## Data and Computer Communications

Tenth Edition by William Stallings

### **CHAPTER 11**

**Local Area Network Overview** 

"There is some evidence that computer networks will have a large impact on society. Likely areas are the economy, resources, small computers, human-to-human interaction, and computer research."



—What Can Be Automated?
The Computer Science and
Engineering Research Study,
MIT Press, 1980

## **Bus Topology**

- Topology
  - Refers to the way in which the endpoints, or stations, attached to the network are interconnected
- Bus topology
  - All stations attach, through a tap, directly to a linear transmission medium, or bus
  - Full-duplex operation between the station and the tap allows data to be transmitted onto the bus and received from the bus
  - A transmission from any station propagates the length of the medium in both directions and can be received by all other stations
  - At each end of the bus is a terminator, which absorbs any signal, removing it from the bus

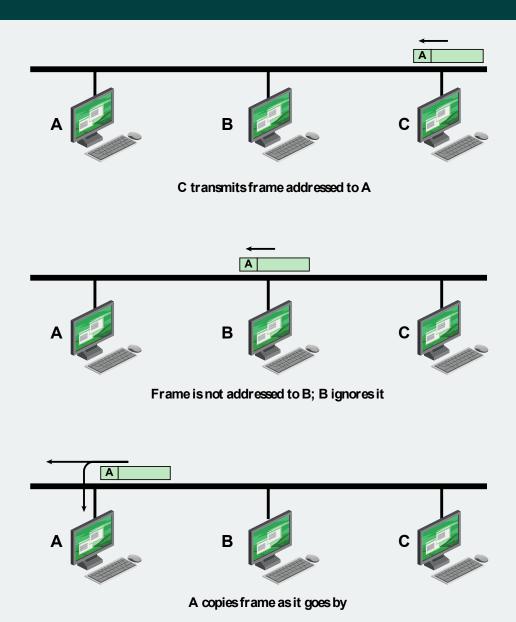


Figure 11.1 Frame Transmission on a Bus LAN

## **Star Topology**

- Each station connects to common central node
  - Usually via two point-to-point links
    - One for transmission and one for reception

#### Central node

- Operate in broadcast fashion
- Physical star, logical bus
- Only one station can transmit at a time (hub)
- Can act as frame switch

