

# Modern Computer Architecture

## Lab Report

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### Abstract

For the practical assignment of the Modern Computer Architecture course we've changed the x264 program. We've profiled x264 and extracted the `x264_pixel_satd_8x4` as kernel function. To get a performance improvement we've moved the execution of the kernel function to the  $\rho$ -VEX, a co-processor of the MicroBlaze. A requirement for this is to load the program into the instruction memory and write and read the data to and from the data memory. We've found that our implementation did not generate a performance increase, but rather a decrease. We think this is due to the communication overhead.

## 1 Introduction

The assignment for this report was basically, find a kernel function from an application and run it on the  $\rho$ -VEX. To do this there are a few things we need to know first:

**Platform** The platform the program will run on is the ERA platform. It consists out of the host processor, the MicroBlaze, accelerated with a VLIW co-processor, the  $\rho$ -VEX. For the co-processor computationally intensive kernels can be extracted to achieve a performance increase.

**Kernel function** The extracted piece of code is called the kernel. The kernel needs to be compiled for the  $\rho$ -VEX so that we can inject the result, the byte-code, in to the co-processor. The rest of the code runs as usual on a regular processor.

**x264** The x264 program is the application we will try to improve using the  $\rho$ -VEX. x264 is a software library for encoding video streams in the H.264 compression format.

The goal for this report is to extract the correct kernel, using profiling and compile this for the  $\rho$ -VEX. To execute the kernel on the  $\rho$ -VEX there should be communication between the processor and the co-processor.

## 2 Profiling

To know what we are going to optimize, we need to profile the application first. Just taking a random function isn't a good idea. Fortunately we can compile the x264 program with the gprof profiling flags enabled. Compiling the program and running it on the Virtual Machine with the `~/Videos/inputs/eledream_640x360_8.y4m` as input video we get the following profiling results:

```
gprof x264 | head -n 10
Flat profile:
```

Each sample counts as 0.01 seconds.

% time	cumulative seconds	self seconds	calls	self ms/call	total ms/call	name
13.70	0.10	0.10	1599044	0.00	0.00	x264_pixel_satd_8x4
13.70	0.20	0.10	570708	0.00	0.00	get_ref
6.85	0.25	0.05	38770	0.00	0.00	x264_pixel_sad_x4_16x16
5.48	0.29	0.04	460484	0.00	0.00	quant_4x4
4.11	0.32	0.03	123076	0.00	0.00	sa8d_8x8

We see that the function `x264_pixel_satd_8x4` is called 1.6 million times during the video conversion. These calls together take up about 14% of the total time. An equal amount of time is spent in `x264_pixel_satd_8x4`.

The video `eledream_640x360_128.y4m` makes the application spend about 18% of the runtime in the pixel processing function and 16% in `get_ref`.

The bigger our input video the more time we will spend processing and the more effect an optimization will have if it targets part of the processing code.

Judging from the profiling information there are two potential places where optimization will be effective: `x264_pixel_satd_8x4` and `get_ref`.

When we looked at the source code of x264 we thought that the nature of `x264_pixel_satd_8x4` was more suitable for optimization because it had some loops and arithmetic in it. The `get_ref` function was a lot more irregular and also harder to understand. On the flip side, the overhead caused by communication between the MicroBlaze processor and the  $\rho$ -VEX would be smaller for `get_ref` because it is called three times less than `x264_pixel_satd_8x4`.

In the end we chose to try and put the computation of `x264_pixel_satd_8x4` on the co-processor and let `get_ref` for what it was. The function `x264_pixel_satd_8x4` was easier to understand and making it work was more important to us than choosing the best optimization area right away.

## 3 Data Layout

The `x264_pixel_satd_8x4` function has the following signature:

```
int x264_pixel_satd_8x4(
    pixel *pix1, int i_pix1,
    pixel *pix2, int i_pix2);
```

This is a function that takes two arrays of pixels, something that has to do with the indices of the pixels in the array, and returns a single integer.

