum of his energy. foretoor.

The recent article [17] talso see. 7[t presents both a (horough of troduction to Heyotield discrete time neural networks and a detailed stochastic analysis of their memone capacity. The analysis in [17] is applicable to both synchronous and asynchronous systems with symmetric interconnection matrices. Section V of that paper mentions the need for methods to test the stability and arractivity of patierns that we desire for the network to store. Such texts are presented in the analysis section of the present paper. The nests are applicable to networks having either the symmetric interconnection matrices of 17 or nonsymmetric interconnection matrices of the present paper. In addition, the present paper presents methtals for estimating the rate of concergence to an equithorons. Resolts similar to those presented herein for miscrete time systems are presented in [9] and [18] for the continuous time Hopfield model.

The application of neural networks to coding theory has been discussed in [16]. This poper considers a twoloyered network, in which the first layer is an Adaline nerwork designed to transfering a set of vectors from a "leature space" to a set of vectors in the "code snace." The new vectors in the contemptor one chosen to be pseudo-orthogonal. The second layer is also an Adahue. network, which has been designed to classify the codes that are output from the first faver. Notice that the network has been assigned and is not adaptive. The examples of the overcut paper are similar to the examples or 114 and may be interpreted in the same manner. However, same complex discuss single laser, tendralek, and auto-associative networks. The addition of a preprocessring layer would read obtainly improve the performance in our examples, because a little patterns to be stored theing outputs of the preprocessort could be chosen as pseudoorthogonal sectors. We did not insestigate this approach

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Moroscopi accioe Marco 8, 1960. The wore of A. S. Michellin III. I. Null was supported by the National Science Foundative inter-trace 17.8 (SCCO). The work of U.A. Virtell was supported to a le kwish 5.1 of the extraction of Note. Dance Applied Mathematics Lander and Synthe Not 1960. Secret FPS 35 07024. This population recommends to the National Applied Mathematics in contract the Associated Lander with the Department of Theory call Language 1.1 octaon of Sopo Dance, Note 16 oct 18 footh, 1.1 a. Language 2.1 octaon of Sopo Dance, Note 16 oct 18 footh, 1.1 a. Language 3.1 octaon of Sopo Dance, Note 16 oct 18 footh, 1.1 a. Language 3.1 octaon of Sopo Dance, Note 16 oct 18 footh, 2.4 a. Language 3.1 octaon of Sopo Dance, Note 16 octaon 3.1 mbhadae, 20 octaon of Sopo Dance, Note 18 octaon 3.1 mbhadae, 20 octaon of Sopo Dance, Note 18 octaon 3.1 mbhadae, 20 octaon of Sopo Dance, Note 18 octaon 3.1 mbhadae, 20 octaon of Sopo Dance, Note 18 octaon 3.1 mbhadae, 20 octaon of Sopo Dance, Note 18 octaon 3.1 mbhadae, 20 octaon of Sopo Dance, Note 18 octaon 3.1 mbhadae, 20 o

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because our main emphasis is on the presentation of general analysis and synthesis procedures that allow networks to be designed with a momore capacity of N, where N is the number of measures at the network.

The present paper does not follow Hopfield's global energy function approach. Instead, we develop in Section IV a local theory for fully intercentracted discrete time neural networks. In our analysis we view the network as an interconnected system of N unity c processors. The local theory includes sufficient constitution for the esymptotic stability of an equilibrium, and bounds on the rate of convergence of system trajectories.

In order to correspond closely to biological neural networks. Hopfieth's discrete time model witches asynchronously. We consider a made that witches synchronously. By making this assuraption, we are able to present deterministic prioris of stability. In addition, the model is suitable for a straightforward implementation via parallel processing computers. Both of these characteristics are important, because we are interested in an engineering application, (pattern recognition) for which we desire guaranteed storage and fast responses

In Section V, using as constraints the stability tests that were determined in Section IV, we deterly a method to design associative merginies. The procedure guarantees to store trace of the reservices I(s, N) the designer wishes to store. The only condition on the sectors is that they must be linearly radependent. The sexign procedure chooses the interconnections and thresholds in a manner to store each desired pattern as an equal brium of the network, and to satisfy an asymptotic stability constraint. This is in contrast to methods that choose the interconnections and thresholds to store the patterns as energy minima. In Section VI we present two design examples.

The designs by the method presented herein will not result in a symmetric interconnection matrix in alost cases. The reader most remember that the symmetry condition is a mathematical assumption foot is necessary for the Hopfield energy turterion approach. Biological networks are not expected to have symmetric interconnections, and stetworks designed to be symmetric chould be expected to line their symmetry property during the implementation process due to parameter inaccuracies. The greatest synthesis procedure is based on the properties of M-matrices. These matrices are robust to parameter inaccuracies.

The present design method does not have a prior of global stability (i.e., oscillatory solutions are not preven to be remeastent) due to the tack of a symmetric interconnection matrix. Instead, we have guaranteed asymptotic stability for the sloved petierns. As our second example illustrates, the number of numeroscipent solutions is very small.

II. Nicion Nickows Moralis

In [1] and [2] Hopfield presents, respectively, models for continuous and discontinuous neural networks. Associated with each of these models is an energy function.

which the trajectories of the system seek to inclinitize. The network is considered to have R neurons. The inputs and outputs of the set of R neurons are described in vector notation, by $\mathbf{a} = [\mathbf{a}_1, \cdots, \mathbf{a}_N]^T$ and $\mathbf{V} = [V_n, \cdots, V_N]^T$, respectively. Here, $\mathbf{a}_i \in R^{\infty} = [0,\infty] \to R \in (-\infty,\infty)$ and $\mathbf{f}_i \in R^{\infty} \to R$ denote the input and output at the oth neuron respectively.

