Lab 3 for uC/OS-II: Ceiling Priority Protocol

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Objective

To implement Ceiling Priority Protocol for ucOS's mutex locks

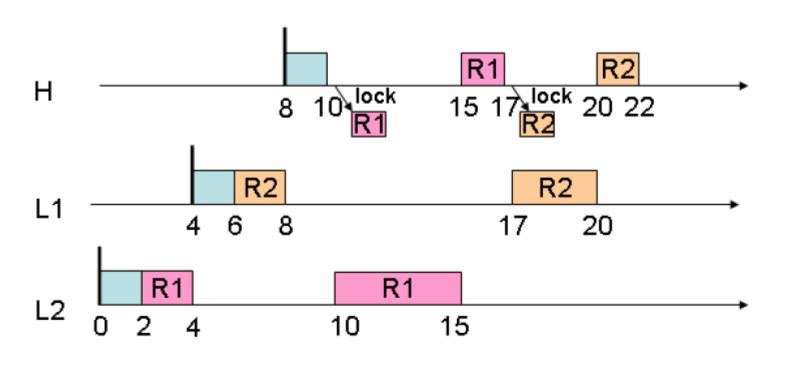
uC/OS Mutex Locks

- A mutex lock is associated with a "priority"
 - Its priority is higher than the highest locker
 - E.g., if T₃ and T₄ share a lock, the priority of the lock should be set to 2
 - When T4 blocks T3, then T4's priority becomes 2

Disadvantages of PIP

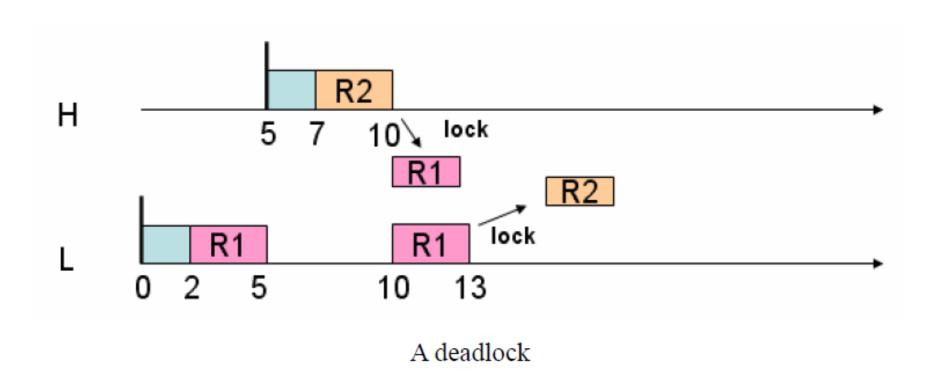
- The "PIP" avoids uncontrolled priority inversion, but it has two disadvantages
 - A high priority task can be blocked multiple times
 - Deadlocks are possible

Scenario 1: Multiple blocking in ucOS2 PIP



Task H is in turn blocked by task L1 and task L2

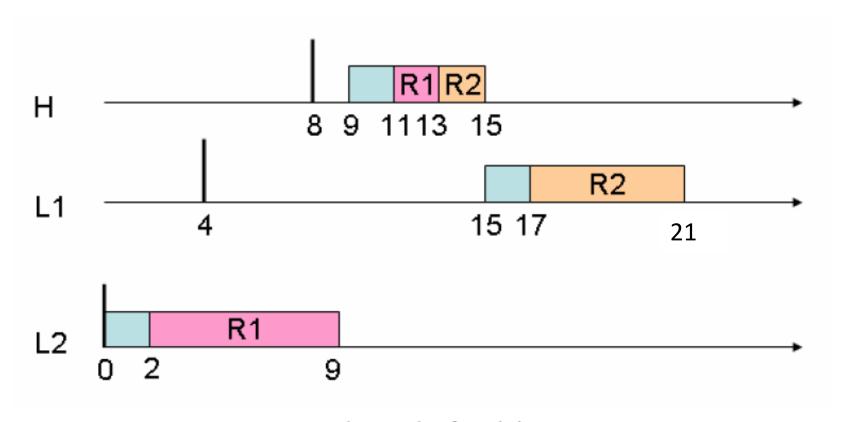
Scenario 2: Deadlock in uCOS-2 PIP



Ceiling Priority Protocol

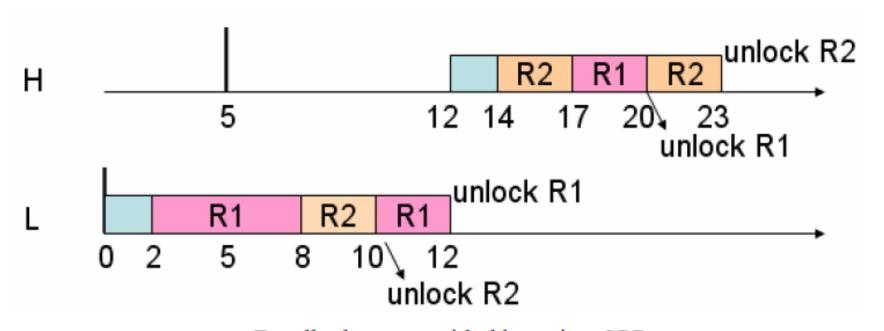
- Highest-Locker Protocol
- When a task acquires a mutex lock, its priority becomes the highest among all lockers' priorities
- In uC/OS, the we use the mutex's priority as the highest-locker's priority

S1 CPP: Removing Multiple Blockings



The result of applying CPP

S2 CPP: Avoiding Deadlocks



Deadlocks are avoided by using CPP

Implementation

- Reuse your code of Lab 1 (do not re-use EDF)
- Modify the following two functions
 - OSMutexPend()
 - If mutex is free, boost the locker's (caller) priority
 - OSMutexPost()
 - Restore the original priority of the locker
- Do not use OSTaskChangePrio()
 - It calls OS_Sched() and results in unexpected behaviors

Implementation

- All tasks should add proper OSTimeDly() at their beginning to emulate their arrival times
- Emulate durations of CPU execution and resource use with your code from Lab 0
 - 2 ticks → lock R1 → 6 ticks → lock R2 → 2 ticks → unlock R2 → 2 ticks → unlock R1

Output

- Similar to those in prior labs, but add lock/unlock events
- Output the results of using CPP for Scenarios 1 and 2

Output Example of S1

Priority initialization:

R1: 1

R2: 2

Task1: 3

Task2: 4

Task3: 5

Task arrival time:

Task1: 8

Task2: 4

Task3: 0

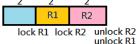
20	lock	R1	(Prio=5	changes	to=1)
90	unlock	R1	(Prio=1	changes	to=5)
90	complete		5	3	
110	lock	R1	(Prio=3	changes	to=1)
130	lock	R2	(Prio=1	changes	to=1)
150	unlock	R2	(Prio=1	changes	to=1)
150	unlock	R1	(Prio=1	changes	to=3)
150	complete		3	4	
170	lock	R2	(Prio=4	changes	to=2)
210	unlock	R2	(Prio=2	changes	to=4)

complete

210

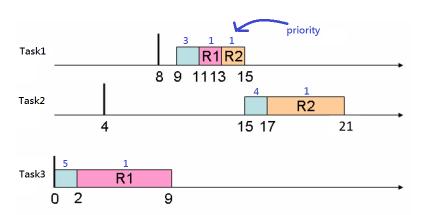
Task execution time and resource used:

Task1:



Task2: R2 unlock R1

Task3: R1 unlock R1



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