

Load Balancing Problem

(Approximation Example)

Input:

m resources with equal processing power

n jobs, where job j takes t_j to process

objective:

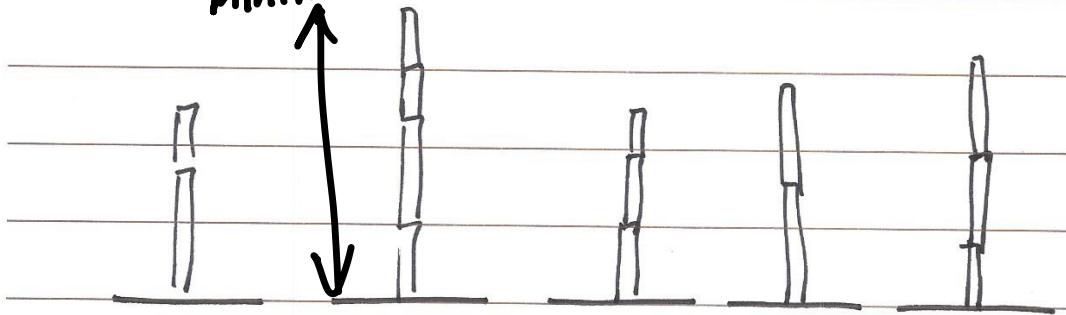
Assignment of jobs to resources such
that the maximum load on any machine
is minimized.

Notation:

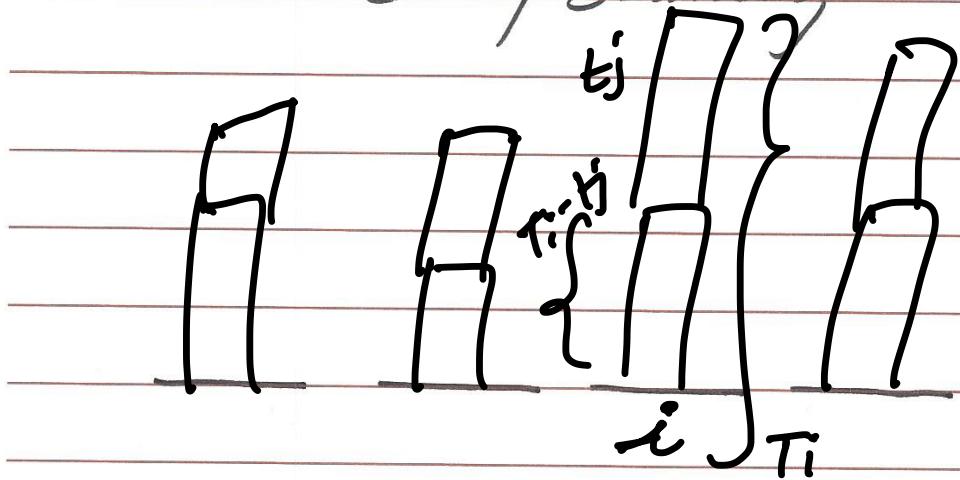
T_i : load on machine/resource i

T^* : value of the optimal solution

minimize this



Greedy Balancing



$\hat{\tau}^* \geq t_j$ load of any process

$$\underline{\tau^* \geq \tau_i - t_j}$$

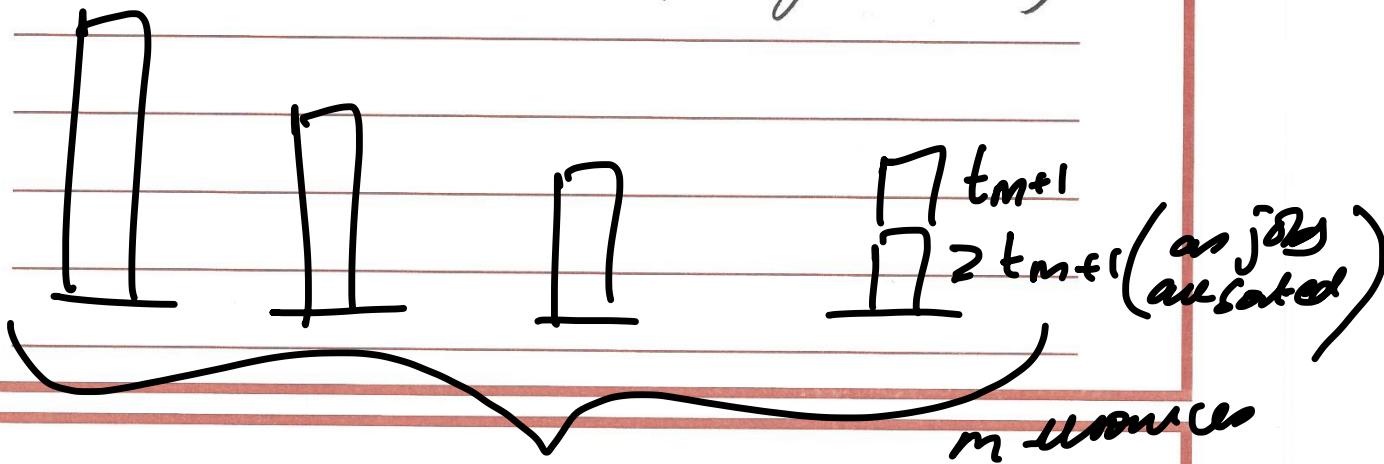
$$2\tau^* \geq \tau_i$$

$$\boxed{\tau_i \leq 2\tau^*}$$

2 - approximation

Improved approximation to greedy balancing

Initially sort jobs in decreasing order of
length then use same greedy balancing



$$\tau^* \geq 2t_{m+1}$$

$$t_j \leq t_{m+1} \rightarrow t_j \leq \frac{1}{2}\tau^*$$

$$t_j \leq \frac{1}{2}\tau^*$$

$$\tau_i - t_j \leq \tau^*$$

$$\boxed{\tau_i \leq \frac{3}{2}\tau^*}$$

1.5 approximation

Vertex Cover Problem

(Approximation Example)

Problem Statement: Find the smallest vertex cover in graph G.

