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**RecSPL is strongly and statically typed.** Its type checking shall be implemented as a *recursive Boolean* procedure that analyses a given syntax tree (from the parser) by "tree crawling", and which

- returns **true** if the syntax tree (with all its branches) was correctly typed
- returns **false** otherwise.

It is thereby assumed that a suitable symbol table is already in place, and that the correctness of the name scopes (for variable names and function names) has already been verified *before* the type checker gets launched.

The type checker's **semantic rules** are *attributed* to their corresponding grammar rules as follows:

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PROG ::= **main** GLOBVARs ALGO FUNCTIONS

*typecheck*(PROG) = ( *typecheck*(GLOBVARs) ^ *typecheck*(ALGO) ^ *typecheck*(FUNCTIONS) )

GLOBVARs ::=

*typecheck*(GLOBVARs) = **true** // base-case of the type-checking-recursion

GLOBVARs<sub>1</sub> ::= VTYP VNAME , GLOBVARs<sub>2</sub>

*typecheck*(GLOBVARs<sub>1</sub>) = { let **T** := *typeof*(VTYP)  
                          let **id** := *symboltable*(VNAME) // access the existing symbol table!  
                          **link** (**T**,**id**) in the symbol table // symbol table now knows the type!  
                          *typeof*(VNAME) = *typeof*(VTYP)  
                          **return** *typecheck*(GLOBVARs<sub>2</sub>) }

VTYP ::= **num**

*typeof*(VTYP) = '**n**' // the auxiliary function *typeof* returns a character that represents the type

VTYP ::= **text**

*typeof*(VTYP) = '**t**' // the auxiliary function *typeof* returns a character that represents the type

VNAME ::= a token of **Token-Class V** from the Lexer

// at this point we assume that the compiler's Scope Analyser has  
// already entered some ID for this variable name into the above-  
// mentioned symbol table, such that *symboltable*(VNAME) will  
// yield use-able information about that node of the syntax tree.

ALGO ::= **begin** INSTRUC **end**

*typecheck*(ALGO) = *typecheck*(INSTRUC)

INSTRUC ::=

*typecheck*(INSTRUC) = **true** // base-case of the type-checking-recursion

INSTRUC<sub>1</sub> ::= COMMAND ; INSTRUC<sub>2</sub>

*typecheck*(INSTRUC<sub>1</sub>) = (*typecheck*(COMMAND) ^ *typecheck*(INSTRUC<sub>2</sub>))

COMMAND ::= **skip**

*typecheck*(COMMAND) = **true** // base-case of the type-checking-recursion

COMMAND ::= **halt**

*typecheck*(COMMAND) = **true** // base-case of the type-checking-recursion

```

COMMAND ::=      print ATOMIC
                  if typeof(ATOMIC)=='n' then typecheck(COMMAND) = true
                  if typeof(ATOMIC)=='t' then typecheck(COMMAND) = true
                  else typecheck(COMMAND) = false

COMMAND ::=      return ATOMIC // must stand 'inside' of a Function-Scope!
typecheck(COMMAND) = { Let a tree-crawler find the FTYP node which belongs
                        to the same Function-Scope inside of which also this
                        COMMAND is standing inside the function's BODY.
                        // We assume that Scope Analysis was already done!
                        Let X be this 'matching' function-type-node in the tree.
                        if typeof(ATOMIC) == typeof(X) == 'n' // functions can return only n
                        then return true
                        else return false }
```

COMMAND ::= ASSIGN    typecheck(COMMAND) = typecheck(ASSIGN)

COMMAND ::= CALL  
           if typeof(CALL) == 'v' // the void-type  
           then typecheck(COMMAND) = **true**  
           else typecheck(COMMAND) = **false**

COMMAND ::= BRANCH    typecheck(COMMAND) = typecheck(BRANCH)

ATOMIC    ::= VNAME  
           typeof(ATOMIC) = typeof(VNAME) // as per symbol table, which gets consulted at this point

ATOMIC    ::= CONST  
           typeof(ATOMIC) = typeof(CONST) // as per symbol table, which gets consulted at this point

CONST    ::= a token of **Token-Class N** from the Lexer  
           typeof(CONST) = 'n' // this is obviously a base-case

CONST    ::= a token of **Token-Class T** from the Lexer  
           typeof(CONST) = 't' // this is obviously a base-case