From the implementation we can see that the data used by the array is always  $4 \times 8 = 32$  pixels per array. This means that the data layout should have at least two times 32 times the size of one pixel. The size of a pixel is defined as `uint_8`. If we choose `i_pix1` and `i_pix2` to always be 8, we don't have to send those values, which saves sending data. Setting those parameters to a fixed value has also the advantage that the array indexes looked up by the function are always predictably 0 to 32

Also for the return value we should reserve 4 bytes, which is the size of an integer. When we have calculated the result, we could write this back to the first position, but we write the result after the data of the two pixel arrays. The advantage for doing this is because it's slightly easier and clearer for future maintainers.

## 4 Implementation

We have access to several (but not enough during the lab) FPGAs which are configured as MicroBlaze processors and run Linux.

The  $\rho$ -VEX is configured as a co-processor that can be controlled using a number of memory mapped files. There is a file for writing the instruction memory, one for reading and writing the data memory, one for reading the status and one for writing control commands.

We abstracted this away to a small interface with the following functionality:

### **rvexInit**

Attempts to open the IMEM, DMEM and SMEM files. You can also specify the bytecode which will be written to the instruction memory.

### **rvexDispose**

Closes all files opened by `rvexInit`.

### **rvexWrite**

Allows you to write to the data memory.

### **rvexRead**

Allows you to read from the data memory.

### **rvexSeek**

Allows you to jump to a given position in the data memory.

### **rvexGo**

Writes the clear and start commands to the control memory and blocks until the status reports that the operation was successful.

## 4.1 Endianness

Since we had to compile for the MicroBlaze using the flag `-DWORD-BIGENDIAN` we figured that the MicroBlaze would be a big endian machine. You can always test it by writing a multibyte value like `0xDEADBEEF` to memory and read the individual bytes. If you read `0xDE 0xAD 0xBE 0xEF` you will know that you have a big endian machine. If you get that sequence but in reverse you know its a little endian machine.

We write the data as big endian to the data memory. For types bigger than one byte we change the endianness on the  $\rho$ -VEX side using a macro. Also before writing data to the memory from the  $\rho$ -VEX the endianness should be changed. For the `x264_pixel_satd_8x4` kernel the size of the pixels is one byte, only the result is an integer, so four bytes, which should be reversed.

## 4.2 Loading Instruction memory

Loading the instruction memory with separate commands before executing the program on the MicroBlaze is pretty inconvenient. That is why we write the bytedata to the instruction memory in the application. The compiled kernel function is actually converted to a piece of C code and saved in a file which can then directly be written to the instruction memory. This solves the inconvenience of multiple files for the same program.

## 5 Results of using the $\rho$ -VEX

We modified the x264 application to log its processing time computed with `clock_gettime`. This required linking the `rt` library.

Then we tried to find a FPGA that was not being used by anyone else and we ran several versions of the x264 application. All of them included the code from Listing 1 which allowed us to get an idea of the total runtime.

Listing 1: Capturing runtime with the monotonic clock

```
struct timespec tss, tse, tsd; // start, end and diff
clock_gettime(CLOCK_MONOTONIC, &tss);

if( !ret )
    ret = encode( &param, &opt );

clock_gettime(CLOCK_MONOTONIC, &tse);
if (tse.tv_nsec > tss.tv_nsec) {
    tsd.tv_sec = tse.tv_sec - tss.tv_sec;
    tsd.tv_nsec = tse.tv_nsec - tss.tv_nsec;
} else {
    tsd.tv_sec = tse.tv_sec - tss.tv_sec - 1;
    tsd.tv_nsec = tse.tv_nsec - tss.tv_nsec + 1000000000;
}
printf("Took %lu.%09lu sec\n", tsd.tv_sec, tsd.tv_nsec);
```

The versions that we created were the following

**vanilla**

The original x264 implementation.

**rvex simple**

Simple implementation where memory mapped files are opened once for stability during execution and only one  $\rho$ -VEX is supported at a time.

**rvex inline**

Based on simple, uses an object to store the file descriptors so allows multiple  $\rho$ -VEX processors to be used and the time critical functions are placed in a header file with the inline annotation.