A 2-approximation algorithm to the vertex cover problem:

Start with $S = \text{Null}$

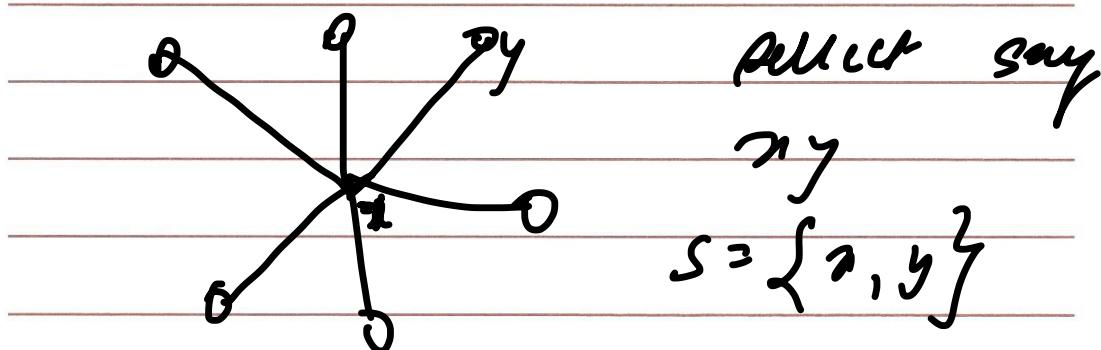
While S is not a vertex cover

Select an edge e not covered by S

Add both ends of e to S

Endwhile

Why is this a 2-approximation?



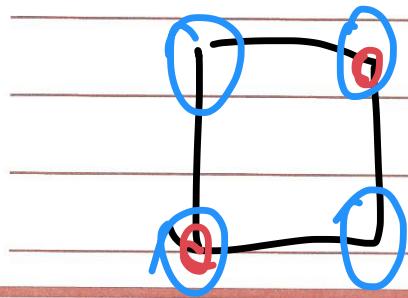
$$\text{Optimal} = \{x\}$$

for every edge we are selecting union of 2 nodes while optimal solution may only have 1 vertex for each edge

Question: Since $\text{Independent Set} \leq_p \text{Vertex Cover}$,

Can I use our 2-approximation algorithm for vertex cover to find a .5-approximation to independent set?

No



Theorem: Unless $P=NP$, there is no $\frac{1}{n^{1-\epsilon}}$ approximation for the Maximum Independent Set problem for any $\epsilon > 0$ where n is the no. of nodes in the graph.

Question: Since $\text{Vertex Cover} \leq_p \text{Set Cover}$

Can I use a 2-approximation algorithm for
set cover to find a 2-approximation to
vertex cover?

Yes

MAX-3SAT

Given a set of clauses of length 3, find a truth assignment that satisfies the largest number of clauses.

A 0.5-approximation to the MAX-3SAT problem:

- set everything to true

If less than 50% of clauses evaluate to true, then

- set everything to False

Worst

More sophisticated approximation methods can get to within a factor of $\frac{8}{9}$ of the optimal sol.

System of Linear Equations

$$[A][x] = [B]$$

Linear Programming

$$[A][x] \leq [B]$$

Objective function: $[C^T][x]$

Goal: Maximize the objective function
subject to the above constraints

Linear Programming Standard Form

- All constraints are of the form \leq
- All variables are non-negative
- Objective function is maximized.