The dynamics of the discontinuous model are described by the following algorithm. Each neutron computes its output according to the formula:

$$p_i = \sum_{i=1}^{N} T_i I_i^* - I_i$$

$$V_{i,k} = \text{sgn}\{n_i\}$$
(DM)

where $I_{\rm q}$, the strength of connection between neutron i and j, belongs to R, and $I_{\rm d}$ an external input to the eth neutron, also belongs to R. The order of which neutrons change states is completely random and asynchronous, but the mean rate of change of the state, $W_{\rm c}$ is fixed. In the above, the function ${\rm kgr}$ is the signorn function defined as

$$\operatorname{sgn}(x) = \begin{cases} -1, & x \ge 0 \\ -1, & x \le 0 \end{cases}$$

In the case where $x = 2 \| V_{-1} \|$ retains its previous value. The equation that describes the dynamics of Hapfield's continuous model can be simplified to

$$\dot{a}_i = \sum_{j=1}^{N} T_j V_j - b_i a_j + I_j$$

$$V_i = g_i(a_i)$$
(CM)

where $a_i = da_i / dt$, $f_i \in R$, $b_i \in R$, and $f_i \in R$

Remark 7: Fire function $g_1 : R \rightarrow R$ is a monotone nondecreasing, continuously differentiable backton with $g_1(0) \rightarrow 0$. Thus $g_1(1)$ satisfies a sector condition: i.e., there are two positive seal numbers g_1 and g_2 such that $g_1 \in g_1(r) / 1 \leqslant g_2$ for all $x \in B(r_1) + \{y \in R^+ | y| \leqslant r\}$ for some $r_1 > 0$. In the above, r_1' denotes absolute value

The design of associative memories (AM's) for the discontinuous model has been considered in several bapers (see (3|-|b|)). The method of (5||b||a||s) discussed in [1]where the outer product method is shown to be capable of storing about 0.15N patterns. A difficulty in designing these networks is due to the discontinuous input-insput relationships. Often, it is possible to state the desired patterns as stable states of (DM); however, the discrete garrie of the model's dynamical equation makes it difficult to grayance that a pattern (stered as an equilibrium point) will attract neighboring states. This model is cittle outs to implement in hardware due to the requirement of a large number of interconnections and the inflorent capacitances and belays. This model is often implemented on a digital compater where the requirement of random asynchronous switching times is approximated by a pseudorandom switching sequence. Synchronous mode a hore recently been discussed (e.g., in [7] and [17]).

AM's con also be implemented by Hopfield's continuous time model, as discussed in [2] [9]-[11], [22], and [22]. The design method of [1] is applicable in the high gain limit where (CMI approaches (DM). This technique does not goarantee to stene cuch of the desired patterns (as asymptotically stolds equilibrium points). The methods discussed in [9]-[11], [22], and [23] guarantee to store each desired pattern. The methods are compared in [10], Again, a problem orises with implementing this circuit in hardwate because of inherent parasitic dynamics, the parameter inaccuration, and the amount of space required to implement a large number of resistors. For a singly of implementation questions see [12], For some interesting implementations that attempt to avoid particles, refer to [24], [25].

Most applications to date have been implemented by special purpose software. If their ICMI was simulated by nameneal routines on a digital computer of IDMI was implemented in a synchronous or pseudoscopiclitonious money. The present paper seeks to avoid the simulation problems of both models, the dosign problems of the discontinuous problem, and the implementation problem of the continuous model by considering a discrete time stocknown model that approximates ICMI. We analyze this model directly, by the method of [13] In addition to these analysis testing, we present in Section V a synthesis technique for the synchronous discrete time model.

By Atler's method, (CMI) can be approximated in discrete time by replacing

$$u(x) = \frac{u((k-1)T) - u(kT)}{T} \qquad \text{for } kT \le t \le (k+1)T.$$

$$(2.5)$$

In the above equation, T is the sampling period, and as T > 0 the accuracy of the approximation invitages. Substitution at 12.10 into (CM) yields

$$\frac{a_i\big((|k-1|)T\big) + a_i\big(kT\big)}{I} = \sum I_i[\mathcal{I}_i(kT) + h_ia_i\big(kT\big) + I$$

Hence.

$$u_i(|k|+1) = \sum T_i V_i(|k|) + (1 - Tr_i(u_i(|k|) + I_i))$$

= $\sum T_i V_i(|k|) + u_i u_i(|k|) + I_i$ (DT)

where $\alpha = 1$. The $\alpha(3\alpha) + (177)\alpha(3\beta)$, and $\alpha(3\alpha) + (187)$. The equations describing ODF) and ODM) are very similar. For both systems, the orbit to each neuron is a weighted sum of the outputs of all the neurons. The features that distinguish model (DT) include its operating in a synchronous minde and its use of a agmoidal input—output function. These features give model (DT) two advantages over (DM); first, (DT) allows direct implementation in software of special pulpose hardware, and second. It allows design techniques that guarantee both sadd to and autial ticire in stored patterns. The dynamics of the discrete time, synchronous model described by (DT) are the subject of the remainder of this paper.

III. Loren Scale Systems Aggreenti

To circumvent the difficulties inherent in the direct analysis of the system (DT) as a whose, we find it convenient to incorporate the method of analysis proposed in [13]. By this approach we will view the system (DT) as the interconnection of tree or isolated subsystems. This section will be decaded into two parts. The first pair transforms (DT) into the form necessary for the second part. The second part explicitly defines the system in a manner consistem with the large-scale systems approach.

3.1. The Translation of the Equalibrium

Defining $T = [T_0] \in R^{N \times N}$, $A = \operatorname{diag}(x_0) \cap \operatorname{diag}(A_1) \in R^{N \times N}$, $A = [t_0] \in R^N$, and $a = [t_0, \cdots, t_N]^T - R^N \to R^N$. Here the equation (DT) in matrix form becomes

$$u(k+1) = I\Gamma(k) + i\omega_0(k) + I$$

 $\Gamma(k) = g(\nu(k)),$ (DF)

In general, the network will have equilibria that may not be logared at the origin $(a_1+b_1)+(a_2+\cdots, b_n)$. The equilibria of (DTI may be desired enchances that we have stored, or may be extraneous equilibria. Each equilibrium must, by definition, satisfy the equation

$$a(k) = T\Gamma(k) + As(k) + I$$
.

