

OS Group Project Docs

ARGUMENT PASSING

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Introduction

The group project of Pintos involves implementing argument passing. At the end of worksheet 3, Pintos can execute based on file name. But the passed arguments are discarded. It is considered argument passing when the arguments extracted, are added to the stack provided by the esp stack pointer. Previously, in worksheet 3, the program name is extracted in the process_execute() function. The function calls are as state below.

process_execute() -> start_process() -> load() -> setup_stack()

The file reading using the program name, is actually handled by the load(). Therefore, I have moved the program name extraction logic into the load function along with argument extraction logic. The setup stack is called; however, we first need to pass the extracted arguments and program name into the function to add them to the stack. The code changed I have brought are described below.

Code Changes

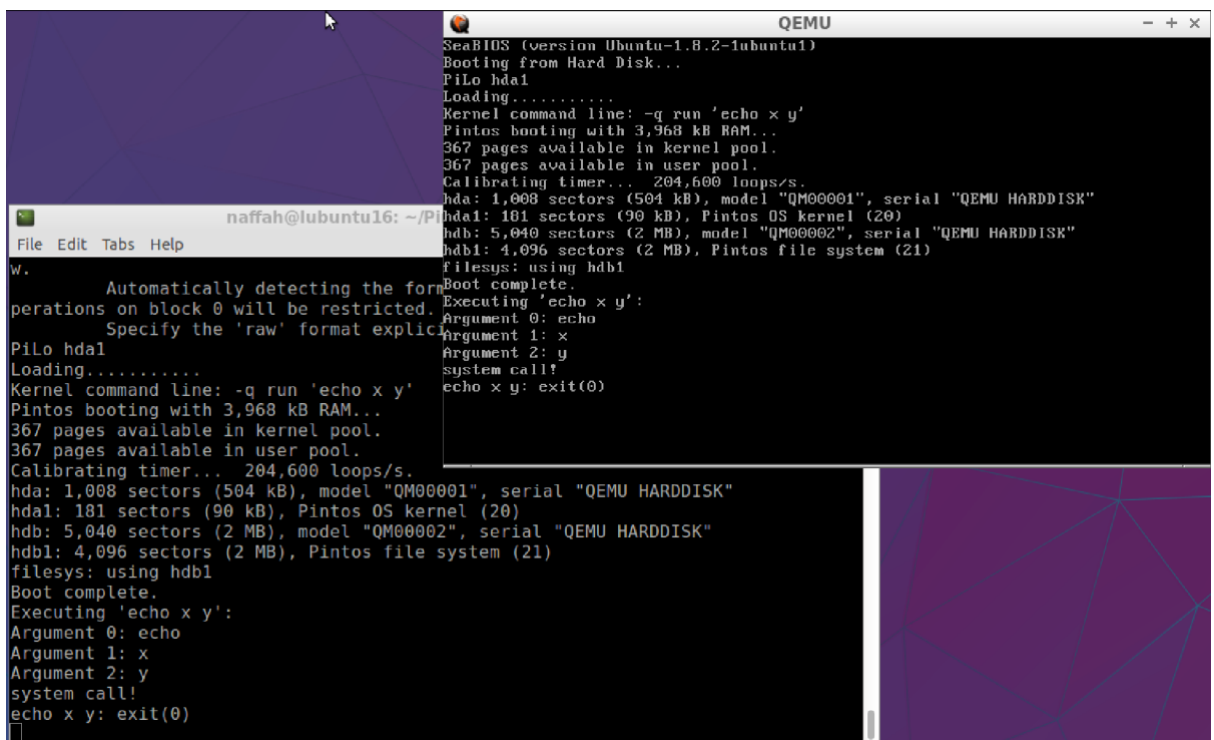
1. The program name is first extracted and added to argv[0]. The remaining pointers in the array is used to store strings of arguments. The argc is simply used to track the number of arguments

```
225  /*
226  |   Extract arguments from file
227  */
228  char *argv[100]; // Maximum number of arguments
229  int argc = 0;
230  char *save_ptr;
231  argv[0] = strtok_r(file_name, " ", &save_ptr);
232  char *token;
233
234  while ((token = strtok_r(NULL, " ", &save_ptr)) != NULL) {
235      argv[++argc] = token;
236  }
237
238  // for (int i = 0; i <= argc; ++i) {
239  //     printf("Argument %d: %s\n", i, argv[i]);
240  // }
241
```

