

## Chapter 10

Mingjia Huo

**Problem 10.4.** Show a concrete attack:

For adversary  $\mathcal{A}$ , he is given  $\text{trans} = (s, u, w)$ , then he can compute  $t = u \oplus s$ . In key-exchange experiment, if  $\hat{k} = w \oplus u \oplus s$ , then  $\mathcal{A}$  outputs  $b' = 0$ ; otherwise,  $\mathcal{A}$  outputs  $b' = 1$ .

So if  $b = 0$ , then  $\mathcal{A}$  always has  $b' = b$ ; and if  $b = 1$ ,  $\mathcal{A}$  guesses right with probability  $1 - 2^{-n}$ .

$$\begin{aligned}\Pr[\text{KE}_{\mathcal{A}, \Pi}^{\text{eav}}(n) = 1] &= \frac{1}{2} \Pr[\text{KE}_{\mathcal{A}, \Pi}^{\text{eav}}(n) = 1 \mid b = 0] + \frac{1}{2} \Pr[\text{KE}_{\mathcal{A}, \Pi}^{\text{eav}}(n) = 1 \mid b = 1] \\ &= \frac{1}{2} + \frac{1}{2}(1 - 2^{-n}) \\ &= 1 - \text{negl}(n),\end{aligned}$$

which is significantly larger than  $\frac{1}{2}$ .