

2-group Belyi maps

Michael Musty

July 9, 2019

Dartmouth College

Acknowledgements

- Dave, Tom, Carl, and John
- Sam, Anna, Jeroen, Edgar, Florian, and Richard
- Mary, Jim, Matt, and Nicole

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Motivation



Galois representations

Let X be an irreducible, smooth projective algebraic curve of genus $g \geq 1$ over a number field K . Let $G_K := \text{Gal}(K^{\text{al}} | K)$ be the absolute Galois group of K and let $\ell \in \mathbb{Z}$ be prime.

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The geometry of X and the arithmetic of ρ are intimately related. For example, if X has good reduction at a prime \mathfrak{p} above $p \neq \ell$, then \mathfrak{p} will be unramified in the **ℓ -torsion field** $K(J[\ell])$.

Belyi's theorem

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Theorem (Belyi 1979)

An algebraic curve (smooth projective) X over \mathbb{C} can be defined over a number field if and only if X admits a Belyi map.

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We can now state Beckmann's theorem to relate Belyi maps to our original motivation.

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Theorem (Beckmann 1989)

*Let $\phi: X \rightarrow \mathbb{P}^1$ be a Galois Belyi map with monodromy group G .
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Let p be a prime not dividing $\#G$.*

Then there exists a number field M satisfying the following properties.

- *p is unramified in M*
- *ϕ is defined over M*
- *X is defined over M*
- *X has good reduction at all primes \mathfrak{p} of M above p*

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Conjecture (Gross 1998)

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The hope is that an explicit nonsolvable field ramified only at 2 can be obtained as $K(\text{Jac}(X)[2])$ where X is the domain of a Galois Belyi map with monodromy group a 2-group.

We call these Belyi maps **2-group Belyi maps**.

Main results

Motivated by the applications of 2-group Belyi maps to arithmetic geometry, we now state the main results.

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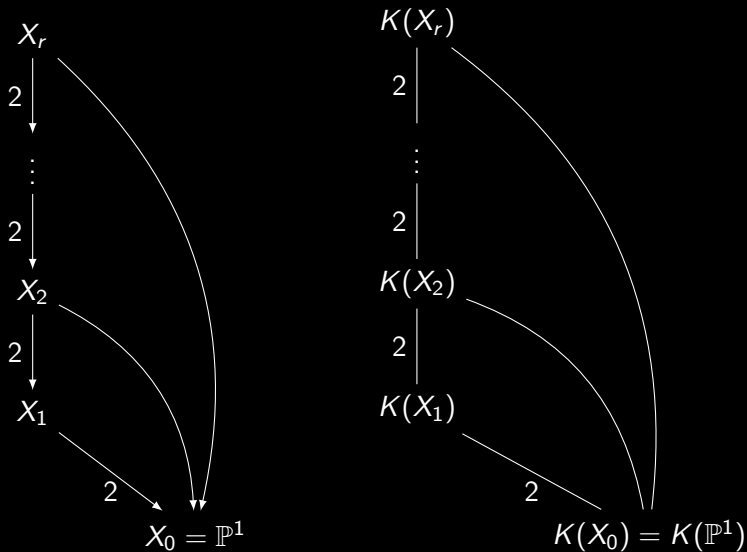
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Motivated by the applications of 2-group Belyi maps to arithmetic geometry, we now state the main results.

- implementation of an algorithm to enumerate isomorphism classes of 2-group Belyi maps
- implementation of an algorithm to compute equations for 2-group Belyi maps over finite fields
- implementation of a *method* to compute equations for 2-group Belyi maps over number fields
- computational and theoretical evidence supporting a conjecture that every 2-group Belyi map is defined over an abelian extension of the rationals

2-group Belyi maps as iterated quadratic extensions



Background



Isomorphism of Belyi maps

Let $\phi: X \rightarrow \mathbb{P}^1$ and $\phi': X' \rightarrow \mathbb{P}^1$ be Belyi maps of degree d .