ASSIGN    ::= VNAME < **input** // we only allow numeric user-inputs in RecSPL  
           if typeof(VNAME) == 'n'  
           then typecheck(ASSIGN) = **true**  
           else typecheck(ASSIGN) = **false**

ASSIGN    ::= VNAME = TERM // texts or numbers can be assigned to variables  
           if typeof(VNAME) == typeof(TERM)  
           then typecheck(ASSIGN) = **true**  
           else typecheck(ASSIGN) = **false**

TERM      ::= ATOMIC    typeof(TERM) = typeof(ATOMIC)

TERM      ::= CALL      typeof(TERM) = typeof(CALL)

TERM      ::= OP        typeof(TERM) = typeof(OP)

CALL ::= FNAME( ATOMIC<sub>1</sub> , ATOMIC<sub>2</sub> , ATOMIC<sub>3</sub> )  
     **if** ( typeof(ATOMIC<sub>1</sub>) == 'n' ^  
         typeof(ATOMIC<sub>2</sub>) == 'n' ^  
         typeof(ATOMIC<sub>3</sub>) == 'n' ) // all three parameters of the function must be numeric  
     **then** typeof(CALL) = typeof(FNAME) // symbol table may be consulted at this point  
     **else** typeof(CALL) = 'u' // undefined which will cause the type-checker to return false

OP ::= UNOP( ARG )  
     **if** typeof(UNOP) == typeof(ARG) == 'b' **then** typeof(OP) = 'b' // bool-type  
     **if** typeof(UNOP) == typeof(ARG) == 'n' **then** typeof(OP) = 'n' // numeric type  
     **else** typeof(OP) = 'u' // undefined which will cause the type-checker to return false

ARG ::= ATOMIC   typeof(ARG) = typeof(ATOMIC)

ARG ::= OP       typeof(ARG) = typeof(OP)

UNOP ::= **not**     typeof(UNOP) = 'b' // bool-type

UNOP ::= **sqrt**    typeof(UNOP) = 'n' // numeric type

OP ::= BINOP( ARG<sub>1</sub> , ARG<sub>2</sub> )  
     **if** typeof(BINOP) == typeof(ARG<sub>1</sub>) == typeof(ARG<sub>2</sub>) == 'b' **then** typeof(OP) = 'b' // bool-type  
     **if** typeof(BINOP) == typeof(ARG<sub>1</sub>) == typeof(ARG<sub>2</sub>) == 'n' **then** typeof(OP) = 'n' // numeric type  
     **if** typeof(BINOP) == 'c' // comparison-type, which "takes" numbers and "yields" a boolean result  
     ^ typeof(ARG<sub>1</sub>) == typeof(ARG<sub>2</sub>) == 'n' **then** typeof(OP) = 'b'  
     **else** typeof(OP) = 'u' // undefined which will cause the type-checker to return false

BINOP       ::=   **or**     typeof(BINOP) = 'b'

BINOP       ::=   **and**   typeof(BINOP) = 'b'

BINOP       ::=   **eq**     typeof(BINOP) = 'c'   // comparison-type

BINOP       ::=   **grt**    typeof(BINOP) = 'c'   // comparison-type

BINOP       ::=   **add**    typeof(BINOP) = 'n'

BINOP       ::=   **sub**    typeof(BINOP) = 'n'

BINOP       ::=   **mul**    typeof(BINOP) = 'n'

BINOP       ::=   **div**    typeof(BINOP) = 'n'

BRANCH      ::=   **if** COND **then** ALGO<sub>1</sub> **else** ALGO<sub>2</sub>  
     // Attention: do not confuse the syntactic if-then-else in the Syntax Tree  
     // with the semantic if-then-else of the recursive type analysis procedure!  
     **if** typeof(COND) == 'b'  
     **then** typecheck(BRANCH) = ( typecheck(ALGO<sub>1</sub>) ^ typecheck(ALGO<sub>2</sub>) )  
     **else** typecheck(BRANCH) = **false**

COND        ::=   SIMPLE    typeof(COND) = typeof(SIMPLE)

COND        ::=   COMPOSIT   typeof(COND) = typeof(COMPOSIT)

```

SIMPLE      ::=    BINOP( ATOMIC1 , ATOMIC2 )
if typeof(BINOP) == typeof(ATOMIC1) == typeof(ATOMIC2) == 'b' then typeof(SIMPLE) = 'b'
if typeof(BINOP) == 'c' // comparison-type
    ^ typeof(ATOMIC1) == typeof(ATOMIC2) == 'n' then typeof(SIMPLE) = 'b'
else typeof(SIMPLE) = 'u' // undefined which will cause the type-checker to return false

