EE3731C Statistical Signal 3.5

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Recap (1)

• Discrete time Markov process

$$p(x_1, \cdots, x_{n-1}, x_{n+1}, \cdots, x_N | x_n) = p(x_1, \cdots, x_{n-1} | x_n) p(x_{n+1}, \cdots, x_N | x_n)$$
PAST FUTURE

• Discrete time, discrete state Markov process: $\pi_{n+1} = \pi_n T$, where $\pi_n =$ probability of different states at time n

$$\pi_{n+1} = [\pi_n(1) \ \pi_n(2) \ \pi_n(3)] \times \begin{bmatrix} 0.9 & 0.075 & 0.025 \\ 0.15 & 0.8 & 0.05 \\ 0.25 & 0.25 & 0.5 \end{bmatrix}$$

Recap (2)

• From previous slide,

$$\pi_{n+1} = \pi_n T$$

$$\implies \pi_{n+2} = \pi_{n+1} T = \pi_n T^2$$

$$\implies \pi_{n+k} = \pi_n T^k$$

- If $\pi^* = \pi^* T$, then π^* is stationary distribution of T.
 - Because $\pi_{n_0} = \pi^* \implies$ probability of any state constant (stationary) for $n \ge n_0$
 - $-\pi^*$ is the left eigenvector of T with eigenvalue 1
- Fundamental Theorem of Markov Chains
 - If there is n_0 , such that $T^n(i,j) > 0$ for all i,j and $n > n_0$, then markov chain has unique stationary distribution π^* .
 - For any π_1 , as $n \to \infty$, $\pi_1 T^n \to \pi^*$

Recap (2)

• From previous slide,

$$\pi_{n+1} = \pi_n T$$

$$\implies \pi_{n+2} = \pi_{n+1} T = \pi_n T^2$$

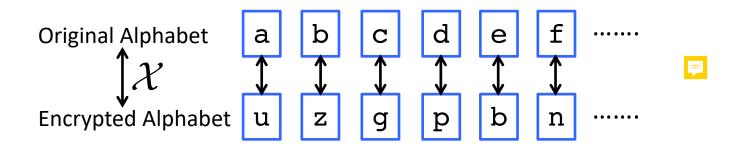
$$\implies \pi_{n+k} = \pi_n T^k$$

- If $\pi^* = \pi^* T$, then π^* is stationary distribution of T.
 - Because $\pi_{n_0} = \pi^* \implies$ probability of any state constant (stationary) for $n \ge n_0$
 - $-\pi^*$ is the left eigenvector of T with eigenvalue 1
- Fundamental Theorem of Markov Chains (in plain English)
 - Suppose we start at an arbitrary state i at time 1. If there is non-zero probability of being in any state any time after finite time n_0 , then Markov chain has unique stationary distribution.
 - Any initial state will reach stationary distribution given sufficient time

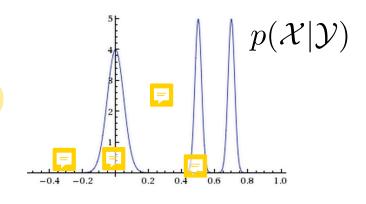
Markov Chain Monte Carlo (MCMC)

Markov Chain Monte Carlo (MCMC)

- Why stationary distribution? MCMC are methods for sampling from p(x) by constructing markov chain whose stationary distribution $\pi^*(x) = p(x)$
- Why sample? Given observations \mathcal{Y} , we often want samples from conditional distribution $p(\mathcal{X}|\mathcal{Y})$
 - e.g., in programming assignment, \mathcal{Y} is encrypted text, \mathcal{X} is mapping between the original and encrypted alphabets.



