
Characterizing Schur Net’s Expressivity Via Homomorphism Expressivity

**NOTICE: THIS IS THE LATE VERSION AND WAS
SUBMITTED AFTER THE DEADLINE**

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Abstract

Graph spectral invariants and group-theoretic constructions each offer complementary pathways toward building expressive Graph Neural Networks (GNNs) [1]. Recent work on higher-order message passing has shown that subgraphs communicating with subgraphs can capture complex local structures (e.g., cycles) by leveraging equivariance to the automorphism group of each local environment [1]. However, explicitly enumerating these automorphism groups is often infeasible. In response to these challenges, Schur Net provides a spectral approach that circumvents direct group enumeration, yet its precise expressive power relative to group-based methods remains an open question [2].

Quite recently, a theoretical framework for measuring the expressivity of GNNs was introduced, dubbed *Homomorphism Expressivity* [3]. Crucially, Homomorphism Expressivity is the first measure of a GNN’s expressivity that does not fall under a Weisfeiler-Lehman test [3]. In this proposal, we layout a plan for theoretically measuring the expressivity of Schur Nets [2] using Homomorphism Expressivity [3]. Specifically, we could characterize how Schur Net’s spectral invariants compare to, and sometimes replicate, the subgraph-counting capabilities typically attributed to automorphism-based frameworks. Our analysis quantifies which families of graphs (and which substructures within them) each paradigm can homomorphism-count, and under what conditions they coincide or diverge. Furthermore, we extend our results to higher-order GNNs, clarifying how spectral techniques and group-theoretic constructions each handle more complex local symmetries. By illuminating these trade-offs, we aim to guide the design of hybrid GNN architectures that harness both the computational simplicity of spectral methods and the robust theoretical grounding of group-theoretic approaches.

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1 Introduction

Placing the task of learning under a mathematical formalism has been the aspiration of different groups of researchers for various reasons: philosophers, logicians, information theorists, researchers studying AI alignment, and others.

Historically, mathematicians used objects from statistics to construct a formal framework for learning. Simultaneously, a handful of mathematicians chose to couch learning in the language of group theory, using representation theory of finite groups [4] to realize models on Turing machines [5], [6].

2 Background

Designers of graph neural networks want their functions to treat a graph with one choice of node labels the same as the graph under another choice of vertex labels. Designers initially achieved this by ensuring that the first $\ell - 1$ layers were equivariant to permutations and then then ensured the last layer was invariant. This because most neural network φ are composed of as a series of layers $\varphi = \varphi^\ell \circ \varphi^{\ell-1} \circ \dots \circ \varphi^1$, and it is known that the composition of two equivariant functions is equivariant, so networks φ invariant to permutations of their input graphs can be constructed a series of equivariant layers, until the last layer is made to be invariant.

However, just as the GNN community realized historically that invariant internal layers could be swapped for more expressive equivariant layers, it has recently been shown [1] that, in order to construct a neural network φ that is *invariant to the entire input graph* \mathcal{G} , internal layers of a GNN φ^i need not be equivariant to the entire input graph \mathcal{G} , but rather need just be equivariant to the automorphism group of a template graph \mathcal{T} .

As we will see in Section 4.1, the template graph \mathcal{T} is chosen based on the problem domain. And so, designers using Autobahn based neural networks need to enumerate and manually choose the proper template graph. Schur Nets [2] was introduced to address this drawback in Autobahn. Schur Nets uses a spectral approach construct a Autobahn-style neural network. This allows Schur Nets to avoid having to explicitly state the template graph automorphism.

While Schur Nets approach to automorphism sub-graph equivariance avoids the overhead imposed by the purely group theoretic approach, Schur Nets' does not necessary yield an irreducible representation for each graph. Therefore, Schur Net-based constructions of equivariance do not necessarily yield the most generalizable forms of automorphism (Autobahn-type) neural networks. Furthermore, the extent of the gap between Schur Nets and group theoretic based automorphism networks remains an open question and area of research.

One of the primary challenges to exploring the expressivity gap between Autobahn and Schur Nets is that Schur Nets is not amenable to Weisfeiler-Lehman [7] tests of expressivity. In fact, the authors of Schur Nets intuit that

“...the ability to learn certain features (such as count cycles of certain sizes or learn some function related to specific functional groups) might be a better criterion ... [than WL type tests]. It'll be an interesting research direction to design suitable criteria for the expressive power of such higher order MPNNs in general.” [2]

This notion, that measuring the ability of GNN to count structures like sub-graph structures, has been recognized recently by other researchers as well, and has been formalized by [8] in a new framework for expressivity measurement, termed *Homomorphism Expressivity*. This initial work established Homomorphism Expressivity for sub-graph counting GNNs. Therefore, such a framework would *not* be directly applicable to Schur Nets, because it is a spectral invariant based GNN. However, following up on that work, Homomorphism Expressivity was extended to spectral invariant GNNs such as Schur Nets [3], giving us the tools necessary to compare pure group theoretic models (Autobahn) with spectral models (Schur Nets) using this framework.

In this paper we seek to kick off the theoretical exploration of this gap. First, we expound both the Autobahn and Schur Nets construction by formulating both approaching on example graphs. Not only does this formulation demonstrate this author's understanding and faculty with both frameworks¹, but also suggests a intuitive understanding of both the characters of the frameworks, which we will enforce by making qualitative observations of both behaviors.

With these preliminaries and intuitions established, we finally get to the main construction of this paper, which is a formal treatment and classification of both frameworks under Homomorphism Expressivity [3]. First, we show that, in the language of *Homomorphism Expressivity*, Schur Nets can be described as a *spectral invariant graph neural network*, and that Autobahn can be classified as a *sub graph counting graph neural network*. Next, using the main result from [3], we compute Schur Nets' Homomorphism Expressivity and compare it to Autobahn.

