

1 Partial Orderings

Definition:

A relation on a set S is called antisymmetric if, whenever $x R y$ and $y R x$ are both true, then $x = y$.

Definition:

A relation R on a set S is called a partial ordering relation if it has the following three properties:
(partial order)

- (a) R is reflexive; that is, $x R x$ is true for every $x \in S$.
- (b) R is antisymmetric; that is, whenever $y R x$ and $x R y$ are true, then $x = y$.
- (c) R is transitive; that is, whenever $x R y$ and $y R z$ are both true, then $x R z$ is true.

Example: Is the relation set equality on some particular set S antisymmetric? Is set equality on some particular set S a partial ordering? Why or why not?

yes: $y = x$ and $x = y \Rightarrow x = y$
for equality

yes - previously said equality is an equivalence relation so reflexive & transitive properties are also satisfied

Example: Suppose S is the set of real numbers. Is \geq an antisymmetric relation on S ? Is \geq a partial ordering on S ? $3 \geq 2$ $2 \geq 3$

antisym: If $y \geq x$ and $x \geq y$ does this mean $x = y$? Yes being both \geq and $\leq \Rightarrow =$.
($4 \geq 2$ but $2 \not\geq 4$ is perfectly fine) (antisym. prop.)

reflexive: Is $x R x$ true for all $x \in S$? Is $x \geq x$ for $x \in$ real numbers? yes

transitive: If $x R y$ & $y R z$ is $x R z$? If $x \geq y$ and $y \geq z$, is $x \geq z$? yes
 \Rightarrow Partial order

Example: Can a relation be both symmetric and antisymmetric? Justify your reasoning.

Yes, equality is both: $y R x \wedge x R y$
 $y = x \wedge x = y \Rightarrow x = y$

(but usually 1 or the other)

Example: Can a partial ordering be an equivalence relation? Must a partial ordering be an equivalence relation? Justify your reasoning.

yes - equality is both

\geq is a partial ordering BUT $3 \geq 2$ while $2 \not\geq 3$ - fails symmetric property
 \Rightarrow not an equivalence relation

\rightarrow but is a partial ordering (passes reflexive, antisymmetric, & transitive properties)

for transitive: $(1,1) \wedge (1,2) \text{ need } (1,2)$ $(1,2) \wedge (2,2) \text{ need } (1,2)$ $(1,3) \wedge (3,3) \text{ need } (1,3)$
 $(1,1) \wedge (1,3) \text{ need } (1,3)$ $(1,2) \wedge (2,3) \text{ need } (1,3)$ $(2,2) \wedge (2,3) \text{ need } (2,3)$
 $(2,3) \wedge (3,3) \text{ need } (2,3)$

Example: Determine whether relation $R = \{(1,1), (1,2), (1,3), (2,2), (2,3), (3,3)\}$ is a partial order on set $S = \{1,2,3\}$ and justify.

reflexive: yes - we have $(1,1), (2,2), \wedge (3,3)$
 satisfy xRx for all $x \in S$

antisymmetric: yes - $(1,1), (2,2), (3,3)$ are fine \wedge if have (x,y) , don't have (y,x)
 $x \neq y$
 (don't have $(2,1), (3,1), \wedge (3,2)$ which is good)

transitive: yes - can chain $xRy \wedge yRz \Rightarrow xRz$ in every case

satisfy all 3 \Rightarrow partial order

Example: Determine whether relation $R = \{(1,1), (1,2), (1,3), (2,2), (2,3), (3,3)\}$ is a partial order on set $S = \{1,2,3,4\}$ and justify.

Now fail reflexive property - don't have $(4,4)$ \wedge need it because $4 \in S$
 (antisymmetric \wedge transitive pass)

Example: Determine whether relation $R = \{(1,1), (1,2), (1,3), (2,2), (2,3), (3,3), (3,1)\}$ is a partial order on set $S = \{1,2,3\}$ and justify.

$(3,1) \wedge (1,2)$ require $(3,2)$ so fails transitive property
 $(x,y) \quad (y,z) \quad (x,z)$

$(1,3) \wedge (3,1)$ both true but $1 \neq 3$ so fails antisymmetric property
 (reflexive still passes)

Definition:

Suppose R_1 is a partial order on set S_1 and R_2 is a partial order on set S_2 . Define relation R by $(a_1, a_2) R (b_1, b_2)$ if and only if one of the following is true:

(a) $a_1 \neq b_1$ and $a_1 R_1 b_1$ OR

(basically do 1 list in order then the other)

(b) $a_1 = b_1$ and $a_2 R_2 b_2$.

