

Memory

In drawings of the ULM the content of memory cells is represented in hexadecimal (skipping the '0x' prefix).

Let $A \in \{0, \dots, 2^{64} - 1\}$ be an unsigned integer:

- $M_1(A)$ denotes the 8-bit pattern in memory cell with address A
- $M_4(A)$ denotes the 32-bit pattern of four consecutive memory cells where the first memory cell has address A .

Example: $M_1(5) = 0x21$, $M_4(4) = 0x20210000$

10	10	00	20	20	21	00	00	14
0x00			0x04				0x08	
(0)			(4)				(8)	

General Purpose Registers

This ULM variant has 16 registers denoted as %0x0, %0x1, ..., %0xF:

- Each register has a width of 64 bits.
- Each register can be addressed with a 4-bit pattern (i.e. one hex digit).
- %0x0 is special. It always contains a bit pattern with only zeros. Writing to it has no effect.

Other Registers

- The instruction pointer (%IP) is 64 bits wide.
- The instruction register (%IR) is 32 bits wide.
- Each of the status flags ZF, CF, OF, SF can store a single bit.

Notation

Let X be a bit pattern:

- $u(X)$ denotes the represented unsigned integer value
- $s(X)$ denotes the represented signed integer value

Let X and Y be a two bit pattern:

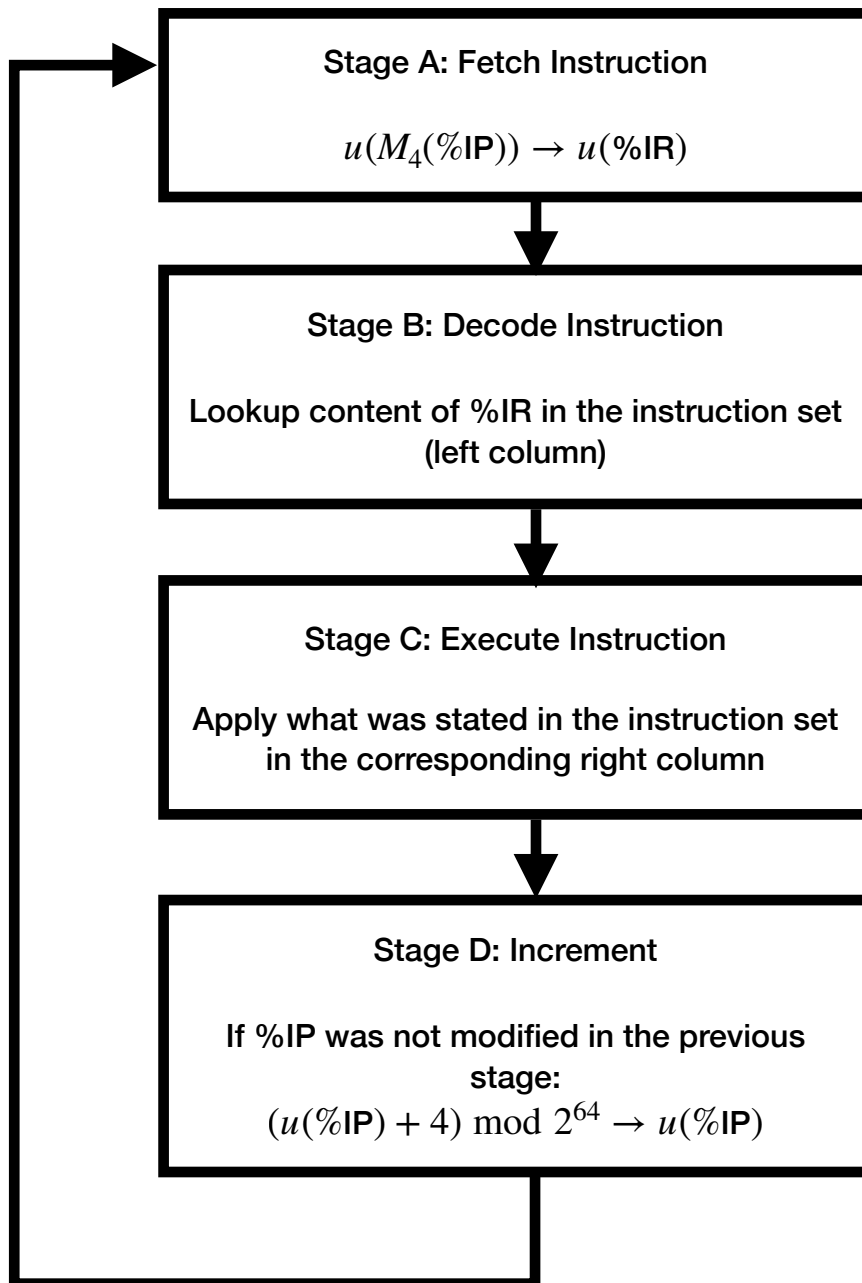
- $u(X) \rightarrow u(Y)$ denotes that Y gets modified such that $u(X) = u(Y)$. This is also called a "Zero extension of X ".
- $s(X) \rightarrow s(Y)$ denotes that Y gets modified such that $s(X) = s(Y)$. This is also called a "Signed extension of X ".

Example:

- $u(0x8F) \rightarrow u(\%0x1)$ writes $0x000000000000008F$ to $\%0x1$
- $s(0x8F) \rightarrow s(\%0x1)$ writes $0xFFFFFFFFFFFF8F$ to $\%0x1$

Von Neumann Cycle

After a program was loaded the cycle described below is applied until in stage C an halt instruction was executed. For educational purposes: After stage D the counter for the instruction cycle gets incremented.



