

2-group Belyi maps

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Outline

Motivation

Background

Computing permutation triples

A refined conjecture

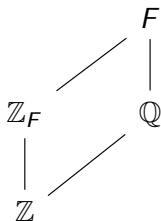
Computing equations

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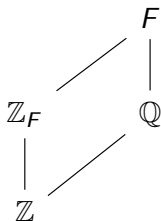
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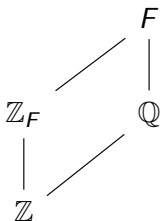
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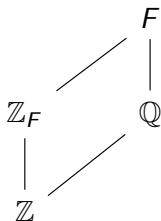


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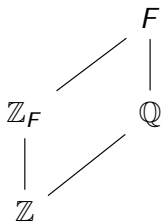
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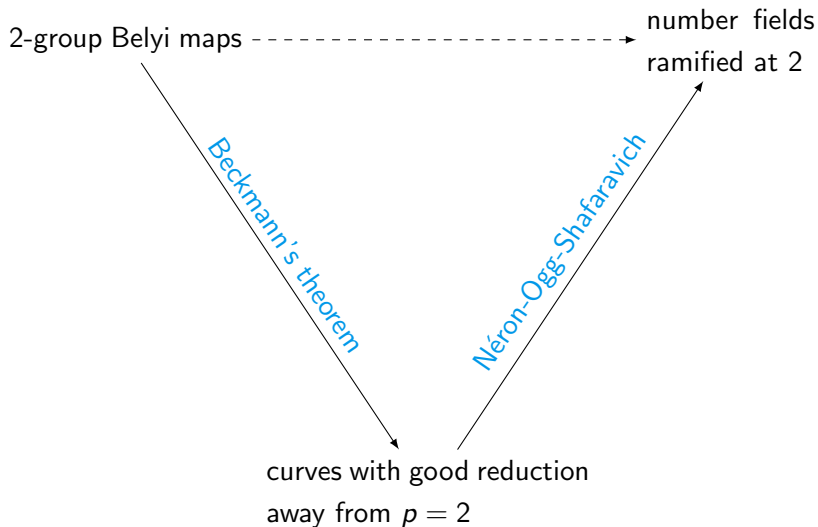
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Nonsolvable?

Why 2-group Belyi maps?



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The hope is that an explicit nonsolvable field ramified only at 2 can be obtained as $K(\text{Jac}(X)[2])$ where X is the domain of a **2-group Belyi map** (which we will define shortly).

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- implementation of an algorithm to enumerate isomorphism classes of 2-group Belyi maps
- implementation of an algorithm to compute equations for 2-group Belyi maps over finite fields
- implementation of a *method* to compute equations for 2-group Belyi maps over number fields

Main results

Motivated by the applications of 2-group Belyi maps to arithmetic geometry, we now state the main results.

- implementation of an algorithm to enumerate isomorphism classes of 2-group Belyi maps
- implementation of an algorithm to compute equations for 2-group Belyi maps over finite fields
- implementation of a *method* to compute equations for 2-group Belyi maps over number fields
- computational and theoretical evidence supporting a conjecture that every 2-group Belyi map is defined over an abelian extension of the rationals

Background



Belyi's theorem

A **Belyi map** is a morphism $\phi: X \rightarrow \mathbb{P}^1$ of smooth projective algebraic curves over \mathbb{C} that is unramified outside of $\{0, 1, \infty\}$.

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Theorem (Belyi 1979)

An algebraic curve (smooth projective) X over \mathbb{C} can be defined over a number field if and only if X admits a Belyi map.

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A **2-group Belyi map** is a Galois Belyi map with monodromy group a 2-group.

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*Let $\phi: X \rightarrow \mathbb{P}^1$ be a Galois Belyi map with monodromy group G .
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Then there exists a number field M satisfying the following properties.

- *p is unramified in M*
- *ϕ is defined over M*
- *X is defined over M*
- *X has good reduction at all primes \mathfrak{p} of M above p*

Permutation Triples

A **transitive permutation triple of degree d** is a triple

$$\sigma = (\sigma_0, \sigma_1, \sigma_\infty) \in S_d^3$$

such that

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The set of degree d Belyi maps up to isomorphism is in bijection with the set of degree d transitive permutation triples up to **simultaneous conjugation** and the group $\langle \sigma \rangle$ is the monodromy group of ϕ .

Passports

A **passport** \mathcal{P} consists of the data (g, G, λ) where $g \geq 0$ is an integer, $G \leq S_d$ is a transitive subgroup, and $\lambda = (\lambda_0, \lambda_1, \lambda_\infty)$ is a triple of partitions of d .

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The **passport of a permutation triple** σ is $(g(\sigma), \langle \sigma \rangle, \lambda(\sigma))$ where

$$g(\sigma) = 1 - d + (e(\sigma_0) - e(\sigma_1) - e(\sigma_\infty))/2$$

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$$e(\tau) = d - \#\text{cycles of } \tau,$$

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We now discuss the importance of organizing triples by passport. 10/43

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The situation improves, however, in the Galois setting.