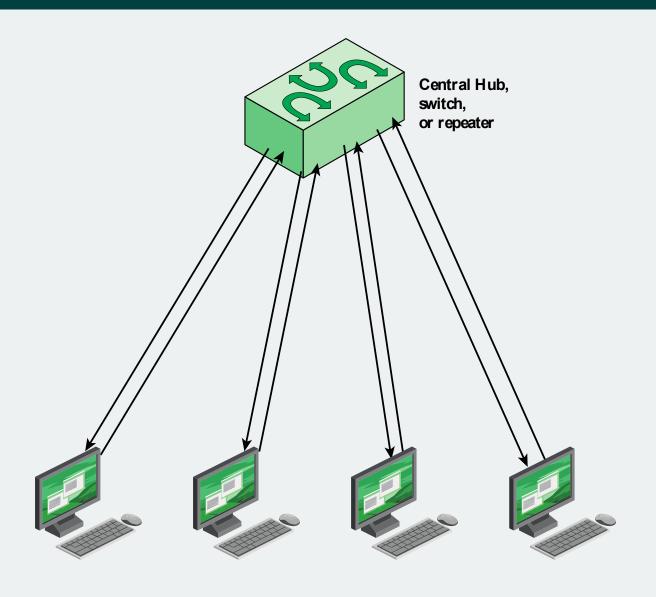


Figure 11.2 Star Topology

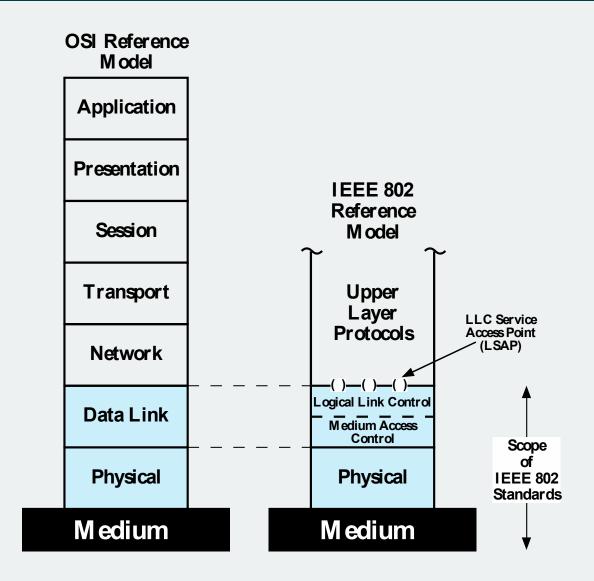
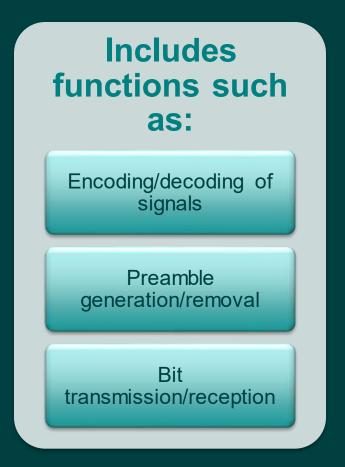


Figure 11.3 IEEE 802 Protocol Layers Compared to OSI Model

### **IEEE 802 Reference Model**

- Lowest layer corresponds to the physical layer of the OSI model
- Includes a specification of the transmission medium and the topology



## **IEEE 802 Layers**

- Logical Link Control Layer (LLC)
  - Provide interface to higher levels
  - Perform flow and error control



#### Media Access Control (MAC)

- On transmit assemble data into frame
- On reception disassemble frame, perform address recognition and error detection
- Govern access to LAN transmission medium

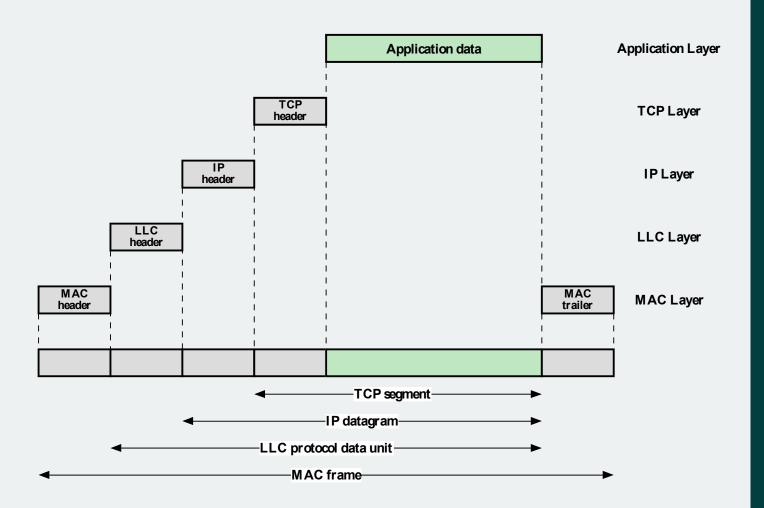


Figure 11.4 LAN Protocols in Context

## **Logical Link Control**

- Transmission of link level PDUs between stations
- Must support multi-access, shared medium
- Relieved of some details of link access by the MAC layer
- Addressing involves specifying source and destination LLC users
  - Referred to as service access points (SAPs)

#### **LLC Services**

#### Unacknowledged connectionless service

- Data-gram style service
- Delivery of data is not guaranteed

#### Connection-mode service

- Logical connection is set up between two users
- Flow and error control are provided

#### Acknowledged connectionless service

 Datagrams are to be acknowledged, but no logical connection is set up

### **LLC Service Alternatives**

#### Unacknowledged connectionless service

- Requires minimum logic
- Avoids duplication of mechanisms
- Preferred option in most cases