- If posterior $p(\mathcal{X}|\mathcal{Y})$ is very "peaky", then samples of \mathcal{X} likely close to peaks (high probability) and good decryption candidates
- MAP \approx sample with largest posterior
- MMSE \approx average of all samples



Sampling $\pi(\mathcal{X})$ with Metropolis Algorithm

1. Start with any $x = x_0$

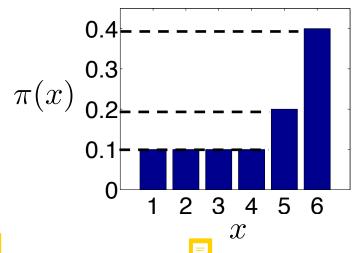
- F
- 2. Sample new x' using any proposal distribution q(x'|x)
 - Only constraint is q(x'|x) = q(x|x')
- 3.

$$\begin{cases} \text{ if } \pi(x') \geq \pi(x) & \text{Replace } x \text{ with } x' \\ \text{if } \pi(x') < \pi(x) & \text{Replace } x \text{ with } x' \text{ with probability } \pi(x')/\pi(x) \\ & \text{Keep } x \text{ with probability } 1 - \pi(x')/\pi(x) \end{cases}$$

- 4. Repeat steps 2 and 3 for a "long" time before stopping
 - Current $x \sim \pi(\mathcal{X})$
- 5. Will explain why this works later. First see example.

Biased Dice Metropolis Example

• Biased dice



• Metropolis:

1. Let
$$x = 1$$

2. Generate x' with q(x'|x) = 1/6 for x' = 1, 2, 3, 4, 5, 6

$$\begin{cases} \text{ if } \pi(x') \geq \pi(x) & \text{Replace } x \text{ with } x' \\ \text{ if } \pi(x') < \pi(x) & \text{Replace } x \text{ with } x' \text{ with probability } \pi(x')/\pi(x) \\ & \text{Keep } x \text{ with probability } 1 - \pi(x')/\pi(x) \end{cases}$$

4. Repeat steps 2 and 3 ten times and return current value of x as one sample from $\pi(x)$

10 iterations of Biased Dice Example

```
Iter 1: Current x = 1, new x' = 5

pi(x') >= pi(x): definitely accept

10.4

Iter 2: Current x = 5, new x' = 3

pi(x') < pi(x): accept with probability pi(x')/pi(x)

Coin toss with p = pi(x')/pi(x) = 0.5; head; accept x'

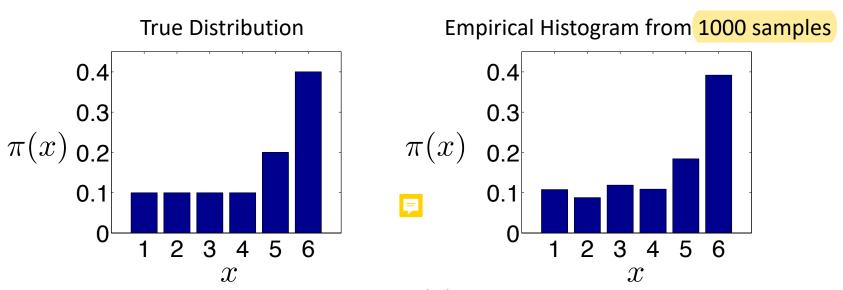
1 2 3 4 5 6
```

Iter 9: Current x = 6, new x' = 6pi(x') >= pi(x): definitely accept

Iter 10: Current x = 6, new x' = 5pi(x') < pi(x): accept with probability pi(x')/pi(x) Coin toss with p = pi(x')/pi(x) = 0.5; tail; reject x'

1st sample of p(x) is 6

Repeat 1000 times to get 1000 samples



- What can we use samples from $\pi(x)$ for?
 - Can estimate E(x) by averaging samples. In dice example, empirical mean = 4.349 (true mean E(x) = 4.4)
 - In real applications, we often sample from p(x|y). Suppose $\pi(x) = p(x|y)$, then if we average samples, we get approximation of E(x|y), i.e., MMSE estimate of x given y

Why do we want to sample from p(x | y)?

