

# MQP TITLE

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by

.....  
MICHAEL FICARRA

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.....  
DANIEL DOUGHERTY  
professor, project advisor

## **Abstract**

This paper describes a method, referred to as the chase, for generating minimal models for a geometric theory. A minimal model for a theory is a model for which there exists a homomorphism to any other model that can satisfy the theory. These models are useful in solutions to problems in many practical applications, including firewall configuration examination and access control evaluation. Also described is a Haskell implementation of the chase and its development process and design decisions.

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# 1 Introduction

Introductory text...

## 1.1 Goals

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## 1.2 The Chase

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## 2 Technical Background

### 2.1 Definitions

In this paper, logic symbols and other possibly ambiguous or uncommon notation will be used extensively, and thusly must be clearly defined.

#### 2.1.1 Models

A *model*  $\mathbb{M}$  is a construct that consists of:

- a set, referenced as  $|\mathbb{M}|$ , called the *universe* or *domain* of  $\mathbb{M}$
- a set of pairings of a *predicate* and a non-negative integral arity
- for each predicate  $R$  with arity  $k$ , a relation  $R_k^{\mathbb{M}} \subseteq |\mathbb{M}|$

It is important to distinguish the predicate, which is just a symbol, from the relation that it refers to when paired with its arity. The relation itself is a set of tuples of elements from the universe.

#### 2.1.2 First-order Logic

#### 2.1.3 Geometric Logic

Geometric logic is first-order logic with constraints on the shape of the expression. Geometric logic formulas are implicitly universally quantified first-order logic expressions of the form

$$A_0 \wedge \dots \wedge A_n \rightarrow E_0 \vee \dots \vee E_m$$

where  $A_0 \dots A_n$  are atomics,  $E_0 \dots E_m$  are first-order logic expressions of the form  $\exists_{x_0 \dots x_k} A_0 \wedge \dots \wedge \exists_{x_0 \dots x_p} A_y$ , and  $n, m, k, p$ , and  $y$  are integers greater than or equal to 0.

(Explain why GL is useful)

## 2.2 Homomorphisms

A homomorphism from  $\mathbb{A}$  to  $\mathbb{B}$  is a function  $h : |\mathbb{A}| \rightarrow |\mathbb{B}|$  such that, for each relation symbol  $R$  and tuple  $\langle a_0, \dots, a_n \rangle$  where  $a_k \in |\mathbb{A}|$  for any  $k$  and  $0 \leq k \leq n$ ,  $\langle a_0, \dots, a_n \rangle \in R^{\mathbb{A}}$  implies  $\langle h(a_0), \dots, h(a_n) \rangle \in R^{\mathbb{B}}$ .

A homomorphism  $h$  is also a *strong homomorphism* if, for each relation symbol  $R$  and tuple  $\langle a_0, \dots, a_n \rangle$  where  $a_k \in |\mathbb{A}|$  for any  $k$  and  $0 \leq k \leq n$ ,  $\langle a_0, \dots, a_n \rangle \in R^{\mathbb{A}}$  if and only if  $\langle h(a_0), \dots, h(a_n) \rangle \in R^{\mathbb{B}}$ .

