

**Due Date: October 20, 2015**

**Problem 1:** Consider the Union-Find data structure discussed in the class. Suppose all *UNION* operations are performed before any of the *FIND* operations, and whenever *UNION*( $x, y$ ) is called,  $x, y$  are the leaders of their sets (so you don't have to find the leaders of  $x$  and  $y$ ). Describe an algorithm that performs *UNION* and *FIND* operations in  $O(1)$  amortized time.

**Problem 2:** Chicago has many tall buildings, but only some of them have a clear view of Lake Michigan. Suppose we are given an array  $A[1..n]$  that stores the height of  $n$  buildings on a city block, indexed from west to east. Building  $i$  has good view of Lake Michigan if and only if every building to the east of  $i$  is shorter than building  $i$ . Here is an algorithm that computes which buildings have a good view of Lake Michigan. What is the running time of this algorithm?

Algorithm 1: GOODVIEW( $A[1..n]$ )
<pre> initialize a stack S for <math>i = 1</math> to <math>n</math>     while (<math>S</math> not empty and <math>A[i] &gt; A[TOP(S)]</math>)         POP(<math>S</math>)     PUSH(<math>S, i</math>) return <math>S</math>                     </pre>

**Problem 3:** A vertex cover of a graph  $G = (V, E)$  is a subset of vertices  $S \subseteq V$  that includes at least one endpoint of every edge in  $E$ . Give a linear-time algorithm for the following task. Given an undirected tree  $T = (V, E)$ , return the size of the smallest vertex cover of  $T$ .

**Problem 4:** Suppose you are given an array  $A[1..n]$  of real numbers.

- Describe and analyze an algorithm that finds the largest sum of elements in a contiguous subarray  $A[i..j]$ .
- Describe and analyze an algorithm that finds the largest product of elements in a contiguous subarray  $A[i..j]$ .

For example, given the array  $A = [-6, 12, -7, 0, 14, -7, 5]$  as input, your first algorithm should return the integer 19 (sum of elements in  $A[2..5]$ ), and your second algorithm should return the integer 504 (product of elements in  $A[1..3]$ ).

**Problem 5:** Given an unlimited supply of coins of denominations  $x_1, x_2, \dots, x_n$ , we wish to make change for a value  $v$ , that is, we wish to find a set of coins whose total value is  $v$ . This might not be possible: for instance, if the denominations are 5 and 10 then we can make change for 15 but not for 12. Given the denominations  $x_1, x_2, \dots, x_n$  and  $v$ , find an  $O(nv)$  dynamic-programming algorithm to decide if it is possible to make change for  $v$  using coins of denominations  $x_1, \dots, x_n$ .