USACO OPEN11 Problem 'mowlawn' Analysis

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First compute partial sums, so that psum $[i] = E_0 + ... + E_{i-1}$. Then, if we let dp $[i]$ be the best possible total over all subsets of the first i cows, since we can choose at most K cows in a row we have the following recurrence:

dp $[i]$ = min over all j such that i - K $\leq$ j $\leq$ i of dp $[j - 1] + E_j + ... + E_{i-1}$ = dp $[j - 1]$ + psum $[i]$ - psum $[j]$.

This immediately gives us an O(NK) solution. However, note that the psum $[i]$ term is independent of j, while dp $[j - 1]$ - psum $[j]$ depends only on j. Thus, we can store paired values of (dp $[j - 1]$ - psum $[j]$, j) in a heap, which allows us to quickly compute the maximum value of dp $[j - 1]$ - psum $[j]$ for j $\geq$ i - K in an amortized fashion. (In particular, we have a max heap sorted by dp $[j - 1]$ - psum $[j]$, and each time we perform a query we pop the heap until the index of the top is at least i - K.) This gives us an O(N log N) solution:

```cpp
#include <cstdio>
#include <queue>
using namespace std;

FILE *fin = fopen ("mowlawn.in", "r"), *fout = fopen ("mowlawn.out", "w");
const int MAXN = 100005;

struct data
{
    int ind;
    long long val;

    inline bool operator < (const data &o) const
    {
        return val < o.val;
    }
};

int N, K;
long long psum [MAXN], dp [MAXN];
priority_queue <data> hp;

int main ()
{
    fscanf (fin, "%d %d", &N, &K);

    for (int cow, i = psum [0] = 0; i < N; i++)
    {
        fscanf (fin, "%d", &cow);
        psum [i + 1] = psum [i] + cow;
    }

    hp.push ((data) {-1, 0});

    for (int i = 0; i <= N; i++)
    {
        while (hp.top ().ind < i - K - 1)
            hp.pop ();

        dp [i] = hp.top ().val + psum [i];
        hp.push ((data) {i, dp [i] - psum [i + 1]});
    }
    fprintf (fout, "%lld\n", dp [N]);
}
```