Ex.: $x_1 - x_2 \geq 0$

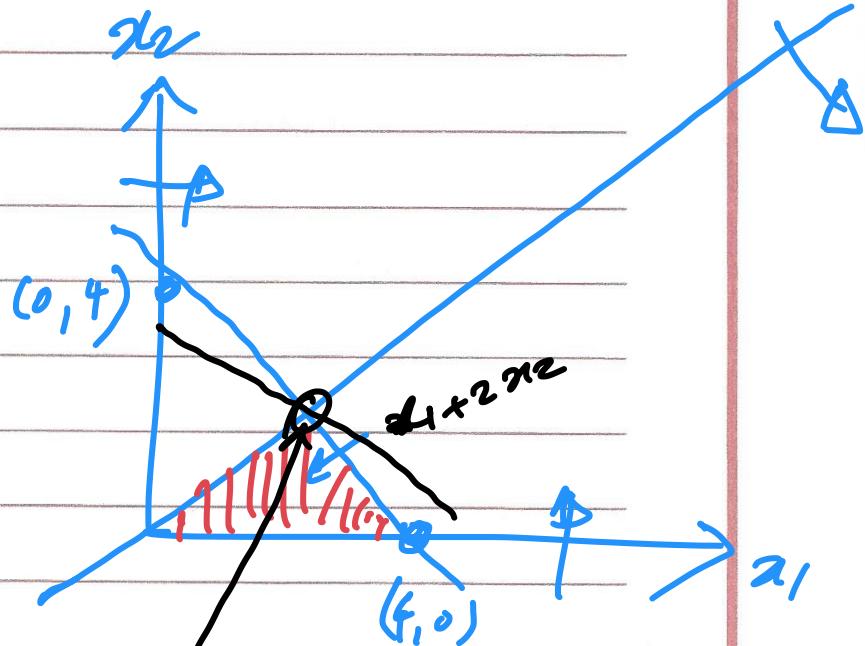
$x_1 \geq 0$

$x_2 \geq 0$

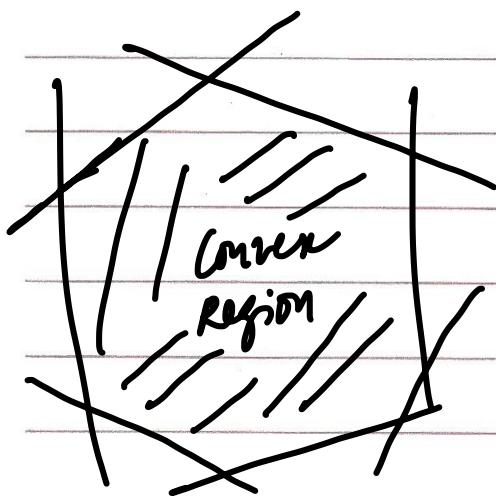
$x_1 + x_2 \leq 4$

objective
maximize

$x_1 + 2x_2$



This point is
in feasible region
and it maximized
the $x_1 + 2x_2$



Any line drawn between
two points of convex
region lie inside the
convex region.

The point that maximizes the objective function is on one of vertices of the region formed.

Simpson Algorithm (worst case performance exponential)

Suppose we are going clockwise we first keep on going till we ↑ the objective function; it starts decreasing so say this is the maximum value of objective function

Linear Programming LP solved

Weighted Vertex Cover Problem

For $G = (V, E)$, $S \subseteq V$ is a vertex cover set such that each edge has at least one end in S

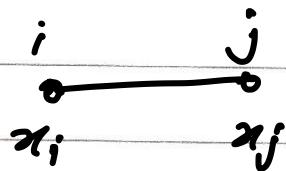
Also, $w_i \geq 0$ for each $i \in V$

So, the total weight of the set = $w(S) = \sum_{i \in S} w_i$

objective: Minimize $w(S)$

x_i is deviation variable for each node $i \in V$

$$\begin{cases} x_i = 0 & \Rightarrow i \notin S \\ x_i = 1 & \Rightarrow i \in S \end{cases}$$



$$x_i + x_j \geq 1$$

Minimize

$$\sum w_i x_i$$

Subject to $x_i + x_j \geq 1$ for $(i, j) \in E$

$$x_i \in \{0, 1\}$$

↳ x_i are discrete values

Since this is not linear programming

It is integer linear programming

belongs to NP

- Linear Programming
Continuous Variables
- Integer Programming
Discrete variables
- Mixed Integer Programming
Both Cont. & Discrete
- Non - Linear Programming
Non Linear Constraints
or Objective Function

finding approximate solution as prob is NP

To find an approximate solution using LP,
drop the requirement that $x_i \in \{0,1\}$ and
solve the LP in polynomial time to
find $\{x_i^*\}$ between 0 & 1.

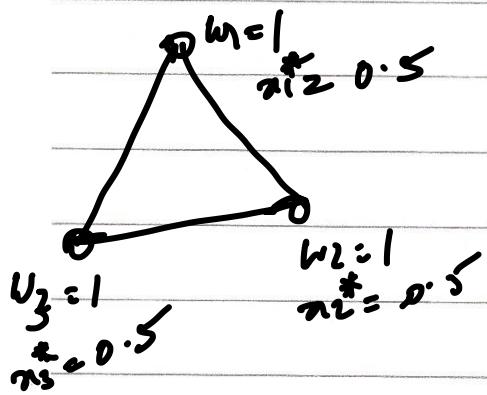
$$W_{LP} = \sum_i w_i x_i^*$$

Assume S' is the optimal vertex cover set
and $W(S')$ is the weight of the opt. solution

weight of our LP solution

$$W_{LP} \leq W(S')$$

↑ weight of
optimal solution



$$\begin{aligned} W_{LP} &= 1 \times 0.5 + 1 \times 0.5 + 1 \times 0.5 \\ &= 1.5 \end{aligned}$$

$$\underline{W(S') = 2}$$

If

$$\pi_i^* = 0 \longrightarrow i \notin S$$

$$\pi_i^* = 1 \longrightarrow i \in S$$

Say $S_{\geq \frac{1}{2}} = \left\{ i \in V : \pi_i^* \geq \frac{1}{2} \right\}$

- check if it is vertex cover (cr. 12)



$$\pi_i^* \quad \pi_j^*$$

edge (i, j) not covered if and only if

$$\pi_i < 0.5 \text{ and } \pi_j < 0.5$$

But this is not possible

$$\Rightarrow \boxed{\pi_i + \pi_j \geq 1}$$

\therefore atleast one of π_i or $\pi_j \geq 0.5$

\therefore It is vertex cover set

$\hat{w}_F \geq 0.5 \rightarrow$ make it 1

Say s is one approx solⁿ

$$w(s) \leq d \cdot w_{LP}$$

$$w_{LP} \leq w(s')$$

$$w(s) \leq 2 \cdot w(s')$$

2-approximation

Maxflow Problem

Variables are flows over edges

objective function: Maximize $\sum_{e \text{ out of } S} f(e)$

Subject to

$$0 \leq f(e) \leq c_e \text{ for each edge } e \in E$$

$$\sum_{e \text{ into } v} f(e) - \sum_{e \text{ out of } v} f(e) = 0 \quad \begin{array}{l} \text{for } v \in V \\ \text{except for } S \text{ & } t \end{array}$$

THIS TYPE

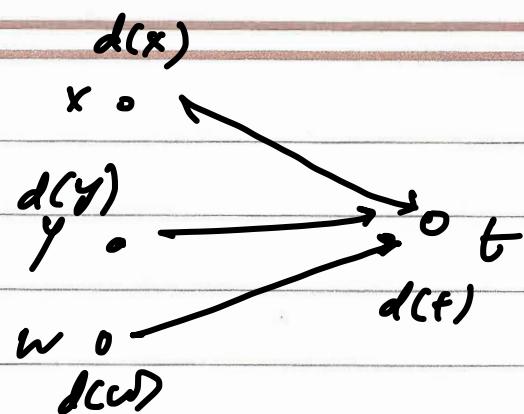
$$A = B \Rightarrow \left\{ \begin{array}{l} A \leq B \\ B \leq A \end{array} \right.$$