3 Preliminaries

In this paper, if we refer to a group action of the symmetric group on a k -ranked tensor without mentioning the particular action, assume it is the action defined below, dubbed *the permutation action*.

Definition 3.1 (*Permutation action on k th order tensors*) Let $\mathbf{T} \in \mathbb{R}^{n^k}$ be k th order (rank) tensor. Let $\sigma \in \mathbb{S}_n$ be a permutation in the symmetric group \mathbb{S}_n . Then we say that the *action of the permutation group \mathbb{S}_n on \mathbf{T}* , also called the permutation action on \mathbf{T} , is

$$\begin{aligned} & \sigma \cdot \mathbf{T} \\ [\sigma \cdot \mathbf{T}]_{i_1, \dots, i_k} &= [\mathbf{T}]_{\sigma^{-1}(i_1), \dots, \sigma^{-1}(i_k)} \end{aligned} \quad (1)$$

Remark 3.2 For a given graph with n nodes, represented by the rank two tensor $A \in \{0, 1\}^{n \times n}$, we say that \mathbb{S}_n permutes the graph by the action defined in Definition 3.1. Conceptually, permuting the graph A by Definition 3.1 can be equivalent to relabelling the nodes of the graph.

Graph neural network designers want their models to be invariant to all permutations on the input graph \mathcal{G} .

Definition 3.3 (*Permutation Invariance*) Let $f : \mathbb{R}^{n^{k_1}} \rightarrow \mathbb{R}^{n^{k_2}}$ be a function. We say that f is *permutation invariant to any permutation* if, for any input $A^{n^{k_1}}$,

$$f(\sigma \cdot A) = \sigma \cdot f(A) \quad (2)$$

for any $\sigma \in \mathbb{S}_n$.

Remark 3.4 Note, I am not sure exactly how you define permutation equivariance if f , represented by a matrix, was not square. For instance if

$$f = L^{m \times n} \quad (3)$$

and $m \neq n$, then $\sigma \cdot L$ cannot be defined using Definition 3.1.

Definition 3.5 (*Permutation Equivariance*) Similarly, we say that $f : \mathbb{R}^{n^{k_1}} \rightarrow \mathbb{R}^{n^{k_2}}$ is *permutation equivariant* We say that f is *permutation invariant to any permutation* if, for any input $A^{n^{k_1}}$,

$$f(\sigma \cdot A) = \sigma \cdot f(A) \quad (4)$$

for any $\sigma \in \mathbb{S}_n$.

¹Facetious

Corollary 3.6 (*An Equivalent Definition of Permutation Equivariance*) Note that the permutation action of \mathbb{S}_n given by Definition 3.1, immediately implies a homomorphism between \mathbb{S}_n and $\text{GL}(n, \mathbb{R})$. This can be used to redefine permutation equivariance when considering linear functions. Let $f : \mathbb{R}^n \rightarrow \mathbb{R}^n$, so f is in $\text{GL}(n, \mathbb{R})$ and can be written as the matrix $L \in \mathbb{R}^{n \times n}$. Next, let $\rho : \mathbb{S}_n \rightarrow \text{GL}(n, \mathbb{R})$, then we can say that $f \in \text{GL}(n, \mathbb{R})$ is permutation equivariant if and only if

$$\rho(\sigma) \star f = f \star \rho(\sigma) \quad (5)$$

for any σ , where \star is the binary operation of the group $\text{GL}(n, \mathbb{R})$. Note that \star is just matrix multiplication if we represent f and $\rho(\sigma)$ with matrices. One can also say that f is permutation equivariant to \mathbb{S}_n if f commutes with every element of \mathbb{S}_n .

In general, I believe a group action $\varphi : G \times \mathcal{X} \rightarrow \mathcal{X}$ always implies a homomorphism between G and the group of functions which map \mathcal{X} to \mathcal{X} under composition (sometimes termed the automorphism group). It is sometimes useful to think of permutation equivariance for a given internal layer φ^i this way as it generalizes its particular input.

Definition 3.7 (*Automorphism of a graph* $A \in \{0, 1\}^{n \times n}$) Let \mathcal{G} be a graph with n nodes. Let the adjacency matrix $A \in M(2, \mathbb{R})$ be \mathcal{G} 's representation. Given the representation A , consider the set of permutations in \mathbb{S}_n , that, when acting on \mathcal{G} via Definition 3.1 leave \mathcal{G} unchanged. That is the set

$$\{\sigma \in \mathbb{S}_n \mid \sigma \cdot A = A\}. \quad (6)$$

We denote this set of permutations by $\text{Aut}(\mathcal{G})$, that is the set of permutations that leave \mathcal{G} unchanged.

Example 3.8 (*4 node graphs*) It was not initially clear to me why one would be interested in the Automorphism group of a graph at all. Furthermore, For any given graph $A \in \{0, 1\}^{n \times n}$, that \mathbb{S}_n was not always essentially $\text{Aut}(A)$. Clearly though $\text{Aut}(A) \neq \mathbb{S}_n$. We can see this by considering the graph A and then considering the permutation $\sigma = (3 \ 1) \in \mathbb{S}_n$.