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Let $\phi: X \rightarrow \mathbb{P}^1$ be a *Galois* Belyi map of degree d with monodromy group G corresponding to the permutation triple σ .

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Then

- ϕ and X are defined over $M(\phi)$,
- $\#G = d$,
- all cycles of σ_s have the same length for $s \in \{0, 1, \infty\}$,
- and if we let a, b, c be the orders of $\sigma_0, \sigma_1, \sigma_\infty$ respectively, we have

$$g(X) = 1 + \frac{\#G}{2} \left(1 - \frac{1}{a} - \frac{1}{b} - \frac{1}{c} \right).$$

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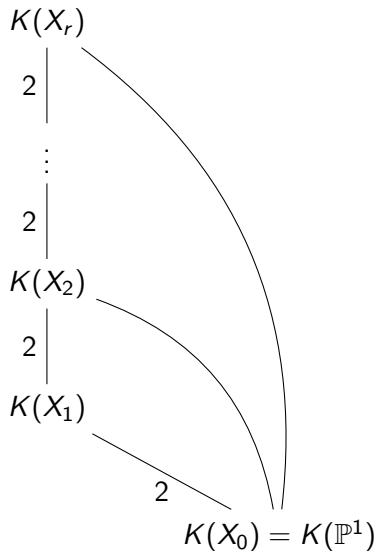
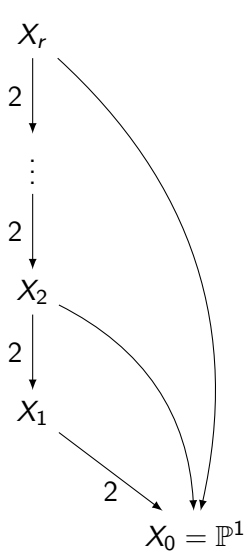
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The monodromy group in this setting corresponds to field automorphisms of the Galois closure of $K(X)$ fixing $K(x)$.

2-group Belyi maps as iterated quadratic extensions



Computing permutation triples



Setup

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A **2-group permutation triple** of degree $d \in \mathbb{Z}_{\geq 1}$ is a triple of permutations $\sigma := (\sigma_0, \sigma_1, \sigma_\infty) \in S_d^3$ satisfying

- $\sigma_\infty \sigma_1 \sigma_0 = \text{id}$;
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We say two degree d 2-group permutation triples σ, σ' are **simultaneously conjugate** if there exists $\tau \in S_d$ such that

$$\sigma^\tau := (\tau^{-1} \sigma_0 \tau, \tau^{-1} \sigma_1 \tau, \tau^{-1} \sigma_\infty \tau) = \sigma'$$

Lifting permutation triples

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A **lift** of σ is a 2-group permutation triple $\tilde{\sigma} \in S_{2d}^3$ such that $\langle \tilde{\sigma} \rangle$ is isomorphic to some extension \tilde{G} of $\mathbb{Z}/2\mathbb{Z}$ by G as in the exact sequence below.

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For a 2-group permutation triple σ , we denote the set of lifts of σ by $\text{Lifts}(\sigma)$ and $\text{Lifts}(\sigma)/\sim$ denotes the set of lifts up to simultaneous conjugation.

Algorithm to compute $\text{Lifts}(\sigma)/\sim$

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$$\text{Lifts}(\sigma) := \bigcup_{f \in H^2(G, A)} \text{Lifts}(\sigma, f)$$

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3. For each extension \tilde{G}_f compute the set $\text{Lifts}(\sigma, f)$ defined by
$$\left\{ \tilde{\sigma} : \tilde{\sigma}_s \in \pi_f^{-1}(\sigma_s) \text{ for } s \in \{0, 1, \infty\}, \tilde{\sigma}_\infty \tilde{\sigma}_1 \tilde{\sigma}_0 = 1, \langle \tilde{\sigma} \rangle = \tilde{G}_f \right\}$$
- 4.

$$\text{Lifts}(\sigma) := \bigcup_{f \in H^2(G, A)} \text{Lifts}(\sigma, f)$$

5. Quotient $\text{Lifts}(\sigma)$ by simultaneous conjugation

Example computing $\text{Lifts}(\sigma)/\sim$: setup

Let $\sigma = ((1\ 2), \text{id}, (1\ 2))$. Then $G = \langle \sigma \rangle = \mathbb{Z}/2\mathbb{Z}$.

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$$\begin{aligned} T_1 = \Big\{ & ((1\ 2)(3\ 4), \text{id}, (1\ 2)(3\ 4)), ((1\ 2)(3\ 4), (1\ 3)(2\ 4), (1\ 4)(2\ 3)), \\ & ((1\ 4)(2\ 3), \text{id}, (1\ 4)(2\ 3)), ((1\ 4)(2\ 3), (1\ 3)(2\ 4), (1\ 2)(3\ 4)) \Big\} \end{aligned}$$

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Let $(\tilde{\sigma}_0, \tilde{\sigma}_1, \tilde{\sigma}_\infty) = ((1\ 2)(3\ 4), (1\ 3)(2\ 4), (1\ 4)(2\ 3))$.

