



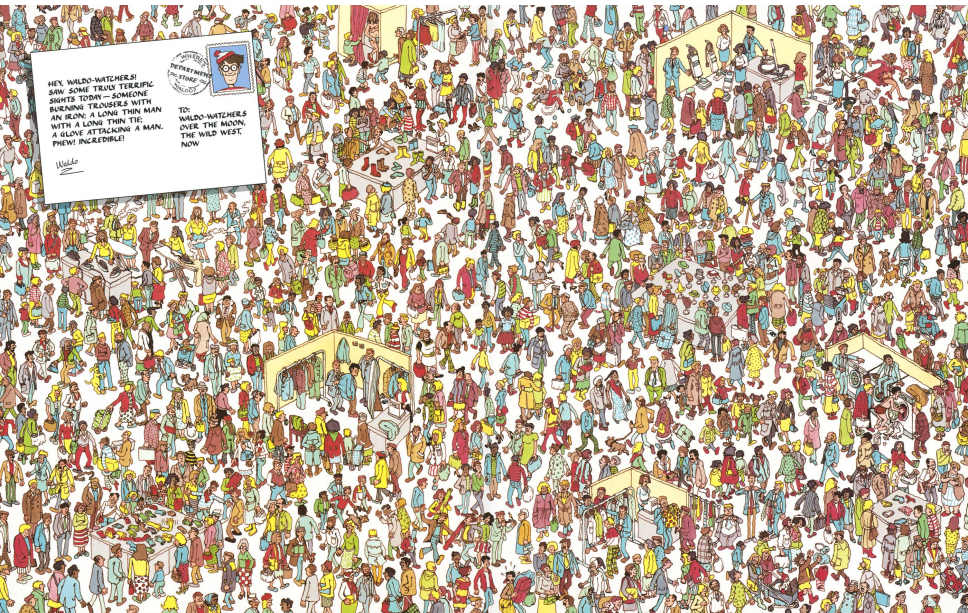
## Pattern Mining on Matrices

Micky Faas

March 14, 2019

With VOUW we propose a method of **mining patterns on matrices**. It encompasses:

- \* A theoretical framework
- \* An optimization problem formulated using MDL
- \* A heuristic algorithm.



Say we have some  $M \times N$  matrix  $A$ , we want to discover and extract recurring structure. Why?

- \* Data analysis
- \* Distance metric for comparison
- \* Clustering



Term?	Form	Encodes...
<b>Pattern</b>	Submatrix of $A$	...relative positions of elements
<b>Offset</b>	$(i, j) \in M \times N$	...position of entire pattern
<b>Instance</b>	$M \times N$ matrix	...absolute positions of elements
<b>Instantiation</b>	$M \times N$ matrix	...what pattern's instance should be at which index

$$\underbrace{\begin{bmatrix} 1 & \cdot \\ \cdot & 1 \end{bmatrix}}_{\text{Pattern}}, \quad \underbrace{\begin{bmatrix} \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & 1 & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & 1 & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \end{bmatrix}}_{\text{Instance with offset (2,2)}}$$

$$\begin{array}{c}
 \text{Original matrix} \\
 A = \begin{bmatrix} 1 & \cdot & \cdot & \cdot & \cdot & 1 \\ \cdot & 1 & \cdot & \cdot & \cdot & \cdot \\ 1 & \cdot & \cdot & \cdot & 1 & \cdot \\ \cdot & 1 & \cdot & \cdot & \cdot & 1 \\ 1 & \cdot & 1 & \cdot & \cdot & \cdot \\ \cdot & 1 & \cdot & 1 & \cdot & \cdot \end{bmatrix}, G = \begin{array}{c} \text{Instantiation} \\ \begin{bmatrix} X & \cdot & \cdot & \cdot & \cdot & Y \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ X & \cdot & \cdot & \cdot & X & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ X & \cdot & X & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \end{bmatrix} \end{array} \\
 \\
 \begin{array}{c} \text{Model} \\ H = \left\{ X = \underbrace{\begin{bmatrix} 1 & \cdot \\ \cdot & 1 \end{bmatrix}}_{\text{Pattern}}, Y = \underbrace{\begin{bmatrix} 1 \end{bmatrix}}_{\text{Pattern}} \right\} \end{array}
 \end{array}$$

We can construct complex patterns by repeatedly combining simpler ones. We use instances for this as they encode the position of one pattern relative to another.

$$\begin{array}{ccccccc}
 \underbrace{X = [0]} & & & & & & \\
 \underbrace{Y = [1]} & \longrightarrow & \bar{X} = \begin{bmatrix} \cdot & \cdot & \cdot \\ 0 & \cdot & \cdot \\ \cdot & \cdot & \cdot \end{bmatrix} & \longrightarrow & \bar{X} + \bar{Y} = \begin{bmatrix} \cdot & \cdot & \cdot \\ 0 & \cdot & \cdot \\ \cdot & 1 & \cdot \end{bmatrix} & \longrightarrow & \underbrace{\begin{bmatrix} 0 & \cdot \\ \cdot & 1 \end{bmatrix}} \\
 \text{Singleton patterns} & & \underbrace{\begin{bmatrix} \bar{Y} = \begin{bmatrix} \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot \\ \cdot & 1 & \cdot \end{bmatrix} \end{bmatrix}}_{\text{Instances}} & & \underbrace{\hspace{1cm}}_{\text{Matrix sum}} & & \text{New pattern}
 \end{array}$$