- In example, I used Metropolis algorithm to generate 1000 samples from biased dice $\pi(x)$, where $\pi(1) = \pi(2) = \pi(3) = \pi(4) = 0.1$, $\pi(5) = 0.2$ and $\pi(6) = 0.4$.
- Since we already know $\pi(x)$, why sample from $\pi(x)$?
- More specifically, suppose $p(x|y) = \pi(x)$ from the dice example. We can compute the MAP estimate of x given y easily. For example, $x_{MAP} = 6$ because $\pi(6)$ has the highest value. Why bother to generate 1000 samples from $\pi(x)$?
- The reason is computation of MAP and MMSE estimates is difficult when x is multi-dimensional.

Why do we want to sample from p(x | y)?

- Let $\vec{x} = \{x_1, \dots, x_N\}$. Let's assume $x_n \in \{1, \dots, K\}$
- Given observation y: $\vec{x}_{MAP} = \operatorname{argmax}_{\vec{x}} p(\vec{x}|y)$
- To find MAP estimate, we can iterate through every possible value of \vec{x} and keep track of the value of \vec{x} with the highest posterior probability.
- However, there are K^N possible values of $\vec{x} \implies$ need to compare K^N possible values of $p(\vec{x}|y)$ to find the best \vec{x} .
- In programming assignment, \vec{x} is mapping between original alphabets and encrypted alphabets. Therefore, N=27 and $K=27 \implies$ "straightforward" approach of computing \vec{x}_{MAP} requires us to compare 27^{27} instances of $p(\vec{x}|y)$.

Why do we want to sample from p(x | y)?

- Instead, we can use Metropolis algorithm to generate M samples from $p(\vec{x}|y)$.
- Denote these samples as $\vec{x}^{(1)}, \dots, \vec{x}^{(M)}$. Then

$$\vec{x}_{MAP} \approx \operatorname*{argmax}_{\vec{x}^{(m)}} p(\vec{x}^{(m)}|y), \sqsubseteq$$

which only requires us to compare M instances of $p(\vec{x}|y)$.

proximate the MAP estimate (even though we could).

• In programming assignment, we run Metropolis algorithm for 15000 iterations to generate 1 sample. As will be seen in assignment, this 1 sample is good enough to unscramble the encrypted message. We do not need to generate multiple samples to ap-

Why does Metropolis work?

Why does Metropolis Algorithm Work?

- 1. Start with any $x = x_0$
- 2. Sample new x' using any proposal distribution q(x'|x)
 - Only constraint is q(x'|x) = q(x|x')

3.

$$\begin{cases} \text{ if } \pi(x') \geq \pi(x) & \text{Replace } x \text{ with } x' \\ \text{ if } \pi(x') < \pi(x) & \text{Replace } x \text{ with } x' \text{ with probability } \pi(x')/\pi(x) \\ & \text{Keep } x \text{ with probability } 1 - \pi(x')/\pi(x) \end{cases}$$

- 4. Repeat steps 2 and 3 for a "long" time before stopping
 - Current $x \sim \pi(\mathcal{X})$

Why does Metropolis Algorithm Work?

- Running Metropolis equivalent to following Markov Chain:
 - Initial state: $x_1 = x_0$
 - Transition probability:

For
$$x \neq x'$$

$$T(x,x') = \begin{cases} q(x'|x) & \text{if } \pi(x') \ge \pi(x) \\ q(x'|x)\pi(x')/\pi(x) & \text{if } \pi(x') < \pi(x) \end{cases}$$

and

$$T(x,x) = C_x$$
 to ensure $\sum_z T(x,z) = 1$

• If T satisfies Fundamental Theorem of Markov Chain (which is quite easy), then Markov chain has stationary distribution $\pi(\mathcal{X})$

Why Stationary Distribution is $\pi(\mathcal{X})$ (1)

- First, show $\pi(x)T(x,x') = \pi(x')T(x',x)$
- Proof is at end of slides
- Property is known as "detailed balance". Markov chain is considered "reversible"

Why Stationary Distribution is $\pi(\mathcal{X})$ (2)