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Let $\phi: X \rightarrow \mathbb{P}^1$ and $\phi': X' \rightarrow \mathbb{P}^1$ be Belyi maps of degree d . ϕ and ϕ' are **isomorphic** (respectively **lax isomorphic**) if the diagrams

$$\begin{array}{ccc} X & \xrightarrow{\sim} & X' \\ & \searrow \phi & \swarrow \phi' \\ & \mathbb{P}^1 & \end{array}, \text{ respectively } \begin{array}{ccc} X & \xrightarrow{\sim} & X' \\ \phi \downarrow & & \downarrow \phi' \\ \mathbb{P}^1 & \xrightarrow[\beta]{\sim} & \mathbb{P}^1 \end{array}$$

commute where $\beta(\{0, 1, \infty\}) = \{0, 1, \infty\}$.

Permutation Triples

A **transitive permutation triple of degree d** is a triple

$$\sigma = (\sigma_0, \sigma_1, \sigma_\infty) \in S_d^3$$

such that

- $\sigma_\infty \sigma_1 \sigma_0 = 1$
- σ generates a transitive subgroup of S_d

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The set of degree d Belyi maps up to isomorphism is in bijection with the set of degree d transitive permutation triples up to **simultaneous conjugation** and the group $\langle \sigma \rangle$ is the monodromy group of ϕ .

Passports

A **passport** \mathcal{P} consists of the data (g, G, λ) where $g \geq 0$ is an integer, $G \leq S_d$ is a transitive subgroup, and $\lambda = (\lambda_0, \lambda_1, \lambda_\infty)$ is a triple of partitions of d .

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The **passport of a Belyi map** $\phi : X \rightarrow \mathbb{P}^1$ is $(g(X), \text{Mon}(\phi), (\lambda_0, \lambda_1, \lambda_\infty))$ with $g(X)$ the genus of X , $\text{Mon}(\phi)$ the monodromy group of ϕ , and the partitions from ramification.

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The **passport of a permutation triple** σ is $(g(\sigma), \langle \sigma \rangle, \lambda(\sigma))$ where

$$g(\sigma) = 1 - d + (e(\sigma_0) - e(\sigma_1) - e(\sigma_\infty))/2$$

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$$e(\tau) = d - \#\text{cycles of } \tau,$$

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We now discuss the importance of organizing triples by passport. 12/46

Fields of moduli, fields of definition, and passports

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The situation improves, however, in the Galois setting...

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Then

- ϕ and X are defined over $M(\phi)$,
- $\#G = d$,
- all cycles of σ_s have the same length for $s \in \{0, 1, \infty\}$,
- and if we let a, b, c be the orders of $\sigma_0, \sigma_1, \sigma_\infty$ respectively, we have

$$g(X) = 1 + \frac{\#G}{2} \left(1 - \frac{1}{a} - \frac{1}{b} - \frac{1}{c} \right).$$

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The monodromy group in this setting corresponds to field automorphisms of the Galois closure of $K(X)$ fixing $K(x)$.

Computing permutation triples



Setup

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- $\sigma_\infty \sigma_1 \sigma_0 = \text{id}$;
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We say two degree d 2-group permutation triples σ, σ' are **simultaneously conjugate** if there exists $\tau \in S_d$ such that

$$\sigma^\tau := (\tau^{-1} \sigma_0 \tau, \tau^{-1} \sigma_1 \tau, \tau^{-1} \sigma_\infty \tau) = \sigma'$$

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For a 2-group permutation triple σ , we denote the set of lifts of σ by $\text{Lifts}(\sigma)$ and $\text{Lifts}(\sigma)/\sim$ denotes the set of lifts up to simultaneous conjugation.

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Choosing

$$\alpha := (1\ d+1)(2\ d+2) \dots (d-1\ 2d-1)(d\ 2d)$$

allows us to label blocks by reducing modulo d .

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Lastly, we quotient by simultaneous conjugation to obtain

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Bipartite graphs of permutation triples

Now that we can lift permutation triples, we now describe some notation for the bipartite graphs that organize these triples.

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For $i \in \mathbb{Z}_{\geq 1}$ we define the bipartite graph denoted \mathcal{G}_{2^i} with the following node sets.