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The induced permutation of $\tilde{\sigma}_0$ on blocks is $(\boxed{1\,3}, \boxed{2\,4})$ which is the same as the permutation $\sigma_0 = (1\,2)$.

Example computing $\text{Lifts}(\sigma)/\sim$: action on blocks

Choose $\alpha = (13)(24)$ to be the generator of $\iota_1(\mathbb{Z}/2\mathbb{Z})$ in \tilde{G}_1 .

Each triple in T_1 must act on the *blocks* $\{\boxed{13}, \boxed{24}\}$ corresponding to the permutations in σ .

Let $(\tilde{\sigma}_0, \tilde{\sigma}_1, \tilde{\sigma}_\infty) = ((12)(34), (13)(24), (14)(23))$.

Note that $\tilde{\sigma}_0(\boxed{13}) = \boxed{24}$ and $\tilde{\sigma}_0(\boxed{24}) = \boxed{13}$.

The induced permutation of $\tilde{\sigma}_0$ on blocks is $(\boxed{13}, \boxed{24})$ which is the same as the permutation $\sigma_0 = (12)$.

Similarly, $\tilde{\sigma}_1$ acts as id on blocks and $\tilde{\sigma}_\infty$ acts as (12) on blocks.

Example computing $\text{Lifts}(\sigma)/\sim$: action on blocks

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Similarly, $\tilde{\sigma}_1$ acts as id on blocks and $\tilde{\sigma}_\infty$ acts as $(1\,2)$ on blocks.

Choosing

$$\alpha := (1\,d+1)(2\,d+2)\dots(d-1\,2d-1)(d\,2d)$$

allows us to label blocks by reducing modulo d .

Example computing $\text{Lifts}(\sigma)/\sim$: conclude

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Lastly, we quotient by simultaneous conjugation to obtain

$$\text{Lifts}(\sigma)/\sim = \left\{ ((1\ 2)(3\ 4), (1\ 3)(2\ 4), (1\ 4)(2\ 3)), \right. \\ \left. ((1\ 4\ 3\ 2), \text{id}, (1\ 2\ 3\ 4)), ((1\ 2\ 3\ 4), (1\ 3)(2\ 4), (1\ 2\ 3\ 4)) \right\}$$

Bipartite graphs of permutation triples

Now that we can lift permutation triples, we now describe some notation for the bipartite graphs that organize these triples.

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For $i \in \mathbb{Z}_{\geq 1}$ we define the bipartite graph denoted \mathcal{G}_{2^i} with the following node sets.

- $\mathcal{G}_{2^i}^{\text{above}}$: the set of isomorphism classes of 2-group Belyi maps of degree 2^i indexed by 2-group permutation triples $\tilde{\sigma}$ up to simultaneous conjugation in S_{2^i}
- $\mathcal{G}_{2^i}^{\text{below}}$: the set of isomorphism classes of 2-group Belyi maps of degree 2^{i-1} indexed by 2-group permutation triples σ up to simultaneous conjugation in $S_{2^{i-1}}$

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For every pair of nodes $(\tilde{\sigma}, \sigma) \in \mathcal{G}_{2^i}^{\text{above}} \times \mathcal{G}_{2^i}^{\text{below}}$ there is an edge between σ and $\tilde{\sigma}$ if and only if $\tilde{\sigma}$ is simultaneously conjugate to a lift of σ .

Algorithm to compute \mathcal{G}_{2^i}

Input: The bipartite graph $\mathcal{G}_{2^{i-1}}$

Output: The bipartite graph \mathcal{G}_{2^i}

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4. For every pair $(\tilde{\sigma}, \sigma) \in \mathcal{G}_{2^i}^{\text{above}} \times \mathcal{G}_{2^i}^{\text{below}}$ place an edge between $\tilde{\sigma}$ and σ if and only if there is a triple in the equivalence class $[\tilde{\sigma}] \in \text{Lifts}(\mathcal{G}_{2^{i-1}})/\sim$ that is a lift of σ

Results : number of triples and passports

Theorem (M.)