```

```

COMPOSIT    ::=    BINOP( SIMPLE1 , SIMPLE2 )
if typeof(BINOP) == typeof(SIMPLE1) == typeof(SIMPLE2) == 'b' then typeof(COMPOSIT) = 'b'
else typeof(COMPOSIT) = 'u' // undefined which will cause the type-checker to return false

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```

COMPOSIT    ::=    UNOP ( SIMPLE )
if typeof(UNOP) == typeof(SIMPLE) == 'b' then typeof(COMPOSIT) = 'b'
else typeof(COMPOSIT) = 'u' // undefined which will cause the type-checker to return false

```

```

FNAME ::= a token of Token-Class F from the Lexer
        // at this point we assume that the compiler's Scope Analyser has
        // already entered some ID for this variable name into the above-
        // mentioned symbol table, such that symboltable(FNAME) will
        // yield use-able information about that node of the syntax tree.

```

```

FUNCTIONS ::=
    typecheck(FUNCTIONS) = true // base-case of the type-checking-recursion

```

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FUNCTIONS1      ::=    DECL FUNCTIONS2
    typecheck(FUNCTIONS1) = ( typecheck(DECL) ^ typecheck(FUNCTIONS2) )

```

```

DECL        ::=    HEADER BODY
    typecheck(DECL) = ( typecheck(HEADER) ^ typecheck(BODY) )

    // Attention! This is exactly the "area" in the tree where the above-mentioned
    // Tree-crawler must find the above-mentioned return ATOMIC command for
    // comparing its type against the function's return-type that is specified in the
    // HEADER!

```

```

HEADER      ::=    FTYP FNAME( VNAME1 , VNAME2 , VNAME3 )
    // Attention! This is exactly the "area" in the tree where the above-mentioned
    // Tree-crawler must find the above-mentioned return ATOMIC command for
    // comparing its type against the function's return-type that is specified in the HEADER!

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typecheck(HEADER) = { let T := typeof(FTYP)
    let id := symboltable(FNAME) // access the existing symbol table!
    link (T,id) in the symbol table // symbol table now knows the type!
    typeof(FNAME) = typeof(FTYP)

    if typeof(VNAME1) ==
        typeof(VNAME2) ==
        typeof(VNAME3) == 'n' // RecSPL allows only numeric arguments
    then return true
    else return false
}

```

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FTYP ::= num    typeof(FTYP) = 'n' // numeric return type
FTYP ::= void    typeof(FTYP) = 'v' // void, for return-less functions

```

```

BODY      ::=    PROLOG LOCVARs ALGO EPILOG SUBFUNCS end
               typecheck(BODY) = ( typecheck(PROLOG)
                                   ^ typecheck(LOCVARs)
                                   ^ typecheck(ALGO)
                                   ^ typecheck(EPILOG)
                                   ^ typecheck(SUBFUNCS) )

PROLOG    ::=    {
               typecheck(PROLOG) = true // base-case of the type-checking procedure

EPILOG    ::=    }
               typecheck(PROLOG) = true // base-case of the type-checking procedure

LOCVARs   ::=    VTYP1 VNAME1 , VTYP2 VNAME2 , VTYP3 VNAME3 ,
               typecheck(LOCVARs) = { let T := typeof(VTYP1)
                                   let id := symboltable(VNAME1) // access the existing symbol table!
                                   link (T,id) in the symbol table // symbol table now knows the type!
                                   typeof(VNAME1) = typeof(VTYP1)

                                   let T := typeof(VTYP2)
                                   let id := symboltable(VNAME2) // access the existing symbol table!
                                   link (T,id) in the symbol table // symbol table now knows the type!
                                   typeof(VNAME2) = typeof(VTYP2)

                                   let T := typeof(VTYP3)
                                   let id := symboltable(VNAME3) // access the existing symbol table!
                                   link (T,id) in the symbol table // symbol table now knows the type!
                                   typeof(VNAME3) = typeof(VTYP3)

                                   return true }

SUBFUNCS  ::=    FUNCTIONS
               typecheck(SUBFUNCS) = typecheck(FUNCTIONS)

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