## **Judgements**

 $\Gamma \vdash_i e \Leftarrow A \dashv \Gamma'$  - in context  $\Gamma$  (which has zero resources), e checks to have type A and consumes resources according to quantities from  $\Gamma'$ .

 $\Gamma \vdash_i e \Rightarrow A \dashv \Gamma'$  - as above, but with type inference.

#### **Notation**

 $\Gamma \vdash_i e \Leftarrow A \dashv \Gamma', s_A^i x : A \text{ is a shorthand for } \Gamma \vdash_i e \Leftarrow A \dashv \Gamma', rx : A \text{ with the condition } s_A^i \sqsubseteq r.$ 

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# Subsumption and annotations

$$\frac{\Gamma \vdash_{i} e \Rightarrow A \dashv \Gamma_{1} \quad \Gamma \vdash_{A} \equiv B \Rightarrow \text{Type}_{r} \dashv \Gamma_{2}}{\Gamma \vdash_{i} e \Leftarrow B \dashv \Gamma_{1} + \Gamma_{2}} \text{SUBSUMPTION}$$

$$\frac{\Gamma \vdash_{A} \Rightarrow \text{Type}_{r} \dashv \Gamma_{1} \quad \Gamma \vdash_{i} e \Leftarrow A \dashv \Gamma_{2}}{\Gamma \vdash_{i} (e : A) \Rightarrow A \dashv \Gamma_{1} + \Gamma_{2}} \text{Annot}$$

# Using variables

$$\frac{\Gamma(x) = A}{\Gamma \vdash_{\mathsf{c}} x \Rightarrow A \dashv |\Gamma|_{x}}$$

$$\frac{\Gamma(x) = A}{\Gamma \vdash_{\text{nc}} x \Rightarrow A \dashv 0 \Gamma}$$

### **Functions**

$$\frac{\Gamma, 0 \times : A \vdash_{i} e \Leftarrow B \dashv \Gamma', 1_{A}^{i} \times : A}{\Gamma \vdash_{i} \lambda \times . e \Leftarrow A \rightarrow B \dashv \Gamma'}$$

$$\frac{\Gamma \vdash_{i} f \Rightarrow A \rightarrow B \dashv \Gamma_{1} \quad \Gamma \vdash_{i} a \Leftarrow A \dashv \Gamma_{2}}{\Gamma \vdash_{i} f \ a \Rightarrow B \dashv \Gamma_{1} + \Gamma_{2}}$$

#### Box

$$\frac{\Gamma \vdash_{i} e \Leftarrow A \dashv \Gamma'}{\Gamma \vdash_{i} \text{box } e \Leftarrow !_{r} A \dashv r \Gamma'}$$

$$\frac{\Gamma \vdash_i e_1 \Rightarrow !_r A \dashv \Gamma_1 \quad \Gamma, 0 \, x : A \vdash_i e_2 \Rightarrow B \dashv \Gamma_2, r_A^i \, x : A}{\Gamma \vdash_i \mathtt{let box} \, x = e_1 \mathtt{ in } e_2 \Rightarrow B \dashv \Gamma_1 + \Gamma_2}$$

## **Empty**

$$\frac{\Gamma \vdash_i e \Leftarrow \texttt{Empty} \dashv \Gamma'}{\Gamma \vdash_i \texttt{Empty-elim} \ e \Leftarrow A \dashv \Gamma'}$$

## Unit

$$\overline{\Gamma \vdash_i \text{unit} \Leftarrow \text{Unit} \dashv 0 \Gamma}$$

$$\frac{\Gamma \vdash_i e_1 \Rightarrow \mathtt{Unit} \dashv \Gamma_1 \quad \Gamma \vdash_i e_2 \Rightarrow A \dashv \Gamma_2}{\Gamma \vdash_i \mathtt{let unit} = e_1 \mathtt{ in } e_2 \Rightarrow A \dashv \Gamma_1 + \Gamma_2}$$

### **Products**

$$\frac{\Gamma \vdash_i e_1 \Leftarrow A \dashv \Gamma_1 \quad \Gamma \vdash_i e_2 \Leftarrow B \dashv \Gamma_2}{\Gamma \vdash_i (e_1, e_2) \Leftarrow A \otimes B \dashv \Gamma_1 + \Gamma_2}$$

$$\Gamma \vdash_{i} e_{1} \Rightarrow A \otimes B \dashv \Gamma_{1} 
\Gamma, 0x : A, 0y : B \vdash_{i} e_{2} \Rightarrow C \dashv \Gamma_{2}, 1_{A}^{i} x : A, 1_{B}^{i} y : B$$

$$\Gamma \vdash_{i} \text{let } (x, y) = e_{1} \text{ in } e_{2} \Rightarrow C \dashv \Gamma_{1} + \Gamma_{2}$$