- $\mathcal{G}_{2^i}^{\text{above}}$: the set of isomorphism classes of 2-group Belyi maps of degree 2^i indexed by 2-group permutation triples $\tilde{\sigma}$ up to simultaneous conjugation in S_{2^i}
- $\mathcal{G}_{2^i}^{\text{below}}$: the set of isomorphism classes of 2-group Belyi maps of degree 2^{i-1} indexed by 2-group permutation triples σ up to simultaneous conjugation in $S_{2^{i-1}}$

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For every pair of nodes $(\tilde{\sigma}, \sigma) \in \mathcal{G}_{2^i}^{\text{above}} \times \mathcal{G}_{2^i}^{\text{below}}$ there is an edge between σ and $\tilde{\sigma}$ if and only if $\tilde{\sigma}$ is simultaneously conjugate to a lift of σ .

Algorithm to compute \mathcal{G}_{2^i}

Input: The bipartite graph $\mathcal{G}_{2^{i-1}}$

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3. Define $\mathcal{G}_{2^i}^{\text{below}} := \mathcal{G}_{2^{i-1}}^{\text{above}}$ and define $\mathcal{G}_{2^i}^{\text{above}}$ by representatives of $\text{Lifts}(\mathcal{G}_{2^{i-1}})/\sim$

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4. For every pair $(\tilde{\sigma}, \sigma) \in \mathcal{G}_{2^i}^{\text{above}} \times \mathcal{G}_{2^i}^{\text{below}}$ place an edge between $\tilde{\sigma}$ and σ if and only if there is a triple in the equivalence class $[\tilde{\sigma}] \in \text{Lifts}(\mathcal{G}_{2^{i-1}})/\sim$ that is a lift of σ

Results : number of triples and passports

Theorem (M.)

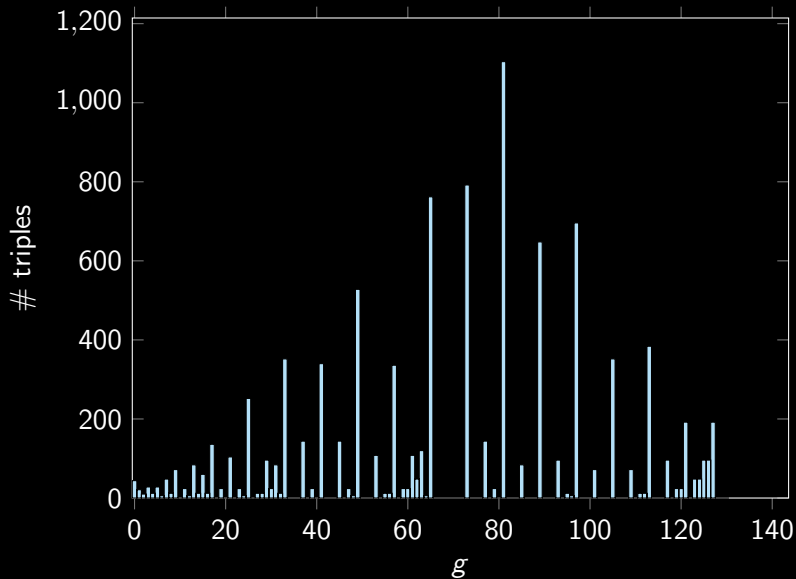
The following tables list the number of isomorphism classes of 2-group Belyi maps, the number of passports, and number of lax passports respectively up to degree 256.

| d | 1 | 2 | 4 | 8 | 16 | 32 | 64 | 128 | 256 |
|-----------|---|---|---|----|----|-----|-----|------|------|
| # triples | 1 | 3 | 7 | 19 | 55 | 151 | 503 | 1799 | 7175 |

