Chapter 10

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Problem 10.4.

Show a concrete attack:

For adversary \mathcal{A} , he is given trans = (s, u, w), then he can compute $t = u \oplus s$. In key-exchange experiment, if $\hat{k} = w \oplus u \oplus s$, then \mathcal{A} outputs b' = 0; otherwise, \mathcal{A} outputs b' = 1.

So if b = 0, then \mathcal{A} always has b' = b; and if b = 1, \mathcal{A} guesses right with probability $1 - 2^{-n}$.

$$\begin{split} \Pr[\mathrm{KE}^{\mathrm{eav}}_{\mathcal{A},\Pi}(n) = 1] &= \frac{1}{2} \Pr[\mathrm{KE}^{\mathrm{eav}}_{\mathcal{A},\Pi}(n) = 1 \mid b = 0] + \frac{1}{2} \Pr[\mathrm{KE}^{\mathrm{eav}}_{\mathcal{A},\Pi}(n) = 1 \mid b = 1] \\ &= \frac{1}{2} + \frac{1}{2} (1 - 2^{-n}) \\ &= 1 - \mathrm{negl}(n), \end{split}$$

which is significantly larger than $\frac{1}{2}$.