What if we have lots of edges but want to use a minimal number of edges to be able to reach everywhere?

Definition:	
A	of a graph ${\mathcal G}$ is a tree formed by using edges and vertices of ${\mathcal G}$ containing all
vertices of	\mathcal{G} .
Example: \mathcal{G} .	Consider the graph below. Remove edges as appropriate to make a spanning tree of
Example: graph below	Use the Breadth-First Search Algorithm to help you create a spanning tree for the

A spanning tree constructed by means of the breadth-first search algorithm is called a

Theorem 5.6:

Definition:

A graph is connected if and only if it has a spanning tree.

Why should this make sense?

Definition: In a weighted graph, the	is the sum of the weights of the edges in the tree.
	thted graph is a spanning tree for which the weight of the
tree is as small as possible. A	in a weighted graph is a spanning tree
for which the weight of the tree is as	
How can we approach creating a minir	nal spanning tree or maximal spanning tree?
_	and a minimal spanning tree and maximal spanning tree fo
the graph below. Give the weight of ea	ach.
Example: Use Prim's algorithm to f	and a minimal spanning tree and maximal spanning tree for
the graph below. Give the weight of ea	
Example: Use Prim's algorithm to fithe graph below. Give the weight of ea	and a minimal spanning tree and maximal spanning tree fo ach.