We ran the script from Listing 2 to get an idea of the runtimes of the different versions. Note that we already did manual test runs to get an idea of how long the runs would take. We noticed that running the same test multiple times produced very similar run times. We did not find it necessary to do multiple runs for the purposes of this lab because the error would be very small in comparison with the run times for different versions of the application:

Took 75.280923348 sec

Took 75.350194581 sec

Took 75.221714038 sec

Took 75.374499116 sec

Listing 2: Test script

```
echo 'eledream 64x36 3 frames with timing' >> group11.log
./x264-timing-sb2 eledream_64x36_3.y4m -o out.mkv | grep Took >> group11.log

echo 'eledream 640x360 8 frames with timing' >> group11.log
./x264-timing-sb2 eledream_640x360_8.y4m -o out.mkv | grep Took >> group11.log

echo 'eledream 64x36 3 frames with rvex interface simple' >> group11.log
./x264-rvex-sb2 eledream_64x36_3.y4m -o out.mkv | grep Took >> group11.log

echo 'eledream 64x36 3 frames with rvex interface struct and inline' >> group11.log
./x264-rvex-struct-inline-sb2 eledream_64x36_3.y4m -o out.mkv | grep Took >> group11.log
```

Listing 3: Test output

```
/ # cat group11.log
eledream 64x36 3 frames with timing
Took 75.699494266 sec
eledream 640x360 8 frames with timing
Took 317.733194410 sec
eledream 64x36 3 frames with rvex interface simple
Took 442.169817935 sec
eledream 64x36 3 frames with rvex interface struct and inline
Took 440.886382661 sec
```

Table 1: Execution times

	64x36 3f	640x360 8f
vanilla	75.7s	317s
rvex simple	442s	>20m
rvex inline	441s	untested

Table 1 summarizes the results found by manual testing and the listed script.

Most importantly, we see that our optimization actually made the application run several times slower. Also, we can conclude that the gcc compiler does a terrific job at inlining functions when `-O3` is enabled, even if you do not ask it to. We couldnt wrap our heads around doing so much communication related tasks in a function that was basically a relatively small fixed number of basic operations (\*, +, -, shifts, comparisons) and achieving a faster result by that. We assumed that the  $\rho$ -VEX was extremely tightly integrated into the MicroBlaze which allowed very good communication speeds but the test results show that this is not the case.

The speedup that we obtained can be calculated by looking at the time spent in the kernel function. This should be about 14% for the input video `eledream_640x360_8.y4m` judging from the profiling results obtained on the VM. We cannot say exactly how much time is spent there because we cannot profile on the MicroBlaze.

Because the  $\rho$ -VEXsimple version took longer than 20 minutes which caused the SSH session to hang we cannot use the exact run time. Let us see what the speedup would be if the run time was 20 minutes exactly.

The vanilla version spends 14% of the 75 seconds that it runs in `x264_pixel_satd_8x4`. So 86% is spent in the rest of the functions. Since we have not changed anything beside the kernel we can assume that the time spent in that part will be equal for vanilla and for rvex simple. This means that  $\frac{1200-0.86\cdot75}{75-0.86\cdot75} = 108$  is the relative speed of the vanilla function in comparison with the rvex version. It actually runs 108 times faster. In other words, the rvex version runs 108 times slower which amounts to a speedup of -10700% obtained by introducing rvex.

Theoretically, the perfect optimization for the selected kernel would be to reduce the computation time to 0. This means that the 14% becomes 0% and so the maximum attainable speedup is  $\frac{100}{100-14} - 100 = 16\%$

## 6 Additional Assignment

Except for the joy of achieving an actual speedup, we did not think that extracting another function would teach us a whole lot more about the process. It would in our eyes teach us a lot about the x264 application but not without severe frustration and confusion as to what is going on.

So we set a challenge for ourselves which was to figure out a way of improving the lab itself. We noticed that some of the people were not that experienced with C and due to the performance heavy x264 application the small number of

$\rho$ -VEX machines, testing during the lab was an absolute pain and took a lot of valuable time away.

The new assignments are placed in four categories:

### **getting your C going**

adder - write a C application that takes two integers as arguments and adds them, compile it for the VM

pow - write a C application that takes two integers  $a$  and  $b$  where  $b \geq 0$  and calculates  $a^b$  without the math library, compile and run it for the Microblaze FPGA

### **using file operations**

bin2c - create a binary to c-code converter, read in a binary file and print a character array with the bytes in it.