#### Connection-mode service

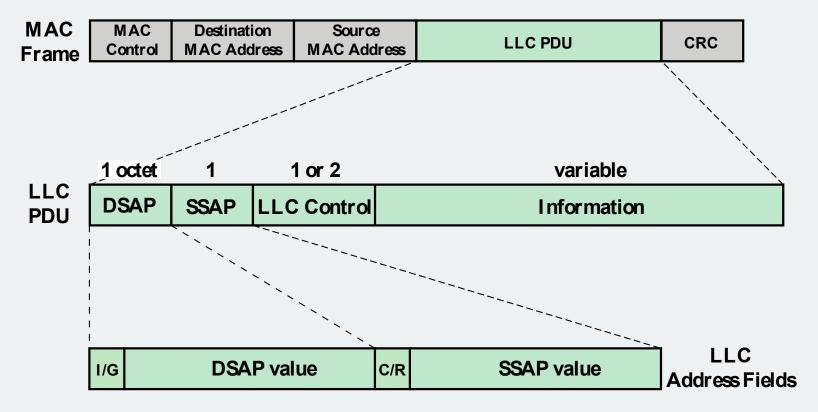
- Used in simple devices
- Provides flow control and reliability mechanisms

#### Acknowledged connectionless service

- Large communication channel needed
- Time critical or emergency control signals

#### **LLC Protocol**

- Modeled after HDLC
- Asynchronous balanced mode
  - Connection mode (type 2) LLC service
- Unacknowledged connectionless service
  - Using unnumbered information PDUs (type 1)
- Acknowledged connectionless service
  - Using 2 new unnumbered PDUs (type 3)
- Permits multiplexing using LSAPs



I/G = Individual/Group C/R = Command/Response

Figure 11.5 LLC PDU in a Generic MAC Frame Format

## Medium Access Control (MAC) Protocol

- Controls access to the transmission medium
- Key parameters:
  - Where
    - Greater control, single point of failure
    - More complex, but more redundant

#### How

- Synchronous
  - Capacity dedicated to connection, not optimal
- Asynchronous
  - Response to demand
  - Round robin, reservation, contention

## **Asynchronous Systems**

#### Round robin

 Each station given turn to transmit data

#### Reservation

- Divide medium into slots
- Good for stream traffic

#### Contention

- All stations contend for time
- Good for bursty traffic
- Simple to implement
- Tends to collapse under heavy load

## **MAC Frame Handling**

- MAC layer receives data from LLC layer
- PDU is referred to as a MAC frame
- MAC layer detects errors and discards frames
- LLC optionally retransmits unsuccessful frames



## **Bridges**

- Connects similar LANs with identical physical and link layer protocols
- Minimal processing
- Can map between MAC formats
- Reasons for use:
  - Reliability
  - Performance
  - Security
  - Geography



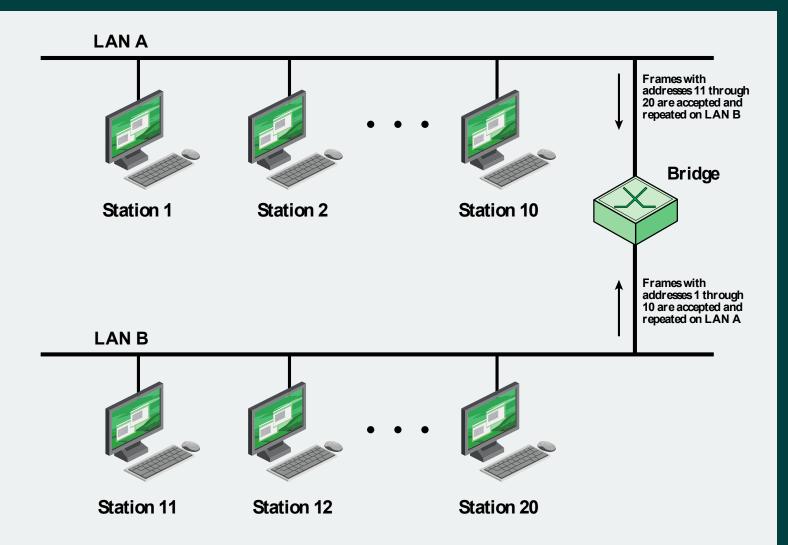


Figure 11.6 Bridge Operation

## **Bridge Design Aspects**

- Makes no modification to the content or format of the frames it receives
- Should contain enough buffer space to meet peak demands
- Must contain routing and addressing intelligence
- May connect more than two LANs
- Bridging is transparent to stations