2. All instances where the file_name was used is replaced with argv[0] containing the extracted program name

```
248      /* Open executable file. */
249      file = filesys_open (argv[0]);
250
251      if (file == NULL)
252      {
253      printf ("load: %s: open failed\n", argv[0]);
254      goto done;
255      }
256
257      /* Read and verify executable header. */
258      if (file_read (file, &ehdr, sizeof ehdr) != sizeof ehdr
259          || memcmp (ehdr.e_ident, "\177ELF\1\1\1", 7)
260          || ehdr.e_type != 2
261          || ehdr.e_machine != 3
262          || ehdr.e_version != 1
263          || ehdr.e_phentsize != sizeof (struct Elf32_Phdr)
264          || ehdr.e_phnum > 1024)
265      {
266      printf ("load: %s: error loading executable\n", argv[0]);
267      goto done;
268      }
```

3. To make sure it works, I will simply print the arguments.



```
QEMU
SeaBIOS (version Ubuntu-1.8.2-1ubuntu1)
Booting from Hard Disk...
Pilo hda1
Loading.....
Kernel command line: -q run 'echo x y'
Pintos booting with 3,968 kB RAM...
367 pages available in kernel pool.
367 pages available in user pool.
Calibrating timer... 204,600 loops/s.
hda: 1,008 sectors (504 kB), model "QM000001", serial "QEMU HARDDISK"
hda1: 181 sectors (90 kB), Pintos OS kernel (20)
hdb: 5,040 sectors (2 MB), model "QM000002", serial "QEMU HARDDISK"
hdb1: 4,096 sectors (2 MB), Pintos file system (21)
filesys: using hdb1
Boot complete.
Executing 'echo x y':
Argument 0: echo
Argument 1: x
Argument 2: y
system call!
echo x y: exit(0)
```

4. Next step is to modify `setup_stack()` function to allow passing of the extracted arguments

```
452  /* Create a minimal stack by mapping a zeroed page at the top of
453  |   user virtual memory. */
454  static bool
455  > setup_stack (void **esp, char **argv, int argc) ...
```

5. The `setup_stack()` is then modified to allow passing arguments to the stack:
 - a. The first loop: This loop iterates through the arguments in reverse order (from the last argument to the first). For each argument, it calculates the length of the argument string (including the null terminator), subtracts the space required from `esp` (the stack pointer), stores the resulting `esp` value in `arr[i]`, and then copies the argument string to the calculated address on the stack.
 - b. The second loop: This loop iterates through the `arr` array in reverse order (from the last argument to the first). For each argument, it subtracts 4 bytes (the size of a 32-bit pointer) from `esp` to create space for a pointer to the argument's address. Then, it assigns the value of `arr[i]` (which is the address of the argument on the stack) to the location pointed to by `esp`.

```
469      // Loop through arguments in reverse order to set up the stack
470      for (int i = argc - 1; i ≥ 0; --i) {
471          // Calculate the space needed for the argument string (including null terminator)
472          size_t arg_size = (strlen(argv[i]) + 1) * sizeof(char);
473
474          // Adjust the stack pointer to make space for the argument string
475          *esp = *esp - arg_size;
476
477          // Store the current stack pointer in the array for future reference
478          arr[i] = (uint32_t *)*esp;
479
480          // Copy the argument string to the stack at the current stack pointer
481          memcpy(*esp, argv[i], arg_size);
482      }
483
484      // Loop through the argument pointers on the stack to set up the stack further
485      for (int i = argc - 1; i ≥ 0; --i) {
486          // Reserve space for a 32-bit pointer on the stack
487          *esp = *esp - sizeof(uint32_t);
488
489          // Store the address of the argument in the array at this location
490          (*(uint32_t **)(*esp)) = arr[i];
491      }
```