We note that $\sigma \cdot A \neq A$, and so $\sigma \notin \text{Aut}(A)$.

$$A = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} \quad \text{vs.} \quad \sigma \cdot A = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix} \quad (7)$$

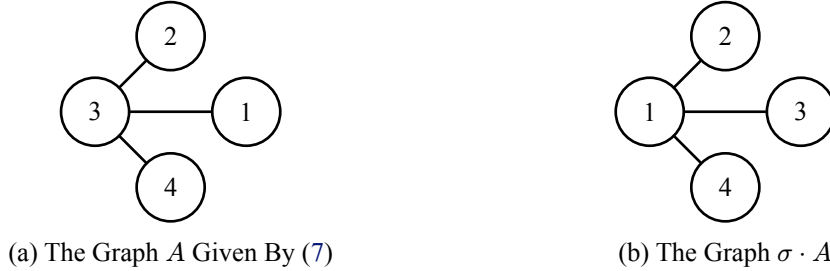


Figure 1: $A \neq \sigma A$ when $\sigma = (3 \ 1)$

Of course, A and $\sigma \cdot A$ look the same, and that is because their node-edge relations are the same, so invariance over \mathbb{S}_n is typically what we want. However, just as \mathbb{S}_n captures symmetries on the entire graph A , automorphisms, capture symmetry up to *sub-graphs*, so it is a finer filter. What I mean is that, notice how the nodes in sub-graph $\{(v_3, v_4), (v_3, v_2)\}$ in A and the nodes in the corresponding sub-graph in $\sigma \cdot A$, i.e. $\{(v_1, v_2), (v_1, v_4)\}$ are not the same across the two sub-graphs (v_3 was swapped out for v_1).

Automorphisms of A , such as $\sigma' = (4 \ 1)$ would ensure that the nodes involved in the sub-graph $\{(v_3, v_4), (v_3, v_2)\}$ in A would remain the same. So the automorphism group of A , $\text{Aut}(A)$ can be thought of as the set of relabellings that leave the sub-graphs of A intact.

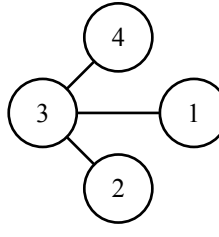
$$\sigma' \cdot A = \begin{pmatrix} 4 & 1 \end{pmatrix} \cdot A = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$


Figure 2: The Automorphism $\sigma' = (4 \ 1)$ Preserves The Sub-Graph Structure of A

Autobahn works by taking the user-given template graph and classifying functions that are equivariant to the entire input graph while judiciously breaking symmetry with the template graph's automorphism group (see Section 4.1.3). This has advantages and drawbacks, as discussed in Section 4.2.

In contrast, Schur Nets take a spectral approach. In particular, Schur Net layers are described by the Corollary proved in their work. This Corollary gives the core characterization of a non-higher order Schur Net.

In the next section we compare and contrast how Autobahn and Schur Nets how the models compute on a pendant graph with a 5-cycle, in precise detail. This will give us the conceptual basis to give a formal analysis of the expressivity of both frameworks. Moreover, it will suggest that the expressivity of both GNNs frameworks is best measured via sub-graph counting, and as a corollary why WL tests are not a satisfactory measure of expressivity for these frameworks. This will lead us to our major contribution, which is providing a precise characterization of Schur Nets and Autobahn within the measurement framework of Homomorphism Expressivity.

4 Exemplifying The Expressivity Gap Between Schur Nets and Autobahn

With its preliminaries established, the first main contribution of this report is to show in full detail how Autobahn [1] and Schur Nets [2] handle a certain graph.

The graph we will be applying Autobahn and Schur Nets to is a labelled 5 cycle graph with a pendant edge.

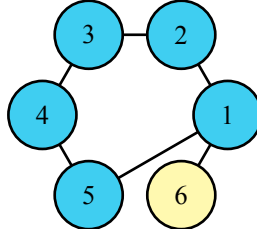
$$A = \begin{bmatrix} 0 & 1 & 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} \quad (8)$$


Figure 3: The Graph Given By The Adjacency Matrix In (8)

Remark 4.1 Before we continue, we make a few key observations about our chosen graph (8).

- The graph has a cycle of length 5, $(v_1, v_2, v_3, v_4, v_5, v_1)$
- v_6 is not in the cycle and the edge (v_6, v_1) is why we call this graph a “pendant”.
- We’ve colored (labelled) the nodes v_i that belong to the cycle in blue, the pendant node (v_6) that does not belong to the cycle in yellow.
- Clearly the colors and the pendant node introduce asymmetry.

4.1 Autobahn’s Behavior On The Pendant Graph (8)

Remark 4.1.1 For simplicity, I may gloss over the promotion and narrowing aspects of Autobahn’s algorithm. Although those are the essential contribution, I am primarily interested in treating group theoretic automorphism based networks, as given by [9]

We now observe how Autobahn works on Figure 3.

4.1.1 Step 1: Choose The Template Graph \mathcal{T}

Our first step is to choose our template graph \mathcal{T} that aligns with our problem domain. Autobahn is designed to recognize the sub-graphs of the input graph A isomorphic to this sub-graph \mathcal{T} and break permutation equivariance across sub-graph boundaries, which gives a more expressive class of neural networks than those whose neurons are permutation equivariant to all of \mathbb{S}_n (i.e MPNNs). Notice that we labelled the nodes with colors corresponding to the selection of our template graph.

Notice the 5-cycle in our node color labelling, so in line with that let us choose the template graph $\mathcal{T} = C_5$.

In particular, the adjacency matrix of \mathcal{T} is given by

$$A_{\mathcal{T}} = \begin{bmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 \end{bmatrix}. \quad (9)$$

The automorphism group of $A_{\mathcal{T}}$ is

$$\text{Aut}(A_{\mathcal{T}}) = D_5 \quad (10)$$

The dihedral group of 5 elements, D_5 is the set of rotations and reflections. Visually this makes sense, as rotating or reflecting this sub-graph should not change it.