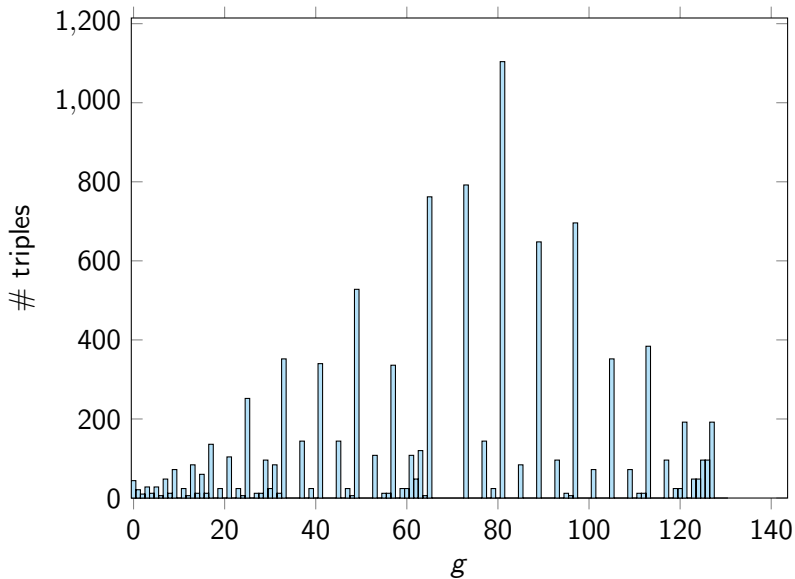
The following tables list the number of isomorphism classes of 2-group Belyi maps, the number of passports, and number of lax passports respectively up to degree 256.

d	1	2	4	8	16	32	64	128	256
# triples	1	3	7	19	55	151	503	1799	7175

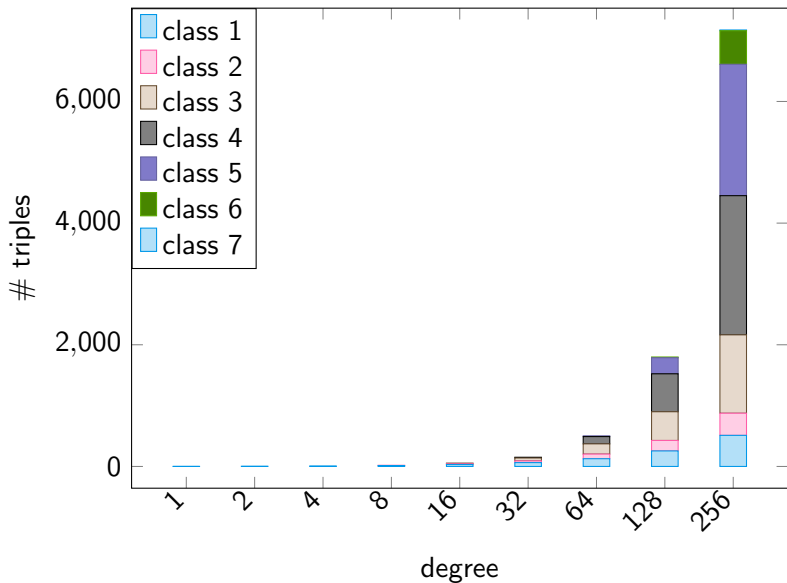
d	1	2	4	8	16	32	64	128	256
# passports	1	3	7	16	41	96	267	834	2893

