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Projet Formel - Calculus of Constructions V4.10 - Vernacular V2.3
                   Reynolds paradox, with the Type: Type axiom
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(* use the Type:Type system, #use"Reyn_tac"; V"Log_Rel" *)
(* this file shows an inconsistency is the logical system U-, defined in Girard's
thesis, and in which the question of the consistency of U- was raised.
The system U- may be described with the axioms
           (star, star), (k, star), (k, k), (k', k)
    in Barendregt-Berardi's GTS. The reader can check that we use only these axioms in the derivation, and not fully the Type:Type axiom.

The lambda-term we get does not loop to itself by reduction *)
(* Reynolds operator *)
Definition PHI.
Body [A:Type] (A->Prop) ->Prop.
(* we extend this map functorialy *)
Definition phi: (A,B:Type)(A->B) -> (PHI A) -> (PHI B) = [A,B:Type][f:A->B][z:(PHI A)][u:B->Prop](z [x:A](u (f x))).
(* preinitial PHI-algebra. We need the axiom (Type, Type) *)
Definition A0.
Body (A: Type) ((PHI A) -> A) -> A.
Definition iter A0.
Body [X:Type] [f:(PHI X)->X] [u:A0] (u X f).
Definition intro : (PHI A0) -> A0 =
[z:(PHI A0)][A:Type][f:(PHI A)->A](f (phi A0 A (iter_A0 A f) z)).
(* extension of PHI to relations. We can thus consider PHI as a functor on sets, that are types with a relations *)
Definition phi2: (A:Type)(Rel A)->(Rel (PHI A)) = [A:Type][R:(Rel A)](power (A->Prop) (power A R)).
(* partial equivalence relation defined on AO, so that the set AO,EO
    is an initial PHI-algebra in the category of sets *)
Definition teta : A0 -> (PHI A0) =
      (iter_A0 (PHI A0) (phi (PHI A0) A0 intro)).
Definition E0 : (Rel A0) = [x1,x2:A0] (E: (Rel A0))
      (per AO E) -> ((z1,z2:(PHI AO))(phi2 AO E z1 z2)->(E (intro z1) (intro z2))) -> ((x1,x2:AO)(E x1 x2)->(E x1 (intro (teta x2)))) ->
Definition F0 : (Rel (A0->Prop)) = (power A0 E0).
Definition G0 : (Rel (PHI A0)) = (power (A0->Prop) F0).
(* the goal of what follows is to show that the set AO,EO is in one-to-one
correspondance with the set (PHI AO), (phi2 AO EO), via intro,teta.
From this, we deduce a contradiction via Cantor-Russell's argument *)
Goal (sym A0 E0).
By reds.
Do res_simpl H0.
Do res_simpl H3.
Do res_simpl H.
Save sym_E0.
Goal (trans A0 E0).
By reds.
Do res_simpl H1.
Do res_exact H5 y.
Do res_simpl H.
Do res_simpl HO.
Save trans_E0.
Goal (per A0 E0).
Resolve per_intro.
Exact sym_E0.
Exact trans_E0.
Save per_E0.
(* intro is a map from (PHI AO), (phi2 AO EO) to AO,EO *)
Goal (z1, z2: (PHI A0)) (GO z1 z2) -> (EO (intro z1) (intro z2)).
By reds.
Do res_simpl H1.
Do res_simpl H.
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Do res_simpl H3.
Do res_simpl H4.
    Save lemmal.
   Goal (per (A0->Prop) F0). Unfold F0.
    Unfold power.
  Do res_simpl per_E0.
Do res_simpl per_equiv.
Resolve per_intro.
By (hyps THEN hyps).
Resolve H1.
Resolve H3
   Resolve H3.
Do res_simpl H.
   By (hyps THEN hyps).
   Do res_exact H2 (y y0).
Do res_simpl H3.
Resolve H4.
   Do res_exact H0 x0.
Do res_simpl H.
   Save per_F0.
   Goal (per (PHI A0) G0).
   Unfold GO.
   Unfold power.
  Do res_simpl per_F0.
Do res_simpl per_equiv.
Resolve per intro.
By (hyps THEN hyps).
Resolve H1.
Resolve H3.
  Do res simpl H.
  By (hyps THEN hyps).
Do res_exact H2 (y y0).
Do res_simpl H3.
Resolve H4.
  Do res_exact HO xO.
  Do res_simpl H.
  Save per_GO.
  Goal (x1,x2:A0) (E0 x1 x2)->(E0 x1 (intro (teta x2))).

