# **CS131: Programming Languages**

Fall 2015 Week #2

## **Today**

- OCaml practice
- HW2

```
head_eq
input: element, list
output: If head of the list is equal to the input element: Some(h)
else None
```

```
let head_eq h list = Do we need all the match
match list with cases?
| [] -> None
| h::_ -> Some h
| _ -> None
```

```
head_eq
input: element, list
output: If head of the list is equal to the input element: Some(h)
else None
```

```
let head_eq h list =
  match list with
  | [] -> None
  | h::_ -> Some h
  | _ -> None
```

Do we need all the match cases?

No. We can remove either the 1st or 3rd match cases. Why?

```
head_eq
input: element, list
output: If head of the list is equal to the input element: Some(h)
else None
```

```
let head_eq h list = let head_eq h =
  match list with function
  | h::_ -> Some h
  | _ -> None | _ -> None
```

Are these the same? A: Yes

```
head_eq
   input: element, list
   output: If head of the list is equal to the input element: Some(h)
          else None
                              What does
let head_eq h list =
  match list with
                              head_eq 3 [1;2;3] output?
  | h::_ -> Some h
  | _ -> None
                              Some 1
```

How do we fix it?

head\_eq

drop

input: list, n

output: the same list but with every nth element removed

let drop list n =

Sample output:

```
drop [] 1 = []
drop [1;2;3] 1 = []
drop [1;2;2;3] 2 = [1;2]
```

```
drop
   input: list, n
   output: the same list but with every nth element removed
let drop list n =
  let rec help i = function
    | [] -> []
    | h :: t -> if i = n
                    then help 1 t
                    else h :: help (i+1) t
  in help 1 list;;
```

## **HW2: Naive Parsing of CFGs**

- A parser generator
- Submission due: Oct 16, 11:55 pm

## **Converting Grammars**

Pair of a nonterminal starting symbol and a list of rules.

Pair of a nonterminal starting symbol and a production function.

The production function is one large pattern match on the nonterminal.

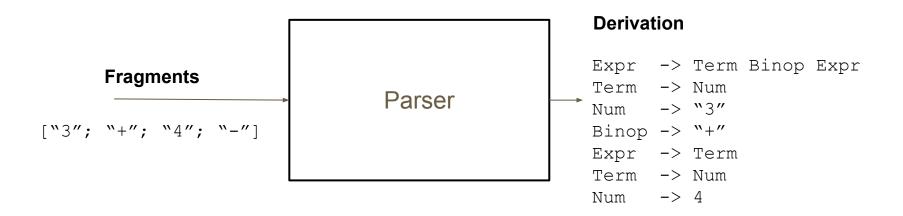
## **Converting Grammars**

Correct.

Careful: RHS of the new grammar is a list of lists

Incorrect!

### **Parser**

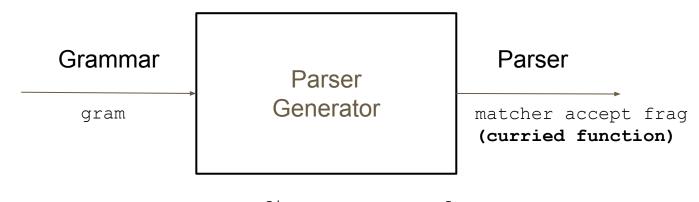


#### **Suffixes**

["-"]

- Input: fragment (i.e. list of tokens)
- Output:
  - the derivation for the given tokens
  - remaining tokens (suffixes)

### The Goal: Parser Generator



- parse\_prefix gram accept frag
- Input: Grammar (starting symbol, production function)
- Output: A parser which has internalized that grammar

## **Acceptor**

- A function that determines whether the given input is "acceptable"
- Input:
  - rules (a derivation)
  - frag (a list of tokens; suffixes)
- Output:
  - Some (rules, frag) if we like the input
  - None if we don't like it

## **Acceptor: examples**

- let accept\_all rules frag = Some (rules, frag)
- let accept\_empty\_suffix rules = function
  | [] -> Some( rules, [] )
  | \_ -> None

let accept\_only\_non\_lvalues = ...
 if a derivation's rules contain 'Lvalue',
 it returns None

### **Matcher**

- A function that matches a prefix of a fragment and checks whether the acceptor passes or not.
- Input: an acceptor and a fragment
- Output: whatever the acceptor returns
  - O Some (rules, frag) | None

### **Basic Matcher**

let match\_num num frag accept =

### **Basic Matcher**

### Hint code

```
let append matchers matcher1 matcher2 frag accept =
 matcher1 frag (fun frag1 -> matcher2 frag1 accept)
let match empty frag accept = accept frag
let make appended matchers make a matcher ls =
  let rec mams = function
    | [] -> match empty
    | head::tail -> append matchers (make a matcher
head) (mams tail)
  in mams ls
```

## **Hint code: make\_appended\_matchers**

- When we define a matcher for the nonterminal symbol expr,
   it can be represented by a combination of three concatenated matchers.
- matcher\_expr acceptor ["3"; "+"; "4"]
  = matcher\_term acceptor ["3"] and
   matcher\_binop acceptor ["+"] and
   matcher expr acceptor ["4"]

## Hint code: make\_or\_matcher

matcher\_combined

Expr -> matcher\_expr

Term Binop Expr matcher\_term 

Binop matcher\_binop 

matcher\_expr

or

Expr -> Term
matcher\_expr matcher\_term

matcher\_expr = or\_matchers
[ matcher\_combined; matcher\_term ]

## **Acceptor: Example Purpose**

Force a partial derivation:

```
let accept_only_non_binop rules frag =
  if contains_binop rules
  then None
  else Some (rules, frag)
```

 Given this acceptor, only accepted derivation of ["1"; "+"; "1"] would be:

```
[Expr, [Term]; Term, [Num]; Num, [1]]
```

and the remaining fragment would be ["+"; "1"]

## **Acceptor: Example Purpose**

```
Expr -> Term | Num
Term -> Num
Num -> 1
```

- Ambiguous grammar: two ways to derive "1"
   What are the ways?
- An acceptor can force one of the derivations to be chosen