Shortest Path using LP

Problem Statement:

Find shortest path from s to t .

shortest distance from s to v is $d(v)$
for each v



$$d(t) \leq d(y) + c_{yt}$$

$$d(t) \leq d(z) + c_{zt}$$

$$d(t) \leq d(w) + c_{wt}$$

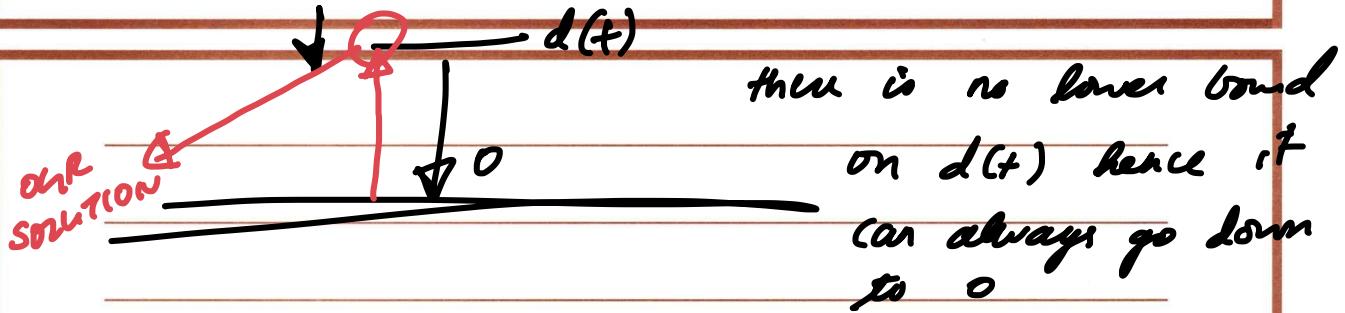
In general

$$\left\{ \begin{array}{l} d(v) \leq d(u) + w(u, v) \\ d(S) = 0 \end{array} \right. \quad \boxed{\text{for each edge } (u, v) \in E}$$

Objective: ~~minimize~~ ^{maximize} $d(t)$

problem

\downarrow $d(f)$
 \downarrow $d(t)$



Discussion 12

1. The *bin packing* problem is as follows. You have an infinite supply of bins, each of which can hold M maximum weight. You also have n objects, each of which has a (possibly distinct) weight w_i (any given w_i is at most M). Our goal is to partition the objects into bins, such that no bin holds more than M total weight, and that we use as few bins as possible. This problem in general is **NP-hard**.

Give a 2-approximation to the *bin packing* problem. That is, give an algorithm that will compute a valid partitioning of objects into bins, such that no bin holds more than M weight, and our algorithm uses at most twice as many bins as the optimal solution. Prove that the approximation ratio of your algorithm is two.

2. Suppose you are given a set of positive integers $A: a_1 a_2 \dots a_n$ and a positive integer B . A subset $S \subseteq A$ is called **feasible** if the sum of the numbers in S does not exceed B .

The sum of the numbers in S will be called the *total sum* of S . You would like to select a feasible subset S of A whose total sum is as large as possible.

Example: If $A = \{8, 2, 4\}$ and $B = 11$ then the optimal solution is the subset $S = \{8, 2\}$.

Give a linear-time algorithm for this problem that finds a feasible set $S \subseteq A$ whose total sum is at least half as large as the maximum total sum of any feasible set $S' \subseteq A$. Prove that your algorithm achieves this guarantee.
You may assume that each $a_i \leq B$.

3. A cargo plane can carry a maximum weight of 100 tons and a maximum volume of 60 cubic meters. There are three materials to be transported, and the cargo company may choose to carry any amount of each, up to the maximum available limits given below.

- Material 1 has density 2 tons/cubic meter, maximum available amount 40 cubic meters, and revenue \$1,000 per cubic meter.
- Material 2 has density 1 ton/cubic meter, maximum available amount 30 cubic meters, and revenue \$1,200 per cubic meter.
- Material 3 has density 3 tons/cubic meter, maximum available amount 20 cubic meters, and revenue \$12,000 per cubic meter.

Write a linear program that optimizes revenue within the constraints. You do not need to solve the linear program.

4. Recall the 0/1 knapsack problem:

Input: n items, where item i has profit p_i and weight w_i , and knapsack size is W .

Output: A subset of the items that maximizes profit such that the total weight of the set $\leq W$.

You may also assume that all $p_i \geq 0$, and $0 < w_i \leq W$.