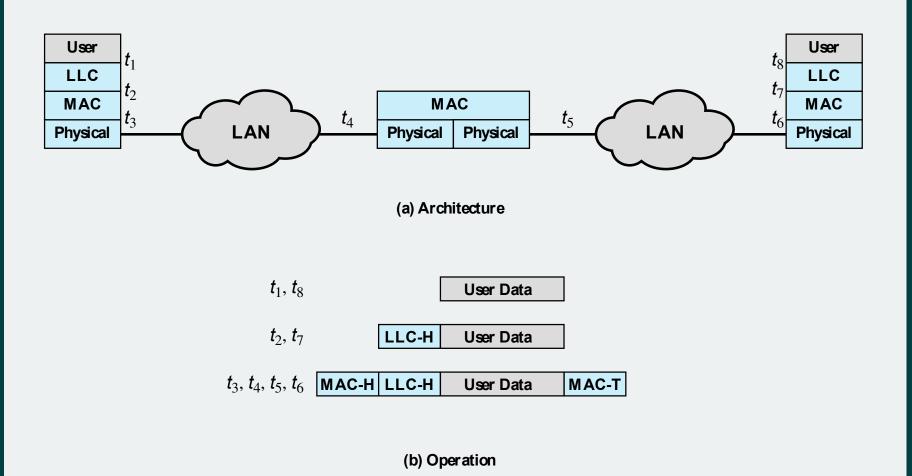


Figure 11.7 Connection of Two LANs by a Bridge

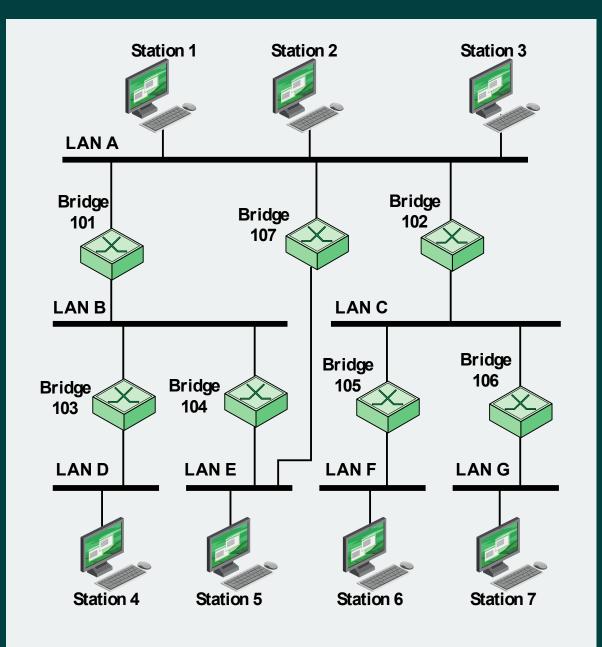


Figure 11.8 Configuration of Bridges and LANs, with Alternate Routes

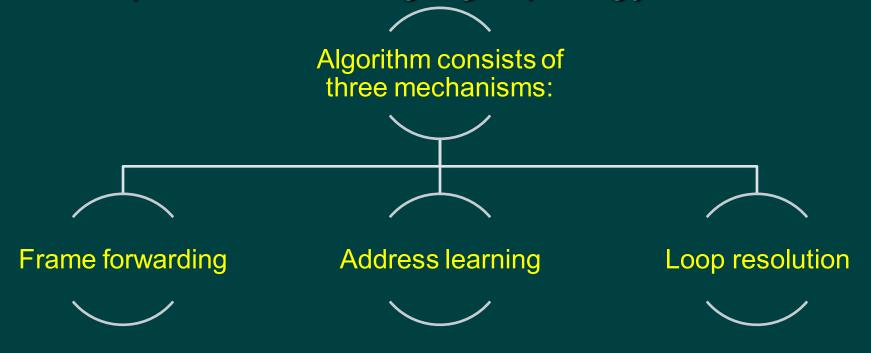
## **Fixed Routing**

- Simplest and most common strategy
- Suitable for small internets and internets that are relatively stable
- Afixed route is selected for each pair of LANs
  - Usually least hop route
- Only change when topology changes
- Widely used but limited flexibility