- From previous slide $\pi(x)T(x,x') = \pi(x')T(x',x)$
- Let's verify $\pi(\mathcal{X})T = \pi(\mathcal{X})$. To evaluate k-th element of $\pi(\mathcal{X})T$, consider

$$T(:,k) = k\text{-th column of } T$$

$$= \sum_{x \in \mathcal{X}} \pi(x) T(x,k)$$

$$= \sum_{x \in \mathcal{X}} \pi(k) T(k,x)$$

$$= \sum_{x \in \mathcal{X}} \pi(k) T(k,x)$$

$$= \pi(k) \sum_{x \in \mathcal{X}} T(k,x)$$

$$= \pi(k)$$

$$\sum_{x \in \mathcal{X}} T(k,x) = \pi(k) T(k,x)$$

$$= \pi(k)$$

$$\sum_{x \in \mathcal{X}} T(k,x) = 1$$

$$= \pi(k)$$

• k-th element of $\pi(\mathcal{X})T$ equals $\pi(k)$, so $\pi(\mathcal{X})$ is stationary distribution

Specifying Proposal Distribution q(x' | x)

Specifying proposal distribution $q(x' \mid x)$

F

- Fundamental Theorem easy to satisfy
 - Exists n_0 , such that $T^n(i,j) > 0$ for all i,j and $n > n_0$
 - Dice example: $q(x'|x) \neq 0$ for all $x' \implies T(x,x') \neq 0$ for all $x' \implies n_0 = 1$
- To sample $\pi(\mathcal{X}) = \frac{1}{Z} f(\mathcal{X})$, observe Z cancels out in the Metropolis algorithm, so replace π with f
 - Useful when Z difficult to compute (e.g., programming assignment)
- Can just specify procedure for generating random x' from x. No need to specify q(x'|x)
- For continuous p(x), replace discrete states with continuous states



Summary

- Monte Carlo Markov Chain
 - Metropolis algorithm
 - Detailed balance: $\pi(x)T(x,x') = \pi(x')T(x',x)$
- Probabilistic signal detection
 - Rather than single estimate (ML, MAP, MMSE), sample posterior distribution p(x|y)
 - Average samples \approx MMSE
 - Sample with biggest probability \approx MAP

Further Optional Readings

- IVLE: Persi Diaconis, The Markov Chain Monte Carlo Revolution (MCMCRev.pdf)
- Another example of Metropolis algorithm: http://www.youtube.com/watch?v=Dzx5xNT79TI
- Search for terms on Wikipedia like "Markov Chain", "Metropolis-Hastings", "MCMC"

Additional Material

Why Stationary Distribution is $\pi(\mathcal{X})$ (1)

- First, show $\pi(x)T(x,x') = \pi(x')T(x',x)$
- If x' = x, obviously true
- If $x' \neq x$, two cases:
- If $\pi(x') > \pi(x)$

$$\pi(x)T(x,x') \qquad T(x,x') \text{ definition} \qquad \pi(x)T(x,x') \qquad T(x,x') \text{ definition}$$

$$= \pi(x)q(x'|x) \qquad q(x'|x) = q(x|x') \qquad = \pi(x)q(x'|x)\frac{\pi(x')}{\pi(x')}$$

$$= \pi(x')\frac{\pi(x)}{\pi(x')}q(x|x') \qquad Multiply\frac{\pi(x')}{\pi(x')} \qquad = q(x|x')\pi(x') \qquad q(x'|x) = q(x|x')$$

$$= \pi(x')T(x',x) \text{ definition} \qquad = \pi(x')T(x',x) \text{ definition}$$

$$= \pi(x')T(x',x') \text{ definition} \qquad = \pi(x')T(x',x) \text{ definition}$$

• If $\pi(x') < \pi(x)$

$$=\pi(x')\frac{\pi(x)}{\pi(x')}q(x|x')$$

$$=q(x|x')\pi(x')$$

$$=q(x|x')\pi(x')$$

$$q(x'|x) = q(x|x')$$

$$T(x'|x) \text{ definition}$$

$$=T(x',x)\pi(x')$$

$$T(x'|x) \text{ definition}$$

$$= \pi(x')T(x',x)$$

Property known as "detailed balance". Markov chain is considered "reversible"