d	1	2	4	8	16	32	64	128	256
# lax passports	1	1	3	6	14	31	85	257	882

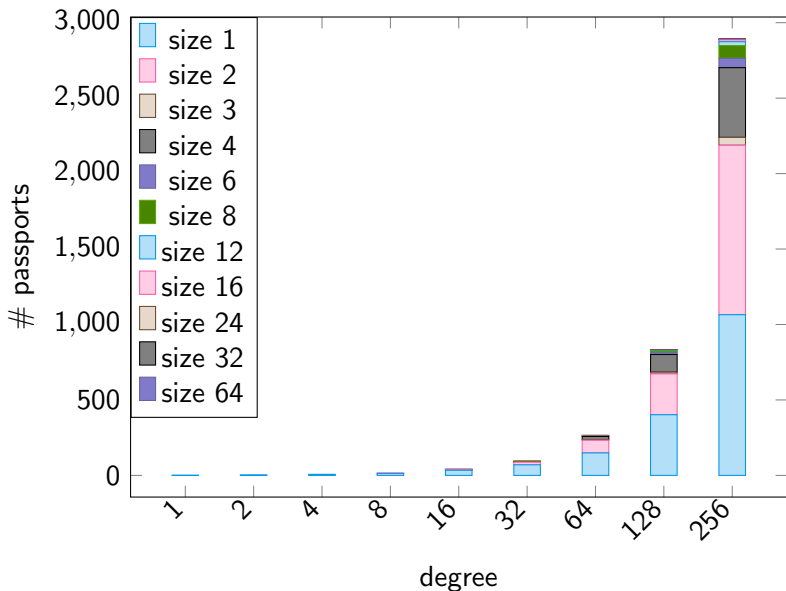
Results : distribution of genera



Results : groups by nilpotency class



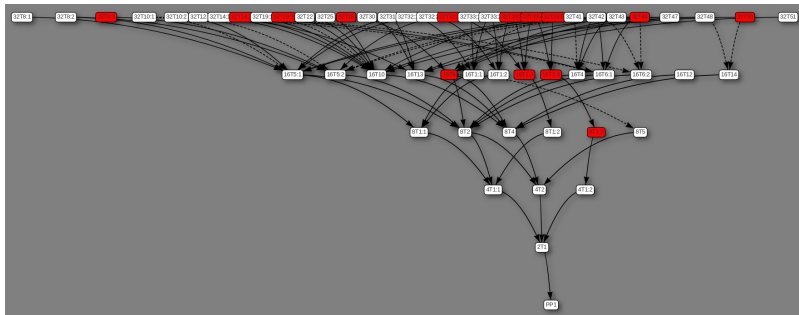
Results : passport sizes



The graph of 2-group Belyi maps

<https://math.dartmouth.edu/~mjmusty/32.html>

<https://math.dartmouth.edu/~mjmusty/32nh.html>



A refined conjecture



Passports

Recall that a passport \mathcal{P} consists of the data (g, G, λ) where $g \in \mathbb{Z}_{\geq 0}$, G is a transitive subgroup of S_d and $\lambda = (\lambda_0, \lambda_1, \lambda_\infty)$ is a triple of partitions of d corresponding to conjugacy classes (C_0, C_1, C_∞) of S_d .

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The size of \mathcal{P} is the cardinality of the set $\Sigma_{\mathcal{P}}$ defined by

$$\left\{ (\sigma_0, \sigma_1, \sigma_\infty) \in C_0 \times C_1 \times C_\infty : \sigma_\infty \sigma_1 \sigma_0 = 1 \text{ and } \langle \sigma_0, \sigma_1 \rangle = G \right\} / \sim$$

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As a result of the action of $G_{\mathbb{Q}}$ on \mathcal{P} , the size of \mathcal{P} bounds the degree of the field of moduli of any Belyi map with passport \mathcal{P} .

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To instead analyze $\text{Gal}(\mathbb{Q}^{\text{al}} | \mathbb{Q}^{\text{ab}})$ we *refine* the notion of a passport.

Refined passports

A **refined passport** \mathcal{P} consists of the data (g, G, c) where $g \in \mathbb{Z}_{\geq 0}$, G is a transitive subgroup of S_d and $c = (c_0, c_1, c_\infty)$ is a triple of conjugacy classes of G .

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where $(\sigma_0, \sigma_1, \sigma_\infty) \sim (\sigma'_0, \sigma'_1, \sigma'_\infty)$ if there exists $\alpha \in \text{Aut}(G)$ with $\alpha(\sigma_s) = \sigma'_s$ for every $s \in \{0, 1, \infty\}$.

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As was the case with passport, every permutation triple σ determines a refined passport $\mathcal{P}(\sigma)$.

A refined conjecture

Theorem (M.)

The size of $\mathcal{P}(\sigma)$ is equal to 1 for every 2-group permutation triple σ with degree ≤ 256 .

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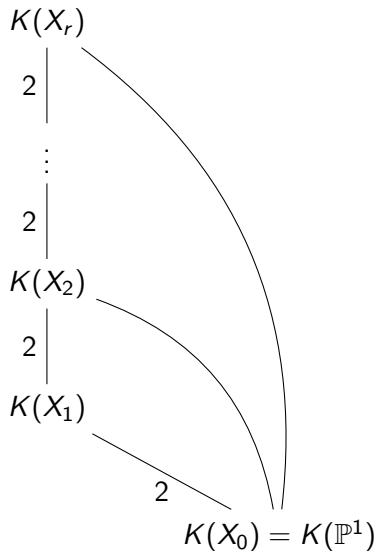
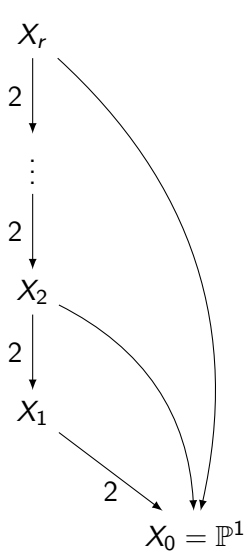
Theorem (M.)

ARC is true for 2-group permutation triples σ with $\langle \sigma \rangle$ dihedral.

Computing equations



2-group Belyi maps as iterated quadratic extensions



A motivating example : setup

Let F be a number field with integers \mathbb{Z}_F . Let $\text{Pl}(F)$ denote the places of F and S_∞ the archimedean places. For $v \in \text{Pl}(F) \setminus S_\infty$ let \mathfrak{p}_v denote the prime ideal of \mathbb{Z}_F corresponding to v .

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Question

How do we construct a quadratic extension of F with ramification prescribed by \mathfrak{a} ?

First, it is possible that no such extension exists.

A motivating example : setup

Let F be a number field with integers \mathbb{Z}_F . Let $\text{Pl}(F)$ denote the places of F and S_∞ the archimedean places. For $v \in \text{Pl}(F) \setminus S_\infty$ let \mathfrak{p}_v denote the prime ideal of \mathbb{Z}_F corresponding to v .

Let $S \subset \text{Pl}(F) \setminus S_\infty$ and let $\mathfrak{a} := \prod_{v \in S} \mathfrak{p}_v$.