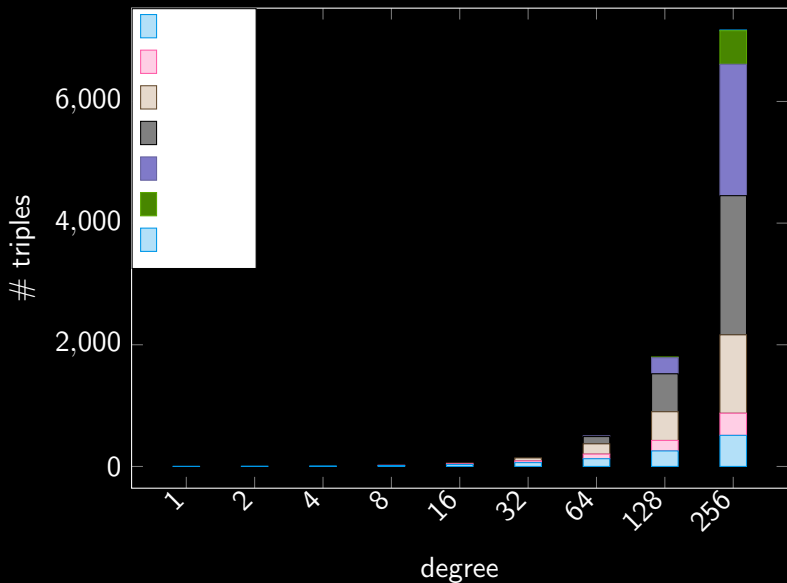
| d | 1 | 2 | 4 | 8 | 16 | 32 | 64 | 128 | 256 |
|-------------|---|---|---|----|----|----|-----|-----|------|
| # passports | 1 | 3 | 7 | 16 | 41 | 96 | 267 | 834 | 2893 |

| d | 1 | 2 | 4 | 8 | 16 | 32 | 64 | 128 | 256 |
|-----------------|---|---|---|---|----|----|----|-----|-----|
| # lax passports | 1 | 1 | 3 | 6 | 14 | 31 | 85 | 257 | 882 |

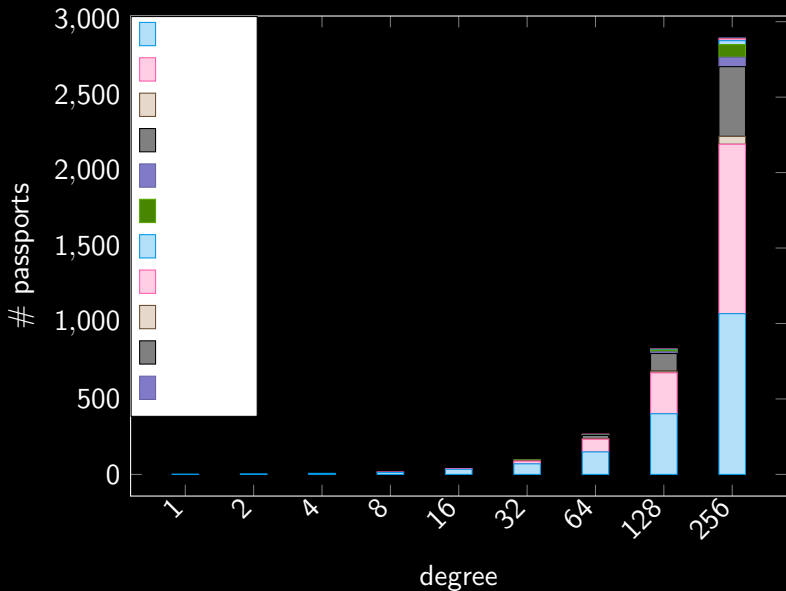
Results : distribution of genera



Results : groups by nilpotency class



Results : passport sizes



A refined conjecture



Passports

Recall that a passport \mathcal{P} consists of the data (g, G, λ) where $g \in \mathbb{Z}_{\geq 0}$, G is a transitive subgroup of S_d and $\lambda = (\lambda_0, \lambda_1, \lambda_\infty)$ is a triple of partitions of d corresponding to conjugacy classes (C_0, C_1, C_∞) of S_d .

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$$\left\{ (\sigma_0, \sigma_1, \sigma_\infty) \in C_0 \times C_1 \times C_\infty : \sigma_\infty \sigma_1 \sigma_0 = 1 \text{ and } \langle \sigma_0, \sigma_1 \rangle = G \right\} / \sim$$

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As a result of the action of $G_{\mathbb{Q}}$ on \mathcal{P} , the size of \mathcal{P} bounds the degree of the field of moduli of any Belyi map with passport \mathcal{P} .

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To instead analyze $\text{Gal}(\mathbb{Q}^{\text{al}} | \mathbb{Q}^{\text{ab}})$ we *refine* the notion of a passport.