## **Spanning Tree**

- Bridge automatically develops routing table
- Automatically updates routing table in response to changing topology



## **Frame Forwarding**

- Maintain forwarding database for each port attached to a LAN
- For a frame arriving on port X:

Search forwarding database to see if MAC address is listed for any port except port X



If destination MAC address is not found, forward frame out all ports except the one from which it was received



If the destination address is in the forwarding database for some port y, check port y for blocking or forwarding state



If port y is not blocked, transmit frame through port y onto the LAN to which that port attaches

## **Address Learning**

- Can preload forwarding database
- When frame arrives at port X, it has come from the LAN attached to port X
- Use source address to update forwarding database for port X to include that address
- Have a timer on each entry in database
- If timer expires, entry is removed.
- Each time frame arrives, source address checked against forwarding database
  - If present timer is reset and direction recorded
  - If not present entry is created and timer set

## **Spanning Tree Algorithm**

- Address learning works for tree layout if there are no alternate routes in the network
  - Alternate route means there is a closed loop
- For any connected graph there is a spanning tree maintaining connectivity with no closed loops
- Algorithm must be dynamic

#### IEEE 802.1 Spanning Tree Algorithm:

- Each bridge assigned unique identifier
- Cost assigned to each bridge port
- Exchange information between bridges to find spanning tree
- Automatically updated whenever topology changes

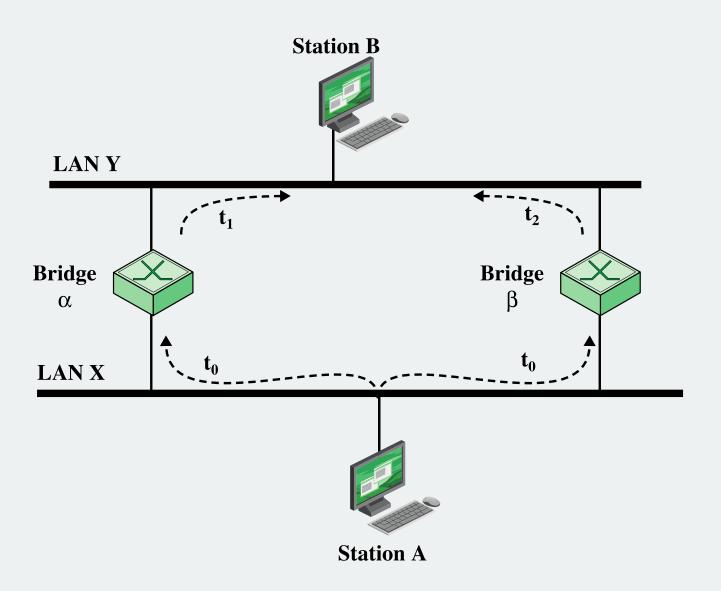


Figure 11.9 Loop of Bridges

#### Hubs

- Activecentral element of star layout
- Each station connected to hub by two lines
- Hub acts as a repeater
- Length of a line is limited to about 100m
- Opticalfiber may be used to about 500m
- Physically a star, logically a bus
- Transmissionfrom any one station is received by all other stations
- Iftwo stations transmit at the same time there will be a collision

