# Integer Programming Sheet 1

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## Exercise B.1

(i) My interpretation of this exercise is to take this extension to mean switching out  $a_k^x x \leq b_k$  in the original formulation for  $A^k x \leq b^k$ . The extended alternative disjunction model then becomes

$$\min_{x \in \mathbb{R}^n} c^T x$$
s.t.  $0 \le x \le u$ 

$$A^1 x \le b^1$$

$$A^2 x \le b^2$$

where we require that at least one of  $A^1x \leq b^1$  or  $A^2x \leq b^2$  is satisfied.

Letting  $M \ge \max_k \max_i \{a_i^k x - b_i^k \mid 0 \le x \le u\}$  we can extend the big-M formulation to become

$$\min_{\substack{(x,y^1,y^2) \in \mathbb{R}^{n+1+1} \\ (x,y^1,y^2) \in \mathbb{R}^{n+1+1}}} c^T x$$

$$A^k x - b^k \le M(\mathbf{1} - \mathbf{1}y^k), \quad (k = 1, 2)$$

$$y^1 + y^2 = 1$$

$$0 \le x \le u$$

$$y^k \in \mathbb{B}^1 \text{ for } k = 1, 2$$

where  $\mathbf{1} = (\underbrace{1,\dots,1}_{m^k \text{ times}})$  and  $m^k$  is the number of rows in  $A^k$ .

(ii) There are two directions to prove.

(⇒): Assume that  $x \in P_1 \cup P_2$ . We can without loss of generality assume that  $x \in P_1$ . We then let

$$y^{1} = 1$$
$$y^{2} = 0$$
$$z^{1} = x$$
$$z^{2} = 0.$$

The constraints then become

$$x = x$$

$$A^{1}x \le b^{1}$$

$$0 \le x \le u$$

$$1 = 1$$

which are obviously all true.

( $\Leftarrow$ ): Assume that the there exists  $x, z^1, z^2, y^1, y^2$  so that the constraints are satisfied. We want to show that  $x \in P_1 \cup P_2$ . Hence, we must show that both  $0 \le x \le u$  and  $A^1x \le b^1$  or  $A^2x \le b^2$ .

For the first requirement, notice that

$$0 \le x = z^1 + z^2 \le u(y^1 + y^2) = u.$$

It thus remains to show either  $A^1x \leq b^1$  or  $A^2x \leq b^2$ . Since  $y^1+y^2=1$  we have either  $y^1=1 \wedge y^2=0$  or  $y^1=0 \wedge y^2=1$ . We can without loss of generality assume that  $y^1=1$  and  $y^2=0$ . The constraints then say that  $A^1z^1 \leq b^1$  and  $0 \leq z^2 \leq u \cdot 0=0$  so that

$$A^1x = A^1z^1 + A^10 \le b^1$$

as desired. This concludes the proof.

# Exercise B.2

The standard form is given by

$$\begin{aligned} \max_{x \in \mathbb{R}^5} &- x_{11} + x_{12} - x_{21} + x_{22} - x_3 \\ \text{s.t.} &- 2x_{11} + 2x_{12} + x_3 \leq 2, \\ &3x_{11} - 3x_{12} + x_{21} - x_{22} + 2x_3 \leq 6 \\ &x_{11} - x_{12} - x_{21} + x_{22} + 3x_3 \leq 3, \\ &- x_{11} + x_{12} + x_{21} - x_{22} - 3x_3 \leq -3 \\ &- x_{21} + x_{22} + x_3 \leq 0 \end{aligned}$$

where  $x_1 = x_{11} - x_{12}$  and  $x_2 = x_{21} - x_{22}$ . This has corresponding dual