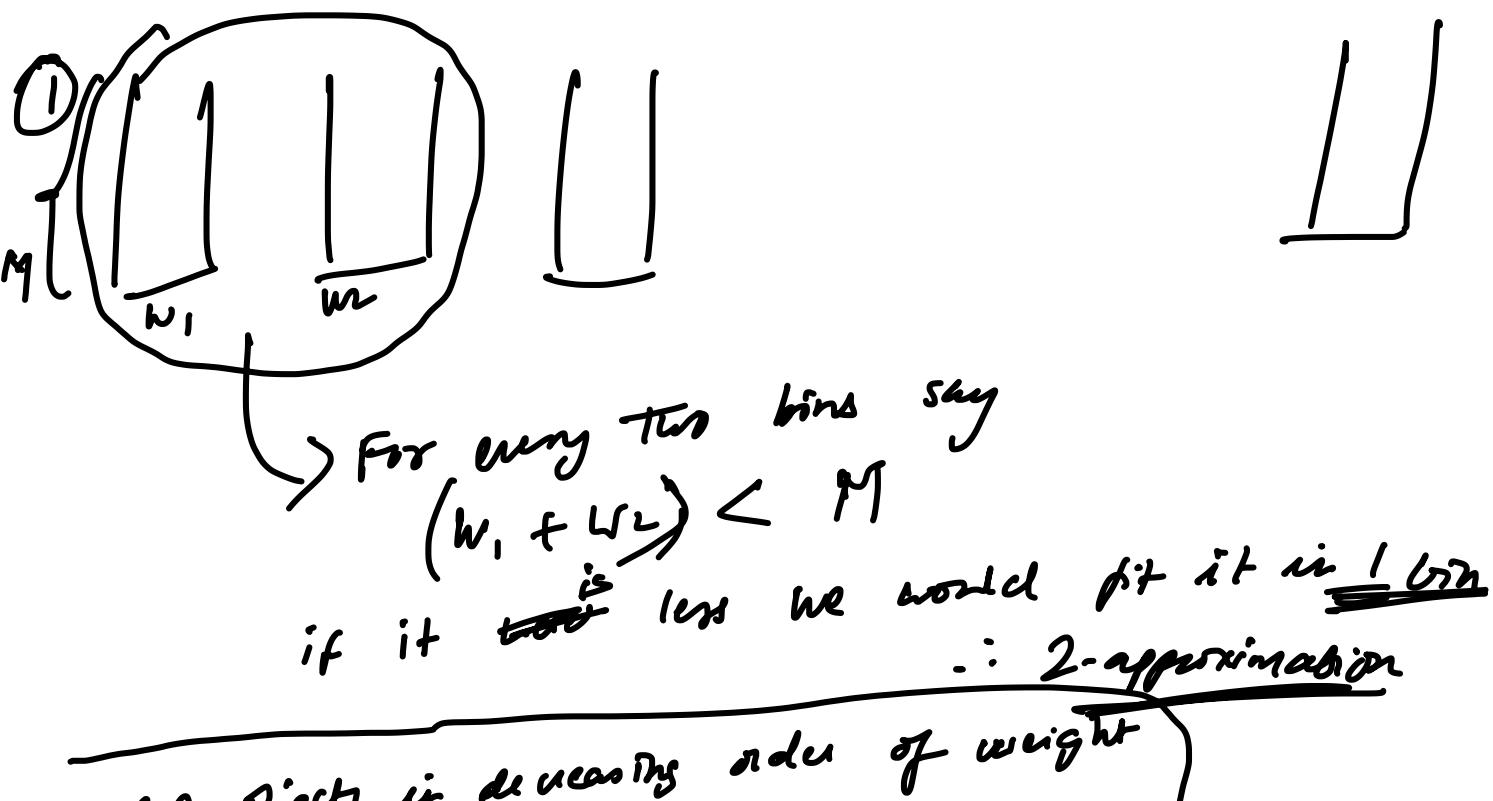
We have created the following greedy approximation algorithm for 0/1 knapsack:

Sort items according to decreasing p_i/w_i (breaking ties arbitrarily).

While we can add more items to the knapsack, pick items in this order.

Show that this approximation algorithm has no nonzero constant approximation ratio.

In other words, if the value of the optimal solution is P^* , prove that there is no constant ρ ($1 > \rho > 0$), where we can guarantee our greedy algorithm to achieve an approximate solution with total profit ρP^* .



sort objects in decreasing order of weight



item 1



item 2

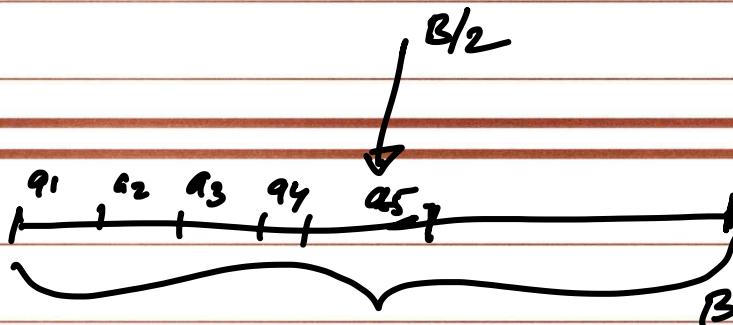
always start to fit items from bin 1

First fit Decreasing

→ 1.22 Approximation

modified FFD → 1.18 Approx

②



either keep going till sum = $B/2$

or we have one item with weight $> B/2$

③ Let M_1, M_2, M_3 denote cubic metres of materials 1, 2, 3

Maximize

$$1000M_1 + 1200M_2 + 1200M_3$$

subject to

$$M_1 + M_2 + M_3 \leq 60$$

$$2M_1 + M_2 + 3M_3 \leq 100$$

$$0 \leq M_1 \leq 40 \quad 0 \leq M_2 \leq 30$$

$$0 \leq M_3 \leq 20$$

④ Item 1 : weight 1 value 2

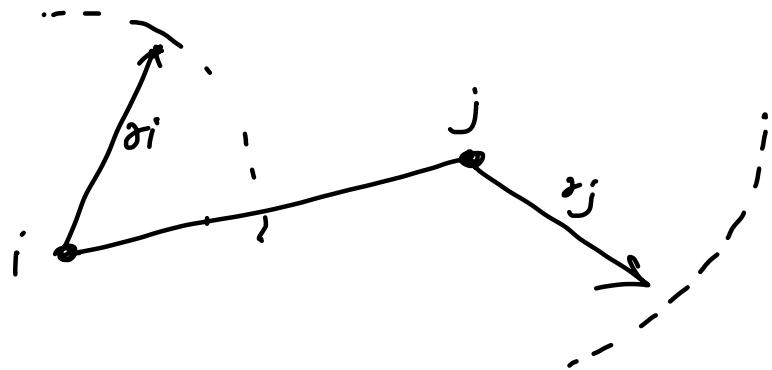
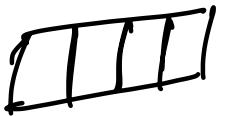
Item 2 : weight w weight w