This equation is equivalent to

$$0 = TV(\lambda) + Ra(\lambda) + I \qquad (3.11)$$

where the matrix H = A - E and E is an M dimensional identity matrix. We are able to translate any equilibrium of (DT) into the origin in the following manner. Let

$$\frac{g(k) - g(k) - g^*}{O(g(k)) - g(g(k)) - g(g^*)}$$

where u^* satisfies (X, I), and $(I = \{(i_1, \cdots, (i_n)\}^T)\}$ Then $p(\hat{x} + I) + u^* = Iq(p(\hat{x}) + u^*) + A(p(\hat{x}) + u^*) + I$ $= I(\{i\}, p(\hat{x})\} + p(u^*)\} + Ap(\hat{x}) + Au^* + I$ $= I(\{i\}, p(\hat{x})\} + Ap(\hat{x}) + Ip(u^*)\} + Au^* + I$

Thus

$$p(k+1) = TG(p(k)) + Ap(k) + Ty(n^*) + Bn^* + I$$

= $TG(p(k)) + Ap(k)$ (2)

which has an equilibrium at p(k) = 0. Henceforth, we will use $1 \ge 1$ to represent the model of the nearst network under study. We will assume that all equilibria of (\ge) are isolated. In the next society, we will address the stability properties of (arbitrary) given equilibrium points. In our approach we will assume arthoral toos of generality that a given equilibrium under its estigation is boated at the one give

Remark 2: Due to the relationship between G(x) and g(x), G(x) is given monomically nondegreesing continuously differentiable function with G(0) = 0. Again, $G_i(0)$ will satisfy a sector condition $x_i \in G(i) \times (x) \in i$, for all $x \in B(x_i)$ with some $x_i + 0$ for $x = 1, \dots, N$. Here, $x_{i,j}$ and x_i , are two positive real numbers that may not be the same as those in Remark.

5.2 Interconnected Systems

We rewrite (\$) into the form

$$p_j(k+1) = \sum_{j=1}^{N} T_{ij} G_j(p_j(k)) + \mathcal{A}_{ij} p_j(k),$$

$$\text{for } j \neq 1, \dots, N = \{\Sigma_i\}$$

System (2.) may be viewed as an interconnection of M free (isolated) subsystems represented by the countries

$$x_i(k-1) + T_n G_i(x_i(k)) + A_n x_i(k), \quad \text{for } i \neq 1, \dots, N$$
(8.)

with the interconnecting structure defined by

$$H_i(\mathbf{x}_1, \cdots, \mathbf{x}_N) \triangleq \sum_{\substack{i=1\\i\neq j}}^N T_{ij}G_j(\mathbf{x}_i(\mathbf{x}_j)), \quad \text{for } i=1, \cdots, N.$$

Following the method of analysis in [13], we will establish stability results for system (Σ) that are phrased in terms of the qualitarive proporties of the free subsystems (S_i) and in terms of the proporties of the interconnecting structure given in (3.2).

IV. STABILITY ANALYSIS

The stability results of the entire neutral network (2) are determined by the stability results of the interconnection structure (3,2) and the properties of the interconnection structure (3,2) of the neutral network. In this section, we first analyze the stability of the free subsystems. Next we establish stability results for the entire neutral network (2) that are phrased in terms of the subship properties of the individual free subsystems (2) and in the terms of the properties of the individual free subsystems (2) and in the terms of the properties of the interconnecting structure given by 13.2).

4.1. Stability Retails of the Free Subsistents

To analyze the stability of a free subsystem (S,t. we need the tollowing assumption

Armonomou (A.D.: For CS.L.

$$0 < \sigma_0 \triangleq ((A_m + T_n)c_{12}) < 1$$

where c_{ij} and c_{ij} are defined in Remark 2

Proposition T if (A-D) is true, then the equilibrium $T_i \in O(nf(S_i))$ is asymptotically stable.

Proof (Thomse $v_i(x_i(k)) = v_i(k)$) as a Lyapunov function for (S.). Then the first furward difference of v_i along the solution of (S.) is given by

$$\begin{split} \nabla c_{i N_{i}}(\mathbf{x}_{i}(k)) & \leq c_{i}(\mathbf{x}_{i}(k)|\mathbf{X}_{i}(k)) = c_{i}(\mathbf{x}_{i}(k)), \\ & = (A_{i} \cdot \mathbf{x}_{i}(k) + I_{n}G_{i}(\mathbf{x}_{i}(k)), = x_{i}(k)), \\ & \approx (cA_{n} - (Y), (k)) + cI_{n}G_{i}(\mathbf{x}_{i}(k)), \\ & \approx [(A_{n}^{i} - (k)|I_{n}^{i}(I_{n}^{i} \mathbf{x}_{i})], \mathbf{x}_{i}(k)). \end{split}$$

Clearly, r_i is positive definite and ∇r_i is negative definite if (A-1) is true. Therefore, the equilibrium $|\mathbf{r}_i| = 0$ is asymptotically stable (cf. $\{\mathcal{H}_i\}_i$).

4.2 Stability Results of the Newall Network

Assumption (A-2). Given σ_{c} of (A-1), the successive principal minors of matrix $D = [D_{c}]$ are all positive, with

$$D_{ij} = \begin{cases} -\{|\boldsymbol{\sigma}_i| \leq 1\}, & i = j \\ -|\boldsymbol{\sigma}_{ij}|, & i \neq j \end{cases}$$

where $\sigma_n = T_{\alpha'}\sigma_{\alpha'}$ ($\sigma_{\alpha'}$ is defined in (A-1)).

Theorem $P: \Pi(A;1)$ and (A;2) are true, then the equilibrium p=0 of (Σ) is asymptotically stable.

Proof — We choose a Lyapunuv function for (Σ) .

$$c(|\rho(k)\rangle = \sum_{i=1}^{N} \lambda_i |\rho_i(k)\rangle$$

where $\lambda_i > 0$ for $i = 1, \dots, N$. Then the first forward difference of ϕ along the solutions of (Σ) is given by

$$\begin{split} &\nabla v_{1,1}A[p(k)] \\ &= v(k+1) - v(k) \\ &= \sum_{i=1}^{N} \lambda_{i} [[p_{i}(x|r)]^{i} - [p_{i}(k)]] \\ &= \sum_{i=1}^{N} \lambda_{i} [[p_{i}(x|r)]^{i} - [p_{i}(k)]] \\ &= \sum_{i=1}^{N} \lambda_{i} \frac{1}{i} [A_{ii}p_{i}(k) + \sum_{i=1}^{N} T_{i}G[(p_{i}(k))]^{i} - [p_{i}(k)]^{i}] \\ &= \sum_{i=1}^{N} \lambda_{i} \frac{1}{i} (A_{ii}^{i}p_{i}(k)) [-\sum_{i=1}^{N} T_{ii}G[(p_{i}kk))]_{1} - [p_{i}(k)]^{i}] \\ &= \sum_{i=1}^{N} \lambda_{i} \left[[A_{ii} + 1] [p_{i}(k)]_{1} + \sum_{i=1}^{N} T_{i} [G_{i}(p_{i}kk)]_{1} \right] \\ &= \sum_{i=1}^{N} \lambda_{i} \left[[A_{i} + 1] [[p_{i}(k)]_{1}^{i} + \sum_{i=1}^{N} T_{i} [G_{i}(p_{i}k)]_{1} \right] \\ &= \sum_{i=1}^{N} \lambda_{i} [A_{i}(-1) + [T_{i}, c_{i,j}]][p_{i}(k)] \\ &+ \sum_{i=1}^{N} \lambda_{i} [A_{i}(-1) + [T_{i}, c_{i,j}]][p_{i}(k)] \\ &= \sum_{i=1}^{N} \lambda_{i} [A_{i}(-1) + [T_{i}, c_{i,j}]][p_{i}(k)] \end{aligned}$$