### 2.2.1 Significance

### 2.2.2 Minimal Models

### 2.2.3 Relation to the Chase



### 3 For later reference:

preceq:  $\mathbb{M} \preceq \mathbb{N}$



Figure 1: Caption in report

## A Chase code

```

1  module Chase where
2  import Parser
3  import Helpers
4  import Debug.Trace
5  import Data.List
6
7  chaseVerify :: [Formula] -> [Formula]
8  — verifies that each formula is in positive existential form and performs some
9  — normalization on implied/constant implications
10 chaseVerify formulae =
11     let isNotPEF = not.isPEF in
12     map (\f -> case f of
13         Implication a b ->
14             if isNotPEF a || isNotPEF b then error ("implication must be in positive existential form")
15             else f
16         - ->
17             if isNotPEF f then error ("formula must be in positive existential form: " ++ showFormula f)
18             else (Implication Tautology f)
19     ) formulae
20
21 chase :: [Formula] -> [Model]
22 — runs the chase algorithm on a given theory and returns a list of models that
23 — satisfy it
24 chase formulae = chase' (chaseVerify formulae) ([], [(mkModel [] [])])
25
26 chase' :: [Formula] -> ([Model], [Model]) -> [Model]
27 — used by the chase function to hide the model identity argument
28 chase' formulae (done, []) = done
29 chase' formulae (done, pending) =
30     let self = chase' formulae in
31     let (p:ending) = pending in
32     trace ("running chase on " ++ show (done, pending)) $
33     if all (\f -> holds p (UniversalQuantifier (freeVariables f) f)) formulae then
34         trace ("all formulae in theory hold for model " ++ showModel p) $
35         trace ("moving model into done list") $
36         self (union done [p], ending)
37     else
38         let possiblySatisfiedModels = attemptToSatisfyFirstFailure p formulae in
39         trace ("at least one formula does not hold for model " ++ showModel p) $
40         trace ("unioning " ++ show ending ++ " with " ++ show possiblySatisfiedModels) $
41         self (done, union ending possiblySatisfiedModels)
42
43 attemptToSatisfyFirstFailure :: Model -> [Formula] -> [Model]
44 — checks if each formula holds, sequentially, until one does not, then tries
45 — to satisfy that formula
46 attemptToSatisfyFirstFailure model (f:formulae) =
47     let self = attemptToSatisfyFirstFailure model in
48     if holds model (UniversalQuantifier (freeVariables f) f) then self formulae
49     else attemptToSatisfy model f
50
51 attemptToSatisfy :: Model -> Formula -> [Model]
52 — returns a model that is altered so that the given formula will hold
53 attemptToSatisfy model formula =
54     let f' = UniversalQuantifier (freeVariables formula) formula in
55     trace ("attempting to satisfy " ++ showFormula formula ++ ")") $
56     attemptToSatisfy' model [] f'
57
58 attemptToSatisfy' :: Model -> Environment -> Formula -> [Model]
59 — hides the environment identity in the 'attemptToSatisfy' function arguments

```

```

60 attemptToSatisfy' model env formula =
61   let (domain,relations) = model in
62   let domainSize = length domain in
63   let self = attemptToSatisfy' model in
64   — trace (" attempting to satisfy (" ++ showFormula formula ++ ") with env " ++ show env)
65   case formula of
66     Tautology -> [model]
67     Contradiction -> []
68     Or a b -> union (self env a) (self env b)
69     And a b -> concatMap (\m -> attemptToSatisfy' m env b) (self env a)
70     Implication a b -> if holds' model env a then self env b else []
71     Atomic predicate vars ->
72       let newRelation = mkRelation predicate (length vars) [genNewRelationArgs env vars (fr
73       let newModel = mkModel (mkDomain domainSize) (mergeRelation newRelation relations) in
74       trace (" adding new relation: " ++ show newRelation) $
75       [newModel]
76     ExistentialQuantifier [] f -> self env f
77     ExistentialQuantifier (v:vs) f ->
78       let f' = ExistentialQuantifier vs f in
79       let nextDomainElement = fromIntegral $ (length domain) + 1 in
80       if any (\v' -> holds' model (hashSet env v v') f') domain then
81         trace (" " ++ showFormula formula ++ " already holds") $
82         [model]
83       else
84         trace (" adding new domain element " ++ show nextDomainElement ++ " for variable
85         attemptToSatisfy' (mkDomain nextDomainElement,relations) (hashSet env v nextDomain
86     UniversalQuantifier [] f -> self env f
87     UniversalQuantifier (v:vs) f ->
88       let f' = UniversalQuantifier vs f in
89       concatMap (\v' -> self (hashSet env v v') f') domain
90     - -> error ("formula not in positive existential form: " ++ showFormula formula)
91
92 genNewRelationArgs :: Environment -> [Variable] -> DomainElement -> [DomainElement]
93 — for each Variable in the given list of Variables, retrieves the value
94 — assigned to it in the given environment, or the next domain element if it
95 — does not exist
96 genNewRelationArgs env [] domainSize = []
97 genNewRelationArgs env (v:ars) domainSize =
98   let self = genNewRelationArgs env in
99   case lookup v env of
100     Just v' -> v' : (self ars domainSize)
101     - -> (domainSize+1) : (self ars (domainSize+1))

```

## References

- [1] A Cottrell, *Word Processors: Stupid and Inefficient*,  
[www.ecn.wfu.edu/~cottrell/wp.html](http://www.ecn.wfu.edu/~cottrell/wp.html)