If we pick item 1 acc to v/w ratio
we cannot pick item 2

But item 2 is optimal solution



A set of n space stations need your help in building a radar system to track spaceships traveling between them. The i^{th} space station is located in 3D space at coordinates (x_i, y_i, z_i) . The space stations never move. Each space station i will have a radar with power r_i , where r_i is to be determined. You want to figure how powerful to make each space station's radar transmitter, so that whenever any spaceship travels in a straight line from one space station to another, it will always be in radar range of either the first space station (its origin) or the second space station (its destination). A radar with power r is capable of tracking spaceships anywhere in the sphere with radius r centered at itself. Thus, a spaceship is within radar range through its trip from space station i to space station j if every point along the line from (x_i, y_i, z_i) to (x_j, y_j, z_j) falls within either the sphere of radius r_i centered at (x_i, y_i, z_i) or the sphere of radius r_j centered at (x_j, y_j, z_j) . The cost of each radar transmitter is proportional to its power, and you want to minimize the total cost of all of the radar transmitters. You are given all of the $(x_1, y_1, z_1), \dots, (x_n, y_n, z_n)$ values, and your job is to choose values for r_1, \dots, r_n . Express this problem as a linear program.

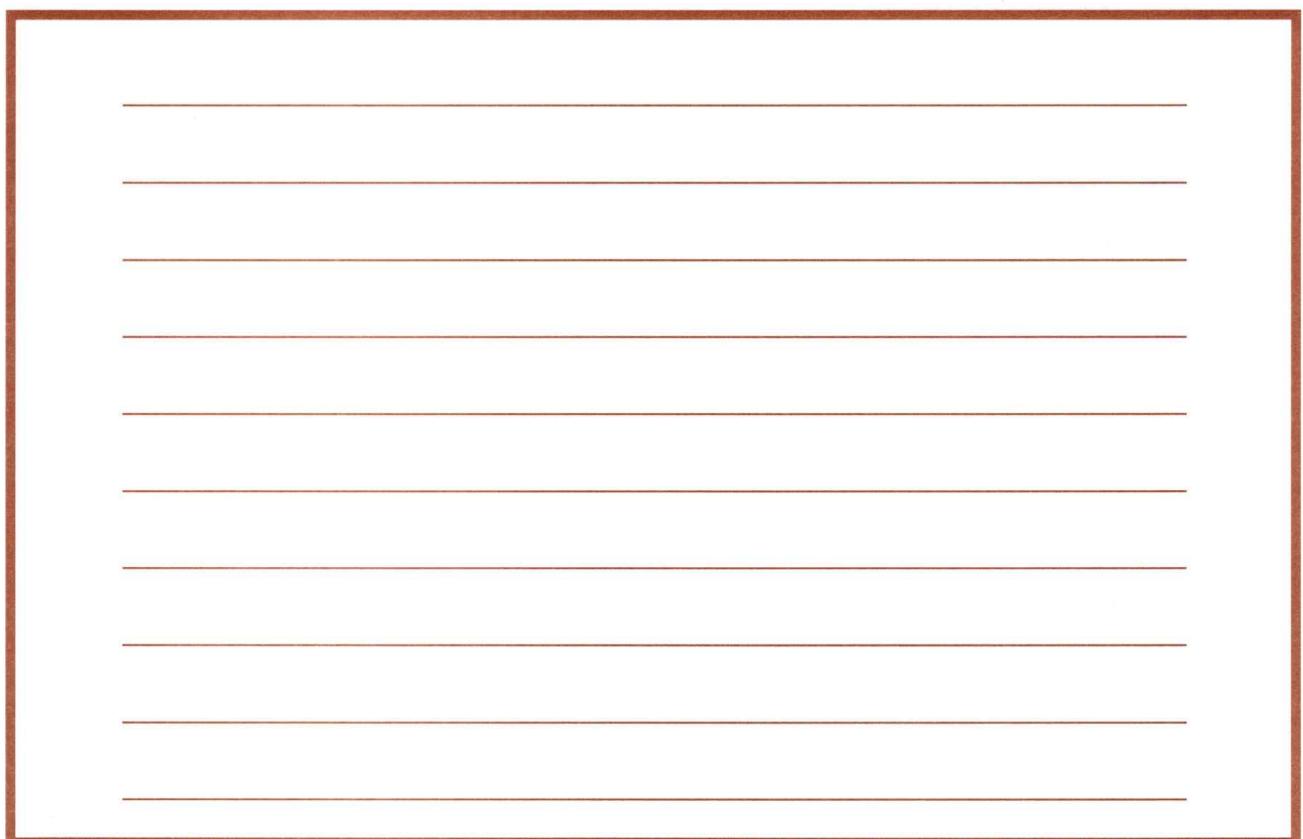


variables: r_i 's

objective : minimize $\sum r_i$'s

subject to

$$r_i + r_j \geq d_{ij} \quad \text{for all } i \text{ and } j$$



NOT IN PORTION

Randomized Algorithms

Two general types of randomization

- 1- Algorithm relies on random nature of input
- 2- Algorithm behaves randomly

Two types of Results

- 1- Some randomized algorithms always produce correct results. Randomization only affects the ran time.
- 2- Some randomized algorithms may produce incorrect solutions, but we try to bound the probability of such an incorrect solution.