where $k = (k_1, \cdots, k_n)^n$ and $\mu = 0$ $p, 1, \cdots, p_n)^n$. Since by (A-2) $|D_{ij}| \leqslant 0$ where $i \neq j$, and the successive principal minors of matrix D are all positive |D| is an M-matrix D the properties of M-matrices, |D| exists and each element of |D| is nonnegative. Hence there exists a vector $p_i = \{y_1, \cdots, y_n\}^n$, with $|y_i > 0|$ for $i = 1, \cdots, N_i$ such that (see [10])

$$y^2u<0$$
, where $y^*=\lambda^*D$, $3\times (D-)^2v>0$.

Therefore, ∇v is negative definite, v is positive definite, and the equilibrium $|\phi| = 0$ of $(\Sigma) > asymptotically stable (c) <math>\{(0\})$

The Lyapunov function in the proof of Theorem 1.5.4 generalized Hamming distance. Thus if for an equilibrium e^* the translation (T) is applied and the assemptions of

Theorems I are found to be true, then Theorem I knows that for initial conditions sufficiently near x^a , the generalized Hamming distance from x^a to solutions of (DTI will asymptotically approach zero. The weighting vector A is included to increase the applicability and to decrease the conservatism of Theorem 1. One method for choosing A is shown in the proof of Theorem I. Since the value of a is not unique, [20] and [21] consider methods for choosing a in an optional fashion.

. Assemblion (4-30) Given ϵ_{ij} and ϵ_{ij} in Remark 2, we define δ_{ij} :

$$\boldsymbol{\delta}_i = \begin{cases} \frac{1}{c_{i,j}} & \text{if } A_{i,j}(p_i) \\ \\ \frac{1}{c_{i,j}}, & \text{if } A_{i,j} \sim 1 \end{cases}$$

Asymptotic (A θ): The successive principal minors of the $A \times B$ test matrix B are all positive where

$$D_{m} = \left\{ \begin{array}{cc} \left(\delta_{0} \left(\mathcal{A}_{m} \right) + 1 \right) = T_{0} : & (-1) \\ \left(D_{m} \right) & (-1) \end{array} \right.$$

Similarly as in Theorem 1, we now obtain the tollowing result.

Corollet 1. The equalibrium $\rho \sim 0$ or (Σ) is asymptotically stable if (A-1), (A-3), and (A-4) are satisfied.

Proof:= We choose a Lyapunov function to (12)

a'

$$v\left(|\rho(X)\right) = \sum_{i=1}^{N} a_i |\rho_i(X)|, \quad \text{with } a_i > 0$$

Lben

$$\begin{split} &\nabla v_{1k}(p(k)) \\ &= v(k+1) - v(k) \\ &= \sum_{i=1}^{N} \lambda_i \{ (p_i(k+1)) - [p_i(k)]_i \} \\ &\leq \sum_{i=1}^{N} \lambda_i \{ (p_i(k+1)) - [p_i(k)]_i + \sum_{i=1}^{N} T_{i,i}(G_i)(p_i(k)) \} \\ &\leq \sum_{i=1}^{N} \lambda_i \{ [(\sigma_i(i))] - [(\frac{p_i(k)}{G_i(p_i(k))}]^2 G_i(p_i(k)) \} \\ &+ \sum_{i=1}^{N} \left[T_i J [G_i(p_i(k))] \right] \\ &\leq \sum_{i=1}^{N} \lambda_i \left\{ \left[\delta_i ((A_{ij} - 1)) - I_{ij} \right] [G_i(p_i(k))] \right\} \\ &\leq \sum_{i=1}^{N} \lambda_i \left\{ \left[\delta_i ((A_{ij} - 1)) - I_{ij} \right] [G_i(p_i(k))] \right\} \\ &= -\lambda^2 D_{ij} \\ &\leq 0. \text{ No proper shores of } \lambda \end{split}$$

where $\lambda = (O_1, \cdots, n_N)^T$ and $w = ((O_1(p_1), \cdots, O_N(p_N))^T)$. Therefore, the equilibrium is asymptotically stable.

4.3. Estimates of Trajectory Bounds

The previous section presented methods for determining the stability properties of the various equilibria of (2). Usually we also desire information about network performance. One critical performance issue concerns the network's rare of convergence from an initial condition to the limit state. The present section develops trajectory found examines that allow the designer to predict the rate of convergence near the equilibria of the network.

In the case where (A-2) is satisfied, D is an M-matrix Properties of M-matrices are discussed in [13]. The particular property of M-matrices that will be used in the sequel is given at (A-5)

Astrophysic (A-9). For the matrix $D = \{D_{ij}\}$ (as defined in (A-20), there exist existents $A_i > 0, i + 1, \cdots, N$, such

$$D_{\alpha} + \sum_{i=1}^{N} \frac{\delta_i}{\delta_i} D_{\alpha} \otimes \epsilon \otimes 0, \qquad i = 1, \cdots, N,$$

The condition for D to be an M-matrix too satisfy IA 20 only requires (A-5) to be satisfied with $\epsilon = 0$. Thus IA-51 is dightly more strongent than (A-2). Using (A-5), we can prove the bollowing.