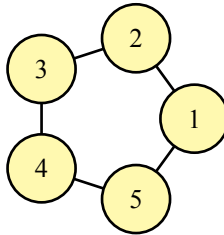


Figure 4: The Template Graph \mathcal{T} , A Sub-Graph of A

Now that we have chosen our template graph, Autobahn starts.

4.1.2 Step 2: Determine The Sub-graphs Isomorphic To C_5

- Objective: To find the sub-graph of A that are isomorphic to the template graph $\mathcal{T} = C_5$. In other words, find all sub-graphs in our input graph that contain 5-cycles.
- How: Find all $\sigma \in \mathbb{S}_6$ so that

$$\sigma \cdot A = A_{\mathcal{T}}. \quad (11)$$

4.1.3 Step 3: Construct A Function That Is Equivariant to $\text{Aut}(C_5) = D_5$

Let $f_i^{\ell-1}$ denote the feature vector from the previous layer ($\ell - 1$) corresponding to the i th vertex. For us, we will consider the first internal layer and let

$$\begin{aligned} f_i^0 &= \deg(v_i) \\ f^0 &= (f_1^0, \dots, f_6^0) \\ &= (3, 2, 2, 2, 2, 1) \end{aligned} \quad (12)$$

Let f'^0 be the feature vector but only on the 5-cycle

$$f_i'^0 = (\deg(v_i)) \quad (13)$$

for $i = \{1, \dots, 5\}$

We construct our sub-graph neuron so that it is pseudo-equivariant $\text{Aut}(C_5)$ if we only compute the features on the sub-graph, but in actuality is not equivariant to $\text{Aut}(C_5)$. We can construct our sub-graph neuron to be do this by having our sub-graph neuron convolve over all permutations in D_5

$$f'^1 = \sum_{\sigma \in D_5} \sigma \cdot w * f'^0 \quad (14)$$

for some learnable weight w .

Notice that, conceptually, if we only compute f'^0 with respect to the 5-cycle sub-graph, then $f_1^0 = \deg(v_1) = 2$ and (14) is indeed equivariant to D_5 . On the other hand, if we compute f'^0 with respect to the entire graph A , then $f_1^0 = \deg(v_1) = 3$, and (14) is not actually equivariant to D_5 . This is the crucial insight. In essence, Autobahn breaks permutation equivariance to D_5 of global features of the graph, while maintaining permutation equivariance to D_5 of local features of the graph.

In other words

$$\begin{aligned} f'^0 &= (2, 2, 2, 2, 2) \text{ When on considered on the isolated subgraph} \\ f'^0 &= (3, 2, 2, 2, 2) \text{ in actuality. I.e. with respect to the entire graph.} \end{aligned} \quad (15)$$

4.1.4 Step 4: Apply Non-Linearity And Add f_6^1

Finally, we construct the output layer

$$f^1 = f'^1 + f_6^1, \quad (16)$$

where f'^1 is computed with $f'^0 = (\deg(v_1), \deg(v_2), \deg(v_3), \deg(v_4), \deg(v_5)) = (3, 2, 2, 2, 2)$

4.2 Summarizing Autobahn

In summary, Autobahn requires you to choose a template graph \mathcal{T} that reflects the input graph's structures and problem domain nicely. For instance, we could have chosen $\mathcal{T} = P_3$, that is, all paths of length 3 in our input graph A .

- Finds the sub-graphs isomorphic to \mathcal{T} on A .
- Constructs neurons that are permutation equivariant to $\text{Aut}(A_{\mathcal{T}})$ with respect the sub-graph
- Incorporates these terms to the neuron's activations on other vertices not in the sub-graph, ensuring global equivariance with respect to the graph.

In summary, by ensuring functions on sub-graphs are permutation equivariant to D_5 **if the features the functions are acting on are computed with respect to the sub-graph**, Autobahn cleverly breaks permutation equivariance to D_5 by applying this function on f'^0 evaluated on the entire input graph.

One apparent drawback of Autobahn is that it requires users to select the template sub-graph. On the other hand, that requirement can actually be considered advantageous as it gives domain practitioners flexibility to choose the template graph that is important to their domain. For instance, chemists applying GNN’s on molecules can choose their template graph that is isomorphic to highly functional sub-structures of the molecule [10].

However, the major drawback of Autobahn is that it requires considering all elements of the automorphism group in order to achieve equivariance, via generalized convolution. Moreover, the complexity that arises due to overlapping sub-graphs isomorphic to the input graph introduces complexities and computational overhead. Autobahn addresses these issues by introducing narrowing and promotion, but it is still a source of complexity. This is discussed in Section A.1.

4.3 How Schur Nets Handles The Pendant Graph

The construction of a Schur Net neuron is given by the following corollary proved in their work.

Corollary 4.3.1 (*Schur Net Neuron (Zhang, Xu & Kondor) [2]*) Consider a GNN on the input graph \mathcal{G} , represented by A . Let n_F be a neuron that operates on sub-graph F of \mathcal{G} . Let input sub-graph F (and its features) be represented by $T \in \mathbb{R}^m$. Let L be the Laplacian of S , U_1, \dots, U_p be the eigenspaces of L , and M_i be an orthogonal basis for the i th eigenspace stacked into a $\mathbb{R}^{n \times \dim(U_i)}$. Then for any collection of learnable weight matrices $W_1, \dots, W_p \in \mathbb{R}^{a \times b}$

$$\varphi : T \longrightarrow \sum_{i=1}^p M_i M_i^T T W_i \quad (17)$$

is a permutation equivariant operation.

