



Errors in Concurrency

Walter Cazzola

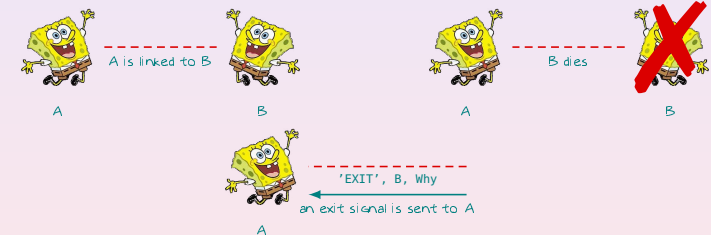
Dipartimento di Informatica
Università degli Studi di Milano
e-mail: cazzola@di.unimi.it
twitter: @w_cazzola



Errors in Concurrent Programs Error Handling on Exit

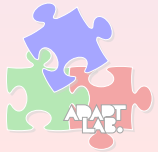
When two processes are related

- the errors of one affect the behavior of the other process;
- the BIF `link` function helps to monitor.



If A is linked to B

- when B dies an exit signal is sent to A;
- the signal is a message like `{'EXIT', Pid, ...}`.



Errors in Concurrent Programs Error Handling on Exit

```
-module(dies).
-export([on_exit/2]).
on_exit(Pid, Fun) ->
  spawn(fun() ->
    process_flag(trap_exit, true),
    link(Pid),
    receive
      {'EXIT', Pid, Why} -> Fun(Why)
    end
  end).
```

```
1> F = fun() -> receive X -> list_to_atom(X) end end.
#Fun<erl_eval.20.67289768>
2> Pid = spawn(F).
<0.35.0>
3> dies:on_exit(Pid, fun(Why) -> io:format("~p died with:~p~n",[Pid, Why]) end).
<0.37.0>
4> Pid ! hello.
<0.35.0> died with:{badarg,[{erlang,list_to_atom,[hello]}]}
=ERROR REPORT==== 9-Nov-2011::17:50:20 ===
Error in process <0.35.0> with exit value: {badarg,[{erlang,list_to_atom,[hello]}]}
hello
```



Errors in Concurrent Programs Details of Error Handling

Links

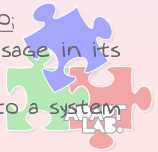
- defines an error propagation path between two processes;
- if a process dies an exit signal is sent to the other process;
- the set of processes linked to a given process is called link set.

Exit Signals

- they are generated by a process when it dies;
- signals are broadcast to all processes in the link set of the dying process;
- the exit signal contains an argument explaining why the process died (`exit(Reason)` or implicitly `set`).
- when a process "naturally dies" the exit reason is `normal`;
- exit signals can be explicitly sent via `exit(Pid, X)`: the sender does not die ("fake death").

System Processes

- a non system process that receives an exit signal dies too;
- a system process receives the signal as a normal message in its mailbox;
- `process_flag(trap_exit, true)` transform a process into a system process.





Errors in Concurrent Programs

Details of Error Handling (Cont'd)

Errors in
Concurrency
Walter Cazzola

Error
Handling
links
monitors
References

Receiver's Behavior

trap_exit	Exit Signal	Action
true	kill	dies & broadcasts it to its link set
true	✗	adds {'EXIT', Pid , X } to the mailbox
false	normal	continues & the signal vanishes
false	kill	dies & broadcasts it to its link set
false	✗	dies & broadcasts it to its link set

Alternatives

- I don't care if a process I create crashes.
Pid = spawn(fun() -> ... end)
- I want to die if a process I create crashes.
Pid = spawn_link(fun() -> ... end)
- I want to handle errors if a process I create crashes
process_flag(trap_exits, true),
Pid = spawn_link(fun() -> ... end).