Theorem 2. If (A-5) is surjefied, there

$$u(k) - h^{\alpha} \Gamma = p(k) \cdot \sin^{\alpha} \left[p(\theta) \right]$$

for all $\|p\| < r + \min\{0, r\}$, where $\|p\| + \sum_{i=1}^{N} \|a_i\| p_i\|_{L^{\infty}}$ is defined in Remark 2, and $|a_i|$ are given by (A-5).

$$\begin{aligned} & Prixf = Choose \ r(p(k)) = \sum_{l=1}^{N} k_l \ p_l^{N} \ Then \\ & \nabla_{\mathcal{L}} r(p(k)) \\ & = r(|p(k+1)| + r(|p(k)|) \\ & = \sum_{l=1}^{N} \lambda_l \Big[|\rho_l(k+1)| - |\rho_l(k)| \Big] \\ & = \sum_{l=1}^{N} \lambda_l \Big[|A_n p_l(k)| + \sum_{l=1}^{N} I_l G(p_l(k)) \Big] - \Big[p_l(k) \Big] \\ & = \sum_{l=1}^{N} \lambda_l \Big[|A_n p_l(k)| + \Big[|\rho_l(k)|_1 + \sum_{l=1}^{N} |I_l \operatorname{det}_2] |p_l(k) \Big] \\ & = -\sum_{l=1}^{N} \lambda_l \Big[|A_n p_l(k)| + \sum_{l=1}^{N} |D_{l_l} p_l(k)| \Big] \\ & = -\sum_{l=1}^{N} \sum_{l=1}^{N} |\lambda_l D_{l_l} |p_l(k)| \\ & = -\sum_{l=1}^{N} |\lambda_l \Big[\sum_{l=1}^{N} \frac{\lambda_l}{\lambda_l} D_{l_l} \Big] |p_l(k)| \\ & \leq -\epsilon \sum_{l=1}^{N} |A_l p_l(k)| = -\epsilon r r r r \left(|p_l(k)| \right) \end{aligned}$$

40.

$$\nabla_{s,t} r(p(k)) \circ - cr(p(k)) \tag{4.1}$$

At this point we invoke the comparison principle (see [(4]), Inequality (4.1) gives rise to the comparison equation

$$\tau_{A+1} = t_2 + - \epsilon \tau_2, \qquad (4.2)$$

The solution of (4.2) is

It now follows from the comparison protectly that

$$(1/p(k)) < \mu'((p(0))$$

where $\mu = (1 + \epsilon)$. From 1A(5), $0 \le \epsilon \le 1$, which intplies $1 \ge a \ge 0$

By use of Theorem 2, we are able to calculate an exponential bound for the tate of convergence from an initial state within the domain of attraction of all.

In the translation (T) and the subsequent analysis, we have assumed that the input I is constant. The model can be generalized by the addition of a time variant input I(k). This model is described by

$$g(k) = IV(k) + Au(k) + I - \tilde{I}(k)$$
 (DT)

With the following assumption. Theorem 2 can be extended to apply to (DT) as is presented in Corollary 2.

Assume that for (DT)

$$\sum_{k=0}^{N} k_k \tilde{I}_k(k) 0 < M, \qquad \text{for all } k \neq 0.$$

for some M > 0.

Corollary, 2: BRAA-5) and CA-6) are title, then

$$\|\varphi(\lambda)\| \leq \left|\alpha - \frac{M}{\epsilon}\right| \mu^{-1} + \frac{M}{\epsilon}$$

provided that $\alpha > M/\epsilon$ and $1000 \le \alpha$.

The proof of this corollary follows the same approach as the proof of Theorem 2.

V. Nexturiors Reserve

As mentioned in the introduction acuta, networks have been shown to effective a implement AMs. This section presents a technique that invokes the stability conditions of the previous sections to writhesize an AM using COO.

5.1 Equiphenan Constraints

For the purpose of designing AM's, we will assume that a set of library concern is given:

$$Y = \{Y^1, Y^2, \dots, Y^n\}, \quad A^n \in \mathbb{R}^N$$

Each element of A is a partern that is to be stored in the AM. In terms of (DT), we will proceed by designing the network to have each element of 15 as an asymptotically stable equilibrium. In this approach leach stored pattern

12 will have an associated contain of autoction (0):

$$(Y + \{V_i \in R^n | V(k_i, 0, V_i)) \rightarrow V) \text{ as } k = *\}.$$

In order too each $V\cap A$ to be an equilibrium of (Δ) . ICCO must be satisfied

$$u^{r} = IV^{rr} + Au^{r} + I,$$
 $(-1, \cdots, r) = (C1)^{r}$

where $V = (V_1^*, \cdots, V_n^*)^T$, $u' = (u_1^*, \cdots, u_N^*)^T$, and $V_1 = G_1(u_1^*)$. Equation (CT) is equivalent to U(1).

Fo compactly present the design algorithm, we define two matrices $\boldsymbol{\ell}$ and \boldsymbol{H} . Let

$$\mathbf{L} = [\mathbf{U} \cap \mathbf{V}^{T}, \cdots, \mathbf{V}^{T}] \tag{5.1}$$

$$H = [u^{\dagger}(u^{\dagger}, \dots ; u^{\dagger})],$$
 (5.2)

where V^* and a^* are the corresponding tth output and input library vectors, respectively. In terms of the matrices defined in (5.1) and (5.2), (CT) can be expressed as

$$H = II$$
, $-AH - I$

where $I = \{I(I) \cdots I\}$ is the $N \times r$ matrix with vector I as each of its columns. Letting E represent the $N \times N$ identity matrix, the above equation is equivalent to

$$0 = Tl_{+} + (A - E)H + I$$

Qur goal is to specify T, A, and I so that the 'inear constraint (CT) is satisfied, where I, and H are given. With this goal in which, we let

$$\begin{aligned} & R_j = \left[|T^f, B_j^f| |Q| \right] \\ & W = \left[|T_{1j}, T_{2j}| | \cdots, T_{N_j}, A_{j-1}, I_j \right] \end{aligned}$$

where H_j is the jth low of H_j and $Q \in R^*$ is a column vector containing all ones. Solving (CT) is equivalent to solving

$$H_i^{F_i} \circ \mathbf{R}^i H^{F_i}$$
, for $j \in \{1, \dots, N_i\}$ (C1)

However, (C.1) has a well known solution (making use of pseudo inverse) after by

$$W^{I} = PH^{I} \tag{5.5}$$

where $P_i = R^2(R,R^3)$: This solution is not, in general, an que. A full discussion of selections of CT1 can be found in [15].

5.2. Assemptione: Stability Constraints

Constraint (CO) allows as as design a network with each element of A as an equilibrium. However, that constraint alone rises not guarantee that each element of A will be asymptotically stable.

Theorem I of the previous section presents a criterion that may be used to guarantee asymptotic stability. Assumption (A-7) below presents a condition equivalent to (A-2). The form of (A-7) is more easily testable than the form of (A-2).