Like Autobahn, Schur Net’s allows the designer to choose the sub-graph F (see Corollary 4.3.1). Therefore, we choose $F = C_5$, inline with the prior section.

Next, we compute the Laplacian of F . For convenience, recall that the adjacency matrix of the input graph and the 5-cycle template sub-graph stated in (8) is given by

$$A = \begin{bmatrix} 0 & 1 & 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} \quad (\text{Input Graph})$$

and,

$$A_{\mathcal{T}} = \begin{bmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 \end{bmatrix} \quad (5\text{-Cycle}).$$

Then the Laplacian of F is given by $L = D_F - A_F$ so

$$L = \begin{bmatrix} 2 & -1 & 0 & 0 & -1 \\ -1 & 2 & -1 & 0 & 0 \\ 0 & -1 & 2 & -1 & 0 \\ 0 & 0 & -1 & 2 & -1 \\ -1 & 0 & 0 & -1 & 2 \end{bmatrix}. \quad (18)$$

Next, we compute the eigenvalues of L . This is known for 5-cycles [11]. Let λ_i denote the i th eigenvalue. Then for $i = k + 1$, compute $i = 1, \dots, 5$

$$\lambda_i = 2 - 2 \cos\left(\frac{2\pi k}{5}\right) \quad (19)$$

We also know the eigenvectors are given by the columns of the discrete Fourier matrix. So the k th eigenvector is given by

$$v_j^{(k)} = \frac{1}{\sqrt{n}} e^{\frac{2\pi i k j}{n}}, \quad j = 0, 1, \dots, n-1, \quad n = 5. \quad (20)$$

The eigenvalues, their multiplicities κ , and the eigenvectors are shown in Figure 1.

Figure 1: The Eigenvalues of The C_5 Laplacian L

i	Eigenvalues (λ_i) of $C(5)$	Eigenvectors (v_i) of C_5	Multiplicity, κ_i
0	0	$\left(\frac{1}{\sqrt{5}}, \frac{1}{\sqrt{5}}, \frac{1}{\sqrt{5}}, \frac{1}{\sqrt{5}}, \frac{1}{\sqrt{5}}\right)$	1
1	1.381966011250105	$\left(\frac{1}{\sqrt{5}}(1, \cos(2*\frac{\pi}{5}), \cos(4*\frac{\pi}{5}), \cos(4*\frac{\pi}{5}), \cos(2*\frac{\pi}{5}))\right)$	2
2	3.618033988749895	$\left(\frac{1}{\sqrt{5}}(0, -\sin(2*\frac{\pi}{5}), -\sin(4*\frac{\pi}{5}), \sin(4*\frac{\pi}{5}), \sin(2*\frac{\pi}{5}))\right)$	2

Therefore, our eigenspaces, expressed as column vectors (with duplicates), are

$$\begin{aligned}
M_1 &= \left[\frac{1}{\sqrt{5}}, \frac{1}{\sqrt{5}}, \frac{1}{\sqrt{5}}, \frac{1}{\sqrt{5}}, \frac{1}{\sqrt{5}}\right]^\top \\
M_2 &= \left[\frac{1}{\sqrt{5}}(1, \cos(2*\frac{\pi}{5}), \cos(4*\frac{\pi}{5}), \cos(4*\frac{\pi}{5}), \cos(2*\frac{\pi}{5}))\right]^\top \\
M_3 &= \left[\frac{1}{\sqrt{5}}(0, -\sin(2*\frac{\pi}{5}), -\sin(4*\frac{\pi}{5}), \sin(4*\frac{\pi}{5}), \sin(2*\frac{\pi}{5}))\right]^\top \\
M_4 &= \left[\frac{1}{\sqrt{5}}(1, \cos(4*\frac{\pi}{5}), \cos(2*\frac{\pi}{5}), \cos(4*\frac{\pi}{5}), \cos(4*\frac{\pi}{5}))\right]^\top \\
M_5 &= \left[\frac{1}{\sqrt{5}}(0, -\sin(\frac{\pi}{5}), \sin(2*\frac{\pi}{5}), -\sin(3*\frac{\pi}{5}), \sin(4*\frac{\pi}{5}))\right]^\top.
\end{aligned} \quad (21)$$

Next we define $T \in \mathbb{R}^{5 \times 2}$. Let $T_{i,j} = \begin{cases} \text{The degree of node } i & j=1 \\ \text{Where node } i \text{ is next to the pendant } & j=2 \end{cases}$

$$T = \begin{bmatrix} 3 & 1 \\ 2 & 0 \\ 2 & 0 \\ 2 & 0 \\ 2 & 0 \end{bmatrix} \quad (22)$$

Finally compute the terms of Corollary 4.3.1. For instance,

$$\begin{aligned}
M_1 M_1^\top &= \frac{1}{5} \begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \end{bmatrix} \\
M_1 M_1^\top T &= \frac{1}{5} \begin{bmatrix} 11 & 1 \\ 11 & 1 \\ 11 & 1 \\ 11 & 1 \\ 11 & 1 \end{bmatrix}
\end{aligned} \quad (23)$$

4.4 Schur Nets Summary

We showed how to construct a Schur Net neuron on the pendant graph. Namely we computed the eigenvalues and eigenspaces of our sub-graph Laplacian via the discrete Fourier basis, which yield orthonormal sub-spaces M_i . The eigenvalues of L had multiplicities, 1, 2, 2, respectively. The corresponding eigenspaces were constructed using the discrete Fourier basis, yielding orthogonal bases M_i . Specifically M_1 corresponded to the eigenvalue of 0 with the eigenvector $\frac{1}{\sqrt{5}}[1 \ 1 \ 1 \ 1 \ 1]^T$, with each M_i being a $5 \times \dim(U_i)$ matrix of our normalized eigenvectors.