This relation is called the lexicographic order on $S_1 \times S_2$ or the dictionary order.

Example: Suppose Mr Webster is scheduled to interview three applicants for a summer internship at 9:00, 10:00, and 11:00 and Ms Collins is to interview three other applicants at the same times. Unfortunately both Mr Webster and Ms Collins have become ill so all 6 interviews are to be conducted by Ms Herrera. She has decided to schedule the applicants to be interviewed in time order, alphabetically by interviewer. Thus $S_1 = \{9:00, 10:00, 11:00\}$ and $S_2 = \{\text{Collins, Webster}\}$. What is the order of applicants and how could R_1 and R_2 be characterized?

$(9:00, \text{Collins}), (9:00, \text{Webster}), (10:00, \text{Collins}), (10:00, \text{Webster}), (11:00, \text{Collins}), (11:00, \text{Webster})$

R_1 is numerical order $\wedge R_2$ is alphabetical order (C before W)
 $(9, 10, 11)$

Note:

The prior example is an intentional variation on Example 2.22 in text. See that example for another way to do a lexicographic ordering.

Theorem 2.5:

If R_1 is a partial order on set S_1 and R_2 is a partial order on set S_2 , then the lexicographic order is a partial order on $S_1 \times S_2$.

Why should this work?

We're maintaining the original orders (which were partial orders) within each set, but just interleaving them (which should not cause problems)

Definition:

A partial order R on a set S is called a total order (or linear order) on S if every pair of elements in S can be compared; that is, for every $x, y \in S$, $x R y$ or $y R x$.

Example: What is an example of a relation on a set that has a total order?

\geq on the real numbers — I know how $2 \neq 1$ relate (have $(2,1)$)
I know how $3 \neq \pi$ relate (have $(\pi, 3)$)
I know how $3 \neq 3$ relate (have $(3,3)$)

For $S = \{1, 2, 3\}$ \geq : $R = \{(1,1), (2,1), (2,2), (3,1), (3,2), (3,3)\}$

Example: What is an example of a relation on a set that only permits partial orders?

$=$ on the real numbers — $(1,2) \neq (2,1)$ not in relation

For $S = \{1, 2, 3\}$ $=$ again $1 \neq 2$ not related either way
 $R = \{(1,1), (2,2), (3,3)\}$ (don't have $(1,2)$ or $(2,1)$)

Definition:

Let R be a partial order on set S . An element $x \in S$ is called a minimal element of S with respect to R if the only element $s \in S$ satisfying $s R x$ is x itself (that is, $s R x$ implies $s = x$).

Let R be a partial order on set S . An element $z \in S$ is called a maximal element of S with respect to R if the only element $s \in S$ satisfying $z R s$ is z itself (that is, $z R s$ implies $z = s$).

Note: No minimal or maximal elements need to exist for a given partial order.

Example: Identify the minimal and maximal elements of S with respect to R , if they exist: S is the set of all integers x such that $0 \leq x < 2$ and $x R y$ if $x \leq y$.

$S = \{0, 1\}$ $R = \{(0,0), (0,1), (1,1)\}$

0 is minimal

1 is maximal

* **Example:** Identify the minimal and maximal elements of S with respect to R , if they exist: S is the set of all real numbers x such that $0 \leq x < 1$, and $x R y$ if $x \leq y$.

minimal: 0 (nothing smaller) $[0, 1)$

maximal: none (never quite reach 1 & could always make something bigger)

* **Example:** Identify the minimal and maximal elements of S with respect to R , if they exist: S is the set of all real numbers x such that $0 \leq x < 1$, and $x R y$ if $x \geq y$.

minimal: none

maximal: 0

Example: Identify the minimal and maximal elements of S with respect to R , if they exist: S is the set of nonempty subsets of $\{1, 2, 3\}$ and $A R B$ if $A \subseteq B$.

minimal: $\{1\}, \{2\}, \{3\}$

$S = \{\{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$

maximal: $\{1, 2, 3\}$

Theorem 2.6:

Let R be a partial order on a finite set S . Then S has both a minimal and a maximal element with respect to R .

Example: Why does this theorem not contradict the example above?

→ notice S is not finite

2 Hasse Diagrams

Description:

A Hasse diagram represents a partial order R on a finite set S by representing each element of S by a point and each pair of distinct points related by R has a line connecting them. Each line is arranged so that the initial point is below its terminal point; that is, an arrow is drawn from the point representing x to the point representing y when $x R y$ and there is no $s \in S$ other than x and y such that $x R s$ and $s R y$. Hasse diagrams are read from bottom to top so all line segments between points are regarded as pointing upward.