As equipment (4.7): For the matrix, $D = [D_{ij}]$ defined in (A.2), assume that there exist constants $A_i \geq 0$. The $1, \dots, N_i$ such that

$$\sum_{i=1}^{N} J_i |D_{ij} > S_i \qquad \text{for } i = 1, \dots, N$$

Assumption (A-7) is called a row-dominance condition. (see [13]). Explicitly, (A-7) requires

$$k_i(1 - |A_i|i) = \sum_{j=1}^{N} k_j T_{ij} k_{jj} > 0$$
 (C2)

for $r = 1, \dots, N$

For the purpose of this design, we will arbitrarily set A = 1, for $r = 1, \cdots, N$. If no solutions exist for this choice of A_i , then other values may be assumed. The optimal choice of A_i has not been existed the because good results have been attained by the present choice.

Due to the signoidal shape of the function μ_i and the fact that the desired patterns are estably comprised of binary elements, the upper sector bound κ usually quite small. The following assumption states this explicitly.

As Sumption (A-8). For $j=1,\cdots,N$, assume that $0 < v_{j2} < 1$.

In the unlikely case when (A-8) is not found to be true, the following analysis can be repeated using the appropriate bound for c_{ij} in place of the value 1.

With the choice of $a_i = 1$ for all $j = 1, \cdots, N$ and (A-8), inequality (C25 can be simplified as follows

$$\begin{split} \lambda_i (\mathbf{I} + [A_n]) &= \sum_{i=1}^N \lambda_i |T_{ij}| c_{ij}, \\ &= 1 - [A_n] + \sum_{j=1}^N (T_n) c_{j2}, \\ &\approx 1 - [A_n] + \sum_{i=1}^N |T_{ij}|, \end{split}$$

This analysis shows that if enequality (C2),

$$F_i = 1 - 1A_{ij} = \sum_{j=1}^{N} T_{ij} > 0, \quad i = 1, \dots, N \quad (C2)$$

is satisfied, then (C2') will be satisfied. Since (C2') is equivalent to (A-7), sotisfaction of constraint (C2) will guarantee asymptotic stability.

5.3. The Synthesis Technique

This section presents a method for designing the desired network. We will consider two cases. In the first case, the gain x of the function y will be considered as variable. The second case considers λ to be constain. In either case the technique will seek to satisfy both constraints (C1) and (C2) for $y = 1, \dots, N$.

Remark 3: Subsequently we will assume that the non-linear function μ_{ij} of the individual neutrons are all the same. Thus we drop the index from g and A_{ij} .

Due to the form of constraints ICD and (C2) Ooth constraints are a function of W_i , which contains one row of T_i , A_i and T_i), the synthesis procedure enables us to design the network one row at a time resulting in a great simplification. For example, the matrix inversion required an (C1) involves an $(r \times r)$ matrix. To solve the entire set

of constraints at once would require the unversion of an $(rN \times rN)$ matrix. The former inversion is much simpler than the latter.

Case I: In this case, the gain of the function g is a variable parameter. To show this explicitly, we will write

$$V = g(\lambda u_i)$$
.

As usual, we will consider the set X to be given. Due to the properties of g discussed in Remark 1, g is invertible Hence, a simple calculation

$$u_i = \frac{1}{2} g \cdot \left(\tilde{\nu}_i \right)$$

will yield the H matrix for any L neatrix. To terther simplify, we set $A_n = 0$, $i = 1, \cdots, N$, in the following procedure. It solutions exist in this special case, then solutions will also exist without the constraint, since the latter case is a generalization of the former

For the present case, the designer chaoses an initial gain λ_n . Using this gain the designer can, in term, find H^n , H^n , and H^n . The initial solution to (5.3) for this choice of λ_n will be denoted H^n , for $i = 1, \dots, N$. In each case, $F_i(H^n)$ is easily calculated. Define

$$s = \min_{i \in \{0,\infty\}} \left(F_i \big(\mathcal{H}_i^{\mathrm{eff}} \big) \right).$$

If $\epsilon > 0$, the constraints (C1) and (C3) are both satisfied. The desired values of T, M, and T are given by the W_i^0 , $i = 1, \cdots, N$

 $|H| = \epsilon > 0$, choose a constant $\delta > 1$, $|\epsilon > 1$, and define

$$egin{aligned} & b = \delta b \ L = L_0 \ & H = H_0 \neq \delta \ . \end{aligned}$$

Ethe to the assumption that A=0, the H_1^0 column may be conceved from R_1 . In this case, neither P_1 nor R_2 are affected by the change in A. Hence, (5.3) shows that decreasing H by the factor δ will decrease each W_1 by a factor of δ . Specifically,

$$\begin{aligned} W_i &= \left[T_{i1}, \cdots, T_{iN}, I_{iJ}\right] \\ &= \left[T_{i1}^n / \delta_i, \cdots, T_{iN}^n / \delta_i, I_i^n / \delta_i\right] = W_i^n / \delta. \end{aligned}$$

The following shows that (C2) is now satisfied: for $i \in \mathbb{N} \to N$,

$$\begin{split} F_i &= 1 - |A_{ij}^0 - \sum_{i=1}^N |T_{ij}^0| = 1 - \sum_{i=1}^N |T_{ij}^0| > \epsilon \\ &\sum_{j=1}^N |T_{ij}^0| < 1 - \epsilon, \epsilon < 0 \\ &\sum_{j=1}^N |T_{ij}| = \sum_{j=1}^N |T_{ij}^0| / \delta < \frac{1 - \epsilon}{\delta} < 1 \\ &F_i = 1 - \sum_{j=1}^N |T_{ij}| > 1 - \left(\frac{1 - \epsilon}{\delta}\right) > 0. \end{split}$$

The above steps are casily translated into a computer algorithm. This algorithm will present the desired F, A,

and I macrices under the conditions that $A \sim 90$ fields and that the assumptions required to solve ICD are satisfied.

In the several specific examples that we considered using this method, the test/is have been very sociessful. Two of mese examples are presented in the heat section of the present paper.

Case 2: In this case, we consider both matrices L and H to be given; thus A is fixed. As was mentioned at Section II. (DT) is meant to be implemented as a computer algorithm. As such, the designer should be able to choose the gain A, and Case 1 should usually apply However for the take of completeness we will present the design algorithm for Case 2. The design approach for Case 1 to very straightforward, while the algorithm for Case 2 may require iteration.

Our goal is to present an algorithm that tests for solutions W, solicting O(1) and O(2). The algorithm presented below solves for T, A, and T, one rew at 5 time to complete a design the algorithm must be repeated. V times torsection each solv!