The key points to distinguish Schur Nets behavior with that of Autobahn is that Autobahn allowed us to specify how to construct our “equivariant” sub-graph neuron (see (14)), while Schur Nets gave us no such flexibility, instead capturing the asymmetry through spectral analysis. This rigidity in Schur Nets ensures a consistent, mathematically grounded approach, relying on the Laplacian’s eigenspaces rather than flexible neuron design, providing a natural characterization of the graph’s structure. Whereas Autobahn computes D_5 directly, allowing more flexibility, Schur Nets uses spectral filters to process the node features, where each orthogonal sub-space corresponds to a part of the cycle.

The main limitation, as shown by the eigenvalue multiplicity Figure 1, is that the Schur Net’s representation of D_5 is not irreducible—namely subspaces M_1 and M_2 .

5 A General Characterization Of The Expressivity of Both Autobahn And Schur Nets

The main contribution of this project is to provide a precise characterization of the expressivity of automorphism based GNNs given by Autobahn and Natural Graph Networks [9], [12] as opposed to Schur Nets. We are able to do this in general thanks to framework introduced by Zhang et al. [8] and in particular for Schur Nets thanks to the recent results from Zhang and Maron et al [3].

Zhang [8] introduces the concept of *Homomorphism Expressivity*.

Definition 5.1 (Zhang et al [3]) Given two graph F and G , a homomorphism from F to G is a mapping $f: V_F \rightarrow V_G$ that preserves edge relations, meaning that, for every $\{u, v\} \in E_F$

$$\{f(u), f(v)\} \in E_G. \quad (24)$$

The set of all homomorphisms from F to G will be denoted as $\text{Hom}(F, G)$

Remark 5.2 A question I had, which I believe is valid, is, can the set of homomorphisms from F to G be counted. The answer turns out to be yes in some cases.

Let φ be graph neural network that outputs some graph invariant. Then, the *Homomorphism Expressivity* of φ can be defined.

Definition 5.3 (Homomorphism Expressivity Zhang et al [8]) Let φ be a GNN (a function on graph G that outputs some invariant). Let $\mathcal{X}_G^\varphi(G)$ be an invariant that φ computes on the graph G . *Homomorphism Expressivity* of φ , denoted by \mathcal{F}^φ , is a family of connected graphs² satisfying the following:

- For any two graphs G and H , $\mathcal{X}_G^\varphi = \mathcal{X}_H^\varphi$ iff $\text{Hom}(F, G) = \text{Hom}(F, H)$ for every $F \in \mathcal{F}^\varphi$.
- For any graph $F \notin \mathcal{F}^\varphi$, there exists a pair of graph G, H so that $\mathcal{X}_G^\varphi = \mathcal{X}_H^\varphi$ and $\text{Hom}(F, G) \neq \text{Hom}(F, H)$.

Let us compute the homomorphism expressivity for a few simple GNNs.

Example 5.4 (*Triangle-Counting GNN*) Let φ be a GNN that outputs the number of triangles in the graph G . For a graph $G = (V_G, E_G)$, $\mathcal{X}_G^\varphi(G) = \|\{\{u, v, w\} \subseteq V_G \mid \{u, v\}, \{v, w\}, \{w, u\} \in E_G\}\|$. This GNN keeps invariance to \mathbb{S}_n by aggregating node features over 3-hop neighborhoods, for example.

²A connected graph is a graph $G = (V, E)$ where there is a path v_i to any given node v_j .

Let G be a complete graph, $G = K_4$. This means it has 4 triangles. Let H be a C_4 graph with one extra edge. Say $H = (v_1, v_2, v_3, v_4, v_1)$ and has the edge (v_1, v_3) .

Therefore, \mathcal{X}^φ is the following over the various graphs

- $\mathcal{X}^\varphi(G) = 4$ (triangles in K_4)
- $\mathcal{X}^\varphi(H) = 1$ (triangle 1-2-3 in C_4)
- $\mathcal{X}^\varphi(G) \neq \mathcal{X}^\varphi(H)$, so we do not need to check if $\text{Hom}(K_3, G) = \text{Hom}(K_3, H)$, by part 1 of Definition 5.3.

Now we consider a new pair G' and H'

- G' Two disjoint triangles $t_1 = (v_1, v_2, v_3)$ and $t_2 = (v_4, v_5, v_6)$, each triangle forming a K_3 .
- H' A cycle C_6 with two triangles. Let us say it has nodes $(v_1, v_2, v_3, v_4, v_5, v_6)$ and edges (v_1, v_3) and (v_4, v_6) .
- $\mathcal{X}^\varphi(G') = 2$ and $\mathcal{X}^\varphi(H') = 2$, so $\mathcal{X}^\varphi(G') = \mathcal{X}^\varphi(H')$. This is a more interesting case because it forces us to consider the homomorphisms from triangle to triangle on G' and H' .

We now consider the homomorphisms $\text{Hom}(K_3, G')$ and $\text{Hom}(K_3, H')$

- A homomorphism from K_3 to a graph maps the triangle to another triangle (see “preserves the edges” in Definition 5.1)
- G' has 2 triangles, so $\text{Hom}(K_3, G') = 2$ as each triangle in G' is distinct: $(v_1, v_2, v_3) \neq (v_4, v_5, v_6)$.

So $K_3 \in \mathcal{F}^\varphi$ because whenever $\mathcal{X}^\varphi(G) = \mathcal{X}^\varphi(H)$, the number of homomorphisms from K_3 , that is the number of triangles counts, is the same across both graphs.

