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The Friendship Theorem

Michelle Hewson, Maxwell Pitney, & Dakota Thompson

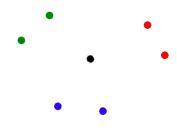
October 13, 2021 Math 479 Project #2



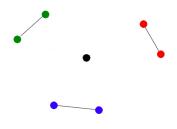
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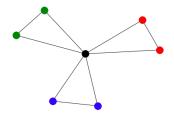
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Introduction

• This concept is the **Friendship Theorem**.

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Theorem

Suppose in a group of people we have the situation that any pair of persons have precisely one common friend. Then there is always a person (the "politician") who is everybody's friend.

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- Although the problem's creator is unknown, there are several famous proof methods.
- We'll tackle this proof using graph theory and linear algebra.

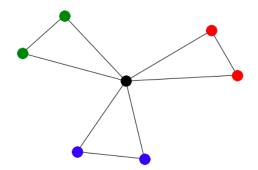


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- An edge can be used to define a relationship between vertices.
- A path is a finite sequence of adjacent vertices and adjacent edges where neither can be repeated.

• An example of a graph:



Linear Algebra

• An identity matrix:

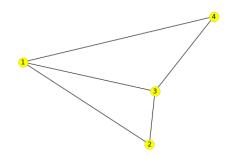
$$I_{1} = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

$$I_{2} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$$

$$I_{3} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$I_{n} = \begin{bmatrix} 1 & 0 & 0 & \cdots & 0 \\ 0 & 1 & 0 & \cdots & 0 \\ 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & 1 \end{bmatrix}$$

• An adjacency matrix:



Linear Algebra

ullet Eigenvalues (λ) are special scalars that have the property,

$$Ax = \lambda x$$
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if x is the eigenvector where $x \in \mathbb{R}^n \neq \mathbf{0}$ and A is a matrix.

Linear Algebra

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if x is the eigenvector where $x \in \mathbb{R}^n \neq \mathbf{0}$ and A is a matrix.

 The trace of a square matrix is the sum of the elements in the main diagonal.

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Theorem

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• The *vertices* represent the people.

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$\mathsf{Theorem}$

Suppose that G is a finite graph in which any two vertices have precisely one common neighbor. Then there is a vertex which is adjacent to all other vertices.

- The vertices represent the people.
- Any given two people that are considered friends will be represented with a connecting edge.



Visualization of The Friendship Theorem

 Graphs that fulfill the properties of the Friendship Theorem are windmill graphs.



Idea of the proof:

• We will prove the Friendship Theorem using a contradiction.

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- To do this, we will create a counterexample graph.

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- Note that *G* still must adhere to the condition of being a friendship graph.

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- Note that G still must adhere to the condition of being a friendship graph.
- Claim: The graph G is a regular graph, or that d(u) = d(v) for all vertices u and v.

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• Suppose d(u) = k, where w_1, w_2, \dots, w_k are the neighbors of u.

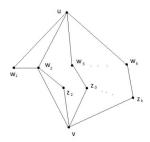
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- Suppose d(u) = k, where w_1, w_2, \dots, w_k are the neighbors of u.
- Without loss of generality, let w_2 be adjacent to v and exactly one other w_i be adjacent to w_2 , we'll say w_1 .

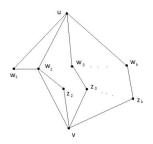
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- Suppose d(u) = k, where w_1, w_2, \dots, w_k are the neighbors of u.
- Without loss of generality, let w_2 be adjacent to v and exactly one other w_i be adjacent to w_2 , we'll say w_1 .
- Then we know that w_i must also have precisely one common neighbor, z_i , where $i \ge 2$, with v.

• This gives us this diagram of G.



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• We can now conclude that $d(v) \ge k = d(u)$, and by symmetry, d(u) = d(v) = k.

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- Note that any vertex, except for w₂, is only adjacent to either u or v, but not both.
- Using what we just concluded, we know that any of these vertices must have degree k.
- Recall that G is a counterexample, so there is no politician. This means there must exist some non-neighbor of w_2 , therefore w_2 must also have degree k.

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$$n=k^2-k+1.$$

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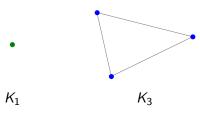
$$n=k^2-k+1.$$

- k^2 is given by summing the degrees of the k neighbors of u.
- -k+1 is given by subtracting all of the k common neighbors of u, except for u.
- This concludes the first part of this proof.

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- If k = 2, we get $k^2 k + 1 = 3 = n$.
- These k and n values result in K_1 and K_3 . Both graphs have a "politician" vertex.



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- Each row has k 1's and each pair of rows has one column where they both have a 1 entry.
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- A must be symmetric because $a_{ij} = a_{ji}$.

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Proof of the Friendship Theorem

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$$A^{2} = \begin{bmatrix} k-1 & 0 & \cdots & 0 \\ 0 & k-1 & & 0 \\ \vdots & & \ddots & \vdots \\ 0 & \cdots & 0 & k-1 \end{bmatrix} + \begin{bmatrix} 1 & 1 & \cdots & 1 \\ 1 & 1 & & 1 \\ \vdots & & \ddots & \vdots \\ 1 & \cdots & 1 & 1 \end{bmatrix}$$

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- Therefore, A^2 has eigenvalues $k-1+n=k^2$ (of multiplicity 1) and k-1 (of multiplicity n 1).

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 - Thus, A has eigenvalues k of multiplicity 1 and $\pm \sqrt{k-1}$ of multiplicity n-1.
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- Note that r + s = n 1.

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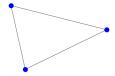
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Recap

• Introduced the *counterexample* graph *G*.

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- Proved that all vertices are of degree k.

• Found *n*, the number of vertices in the graph.

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- Showed that the adjacency matrix A is symmetric and diagonalizable.
- Used the eigenvalues of A to produce a contradiction.

Further Discussion

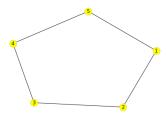
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Further Discussion

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- Windmill graphs are only possible with the presence of a politician.

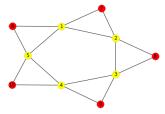
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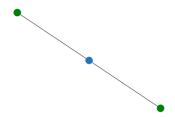
- Note that the Friendship Theorem only works with finite graphs.
- Windmill graphs are only possible with the presence of a politician.
- A counterexample can be constructed beginning with a 5 cycle:



• What if we rephrased the Friendship Theorem to say *G* is a graph with the property that there is exactly one path of length 2 between any two vertices?

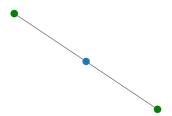


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- This reconstruction is another way to present the friendship condition.
- What if the path length is greater than 2?



Theorem

Let $\ell > 2$, then there are no finite graphs with the property that between any two vertices there is precisely one path of length ℓ .

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- This has been verified for all $\ell \leq 33$.
- A general proof of this conjecture has yet to be formulated.

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- This proof, which was presented by Paul Erdös, is the first of several different proofs for this theorem.
- It uses applications of both graph theory and linear algebra.
- The Friendship Theorem is a fascinating example of a real life problem that can be translated to and solved with mathematics.

References

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