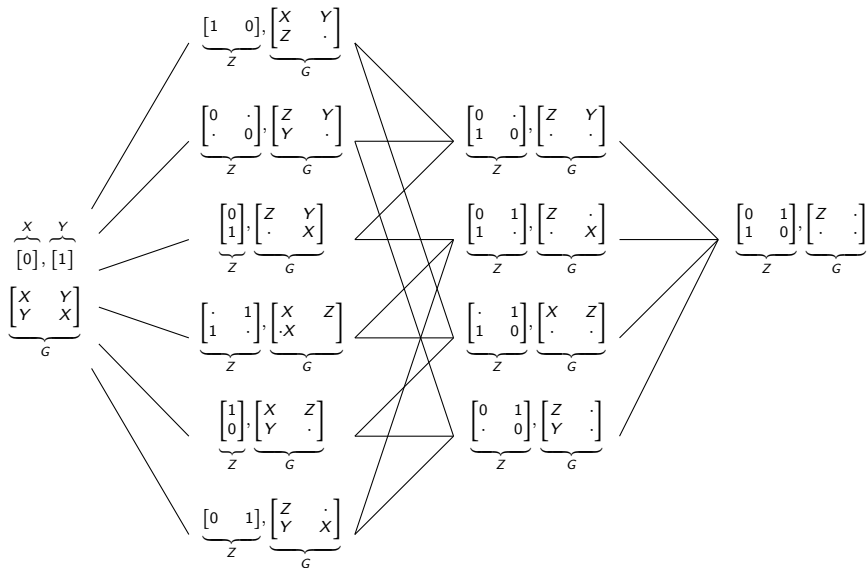


Idea: we inductively define the model space by starting with singletons for each element of  $A$ . Then we repeatedly merge instances until we finally obtain a pattern that equals  $A$ .

- \* At the beginning we have the completely underfit model
- \* We end up with a completely overfit model

Everything in between is a possible solution, but *we want the solution that describes  $A$  best.*

For example  $A = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$ : what is the space of possible solutions?



How complex is this lattice? Very complex!

Idea: pick two instances to combine  $MN - 1$  times.

$$\prod_{n=0}^{MN-2} \binom{MN-n}{2} = \prod_{n=0}^{MN-2} \frac{MN-n}{2(MN-n-2)!} = \frac{(MN)!(MN-1)!}{2^{MN-1}}.$$

We will need to use heuristics.

Idea: the best solution is a balance of model and instantiation complexity. We use two-part MDL:

$$\underbrace{L(H)}_{\text{Model}} + \underbrace{L(A|H)}_{\text{Data given model}}$$

In this case:

$$L\left(\underbrace{\begin{pmatrix} \text{Pattern} & \text{Pattern} \\ \begin{bmatrix} 1 & \cdot \\ \cdot & 1 \end{bmatrix} & \begin{bmatrix} 1 \end{bmatrix} \end{pmatrix}}_{\text{Model}}\right) + L\left(\underbrace{\begin{pmatrix} \text{Instantiation Matrix} \\ \begin{matrix} X & \cdot & \cdot & \cdot & \cdot & Y \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ X & \cdot & \cdot & \cdot & X & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ X & \cdot & X & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \end{matrix} \end{pmatrix}}_{\text{Data given model}}\right)$$

Why does this work? MDL is founded on these principles:

**Kolmogorov Complexity** of given data is the shortest computer program to produce that data.

**Occam's Razor** . Given competing hypotheses, pick the one with the fewest assumptions.

**No-hypercompression theorem** . Data with no inherent structure, cannot be compressed.

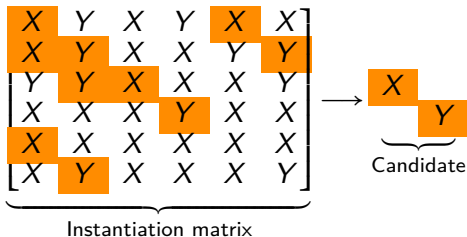
**Kraft Inequality** . One-to-one correspondence between code lengths and probabilities.

Idea: look at  $A$  once to build an instantiation matrix  $G$ . Make a model  $H$  using only singleton patterns, one for each possible value in  $A$ .

Then we constantly refine our description by merging instances to form patterns. Once we merge instances, we never reconsider (greedy approach).

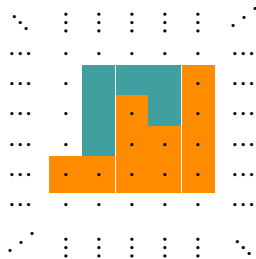
$$\begin{array}{c} \overbrace{[0]}^X, \overbrace{[1]}^Y \\ \underbrace{\begin{bmatrix} X & Y \\ Y & X \end{bmatrix}}_G \end{array} \rightarrow \underbrace{\begin{bmatrix} 1 & 0 \end{bmatrix}}_Z, \underbrace{\begin{bmatrix} X & Y \\ Z & \cdot \end{bmatrix}}_G \rightarrow \underbrace{\begin{bmatrix} 0 & \cdot \\ 1 & 0 \end{bmatrix}}_Z, \underbrace{\begin{bmatrix} Z & Y \\ \cdot & \cdot \end{bmatrix}}_G$$

Which pair of instances do we combine? We will first derive a list of candidates from the instantiation matrix:



The candidate that minimizes the MDL equation best is picked.

We constrain the idea of the last slides by only looking at **adjacent** instances.



This gives two benefits: (1) visits all candidates just once, (2) reduces candidate space.



This simple algorithm gives an acceptable complexity:

- \* Each iteration (find candidates and merge best) is (almost) linear.
- \*  $MN - 1$  iterations worst-case, but we need only a fraction in any practical case.

However:

- \* Greedy approach sometimes gives completely wrong results.
- \* No tolerance to noise.

A 'wish list' for the future:

- \* Actual data and use-cases (!)
- \* Noise robustness.
- \* Multi-instance candidates (more than 2).
- \* Hierarchical encoding of patterns.
- \* Improvement of the current (prequential) encoding.