Conclusion: So triangles $K_3 \in \mathcal{F}^\varphi$. In lay terms, our GNN φ can recognize triangles.

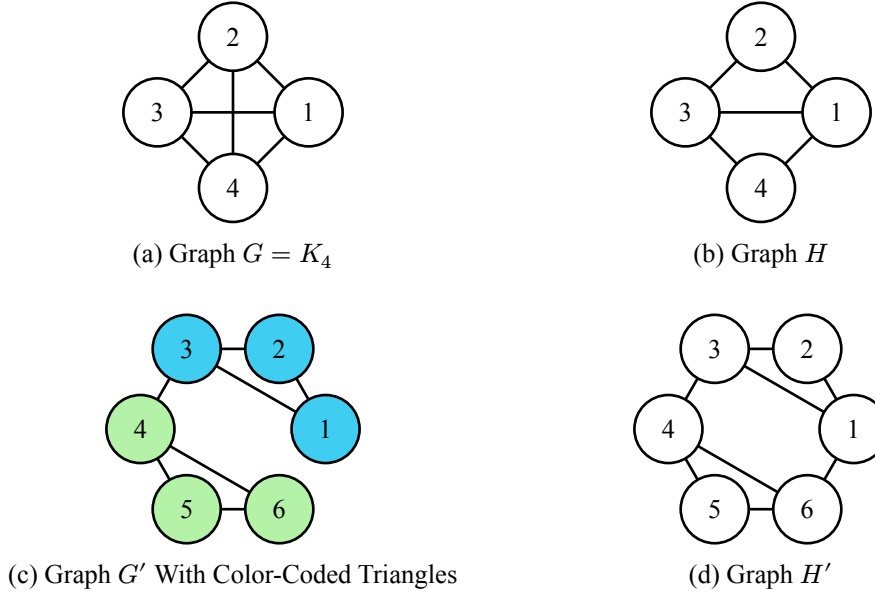


Figure 5: The Graphs of G, H, G' , and H'

So our GNN φ is expressive for triangles. Let us now see how it does with 3-paths.

Example 5.5 (*Homomorphism Expressivity of Triangle GNN for P_3*) We will compute the homomorphism expressivity using our triangle GNN φ , which we defined in the last example.

- Let $F = P_3$, where F is given by Definition 5.1.
- So F is a 3-path (v_1, v_2, v_3) .
- Let G be a 4-cycle. So G has no triangles, and $\mathcal{X}^\varphi(G) = 0$
- Let H be a star graph \mathbb{S}_4 . No triangles also, so $\mathcal{X}^\varphi(H) = 0$
- $\mathcal{X}^\varphi(G) = \mathcal{X}^\varphi(H) = 0$

Next consider the Homomorphisms $\text{Hom}(P_3, G)$ and $\text{Hom}(P_3, H)$

- $\text{Hom}(P_3, G)$: Map P_3 to a path of length 2 in C_4 . C_4 has 4 such paths, and for each path, a P_3 can be mapped in 2 directions: $u \rightarrow 1, v \rightarrow 2, w \rightarrow 3$ or $u \rightarrow 3, v \rightarrow 2, w \rightarrow 1$. So There are $4 \times 2 = 8$ homomorphisms.
- $\text{Hom}(P_3, H)$: In \mathbb{S}_4 map v to the center and u, w to two distinct leaves. So we can choose 2 leaves out of 3, and for each choice, we can map u, w in 2 ways ($u \rightarrow 1, w \rightarrow 2$ or $u \rightarrow 2, w \rightarrow 1$). Giving us $3 \times 2 = 6$ homomorphisms.
- $\text{Hom}(P_3, G) = 8, \text{Hom}(P_3, H) = 6$, so $\text{Hom}(P_3, G) \neq \text{Hom}(P_3, H)$.

Conclusion: Since $\mathcal{X}^\varphi(G) = \mathcal{X}^\varphi(H)$ but $\text{Hom}(P_3, G) \neq \text{Hom}(P_3, H)$, this implies that $P_3 \notin \mathcal{F}^\varphi$ by Definition 5.1.

5.1 Aside: Homomorphism Expressivity Computation of A Basic Schur Nets

To recapitulate the examples given in Section 5 using Homomorphism Expressivity Definition 5.1, let us compute the Homomorphism Expressivity Definition 5.1 for a very simple Schur Net GNN.

Example 5.1.1 (*Homomorphism Expressivity For A Simple Schur Net*) Let us define φ to a simplified Schur Net that outputs the smallest non-zero eigenvalue of the Laplacian as its invariant $\mathcal{X}^\varphi(\mathcal{G})$.

$$\mathcal{X}^\varphi(\mathcal{G}) = \lambda_1(\mathcal{G}) \quad (25)$$

- Let G be a 5-cycle, $G = C_5$
- Let H be the pendant in Section 4.3

Using our computations of eigenvalues of the 5-cycle in Figure 1, we have that $\lambda_1 = 2 - 2\cos(2\pi/5) \approx 1.38$ for \mathcal{G} , a 5-cycle (see Section 4.3). So $\mathcal{X}^\varphi(G) = \lambda_1 \approx 1.38$

On the other hand, I still need to write code to compute the spectrum of H , our pendant in Section 4.3, but I will assume is close to 1.38. So $\mathcal{X}^\varphi(H) = \lambda_1(H) \approx 1.38$.

With this assumption, we have that

$$\mathcal{X}^\varphi(G) = \mathcal{X}^\varphi(H) \approx 1.382, \quad (26)$$

allowing us to applying the first criteria of Definition 5.1.