### **introducing rvex**

rvex-adder - write an rvex kernel that adds two integers

rvex-pow - write an rvex kernel that multiplies two integers, use it to calculate the power

### **abstracting rvex**

rvex-lib - write a simple interface for rvex.c and rvex.h containing the following functions

1. rvexInitialize: initialize rvex object
2. rvexDispose: clean up rvex object
3. rvexBytecode: set the bytecode for the rvex
4. rvexSeek: set the data cursor position
5. rvexRead: read from data memory
6. rvexWrite: write to data memory
7. rvexGo: start the rvex and block until the operation is finished

They should operate on an Rvex object that you must define: struct, extern.

Apply the rvex-lib to your rvex-adder and then to your rvex-pow application. Test if everything still is working as intended.

### real-life application

rvex-x264 - modify x264 to use a kernel

You can supply a basic project setup with Makefiles. Each category can use its own introduction giving hints such as use open instead of fopen for memory mapped files.

The adder application is useful for its extreme simplicity. It is easy to test and discover argument passing and endianness handling (for the rvex version).

The pow application is useful because it requires multiple calls to rvexGo. This will remind students of resetting the data memory cursor for the rvex version.

Listing 4: pow - Microblaze part

```
#include <stdio.h>
#include <stdlib.h>
#include "rvex.h"
#include "bytecodes/mult.h"

int main(int argc, char **argv) {

    if (argc != 3) {
        printf("Use like this: '%s n c' to let the "
               "program calculate n^1..c\n", argv[0]);
        return -1;
    }

    rvexInit(&rvex0, bytecode, sizeof(bytecode),
            RVEX_O_INSTRUCTION_MEMORY_FILE,
            RVEX_O_DATA_MEMORY_FILE,
            RVEX_O_CORE_CTL_FILE,
            RVEX_O_CORE_STATUS_FILE);

    int n = atoi(argv[1]);
    int c = atoi(argv[2]);
    int temp = 1;
    while(c > 0) {
        rvexSeek(&rvex0, 0);
        rvexWrite(&rvex0, &n, sizeof(int));
        rvexWrite(&rvex0, &temp, sizeof(int));
        rvexGo(&rvex0);
        rvexRead(&rvex0, &temp, sizeof(int));
        c--;
    }

    printf("n^c = %d\n", temp);

    rvexDispose(&rvex0);

    return 0;
}
```



It is a simple program but uses all the functionality with rvex that you need for x264.

Listing 5: pow - kernel part

```
int a = 0xAAA00A00, b = 0xBBB00B00, result = 0xCCC00C00;

#define FLIP_ENDI_32(a) ( \
    (((a)&0x000000FF) << 24) | \
    (((a)&0x0000FF00) << 8) | \
    (((a)&0x00FF0000) >> 8) | \
    (((a)&0xFF000000) >> 24) \
)

int main () {
    a = FLIP_ENDI_32(a);
    b = FLIP_ENDI_32(b);
    result = a * b;
    result = FLIP_ENDI_32(result);
    return 0;
}
```

It is possible to distribute an already filled in rvex.h and leave the implementation to the students if they are running out of time. At first it is not even necessary to give that much detail about a possible implementation.

## 7 Conclusion

We've identified a kernel function and changed the x264 program so the kernel is executed on the  $\rho$ -VEX. The kernel we used is the `x264_pixel_satd_8x4` function. Using our abstracted `rvex*` functions you can write the instruction memory, write and read the data memory and start the  $\rho$ -VEX. Unfortunately we didn't see a speedup, instead we even saw a decrease in performance. Assumably this is because of the communication overhead. The `x264_pixel_satd_8x4` function is called many times. Maybe by sending more data at once, and thus executing the  $\rho$ -VEX less times might improve the performance.

Using our `rvex*` functions, we were able to create different test programs which helped us testing the communication and get to know the platform, getting the endianness right and fixing multiple executions of the  $\rho$ -VEX.

All in all the lab was definitely fun. It felt really rewarding to hack something together, inject some bytecode and see it working.

Regarding the results, we have some suggestions that we would have tried out given enough time.

We still believe in the  $\rho$ -VEX so we would try to reduce the performance loss of the current bottleneck: communication. This can be done by reducing the number of reads and writes to the  $\rho$ -VEX. This means selecting a bigger part of the application and send it over to the  $\rho$ -VEX in its entirety.

We are interested in if the  $\rho$ -VEX performance is dependent on the type of program that is run. Is it more efficient for highly repeating code or does it perform just as well for dynamic code with little repetition compared to the normal CPUs?