## Division

$$r_1/r_2 = \sup\{s \in \mathcal{Q} \mid s \cdot r_2 \sqsubseteq r_1\}$$

$r_1/r_2$	0	1	?	+	*
0	0	0	0	0	0
1		1	1	1	1
?		?	1	?	1
+		+	+	1	1
*		*	+	?	1

### Sums

$$\frac{\Gamma \vdash_{i} e \Leftarrow A \dashv \Gamma'}{\Gamma \vdash_{i} \text{ inl } e \Leftarrow A \oplus B \dashv \Gamma'} \frac{\Gamma \vdash_{i} e \Leftarrow B \dashv \Gamma'}{\Gamma \vdash_{i} \text{ inr } e \Leftarrow A \oplus B \dashv \Gamma'}$$

$$\frac{\Gamma, 0 \times : A \vdash_{i} e_{1} \Leftarrow C \dashv \Gamma_{2}, r_{1} \times : A}{\Gamma \vdash_{i} e \Rightarrow A \oplus B \dashv \Gamma_{1}} \frac{\Gamma, 0 \times : A \vdash_{i} e_{2} \Leftarrow C \dashv \Gamma_{3}, r_{2} \times : A}{\Gamma \vdash_{i} e \Rightarrow A \oplus B \dashv \Gamma_{1}} \frac{\Gamma, 0 \times : A \vdash_{i} e_{2} \Leftarrow C \dashv \Gamma_{3}, r_{2} \times : B}{r = (r_{1}/1_{A}^{i}) \sqcap (r_{2}/1_{B}^{i})}$$

$$\frac{\Gamma \vdash_{i} \text{ case } e \text{ of } \{x.e_{1}; y.e_{2}\} \Leftarrow C \dashv r \Gamma_{1} + (\Gamma_{2} \sqcap \Gamma_{3})}{\Gamma \vdash_{i} e \Rightarrow A \oplus B \dashv \Gamma_{1}} \frac{\Gamma, 0 \times : A \vdash_{i} e_{1} \Rightarrow C \dashv \Gamma_{2}, r_{1} \times : A}{\Gamma \vdash_{i} e \Rightarrow A \oplus B \dashv \Gamma_{1}} \frac{\Gamma, 0 \times : A \vdash_{i} e_{1} \Rightarrow C \dashv \Gamma_{3}, r_{2} \times : B}{r = (r_{1}/1_{A}^{i}) \sqcap (r_{2}/1_{B}^{i})}$$

$$\Gamma \vdash_{i} \text{ case } e \text{ of } \{x.e_{1}; y.e_{2}\} \Rightarrow C \dashv r \Gamma_{1} + (\Gamma_{2} \sqcap \Gamma_{3})$$

### Let

$$\frac{\Gamma \vdash_{i} e_{1} \Rightarrow A \dashv \Gamma_{1} \quad \Gamma, 0 \times A \vdash_{i} e_{2} \Rightarrow B \dashv \Gamma_{2}, r \times A \quad r' = r/1_{A}^{i}}{\Gamma \vdash_{i} \text{let } x = e_{1} \text{ in } e_{2} \Rightarrow B \dashv r' \Gamma_{1} + \Gamma_{2}}$$

### Natural numbers

$$\frac{\Gamma \vdash_{i} 0 \Leftarrow \mathbb{N} \dashv 0\Gamma}{\Gamma \vdash_{i} n \Leftarrow \mathbb{N} \dashv \Gamma'}$$

$$\frac{\Gamma \vdash_{i} n \Leftarrow \mathbb{N} \dashv \Gamma'}{\Gamma \vdash_{i} \operatorname{succ} n \Leftarrow \mathbb{N} \dashv \Gamma'}$$

$$\frac{\Gamma \vdash_{i} n \Rightarrow \mathbb{N} \dashv \Gamma_{1} \quad \Gamma \vdash_{i} z \Rightarrow A \dashv \Gamma_{2} \quad \Gamma \vdash_{i} s \Rightarrow A \to A \dashv \Gamma_{3}}{\Gamma \vdash_{i} \operatorname{elim}_{\mathbb{N}} z \ s \ n \Rightarrow A \dashv \Gamma_{1} + \Gamma_{2} + *\Gamma_{3}}$$