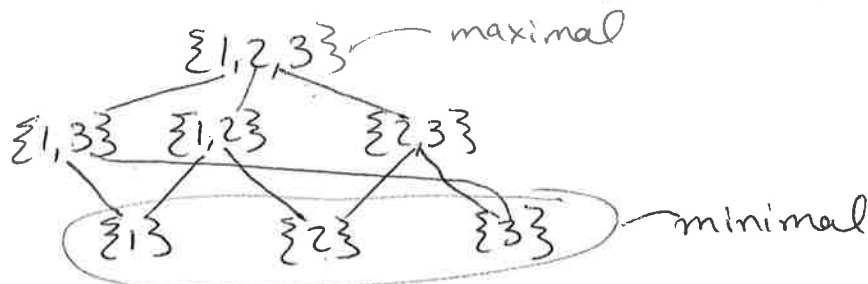
Note: segments need not be drawn with arrows. The direction is implicit in how the Hasse diagram is drawn.

Example: Where will minimal and maximal element(s) of partial order R on set S appear in a Hasse diagram?

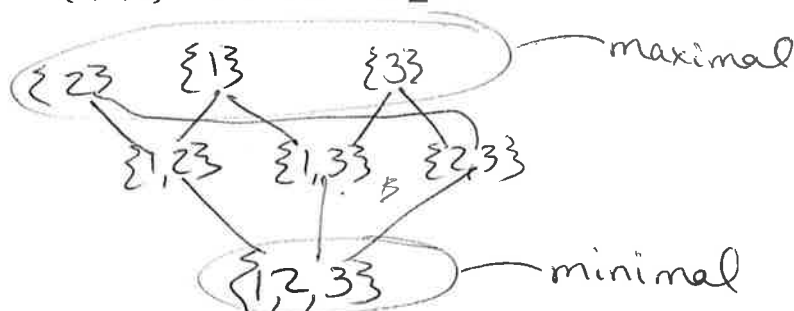
at the bottom

at the top

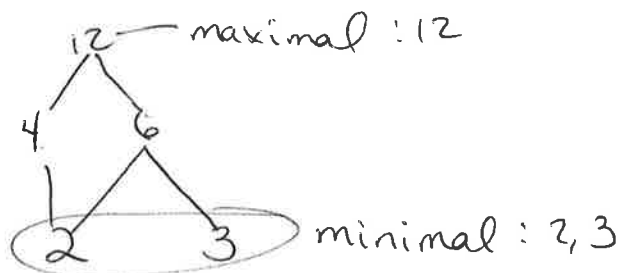
Example: Create a Hasse diagram for relation R on set S : S is the set of nonempty subsets of $\{1, 2, 3\}$ and $A R B$ if $A \subseteq B$. $S = \{\{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$



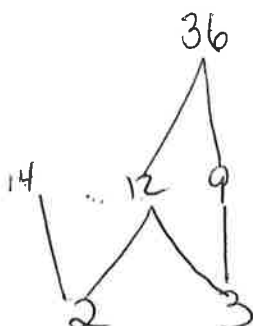
Example: Create a Hasse diagram for relation R on set S : S is the set of nonempty subsets of $\{1, 2, 3\}$ and $A R B$ if $B \subseteq A$.



Example: Create a Hasse diagram for relation R on set S : let R be the partial order "divides" on the set $S = \{2, 3, 4, 6, 12\}$. What does this Hasse diagram indicate about the minimal and maximal element(s) of S ? (i.e. x divides y)



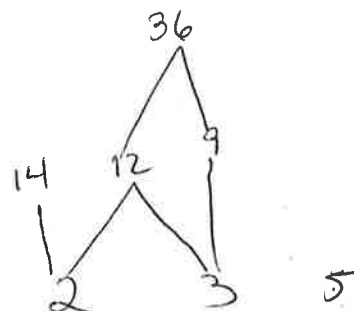
Example: Create a Hasse diagram for relation R on set S : let R be the partial order "divides" on the set $S = \{2, 3, 9, 12, 14, 36\}$. What does this Hasse diagram indicate about the minimal and maximal element(s) of S ?



maximal: 14, 36

minimal: 2, 3

if changed $S = \{2, 3, 5, 9, 12, 14, 36\}$



max: 5, 14, 36
min: 2, 3, 5