Ex1: Problem Statement:

Find the medians in an unsorted array.

sort $\rightarrow O(n \log n)$
pick middle $\rightarrow O(1)$

Another Approach



look for median on right

If looking for the median, we then only need to look into either the right half or the left half. (Not necessarily equal halves)

Solutions :

Use the random nature of data / input
to find the medians:

Divide: Pick a pivot value (last term
is the array) and split the array in two.
Terms in left half are all less than
the pivot value
Terms in right half are all greater
than the pivot value.

Conquer :

if X is the median term we are
looking for return it
else if the median is in the left
half then explore the left half recursively
else explore the right half recursively
end if

A Simplistic Analysis of Runtime

Assume array size is only reduced by 10% each time.

$$T(n) = Cn + \left(\frac{9}{10}\right)Cn + \left(\frac{9}{10}\right)^2Cn + \dots$$

$$= O(n) \leftarrow \text{Expected Runtime}$$

Ex.2: Global Min-Cut Problem

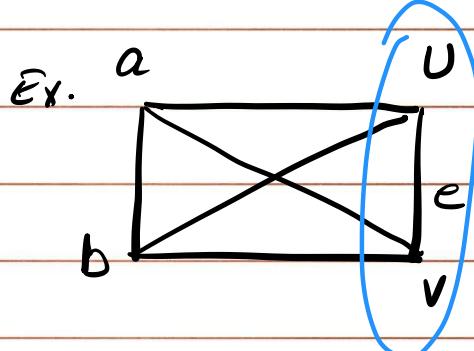
Given a connected unweighted multigraph, partitions vertices into two disjoint sets with the minimum no. of edges connecting the two sets.

Deterministic Solutions

- Pick any vertex as a source
- For each other vertex in V
Run Max Flow and find mincut.
- Return smallest min cut.

Randomize Solutions Using Edge Contractions

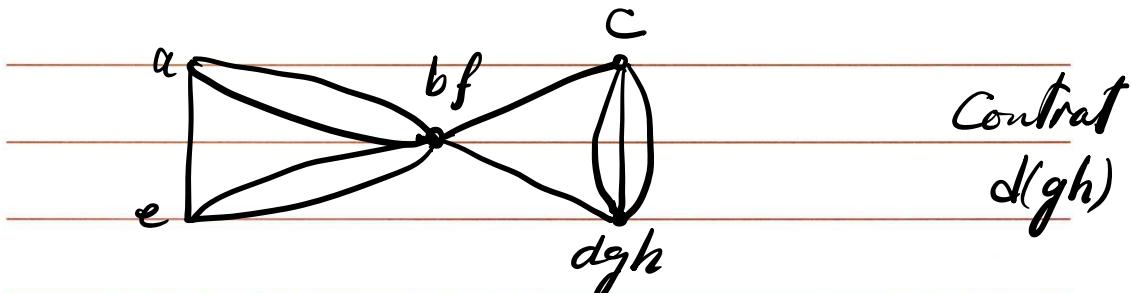
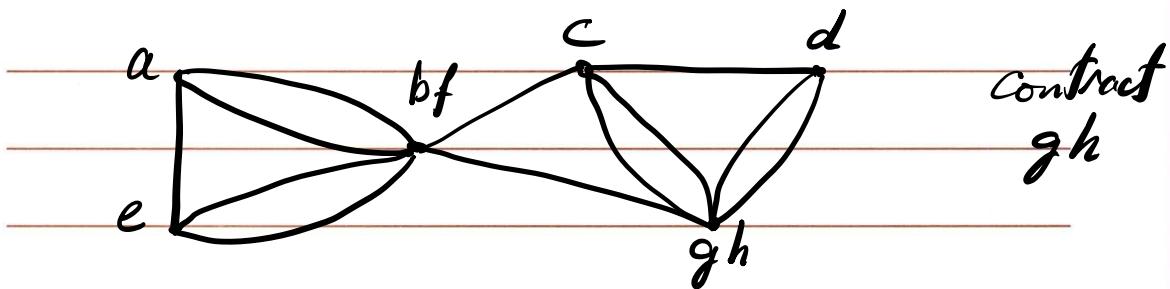
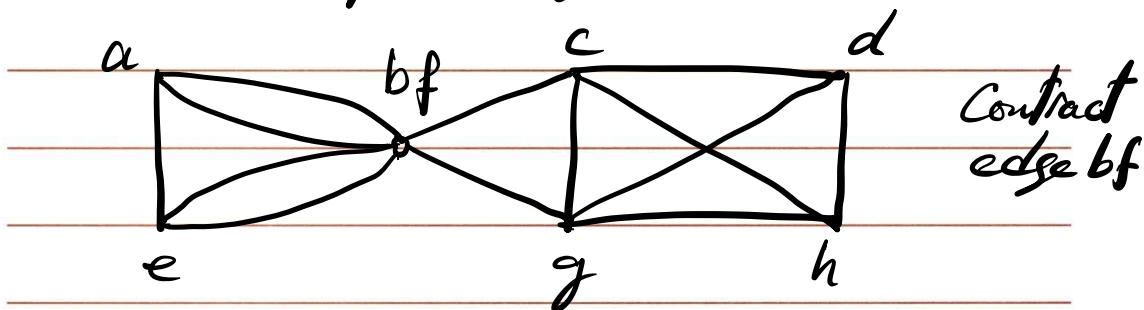
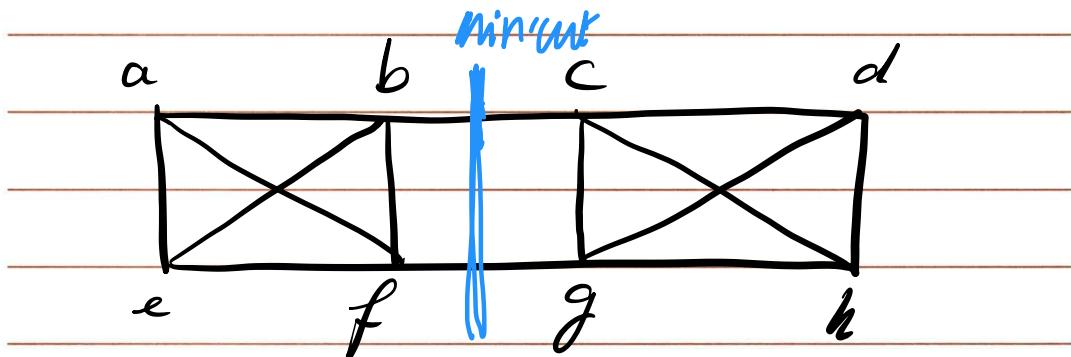
Edge Contraction: Contraction of an edge e with end points U & V is the replacement of U and V with a single vertex such that edges incident to the new vertex are the edges (other than e) that were incident on U or V .