Question

How do we construct a quadratic extension of F with ramification prescribed by \mathfrak{a} ?

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If \mathfrak{a} is not principal, then the question requires more care.

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To summarize, in the case where \mathfrak{a} is not principal but there exists \mathfrak{b} with $\mathfrak{a}\mathfrak{b}^2$ principal we have $[\mathfrak{a}] \in \text{Cl}_F^2$ and $[\mathfrak{b}]$ is unique up to multiplication by $[\mathfrak{c}] \in \text{Cl}_F[2]$.

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The algorithms in this section rely on transporting this technique to the function field setting.

Algorithm in characteristic $p \geq 3$: setup

We want to do computations modulo primes, so we need to think about Belyi maps not over \mathbb{C} or \mathbb{Q}^{al} , but over finite fields in an analogous way.

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A **2-group Belyi map modulo q** is a Galois extension of function fields $\mathbb{F}_q(x) \hookrightarrow F$ with $[F : \mathbb{F}_q(x)]$ a power of 2 unramified outside $\{0, 1, \infty\}$ with $p \geq 3$.

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7. Lift automorphisms to identify the passports of the Galois extensions $F(\sqrt{f_a})$.

Notation for examples

DNG-a,b,c-gE-H

D : degree in $\{2, 4, 8, 16, 32, 64, 128, 256\}$

N : either T or S identifying group database

G : a positive integer identifying the group

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To identify a passport we omit the hash.

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$F(\sqrt{f}) = \mathbb{F}_3(x)(\alpha)$ where α satisfies

$$\begin{aligned} & y^{16} + (x^7 + 2x^6 + x^5)y^{14} + (x^{16} + 2x^{14} + x^{13} + 2x^{11})y^{12} \\ & + (x^{23} + x^{21} + 2x^{20} + 2x^{18} + x^{17} + x^{15})y^{10} \\ & + (x^{31} + 2x^{30} + 2x^{29} + 2x^{28} + x^{27} + 2x^{26} + x^{25} + 2x^{24} + x^{23} + x^{20})y^8 \\ & + (2x^{38} + 2x^{36} + x^{35} + 2x^{29} + 2x^{27} + x^{26})y^6 \\ & + (x^{48} + 2x^{47} + 2x^{46} + 2x^{44} + 2x^{43} + 2x^{42} \\ & \quad + x^{39} + 2x^{38} + 2x^{37} + 2x^{35} + 2x^{34} + 2x^{33})y^4 \\ & + (x^{54} + x^{52} + x^{51} + x^{50} + x^{49} + x^{47} \\ & \quad + x^{45} + x^{43} + x^{42} + x^{41} + x^{40} + x^{38})y^2 \\ & + x^{64} + 2x^{62} + x^{61} + 2x^{59} + 2x^{58} + x^{56} \\ & + x^{52} + 2x^{50} + 2x^{49} + x^{47} + 2x^{46} + x^{44} \end{aligned}$$

Algorithm in characteristic $p \geq 3$: compute entire passport

One issue with our technique to compute 2-group Belyi maps is that it only guarantees that the resulting Belyi map has the desired *passport* and does not allow us to control which *isomorphism class* we get.

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To recover from this we use isomorphism testing of function fields to determine if we have redundant Belyi maps with a given passport.

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Since we know the sizes of passports from our work with permutation triples, we know that we have representatives from every isomorphism class even if we cannot match the Belyi maps to their corresponding permutation triples.

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Although this implementation does not allow us to compute all 2-group Belyi maps for a given degree, it does work well in practice.

<https://github.com/michaelmusty/2GroupDessins>

- *all* 2-group Belyi maps modulo 3 up to degree 32
- hundreds of 2-group Belyi maps up to degree 256

Future work

- higher degree over \mathbb{F}_3
- non-Galois setting
- an algorithm in characteristic zero
- prove ARC for other families of 2-groups
- p -group Belyi maps for p odd
- compute torsion fields

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- Sam, Jeroen, Edgar, Florian, and Richard
- Mary, Jim, Matt, and Nicole

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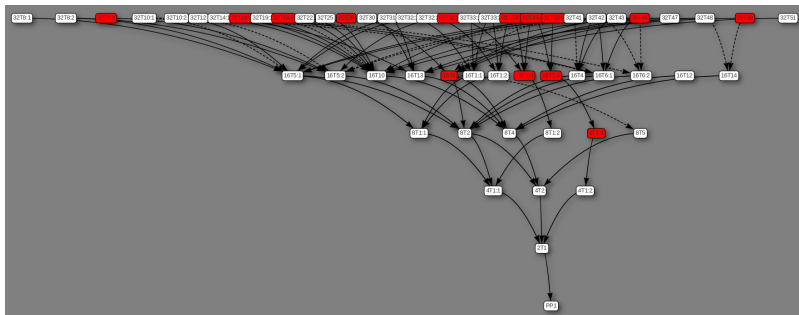
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Thanks for listening!

