## **Category Theory**

### Summer of Science 2025 Project

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### **Contents**

I.	. Category Theory in Context		
1.	Categories, Functors, Natural Transformations	7	
	1.1. Abstract and concrete categories	7	
	1.2. Duality	7	
11.	Tom Leinster	9	

# Part I. Category Theory in Context

# 1. Categories, Functors, Natural Transformations

#### 1.1. Abstract and concrete categories

#### 1.1.i

*Proof.* Observe that  $gfh: y \to x$ . Also gfh = (gf)h = h and gfh = g(fh) = g, thus g = h. This means that f has an inverse morphism g such that  $fg = 1_x$  and  $gf = 1_y$ , so f must be an isomorphism.

#### 1.1.ii

Proof. First we show that the collection of isomorphisms in  $\mathcal{C}$  with the objects of  $\mathcal{C}$  forms a category. The identity morphism exists in the new category since it is its own inverse and therefore an isomorphism. Associativity is inherited from morphisms on  $\mathcal{C}$ . To prove composability, say  $f: x \to y$  and  $g: y \to z$  are isomorphisms, i.e., there exist  $f_*$  and  $g_*$  such that  $ff_* = 1_y$ ,  $f_*f = 1_x$  and  $gg_* = 1_z$ ,  $g_*g = 1_y$ . Clearly gf is a morphism since so are f and g. We claim that  $f_*g_*$  is its inverse morphism, which is trivial to verify. Hence gf is an isomorphism. It is clear that the resultant category is a groupoid so let us prove its maximality. Suppose to the contrary that there is a larger groupoid  $\mathcal{C}'$  containing it. Since they have the same objects,  $\mathcal{C}'$  must contain an isomorphism not in  $\mathcal{C}$ . This is a contradiction because  $\mathcal{C}$  contains all the isomorphisms between its objects.

#### 1.1.iii

*Proof.* Associativity follows in either case because the maps in  $\mathcal{C}$  are associative.

- (i) First note that there is an identity morphism for every  $f: c \to x$  because we can take g = f. To show composability, let  $h_1: a \to b$  be from  $f: c \to a$  to  $g: c \to b$  and  $h_2: b \to d$  be from  $g: c \to b$  to  $h: c \to d$ . Note that  $g = h_1 f$  and  $h = h_2 g$  implies  $h = h_2(h_1 f) = (h_2 h_1) f$  and so the diagram with  $h_2 h_1$  also commutes.
- (ii) Identity morphism for f is found by taking g = f. To show composability, proceed as before  $f = gh_1 = (hh_2)h_1 = h(h_2h_1)$ .

#### 1.2. Duality

# Part II. Tom Leinster