# IO Problem Set 1

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# 1 Question 1

#### 1.1 Part (a)

The demand curve with constant elasticity can be written as  $Q=aP^{-c}$ . Rewriting, the corresponding inverse demand function is  $P(Q)=a^{1/c}Q^{-1/c}$ . Then,  $P'(Q)=(-1/c)a^{1/c}Q^{-(1+c)/c}$ ,  $P''(Q)=(-1/c)(-(1+c)/c)a^{1/c}Q^{-(1+2c)/c}=((1+c)/c^2)a^{1/c}Q^{-(1+2c)/c}$  We then have the following:

$$P'(Q) + QP''(Q) = (-1/c)a^{1/c}Q^{-(1+c)/c} + Q(((1+c)/c^2)a^{1/c}Q^{-(1+2c)/c})$$
$$= (-1/c)a^{1/c}Q^{-(1+c)/c} + (((1+c)/c^2)a^{1/c}Q^{-(1+c)/c})$$
$$= (1/c^2)a^{1/c}Q^{-(1+c)/c} > 0.$$

#### 1.2 Part (b)

Next, let N firms be competing a la Cournot.

Assumption (A1) is that  $0 \ge P''(Y)y_i + P'(Y)\forall y_i < Y$ .

Assumption (A2) states that  $0 \ge P'(Y) - C_i''(y_i) \forall y_i < Q$ . We are given that each firm has identical cost functions, so  $C_i(y) = C(y)$ . Note that (A2) therefore states that  $C''(y_i) \ge P'(Y)$ . With identical costs, in equilibrium  $y_i = y = Y/N$ . Using this, (A1) becomes  $0 \ge P''(Y)Y/N + P'(Y)$ 

Let us set up the maximization problem for each firm:

$$\max_{y_i} P(y_i + Y_{-i})y_i - C(y_i)$$

$$\Rightarrow P'(Y)y_i + P(Y) - C'(y_i) = 0$$

$$\Rightarrow P(Y) = C'(Y/N) - P'(Y)Y/N$$
(1)

Differentiating both sides with respect to N,

$$\frac{\partial P(Y)}{\partial N} = -C''(Y/N)(Y/N^2) + P'(Y)(Y/N^2)$$
$$= (Y/N^2)(P'(Y) - C''(y))$$
$$< 0.$$

Now we can rewrite (1) for y and differentiate with respect to N:

$$\begin{split} y &= \frac{C'(y) - P(Ny)}{P'(Ny)} \\ \frac{\partial y}{\partial N} &= -\frac{C'(y)}{(P'(Ny))^2} P''(Ny) y - y + \frac{P(Ny)}{(P'(Ny))^2} P''(Ny) y \\ &= -y - \frac{P''(Y)y}{P'(Y)^2} (P(Y) - C'(y)) \\ &= -y - \frac{P''(Y)y}{P'(Y)^2} P'(Y) y \\ &= y \left( -1 - \frac{P''(Y)y}{P'(Y)} \right) &\leq 0, \end{split}$$

where the final inequality uses our rewritten form of (A1).

### 2 Question 2

Each bidder chooses a bid  $b_i \in \mathbb{R}$  to maximize their payoffs:

$$b_i = \arg\max_b \pi(b, b_{-1}),$$

where the payoff  $\pi(b, b_{-i})$  is:

$$\pi(b, b_{-i}) = \begin{cases} V - b, b > b_{-i} \\ 0, b < b_{-i} \\ (1/2)(V - b), b = b_{-i} \end{cases}.$$

The equilibrium is  $b_i = V \forall i$ . Why is this the case? Suppose instead player i bid  $b_i > V$ . Then, their payoff would be  $V - b_i < 0$ . Moreover, suppose  $b_i < V$ . Then, person i still only gets 0 payoff. So,  $b_i = V \forall i$  is an equilibrium. No other equilibrium can exist. To see why, first note that equilibria can only exist with  $b_i = b_{-i}$  as otherwise the bidder with the largest bid is strictly better off reducing their bid by some  $\epsilon$ . If  $b_i = b_{-i} < V$  then bidder i is better off increasing their bid by an epsilon and winning positive payoff. If  $b_i = b_{-i} > V$  then the best response for i is to reduce their bid by an  $\epsilon$  such that they are sure to receive 0 payoff instead of negative expected payoff. So, the only equilibrium is  $b_i = b_{-i} = V$ .