Slide 5 of 14



Errors in Concurrent Programs

Going into Details of Error Handling

Errors in
Concurrency
Walter Cazzola

Error
Handling
links
monitors
References

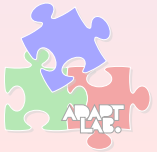
```
-module(edemol).
-export([start/2]).

start(Bool, M) ->
  A = spawn(fun() -> a() end),
  B = spawn(fun() -> b(A, Bool) end),
  C = spawn(fun() -> c(B, M) end),
  sleep(1000), status(b, B), status(c, C).

a() -> process_flag(trap_exit, true), wait(a).
b(A, Bool) -> process_flag(trap_exit, Bool), link(A), wait(b).
c(B, M) -> link(B),
  case M of
    {die, Reason} -> exit(Reason);
    {divide, N} -> 1/N, wait(c);
    normal -> true
  end.
```

This starts 3 processes: **A**, **B** and **C**

- **A** will trap exits and watch for exits from **B**,
- **B** will trap exits if **Bool** is true and
- **C** will die with exit reason **M**



Slide 6 of 14



Errors in Concurrent Programs

Going into Details of Error Handling (Cont'd)

Errors in
Concurrency
Walter Cazzola

Error
Handling
links
monitors
References

```
wait(Prog) ->
  receive
  Any ->
    io:format("Process ~p received ~p~n", [Prog, Any]),
    wait(Prog)
  end.

sleep(T) ->
  receive
  after T -> true
  end.

status(Name, Pid) ->
  case erlang:is_process_alive(Pid) of
    true -> io:format("process ~p (~p) is alive~n", [Name, Pid]);
    false -> io:format("process ~p (~p) is dead~n", [Name, Pid])
  end.
```

This starts 3 processes: **A**, **B** and **C**

- wait/1 just prints any message it receives;
- sleep/1 awakes the invoking process after a period of time;
- status/2 prints the aliveness of the invoking process.



Slide 7 of 14



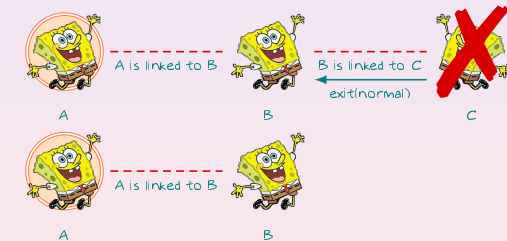
Errors in Concurrent Programs

Going into Details of Error Handling (Cont'd)

Errors in
Concurrency
Walter Cazzola

Error
Handling
links
monitors
References

```
1> edemol:start(false, {die,normal}).
process b (<0.48.0>) is alive
process c (<0.49.0>) is dead
ok
```



- **B** is not a system process;
- when **C** dies with normal signal, **B** doesn't die.



Slide 8 of 14

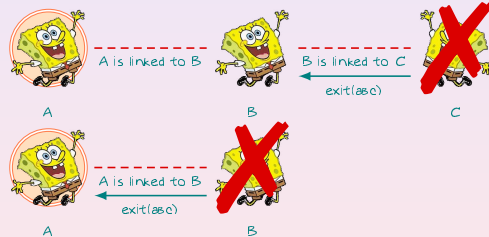


Errors in Concurrent Programs Going into Details of Error Handling (Cont'd)

Errors in
Concurrency
Walter Cazzola

Error
Handling
links
monitors
References

```
1> edemol:start(false, {die, abc}).
Process a received {'EXIT',<0.40.0>,abc}
process b (<0.40.0>) is dead
process c (<0.41.0>) is dead
ok
```



- B is not a system process;
- when C evaluates `exit(abc)`, process B dies;
- when B exits rebroadcasts the unmodified exit signal to its link set
- A traps the exit signal and convert it to the error message



Slide 9 of 14

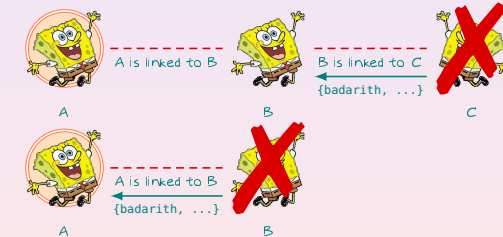


Errors in Concurrent Programs Going into Details of Error Handling (Cont'd)