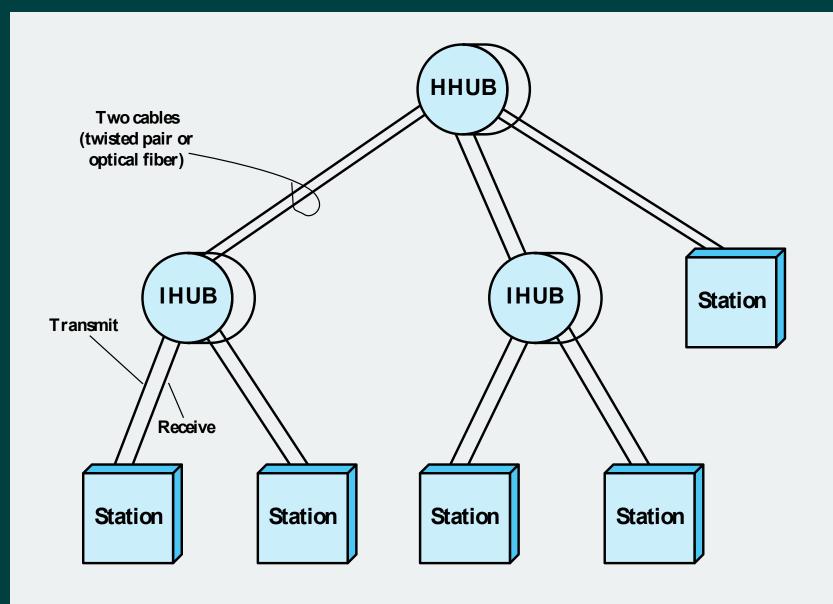


Figure 11.10 Two-Level Star Topology

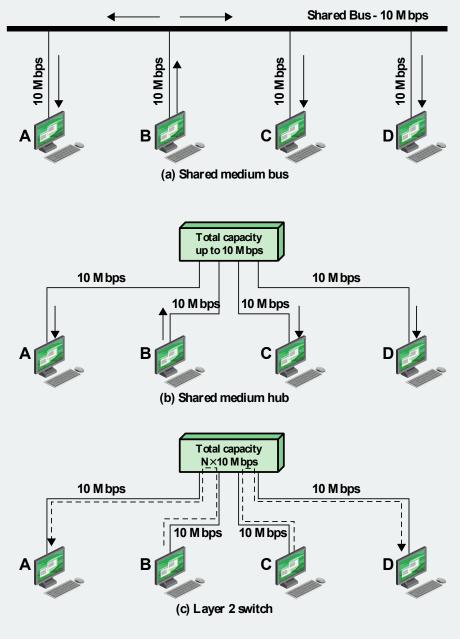


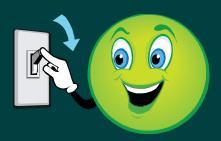
Figure 15.11 LAN Hubs and Switches

## **Layer 2 Switch Benefits**

- No change is required to the software or hardware of the attached devices to convert a bus LAN or a hub LAN to a switched LAN
- Have dedicated capacity equal to original LAN
  - Assuming switch has sufficient capacity to keep up with all devices
- Scales easily
  - Additional devices attached to switch by increasing capacity of layer 2

## Types of Layer 2 Switches

- Store-and-forward switch
  - Accepts frame on input line, buffers briefly, routes to appropriate output line
  - See delay between sender and receiver
  - Boosts overall integrity



- Cut-through switch
  - Use destination address at beginning of frame
  - Switch begins repeating frame onto output line as soon as destination address is recognized
  - Yields highest possible throughput
  - Risk of propagating bad frames

## Layer 2 Switch vs. Bridge

Differences between switches and bridges:

Bridge

Frame handling done in software

Analyzes and forwards one frame at a time

Uses store-andforward operation **Switch** 

Performs frame forwarding in hardware

Can handle multiple frames at a time

Can have cutthrough operation

- Layer2 switch can be viewed as full-duplex hub
- Incorporates logic to function as multiport bridge
  - New installations
    typically include layer
    2 switches with bridge
    functionality rather
    than bridges