By (hyps THEN hyps).

Do res_simpl H2.

Do res_simpl H.

Save id_intro_teta.
  Goal (z1, z2: (PHI A0)) (G0 z1 z2) \rightarrow (G0 z1 (teta (intro z2))).
 Goal (Z1, Z2: (PHI AU)) (GU Z1 Z2) -> (GU Z1 (teta (intro Z. By (hyps THEN hyps).

Change (equiv (z1 x) (z2 [x:A0] (y (intro (teta x))))).

Do res_simpl H.

Do res_intro HO.

Do res_simpl id_intro_teta.

Save id_teta_intro.
  Goal (x1, x2:A0) (E0 x1 x2) \rightarrow (G0 (teta x1) (teta x2)).
  By hyps.
Change ([u, v: A0] (G0 (teta u) (teta v)) x1 x2).
 By hyps.
Do res_simpl id_teta_intro.
 Save lemma_teta.
 Definition psi : (A0->Prop)->A0 =
[u:A0->Prop](intro (F0 u)).
 Definition inter : (PHI A0) -> A0-> Prop =
 [C: (PHI A0)][x:A0] (P:A0->Prop) (FO P P)-> (C P)-> (P x).
 Goal (z1, z2: (PHI A0)) (GO z1 z2) \rightarrow (FO (inter z1) (inter z2)). By (hyps THEN hyps). Do res_intro equiv_intro.
 By hyps.
 Do cut (equiv (P x) (P y)).
Do res_intro H4.
 Resolve H5.
 Do res_simpl H1.
Do cut (equiv (z1 P) (z2 P)).
Do res_intro H7.
 Exact (H9 H3).
 Do res_simpl H.
Do res_simpl H2.
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By hyps.
Do cut (equiv (P x) (P y)).
Do res_intro H4.
 Resolve H6.
Do res_simpl H1.
Do cut (equiv (z1 P) (z2 P)).
Do res_intro H7.
Exact (H8 H3).
Do res_simpl H.
Do res_simpl H2.
 Save lemma inter.
 (* we follow Cantor-Russell's paradox *)
Definition khi : A0 -> (A0->Prop) = [x:A0] (inter (teta x)).
 Section paradox.
 Variable p:Prop.
 Definition u0: A0 \rightarrow Prop = [x:A0](khi x x) \rightarrow p.
Definition x0: A0 = (psi u0).
 Goal (E0 x0 x0).
 Unfold x0.
 Unfold x0.
Unfold psi.
Unfold psi.
Do res simpl lemmal.
Do res_simpl per_F0.
Do res_simpl equiv_intro.
Do res_exact H1 x.
Do res_exact H0.
 Save lemma4.
 Goal (F0 u0 u0).
 By hyps.
Do cut (F0 (khi x) (khi y)).
Do cut (equiv (khi x x) (khi y y)).
Do res_simpl H1.
Unfold u0.
 Unfold u0.
Do res_simpl equiv_intro.

Exact (H4 (H3 H5)).

Exact (H4 (H2 H5)).

Do res_simpl H0.

Unfold khi.

Unfold khi.
Resolve lemma_inter.
Do res_simpl lemma_teta.
Save lemma3.
Goal (F0 u0 (khi x0)).
Unfold x0.
Unfold khi.
Unfold psi.
Do apply per F0.
Do res_simpl H.
Do res_exact H1 (inter (F0 u0)).
By hyps.
Do res_simpl equiv_intro.
By hyps.
Do cut (equiv (u0 x) (P y)).
Do res_simpl H6.
Exact (H7 H3).
Do res_simpl H5.
Do apply lemma3.
Do cut (equiv (u0 x) (u0 y)).
Do res_simpl H5.
Resolve H7.
Do res_simpl H3.
Do res_simpl H4.
Do res_simpl lemma_inter.
Do cut (G0 (F0 u0) (teta (intro (F0 u0)))).
Do res_simpl H3.
Do res_simpl id_teta_intro.
Do res_simpl equiv_intro.
Do res_exact H1 x0.
Do res_exact H1 y0.
Do res_simpl H0.
Save lemma5.
Goal (equiv (u0 x0) (khi x0 x0)).
Do apply lemma5. Resolve H.
Exact lemma4.
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Save lemma6.

Goal (u0 x0).
By hyps.
Do res simpl lemma6.
Exact (H1 H H).
Save lemma7.

Goal p.
Do res simpl lemma6.
Exact (lemma7 (H lemma7)).
Save Reynolds.

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