Instruction Set (for decoding and executing instructions)

Content of % IR (instruction register)	Effect in execution stage
<div> <div>32</div> <div>24</div> <div>16</div> <div>0</div> <div>0x01</div> <div>imm</div> <div></div> </div>	Halt program execution with exit code $u(imm)$
<div> <div>32</div> <div>24</div> <div>0</div> <div>0x04</div> <div>offset</div> </div>	if ZF = 1 then $(u(\%IP) + 4 \cdot s(offset)) \bmod 2^{64} \rightarrow u(\%IP)$
<div> <div>32</div> <div>24</div> <div>0</div> <div>0x05</div> <div>offset</div> </div>	$(u(\%IP) + 4 \cdot s(offset)) \bmod 2^{64} \rightarrow u(\%IP)$
<div> <div>32</div> <div>24</div> <div>20</div> <div>0</div> <div>0x10</div> <div>dest</div> <div>imm</div> </div>	$u(imm) \bmod 2^{64} \rightarrow u(\%dest)$
<div> <div>32</div> <div>24</div> <div>20</div> <div>16</div> <div>0</div> <div>0x12</div> <div>z</div> <div>y</div> <div>imm</div> </div>	$(u(\%y) + u(imm)) \bmod 2^{64} \rightarrow u(\%z)$ Update status flags: <div> <div>Flag</div> <div>Condition</div> <div>ZF</div> <div>$u(\%y) + u(\%x) = 0$</div> <div>CF</div> <div>$u(\%y) + u(\%x) \geq 2^{64}$</div> <div>OF</div> <div>$s(\%y) + s(\%x) \notin \{-2^{63}, \dots, 2^{63} - 1\}$</div> <div>SF</div> <div>$s(\%y) + s(\%x) < 0$</div> </div>
<div> <div>32</div> <div>24</div> <div>20</div> <div>16</div> <div>0</div> <div>0x14</div> <div>z</div> <div>y</div> <div>imm</div> </div>	$(u(\%y) - u(imm)) \bmod 2^{64} \rightarrow u(\%z)$ Update status flags: <div> <div>Flag</div> <div>Condition</div> <div>ZF</div> <div>$u(\%y) - u(imm) = 0$</div> <div>CF</div> <div>$u(\%y) - u(imm) < 0$</div> <div>OF</div> <div>$s(\%y) - s(imm) \notin \{-2^{63}, \dots, 2^{63} - 1\}$</div> <div>SF</div> <div>$s(\%y) - s(imm) < 0$</div> </div>
<div> <div>32</div> <div>24</div> <div>20</div> <div>16</div> <div>0</div> <div>0x20</div> <div>data</div> <div>addr</div> <div>offset</div> </div>	$u(M_1(A)) \rightarrow u(\%data)$ where $A = (u(\%addr) + s(offset)) \bmod 2^{64}$
<div> <div>32</div> <div>24</div> <div>20</div> <div>0</div> <div>0x30</div> <div>x</div> <div></div> </div>	Print the character with ASCII code $u(\%x) \bmod 2^8$

ASCII Table

Dec	Hex	Char	Dec	Hex	Char	Dec	Hex	Char	Dec	Hex	Char
0	00	Null	32	20	Space	64	40	@	96	60	`
1	01	Start of heading	33	21	!	65	41	A	97	61	a
2	02	Start of text	34	22	"	66	42	B	98	62	b
3	03	End of text	35	23	#	67	43	C	99	63	c
4	04	End of transmit	36	24	\$	68	44	D	100	64	d
5	05	Enquiry	37	25	%	69	45	E	101	65	e
6	06	Acknowledge	38	26	&	70	46	F	102	66	f
7	07	Audible bell	39	27	'	71	47	G	103	67	g
8	08	Backspace	40	28	(72	48	H	104	68	h
9	09	Horizontal tab	41	29)	73	49	I	105	69	i
10	0A	Line feed	42	2A	*	74	4A	J	106	6A	j
11	0B	Vertical tab	43	2B	+	75	4B	K	107	6B	k
12	0C	Form feed	44	2C	,	76	4C	L	108	6C	l
13	0D	Carriage return	45	2D	-	77	4D	M	109	6D	m
14	0E	Shift out	46	2E	.	78	4E	N	110	6E	n
15	0F	Shift in	47	2F	/	79	4F	O	111	6F	o
16	10	Data link escape	48	30	0	80	50	P	112	70	p
17	11	Device control 1	49	31	1	81	51	Q	113	71	q
18	12	Device control 2	50	32	2	82	52	R	114	72	r
19	13	Device control 3	51	33	3	83	53	S	115	73	s
20	14	Device control 4	52	34	4	84	54	T	116	74	t
21	15	Neg. acknowledge	53	35	5	85	55	U	117	75	u
22	16	Synchronous idle	54	36	6	86	56	V	118	76	v
23	17	End trans. block	55	37	7	87	57	W	119	77	w
24	18	Cancel	56	38	8	88	58	X	120	78	x
25	19	End of medium	57	39	9	89	59	Y	121	79	y
26	1A	Substitution	58	3A	:	90	5A	Z	122	7A	z
27	1B	Escape	59	3B	;	91	5B	[123	7B	{
28	1C	File separator	60	3C	<	92	5C	\	124	7C	
29	1D	Group separator	61	3D	=	93	5D]	125	7D	}
30	1E	Record separator	62	3E	>	94	5E	^	126	7E	~
31	1F	Unit separator	63	3F	?	95	5F	_	127	7F	□