The graph of 2-group Belyi maps

<https://math.dartmouth.edu/~mjmusty/32.html>

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Galois representations

Let X be an irreducible, smooth projective algebraic curve of genus $g \geq 1$ over a number field K . Let $G_K := \text{Gal}(K^{\text{al}} | K)$ be the absolute Galois group of K and let $\ell \in \mathbb{Z}$ be prime.

Let $J := \text{Jac}(X)$ be the **Jacobian variety** of X . J is an abelian variety of dimension g .

G_K acts on the ℓ -torsion points $J[\ell](K^{\text{al}}) \cong (\mathbb{Z}/\ell\mathbb{Z})^{2g}$ of X .

This action determines a **mod- ℓ Galois representation**

$$\rho: G_K \rightarrow \text{Aut}(J[\ell]) \cong \text{GL}_{2g}(\mathbb{Z}/\ell\mathbb{Z}).$$

The geometry of X and the arithmetic of ρ are intimately related. For example, if X has good reduction at a prime \mathfrak{p} above $p \neq \ell$, then \mathfrak{p} will be unramified in the **ℓ -torsion field** $K(J[\ell])$.

Isomorphism of Belyi maps

Let $\phi: X \rightarrow \mathbb{P}^1$ and $\phi': X' \rightarrow \mathbb{P}^1$ be Belyi maps of degree d . ϕ and ϕ' are **isomorphic** (respectively **lax isomorphic**) if the diagrams

$$\begin{array}{ccc} X & \xrightarrow{\sim} & X' \\ & \searrow \phi & \swarrow \phi' \\ & \mathbb{P}^1 & \end{array}, \text{ respectively } \begin{array}{ccc} X & \xrightarrow{\sim} & X' \\ \phi \downarrow & & \downarrow \phi' \\ \mathbb{P}^1 & \xrightarrow[\beta]{\sim} & \mathbb{P}^1 \end{array}$$

commute where $\beta(\{0, 1, \infty\}) = \{0, 1, \infty\}$.

Algebraic function fields : setup

Let K be a perfect field. An **algebraic function field in one variable over K** is a field extension F over K such that there exists $x \in F$ transcendental over K and $[F : K(x)]$ is finite.

K is the **constant field** of F and the **exact constant field** of F is the algebraic closure of K in F .

As an example, let X be an irreducible affine plane curve (possibly singular) defined by the equation $f(x, y) = 0$ with $f \in K[x, y]$.

Then the **function field of X** , denoted $K(X)$ is the field of fractions of the coordinate ring $K[x, y]/(f(x, y))$ of X .

A **place** of F is the maximal ideal of some DVR \mathcal{O}_P of F . The valuation on F corresponding to P is denoted ord_P .

The set of places of F is denoted $\text{Pl}(F)$ and the **degree** of P is the index $[\mathcal{O}_P/P : K]$ of the **residue class field**.

Algebraic function fields : Picard group and $\mathcal{L}(D)$

The **divisor class group** $\text{Div}(F)$ of F is the free abelian group generated by the places of F . A **divisor** $D \in \text{Div}(F)$ is represented by a sum of places $\sum_P a_P P$ and the **degree** of D is $\sum_P a_P \deg(P)$. The set of **degree zero divisors** is denoted $\text{Div}^0(F)$.

The image of $\text{div}: F^\times \rightarrow \text{Div}(F)$ defined by $\text{div}(f) = \sum_P \text{ord}_P(f) P$ is the subgroup of **principal divisors** of F denoted $\text{Princ}(F)$.

The **Picard group** of F is $\text{Pic}(F) := \text{Div}(F) / \text{Princ}(F)$.

The **Jacobian** of F is $\text{Pic}^0(F) := \text{Div}^0(F) / \text{Princ}(F)$.

There is a partial order on $\text{Div}(F)$ defined by $D_1 \geq D_2$ if $\text{ord}_P(D_1) \geq \text{ord}_P(D_2)$ for every $P \in \text{Pl}(F)$.

The **Riemann-Roch space** of a divisor $D \in \text{Div}(F)$ is defined by $\mathcal{L}(D) := \{f \in F : \text{div}(f) + D \geq 0\} \cup \{0\}$.

Algebraic function fields : quadratic extensions

Lemma

Let $aF^{\times 2}$ be a nontrivial coset of $F^{\times}/F^{\times 2}$ and consider the extension $L := F(\sqrt{a})$. Then a prime P of F is ramified in L if and only if $\text{ord}_P(a)$ is odd.

We now revisit the question of finding a quadratic extension of $F(\sqrt{a})$ of F with ramification prescribed by $R \in \text{Div}(F)$.

The simple cases are when no such a exists or when R is principal. The last case occurs when there exists D such that $R - 2D = \text{div}(a)$ for some $a \in F$.

As in the number field setting, this implies $R \in 2\text{Pic}(F)$ and D is unique up to addition by $T \in \text{Pic}^0(F)[2]$.

