

DIS 4A

EUCLID'S ALGORITHM

goal: calculate $\gcd(x, y)$

theorem: $\gcd(x, y) = \gcd(y, x \bmod y)$, $x \geq y$

↳ we can iteratively reduce the numbers we are calculating the gcd for until one of them is zero.

↳ intuitively: a number d divides x and y



d divides y and the remainder when dividing x by y

example: $\gcd(10, 6) = \gcd(6, 4)$
 $\uparrow \quad \uparrow$ $\uparrow \quad \uparrow$
 $x \quad y$ $y \quad x \bmod y$

$$\gcd(10, 6) = \gcd(6, 4) = \gcd(4, 2) = \gcd(2, 0) = 2$$

EXTENDED EUCLID'S ALGORITHM

goal: in addition to calculating $\gcd(x, y)$,
find a, b such that $\gcd(x, y) = ax + by$

how? q1 on the worksheet :)

note: when $\gcd(x, y) = 1$, we can solve for
the inverses $x^{-1} \pmod{y}$ and $y^{-1} \pmod{x}$

ex. $ax + by = 1 \pmod{y}$

$$ax \equiv 1 \pmod{y}$$

$$x^{-1} \equiv a \pmod{y}$$

(and similarly for y^{-1})

CHINESE REMAINDER THEOREM

for coprime positive integers n_1, n_2, \dots, n_k , there is a unique solution $(\text{mod } n_1 n_2 \dots n_k)$ to the system of equations.

$$x \equiv a_1 \pmod{n_1}$$

$$x \equiv a_2 \pmod{n_2}$$

\vdots

$$x \equiv a_k \pmod{n_k}$$

the solution:

don't memorize this!
understand each of the terms (q2)

$$x = \left[\sum_{i=1}^k a_i \left(\frac{N}{n_i} \right) \left(\left(\frac{N}{n_i} \right)^{-1} \pmod{n_i} \right) \right] \pmod{N}$$

where $N = n_1 n_2 \dots n_k$