An initial solution to (CI) is given by (5.3) as

$$0.5 - P.H/$$

where $P = R_i^T(R_iR_i^T)$ is and $W_{in} = (F_{in}, T_{i+1})$. $T_{\infty} = A_{in}(F_i)$ for this initial solution, $F_i(W_i)$ is calculated. (1) If $F_i(W_i)$ is positive, then a solution has been found for the current row. In this case, the solution can either be further optimized by use of the remaining degrees of freedom or it can be accepted. In the latter case the algorithm will proceed to some for the next row, $W_{i-1,0}^T$ until all N cows have been determined.

29 If $F(W_m)$ is negative, then we apply the following algorithm. The gradient of $F(W_n)$ is

$$\nabla F_{i}(B_{i}) = \langle -\operatorname{sgn}(T_{i}), \cdots, -\operatorname{sgn}(T_{i}), -\operatorname{sgn}(A_{i}), 9 \rangle,$$

The direction of the gradient for a given value of B_i will be denoted by A. The vector A^i is the direction of maximum because of $F(B_i)$.

Since H₂, tails to satisfy $\{C2\}$, but does satisfy $\{C1\}$, the algorithm will produce a new vector W_1 , by making from W_2 , and direction that both increases F_1 and forces W_2 , to satisfy $\{C1\}$. The required direction is given by

$$a \sim P^* d$$

where $P^+ = (E - R_i^T (R_i R^T)^{-1} R_i)$ and E is the N + 2 dimensional identity matrix. The matrix P^+ projects any vector, $S \in R_i^{N+2}$ into the null space of R_i (the null space will be denoted by N_N). Since $d \in N_n$.

$$R_i(\mathbf{H}_m + \alpha d) + R_i W_m + \alpha R_i d + R_i W_m + H_i$$
.

Thus by defining $A_1(\alpha) \in W_{\alpha}(1, \alpha)$, A_1 , will satisfy (C1). Due to the fact that F is linear when restricted to any one of the $2^{\alpha+2}$ α -tents of $R^{\alpha+2}$, if the angle between d and d is less than $\theta(0)$, then d will be a direction from W_{α} that increases F_{α} and angle is less than $\theta(0)$ by the projection theorem. Since the gradient is constant until

one of the components of W(p, a) becomes zero, F(M, p, a) is the smallest positive matner that sets one of the components of W(p, a) to zero. When the oth component of W(p, a) is set equal to zero, the corresponding column of R, should be removed in this manner, $W^{(a)}$ will remain zero for all subsequent iterations, and at each iteration the dimension of A_R is reduced by one

The above algorithm can be used iteratively. At ℓ ich step $(\ell, 1)$ as satisfied and J_1 is arcressed. The algorithm is guarantees to sorp after no more than $(N + 2 + \ell)$ denotions: since N_{N_1} begans with all more long $(N + 2 + \ell)$ and the dimension decreases by one at each iteration.

At the algorithm site minimation, either (CO) will be satisfied and the perwork will have stored each element of L as an asymptotically stable equilibrium, or the algorithm will have failed to design the required network. In the latter case, the choice of $A_{ij}(z+1) \leq \Delta$ (from CO)) may be altered or new patterns may be chosen.

VI SIMULATIONS

In the current wation, we present two simulation exomples using the synthesis method which is given in Section V. The purpose of the first example is to esomine the average performance of 1017 based on a total of 130 13-region networks that were designed. The second example demonstrates an application of (DC) to pattern recognition. A neural network was designed to recognize 27 key patterns tibe 28 appear case letters and a blank!, each pattern being represented by a 9 × 9 a selarities.

Remark 4. In both of the following examples, the designation. Thirary? will mean that the only possible values that a variable may assume is 1-1 on 1.

Example 3: A 13 neuron system has \$192 (in [3,7) prosible briary inputs. For each design each of the possitite birmry inputs will, in turn, be applied as an in-tail condition for the system. The network will be adopted to godye from its initial condition to a final state. We will autorprot this final state as the network's briary response in the given matrix expedition. The Hamming distance, from the inmal condition to the fand butary response, will be used as a passy for estimating basins of attraction of the stored patterns. Specifically, we would like to detail mine how the network's performance will be affected by the number of patterns (all to be speed) for each value of r_i between 1 and 1.9 - 8000 made were made, Each Calconsisted of choosing a set of a Lipporte independent, but otherwise random, binary pasterny of length NU 130. For each set of a potterns, a perwork was designed and simulated. The results of the ten totals for each cake of a were then averaged. The results are presented graphically in Engs 1-9

In Figs. 5-8 the number of patterns to be stored in the judependent variable. The dependent variable is the everage momber of binary inputs converging to a desired pattern from the Standard distance (14D) noted in the caption of each figure.

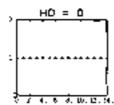


Fig. 1.—As a flatetion of the number of patterns stored thoughough, the actuage stunder of timory input patterns converging from Hamming distance. HD territoral to a desired subject.

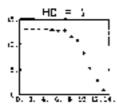


Fig. 2. As a lanction of the number of patterns stored than isostal) the increase number of binary input patterns converging from Hamming distance. HD (certical), to a desired curput.

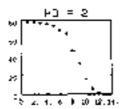


Fig. 3.—As a function of the tuniber of patients scoted thin cortial CPs: average natibe, of binary input patients universing from Hamming Cistance, HD (vertically to a desired output.)

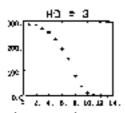


Fig. 6.—Avia function of the number of patterns storrel (bordennish, the average number or binary input patterns converging from Hamming distance. 4ID (vertical), to a desired mappy

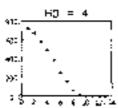


Fig. 7. As a function of the nursber of gatterns stored Promotoral), the average number of binary diput gatterns converging from Hammina distance. FED Gentical) to a desired output.

Fig. 1 shows that each stored pattern attracts the binary input that is at a **Hernning** distance of zero. In addition to verifying that each desired pattern is in fact stored, this verifies that each stored pattern does have a nonempty basin of attraction.

Fig. 2 shows that for z = 1 to 4.05 z (0.3 M), in each of the ten irrals, cack of the stored patterns attracted a liot

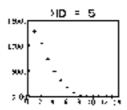


Fig. 6.—As a function of the number of paperts stored beautomate the assume number of broady input patients consenging how. Hastining costumer, HD test cuts was drapped coupur.

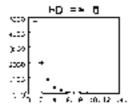


Fig. 3. As a function of the number of patterns stoned thericancer the ascence models of brokes input patterns converging from Hamming Ostance, HD Gert call Co., desired output.