Algorithm in characteristic $p \geq 3$: Galois test

Input:

- F a Galois extension of $\mathbb{F}_q(x)$
- $\text{Gal}(F | \mathbb{F}_q(x))$ explicitly given as automorphisms of F
- $a \in F$

Output: True if $F(\sqrt{a})$ is Galois over $\mathbb{F}_q(x)$ and False otherwise

- For each generator $\sigma \in \text{Gal}(F | \mathbb{F}_q(x))$ test if $\sigma(a)/a$ is a square in F
- Return True if $\sigma(a)/a$ is a square in F for all generators σ and otherwise return False

Similarly, we can apply the same test after extending the constant field from \mathbb{F}_q to \mathbb{F}_{q^2} .

Algorithm in characteristic $p \geq 3$: get candidates

Input:

- F a 2-group Belyi map modulo q of degree $d = 2^m$ corresponding to a 2-group permutation triple σ
- A passport $\mathcal{P} = (\tilde{G}, (a, b, c))$ with \tilde{G} a 2-group of order $2d$ such that there exists a 2-group permutation triple $\tilde{\sigma}$ with passport \mathcal{P} that is a lift of σ
- $\text{Gal}(F \mid \mathbb{F}_q(x)) \cong \langle \sigma \rangle$ explicitly given as automorphisms of F

Output: A list of candidate functions $\{f_i\}$ with each $f_i \in F$ such that $F(\sqrt{f_i})$ is a 2-group Belyi map modulo q with passport \mathcal{P} .

Algorithm in characteristic $p \geq 3$: get candidates (steps 1-4)

1. For $s \in \{0, 1, \infty\}$ compute

$$r_s := \begin{cases} 0 & \text{if } \text{order}(\sigma_s) = \text{order}(\tilde{\sigma}_s) \\ 1 & \text{if } \text{order}(\sigma_s) < \text{order}(\tilde{\sigma}_s) \end{cases}$$

2. Compute

$$R := \sum_{s \in \{0, 1, \infty\}} r_s R_s \in \text{Div}(F)$$

where R_0, R_1, R_∞ are defined to be the supports of $\text{div}(x)$, $\text{div}(x - 1)$, and $\text{div}(1/x)$ respectively.

3. Compute the abelian group $\text{Pic}(F) = T \oplus \mathbb{Z}$ (with T a finite abelian group) along with a map $\psi: \text{Pic}(F) \rightarrow \text{Div}(F)$.
4. Compute $[R] := \psi^{-1}(R)$.

Algorithm in characteristic $p \geq 3$: get candidates (step 5)

5. For each $a \in \text{Pic}(F)[2]$ compute the following:
 - (a) Let $D_a := \psi(a + [R]/2) \in \text{Div}(F)$.
 - (b) Compute $\mathcal{L}(R - 2D_a)$.
 - (c) If $\mathcal{L}(R - 2D_a)$ has dimension 1, then compute $f_a \in F$ with $\text{div}(f_a)$ generating $\mathcal{L}(R - 2D_a)$ and go to Step 5d Otherwise go to the next $a \in \text{Pic}(F)[2]$.
 - (d) Apply Galois test to F , $\text{Gal}(F | \mathbb{F}_q(x))$, and f_a from Step 5c to see if $F(\sqrt{f_a})$ generates a Galois extension. If $F(\sqrt{f_a})$ is Galois over $\mathbb{F}_q(x)$ then save f_a and go to the next $a \in \text{Pic}(F)[2]$. If $F(\sqrt{f_a})$ is not Galois over $\mathbb{F}_q(x)$, then go to Step 5e.
 - (e) Let F' be the function field F after extending the field of constants \mathbb{F}_q to \mathbb{F}_{q^2} . Apply Galois test to F' , $\text{Gal}(F' | \mathbb{F}_{q^2}(x))$, and f_a (viewed as an element of F') from Step 5c to see if $F'(\sqrt{f_a})$ generates a Galois extension. If $F'(\sqrt{f_a})$ is Galois over $\mathbb{F}_{q^2}(x)$ then save f_a . Go to the next $a \in \text{Pic}(F)[2]$.

Algorithm in characteristic $p \geq 3$: get candidates (steps 6-8)

6. Let S be the set of f_a saved in Step 5d. Let S' be the set of f_a saved in Step 5e.
7. • If S is nonempty, then for each $f_a \in S$ compute $F(\sqrt{f_a})$,

$$G_a \cong \text{Gal}(F(\sqrt{f_a}) | \mathbb{F}_q(x)),$$

and let $S'' = \{f_a \in S : G_a \cong \tilde{G}\}$.

- If S is empty, then for each $f_a \in S'$ compute $F'(\sqrt{f_a})$,

$$G_a \cong \text{Gal}(F'(\sqrt{f_a}) | \mathbb{F}_{q^2}(x)),$$

and let $S'' = \{f_a \in S' : G_a \cong \tilde{G}\}$.

8. Return the list S''