Now consider the all-pay auction. The expected payoff  $\pi(b, b_{-i})$  is given by the following:

$$\pi(b, b_{-i}) = \begin{cases} V - b, b > b_{-i}, \\ -b, b < b_{-i} \\ (1/2)V - b, b = b_{-i} \end{cases}$$

First we will show that a pure strategy Nash equilibrium does not exist. Suppose it does. Then,  $b_i = b_{-i}$  because otherwise the highest bidder would be strictly better off by reducing their bid by an  $\epsilon$ . Consider, then,  $b_i = b_{-i} = b$ . If b < V then either bidder would be better off increasing their bid by an  $\epsilon$  and winning V surely. If  $b \ge V$  then either bidder would be better off not bidding (or bidding zero). Therefore no pure strategy Nash equilibrium can exist.

We will now search for a mixed-strategy Nash equilibrium. The equilibrium solves for a pair of distributions from which bids are drawn,  $(F_1(b), F_2(b))$  are the corresponding CDFs. We know that the distributions satisfy boundary conditions  $F_i(0) = 0$ ,  $F_i(V) = 1$  by nonnegativity of our bids and from the value of the object being V. Given these distributions, and assuming differentiability, the maximization problem and first order conditions are the following:

$$\max_{b} F_{-i}(b)V - b$$
$$V f_{-i}(b) = 1$$

From this first order condition we see that the pdf  $f_j(b)$  is constant across the support of the distribution, implying a uniform distribution across our support. Our CDFs then take the functional form  $F_i(b) = b/V$ . The seller's expected revenue is  $E[R] = 2E[b_i] = 2V/2 = V$ .

It is important to note that the bidders are just indifferent between mixing in this way and not bidding (or bidding zero surely). That is, each individual's  $E[\pi] = P(\text{winning}_i)V - E[b_i] = (1/2)V - V/2 = 0$ . If the seller sets a reserve price then (applying symmetry) the probability of winning for each bidder will not change, so their expected bid will not increase (or they would be strictly better off by not bidding at all). Since each bidder's expected bid is will not change, the expected revenue for the seller will not increase.

# 3 Question 3

Consumers derive v = 3 from coffee. Starbucks coffee locations are at locations 0,1 and have prices  $p_0, p_1$  respectively. Esquire is at location 0.5 and has price q. Marginal cost is 0. Travel cost of the consumer is  $d^2$  where d is the distance (in miles) from the consumer to the coffee shop. We first solve for the indifferent consumers, under the assumption of existence. The indifferent consumer between Starbucks 0 and Esquire has the following expressions:

$$v - p_0 - x^2 = v - q - (0.5 - x)^2$$
$$q - p_0 = x^2 - x^2 + x - 0.25$$
$$x = (q - p_0) + (1/4)$$

The indifferent consumer between Starbucks 1 and Esquire has the following similar expression:

$$v - q - (0.5 - y)^{2} = v - p_{1} - (1 - y)^{2}$$
$$p_{1} - q - y^{2} + y - 0.25 = -y^{2} + 2y - 1$$
$$(3/4) + p_{1} - q = y.$$

Given  $p_0, p_1$ , Esquire sets q to maximize profits:

$$\max_{q} q(y - x)$$

$$\max_{q} q((1/2) + p_1 - 2q + p_0)$$

$$\Rightarrow (1/2) + p_1 + p_0 = 4q$$

$$\Rightarrow q = 1/8 + (1/4)(p_1 + p_0)$$

Given q, Starbucks sets  $p_1, p_0$  to maximize profits:

#### 4 Question 4