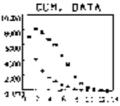


Fig. 4.— As a function of the perform Authority stored Heavenfully the assumption of binary deput patterns converging from Hazming a status. Hit term of this adjusted output.

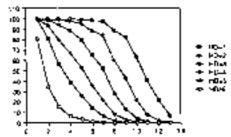
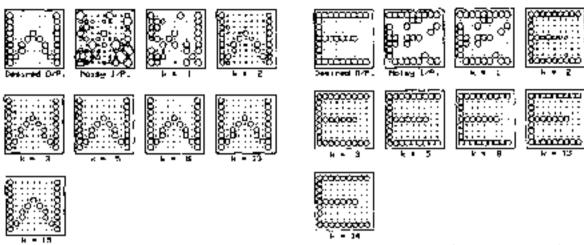


Fig. 9.— The completed property (1999) of the state from Figs. 1–2. Farhold options, who is noticed for in, the source property and all the imprisy at Hammong distance HIII from a one two pattern that converge to the Desired partier.

the most patterns that only differed by a single bit. As a mercused further, the arraction of patterns from a Hamming distance of one remained ovar the maximum of 33 up if $\gamma=8$

Each of the plots is expected to decrease monotonically as a approaches. A Task is because the number of possible binary inputs 18192) is constant. Hence, on the average, each stored pattern should attract 819277 of the possible inputs. Fig. 8 plots the average total number of inputs attracted per output pattern (indicated by "\="") and the average total number of the input patterns findicated by "\="") attracted to all or the 7 stored vectors. Interestingly, the lower plot of Fig. 8 does appear to decrease approximately as 177. The upper plot remains



(9) It is I for evolution at a found retwink starting 27 p. fleres. The network monogeness the latter "W" from an initial and two that has belief former an even with servencian (see 150 standard devolution to the "W" pattern.

Fig. 17. Time evolution of a normal network incommon the father than a formal independent for the Hamming distance of 43 from the above comparison for the "Fig.".

above 8192/2 until r = 6. At this point, the extraneous equalibria are arrancing almost half of the input patients. The usefulness of networks with r near X is application-dependent. For example, in a secure voice recognition application, it is desirable for the majority of inputs to be attracted by extraneous equilibria. It should be nested that for each of the 19 mals with r = 2, all 8192 monts were attracted to the two stored patterns.

Finally, Fig. 9 shows the percentage of imput patterns that are at the Harming distance of IID from a stored pattern, which are attracted to the steroid pattern for III = 1.5 and IID =0. As is expected, the percentage of patterns converging from large Harming distances drops off faster than the percentage from a smaller Enabring distance. The form of this graph is similar to the "waterfall" graphs common in cashing theory and agriculture assign. Waterfall graphs are used to display the degradation of the system perioditative as the input race increases. Using this type of interpretation, Fig. 9 displays that the ability of the network is handle small agrid to make ratios (large Standing distances) decreases as the number of patterns stored (r) increases.

Arampho 2: In this example, we design an SI-neuron network. The given set of library vectors corresponds to the 26 appear case letters and a blank pattern $(a_0, N = 81)$ and x = 27. The corpose of the perwork x to recognize the key mattern truin a reosy unput pattern. The two different types of noise that will concern us in the semulations are the Gaussian noise and the binary noise x kind of noise that will randomly flip binary buts. For the termer, the noisy input pattern is generated by adding zero mean Gaussian noise with a given specific standard deviation to the key pattern (Fig. 10). For the latter, the noisy input pattern is generated by readomly flipping a certain noise of the hintern bits of the key pattern to produce an input pattern at the desired Christian deviation the desired nutring distance from the desired nutring overlap (Fig. 11).

. **3**,011-1 Zero регар Сероран. Shary Noise 18,830 'yps AMP IN Appropriate Common/AUI Harring Datasets (F.) 7managi 165 42/ T 00 13 Conserer to -44 45.7 E.. 33.1 454 95.5 Converge to 17 1" ! ٠, 148 uca 10.7 Spanoučisau ů. Not converged Date 3.0 : 4 ΕU

³ Beorge the system dall not converge with a 1-0 four units (1.2) game ground to SNR of 7 dB and 1 volume 10 mg/. [MS] 0, respectively.

The typical time goodarion of the neural network is shown in Figs. 10 and 11 for the Grassion pose and the bursty mass, respectively. For each figure, the plot in the left appear contains the desired catput pattern and the next camera to the right is the masy input pattern. The rest of the plots are the time evolution results like caption on the bottom of each plot indicates the current value of the time ray each of the successive plots). The final result (i.e., the pattern to which the requestion of capture of the successive plots), the final result (i.e., the pattern to which the requestions of the last row, which is furthes, to the right.

Remark 5. Figs. 10 and 11 ase a small dot and a longe spin's to represent the catput of netrops with values of 1 and 1 respectively. Output values between these two-corones are sepresented by cacles whose size is scaled between the large circle and the small dot. This is the continuous nature of G(c) (discussed in Remark 1), we see that the final netpat state can never assume the desired values of 1 and 1 (instead, we recognize values sufficiently near the binary values as the binary values.

We lined made six different sets of smioletions, three for Gaussian noise with different standard deviations and three for binary noise with different Hamming distances.

TABLE III

1 7114	Zero mean Georgia Turk			Bran Nove		
i Turk Mar	Planning Personn			-Distances		
Table /	. 0.7]	9,45	1.00	٠.	10	13
Aserage table to con-ergo	154	:02	241	157	153	27.7

that including the renominational cases.

For each case, we test all of the key patterns twice free. $2 \times 27 = 54$ trials). The results are summarise a in Tarkes I. and H

From Table I we line that the performance of the system decreases as expected, as the name level increases. However, about 950 of the mosy input patterns. with a 2-dB SNR (or a Barmming distance of is away from the key pattern) can be successfully recognized. Table II. shows that the convergence time of the network ipageases. as the input tooks increases. This is expected because the mercased noise moves the input pattern faither away. from the equilibrium positi Chemory C

VII. CONCLUSIONS

In this paper we have presented a detailed stability analysis of and design algorithm for a class of neural hetworks. We offize the standity sanalysis to derive the techniques for synthesizing a discrete time neural metwork. Compared to the already eyesing reconiques for designing neural networks, the algorithm of the present paper is very easy to implement and yields guaranteed stability properties. To demonstrate the applicability of our results, we have presented two specific examples Both examples are simulated on a digital computer using a Fortrap program requiring less than 100 lines of code

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