Given $\mathcal{X}^\varphi(G) = \mathcal{X}^\varphi(H)$. We compute $\text{Hom}(F, G)$ and $\text{Hom}(P_2, H)$. Of course the question is, which graph F do we want to choose? Let us choose $F = P_2$.

6 Acknowledgements

This work is incomplete. I will send the completed version this weekend, but I acknowledged that this submission will be the one that is graded.

6.1 TODOs

- Transcribe C4 example of 5 node (pendant node makes it clear, but I like comparing it to \mathbb{S}_4 better) from scratch notes but already 1/3 way.
- Transcribe C3 example (my first attempt earlier this week, include it ?)

- Now, gather parallel trees analysis notes for $d = 2$ and 3. (Too many examples)?
- Now, show exactly how Schur Nets is good at capturing parallel trees up to d , show autobahn captures symmetries more generally. **Schur Nets may miss symmetries in parallel trees after a certain depth.**

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A Appendix

Definition A.1 (*Narrowing In Autobahn [1]*) Let (i_1, \dots, i_k) be an ordered subset of $\{1, 2, \dots, m\}$. The set $\{1, \dots, m\}$ corresponds to the sub-graph isomorphic to the template graph \mathcal{T} , and the ordered subset (i_1, \dots, i_k) corresponds to k “ambassador nodes” that overlap with a different sub-graph.

Now, for all $u \in \mathbb{S}_k$, let $\tilde{u} \in \mathbb{S}_k$ be the permutation that u to the first k elements and for all $s \in \mathbb{S}_{m-k}$, let $\tilde{s} \in \mathbb{S}_m$ be the permutation that applies s to the last $m - k$ elements. Then, given our activation function $f : \mathbb{S}_m \rightarrow \mathbb{R}^d$, the *narrowing* of f to (i_1, \dots, i_k) is

$$f \downarrow_{i_1, \dots, i_k} = (m - k)!^{-1} \sum_{s \in \mathbb{S}_{m-k}} f(\tilde{u} \tilde{s} t) \quad (27)$$

Remark A.2 This was taken from [1]. At a high-level narrowing produces a function $f \downarrow$ that is only dependent on the important nodes $(v_{i_1}, \dots, v_{i_k})$.

A.1 How Narrowing And Promotion Work

Remark A.1.3 We omitted details of Autobahn’s primary contribution, that is, the formalized notions of *narrowing* and *promotion*. This is because we are concerned with Autobahn’s sub-graph counting abilities, and narrowing and promotion are techniques to reduce the computation complexity of achieving automorphism based neurons when graphs overlap [9].

We constructed our example to be particularly simple, and showed such an automorphism based neuron [9]. The point of the section was to illustrate, through elementary computation and novice language, the sub-graph counting abilities of automorphism based neuron networks.

See the appendix Section A.1 for an explanation of how Autobahn uses narrowing and promotion to reduce the computation overhead, which is important when the input graph contains more than 1 sub-graph isomorphic to the template graph.

At any given layer i , Autobahn first identifies all the sub-graphs of A that are isomorphic to $\mathcal{A}_{\mathcal{T}}$. In our case, $(v_1, v_2, v_3, v_4, v_5)$ is the only such sub graph of A .

Next, given we are looking at the first layer, we need to define an input domain. Let the input to the first layer be the degree of each node on the feature graph.

Next, for each sub-graph isomorphic to A , Autobahn performs *narrowing*. *Narrowing* will be described plainly here. The formal definition is given in the appendix. *Narrowing* takes the sub-graph $(v_{i_1}, \dots, v_{i_m})$ and partitions the nodes of the graph into two groups: the important group $\mathcal{I}_1 = \{i_1, \dots, i_k\}$ and the not important group consisting of the remaining $k - m$ nodes, $\mathcal{I}_2 = \{i_{k+1}, \dots, i_m\}$. Note that I assumed that the first 1 through k indices were the important nodes, but typically narrowing adds the indices to the important group which *overlap* with a another sub-graph. Let us call such nodes the “ambassador” nodes.

In our example, narrowing operates on our sub-graph $\{v_1, \dots, v_5\}$. It then adds node v_1 to the important group $\mathcal{I}_1 = \{i_1\}$ and puts the remaining 2, ..., 4 indices in the not important group $\mathcal{I}_2 = \{i_2, \dots, i_4\}$. Finally, the narrowing procedure is applied to our sub-graph neuron, which makes our sub-graph neuron only depend on the important nodes in \mathcal{I}_1 . See the appendix Definition A.1 for the formal definition of narrowing.

Similarly to how narrowing is done, then *promoting* is done to the neuron, which takes the narrowing activations and spreads them out to all the m nodes in the sub-graph that this neuron is working on.

We won’t go into narrowing or promotion in too much detail, but we can provide a sketch of how to ensure that our neuron \mathcal{n} is equivariant to $\text{Aut}(A_{\mathcal{T}})$ and also equivariant to all of A in general.

For example, in our case let the input feature vectors be $f_1 = \deg(v_1) = 2, f_2 = \deg(v_2) = 2 = \dots = f_5$. Notice that $f_1 = 2$, because we are only considering degrees within the sub graph v_1, \dots, v_5 .

Now, we can ensure equivariance of our function \mathcal{n} on this 5-tuple with respect to D_5 , by convolving over the element of D_5 .

$$D_5 = \sum_{\sigma \in D_5} f \quad (28)$$

Now, in order to account for v_1 ’s connection to v_6 , we redefine the feature vectors f_i to include a indicator variable saying if the node is adjacent to the pendant node v_6 .

$$f_1 = (\deg(v_1), 1) = (2, 1) \quad f_2 = (\deg(v_1)), 0) = (2, 0) = f_2 \dots = f_5 \quad (29)$$