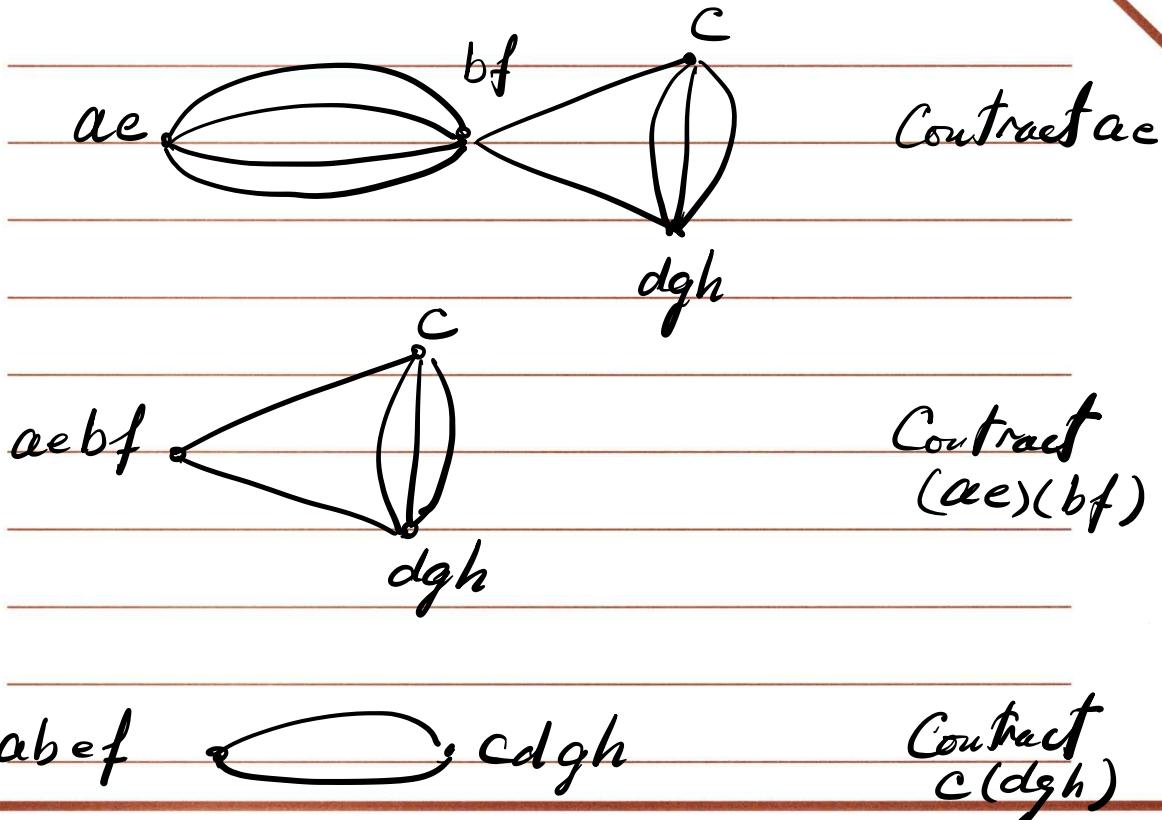


aFTER CONTRACTION
OF edge uv



Find Global Min-Cut in the graph below:





Randomized Algorithms

while there are more than ≥ 2 nodes left
 pick an edge at random
 and contract the edge
 (Remove any self loops)
 end while

For the two remaining nodes U_1 & U_2
 set $V_1 = \{ \text{nodes that went into } U_1 \}$
 $V_2 = \{ \dots \sim \dots \sim U_2 \}$

Analysis

Assume that a graph has a single min cut C . If we never contract edges from that min cut, then the algorithm will end up with the correct min cut.

How likely is this?

Say size of min cut is k .

So each ~~edge~~^{node} must have at least degree k .

\Rightarrow Total no. of edges in $G \geq k^{n/2}$

So, the probability that the 1st contracted edge is not on our mincut is at least $1 - \frac{k}{k^{n/2}}$

The probability that all contracted edges are Not on our min cut is at least

$$\left(1 - \frac{2}{n}\right) \cdot \left(1 - \frac{2}{n-1}\right) \cdot \left(1 - \frac{2}{n-2}\right) \cdots \left(1 - \frac{2}{3}\right)$$
$$= \Omega\left(\frac{1}{n^2}\right)$$

The probability of the algorithm failing N times is

$$\left(1 - \frac{1}{n^2}\right)^N$$

If we choose $N = 100n^2$ we can make it close to zero! Since $\left(1 - \frac{1}{n^2}\right)^{100n^2} \leq e^{-100}$

$$\left(\lim_{x \rightarrow \infty} \left(1 - \frac{1}{x}\right)^{ax} = e^{-ax}\right)$$