Errors in
Concurrency
Walter Cazzola

Error
Handling
links
monitors
References

```
6> edemol:start(false, {divide,0}).
Process a received {'EXIT',<0.56.0>,{badarith,[{edemol,c,2}]}}
=ERROR REPORT==== 11-Nov-2011::18:03:29 ===
Error in process <0.57.0> with exit value: {badarith,[{edemol,c,2}]}}
process b (<0.56.0>) is dead
process c (<0.57.0>) is dead
ok
```



- B is not a system process;
- when C tries to divide by zero an error occurs and C dies with a `{badarith, ...}` error;
- B receives this and dies and the error is propagated to A



Slide 10 of 14

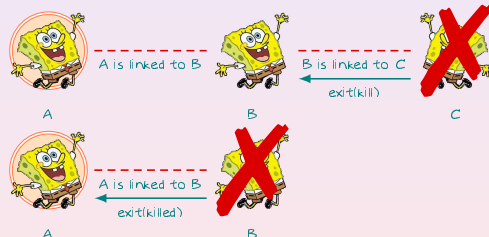


Errors in Concurrent Programs Going into Details of Error Handling (Cont'd)

Errors in
Concurrency
Walter Cazzola

Error
Handling
links
monitors
References

```
1> edemol:start(false, {die, kill}).
Process a received {'EXIT',<0.60.0>,killed}
process b (<0.60.0>) is dead
process c (<0.61.0>) is dead
ok
```



- B is not a system process;
- the exit reason `kill` causes B to die, and the error is propagated to its link set.



Slide 11 of 14

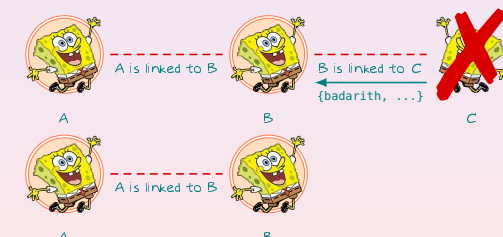


Errors in Concurrent Programs Going into Details of Error Handling (Cont'd)

Errors in
Concurrency
Walter Cazzola

Error
Handling
links
monitors
References

```
8> edemol:start(true, {divide,0}).
Process b received {'EXIT',<0.65.0>,{badarith,[{edemol,c,2}]}}
=ERROR REPORT==== 11-Nov-2011::18:16:47 ===
Error in process <0.65.0> with exit value: {badarith,[{edemol,c,2}]}}
process b (<0.64.0>) is alive
process c (<0.65.0>) is dead
ok
```



- B is a system process;
- in all cases, B traps the error;
- the error is never propagated to A



Slide 12 of 14



Errors in Concurrent Programs

Monitors: Unidirectional Links

Errors in
Concurrency
Walter Cazzola

Error
Handling
links
monitors
References

Links are **symmetric**

- i.e., if **A** dies, **B** will send an exit signal and vice versa;
- to prevent a process from dying, we have to make it a system process that is not always desirable.

A monitor is an **asymmetric** link

- if **A** monitors **B** and **B** dies **A** will be sent an exit signal but
- if **A** dies **B** will **not** be sent a signal.

A can create a monitor for **B** calling `erlang:monitor(process, B)`

- if **B** dies with exit reason **Reason** a `'DOWN'` message `{'DOWN', Ref, process, B, Reason}` is sent to **A** (**Ref** is the reference to the monitor).
- the monitor is unidirectional:
 - to repeat the above call will create several, independent monitors and each one will send a `'DOWN'` message when **B** terminates.



Slide 13 of 14



References

Errors in
Concurrency
Walter Cazzola

Error
Handling
links
monitors
References

- Gul Agha.
Actors: A Model of Concurrent Computation in Distributed Systems.
MIT Press, Cambridge, 1986.
- Joe Armstrong.
Programming Erlang: Software for a Concurrent World.
The Pragmatic Bookshelf, fifth edition, 2007.
- Francesco Cesarini and Simon J. Thompson.
Erlang Programming: A Concurrent Approach to Software Development.
O'Reilly, June 2009.



Slide 14 of 14