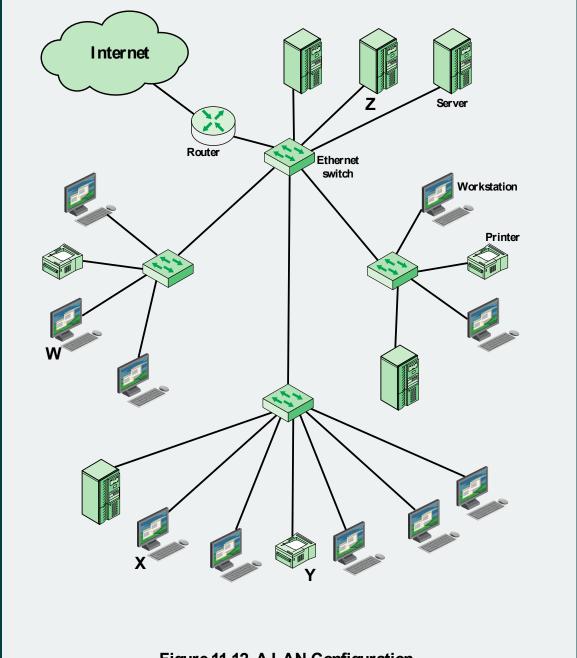


Figure 11.12 A LAN Configuration

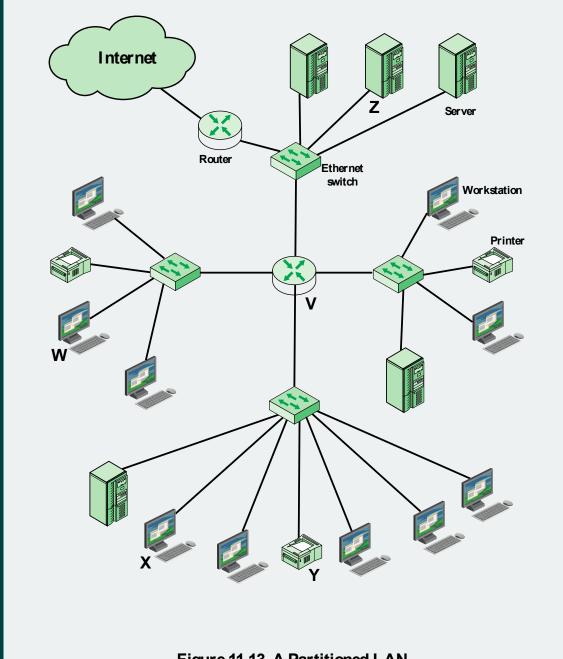


Figure 11.13 A Partitioned LAN

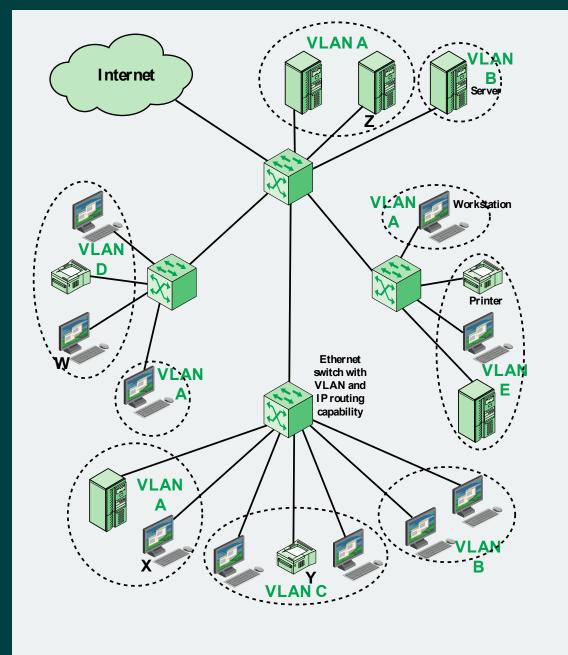


Figure 11.14 A VLAN Configuration

## **Defining VLANs**

- Broadcast domain consisting of a group of end stations not limited by physical location and communicate as if they were on a common LAN
- Membership by:
  - Port group
  - MAC address
  - Protocol information



## Communicating VLAN Membership

Switches need to know VLAN membership

- Configure information manually
- Network management signaling protocol
- Frame tagging (IEEE802.1Q)

# Summary

- Bus and tree topologies and transmission media
  - Topologies
  - Choice of topology
  - Choice of transmission medium
- LAN protocol architecture
  - IEEE 802 reference model
  - Logical link control
  - Medium access control
- Hubs and switches
  - Hubs
  - Layer 2 switches

- Bridges
  - Functions of a bridge
  - Bridge protocol architecture
  - Fixed routing
  - The spanning tree approach
- Virtual LANs
  - The use of virtual LANs
  - Defining VLANs
  - Communicating VLAN membership