Refined passports

A **refined passport** \mathcal{P} consists of the data (g, G, c) where $g \in \mathbb{Z}_{\geq 0}$, G is a transitive subgroup of S_d and $c = (c_0, c_1, c_\infty)$ is a triple of conjugacy classes of G .

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As was the case with passport, every permutation triple σ determines a refined passport $\mathcal{P}(\sigma)$.

A refined conjecture

Theorem (M.)

The size of $\mathcal{P}(\sigma)$ is equal to 1 for every 2-group permutation triple σ with degree ≤ 256 .

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Theorem (M.)

ARC is true for 2-group permutation triples σ with $\langle \sigma \rangle$ dihedral.

Computing equations



A motivating example : setup

Let F be a number field with integers \mathbb{Z}_F . Let $\text{Pl}(F)$ denote the places of F and S_∞ the archimedean places. For $v \in \text{Pl}(F) \setminus S_\infty$ let \mathfrak{p}_v denote the prime ideal of \mathbb{Z}_F corresponding to v .

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Question

How do we construct a quadratic extension of F with ramification prescribed by \mathfrak{a} ?

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How do we construct a quadratic extension of F with ramification prescribed by \mathfrak{a} ?

First, it is possible that no such extension exists.

A motivating example : setup

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If \mathfrak{a} is not principal, then the question requires more care.

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To summarize, in the case where \mathfrak{a} is not principal but there exists \mathfrak{b} with $\mathfrak{a}\mathfrak{b}^2$ principal we have $[\mathfrak{a}] \in \text{Cl}_F^2$ and $[\mathfrak{b}]$ is unique up to multiplication by $[\mathfrak{c}] \in \text{Cl}_F[2]$.

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The algorithms in this section rely on transporting this technique to the function field setting.

Algebraic function fields : setup

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Then the **function field of X** , denoted $K(X)$ is the field of fractions of the coordinate ring $K[x, y]/(f(x, y))$ of X .

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The set of places of F is denoted $\text{Pl}(F)$ and the **degree** of P is the index $[\mathcal{O}_P/P : K]$ of the **residue class field**.

Algebraic function fields : Picard group

The **divisor class group** $\text{Div}(F)$ of F is the free abelian group generated by the places of F . A **divisor** $D \in \text{Div}(F)$ is represented by a sum of places $\sum_P a_P P$ and the **degree** of D is $\sum_P a_P \deg(P)$.

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The **Picard group** of F is $\text{Pic}(F) := \text{Div}(F) / \text{Princ}(F)$.

The **Jacobian** of F is $\text{Pic}^0(F) := \text{Div}^0(F) / \text{Princ}(F)$.

Algebraic function fields : Riemann-Roch spaces

Algebraic function fields : quadratic extensions

Algorithm in characteristic $p \geq 3$: setup

Algorithm in characteristic $p \geq 3$: Galois test

Algorithm in characteristic $p \geq 3$: get candidates

Algorithm in characteristic $p \geq 3$: lift Belyi map

Implementation in characteristic zero

<https://github.com/michaelmusty/2GroupDessins>

- *all* 2-group Belyi maps modulo 3 up to degree 32
- hundreds of 2-group Belyi maps up to degree 256

Examples



Notation

DNG-a, b, c-gE-H

D : degree in $\{2, 4, 8, 16, 32, 64, 128, 256\}$

N : either T or S identifying group database

G : a positive integer identifying the group

a : ramification index of 0 in $\{2, 4, 8, 16, 32, 64, 128, 256\}$

b : ramification index of 1 in $\{2, 4, 8, 16, 32, 64, 128, 256\}$

c : ramification index of ∞ in $\{2, 4, 8, 16, 32, 64, 128, 256\}$

g : just the letter g

E : the genus in $\mathbb{Z}_{\geq 0}$

H : the hash of the 2-group permutation triple a positive integer

An interesting example

<https://michaelmusty.github.io/d3ssins>

Future work

- higher degree over \mathbb{F}_3
- an algorithm in characteristic zero
- prove ARC for other families of 2-groups
- p -group Belyi maps for p odd
- compute torsion fields

Backup slides