$$\min_{y \in \mathbb{R}^5} 2y_1 + 6y_2 + 3y_3 - 3y_4$$

$$\text{s.t } -2y_1 + 3y_2 + y_3 - y_4 = -1$$

$$y_2 - y_3 + y_4 - y_5 = -1$$

$$y_1 + 2y_2 + 3y_3 - 3y_4 + y_5 = -1$$

$$y_1, y_2, y_3, y_4, y_5 \ge 0.$$

#### Exercise B.3

First, to see that the new system doesn't involve  $x_k$  is fairly straightforward. As  $a_{ij} = 0$  for  $i \in M_0^k$  we have that

$$\sum_{j=1}^{n} a_{ij} x_j \le b_i, \quad (i \in M_0^k)$$

clearly doesn't involve  $x_k$ . Moreover, if we look at the kth term of the sum

$$\sum_{j=1}^{n} (a_{ik}a_{lj} - a_{lk}a_{ij})x_j$$

we see that it equates to

$$a_{ik}a_{lk} - a_{lk}a_{ik} = 0$$

so that the coefficient of  $x_k$  is zero. Hence this sum also does not involve  $x_k$  and so the new system does not involve  $x_k$ .

For the second statement there are two directions to be proven.

 $(\Rightarrow)$ : Assume that  $(x_1, \ldots, x_{k-1}, x_{k+1}, \ldots, x_n)$  satisfies the new system. We then want to find a value of  $x_k$  so that  $(x_1, \ldots, x_k, \ldots, x_n)$  satisfies the original system.

Since all the  $x_j$ 's are given, with the exception of  $x_k$  of course, this places restrictions on what values  $x_k$  can have. Solving the original system with respect to  $x_k$  we get the following constraints:

$$x_k \le \frac{b_i - \sum_{j \ne k} a_{ij} x_j}{a_{ik}} \quad a_{ik} \in M_+^k$$
$$x_k \ge \frac{b_i - \sum_{j \ne k} a_{ij} x_j}{a_{ik}} \quad a_{ik} \in M_-^k.$$

Putting this information together tells us that a value  $x_k$  exists if and only if

$$\max_{i \in M_-^k} \left( \frac{b_i - \sum_{j \neq k} a_{ij} x_j}{a_{ik}} \right) \leq \min_{i \in M_+^k} \left( \frac{b_i - \sum_{j \neq k} a_{ij} x_j}{a_{ik}} \right).$$

Letting  $i^*$  and  $l^*$  be the corresponding maximizing and minimizing indices we get, after rearranging terms, the requirement

$$\sum_{j \neq k} (a_{i^*k} a_{l^*j} - a_{l^*k} a_{i^*j}) x_j \le a_{i^*k} b_{l^*} - a_{l^*k} b_{i^*}$$

which holds by assumption.

 $(\Leftarrow)$ : Suppose  $(x_1, \ldots, x_k, \ldots x_n)$  satisfies the original system. By assumption we then have

$$\sum_{j=1}^{n} a_{ij} x_j \le b_i, \quad (i \in M_0^k).$$

We also have for all  $(i, l) \in M_+^k \times M_-^k$  that

$$\sum_{j=1}^{n} (a_{ik}a_{lj} - a_{lk}a_{ij})x_j = a_{ik} \sum_{j=1}^{n} a_{lj}x_j - a_{lk} \sum_{j=1}^{n} a_{ij}x_j$$

$$\leq a_{ik}b_l - a_{lk}b_i$$

which concludes the proof.

## Exercise B.4

If

$$\sum_{i=1}^{n} a_{ij} x_j \le b_i, \quad (i = 1, \dots, m)$$

is infeasible, then by the theorem of alternatives for linear inequalities (Farkas' lemma) we get that there exist a  $u \in \mathbb{R}^m$  such that

$$\sum_{i=1}^{m} u_i a_{ij} = 0, \quad (j = 1, \dots, n)$$
$$u_i \ge 0, \quad (i = 1, \dots, m)$$
$$\sum_{i=1}^{m} u_i b_i < 0.$$

The objective function of (D) is  $\sum_{i=1}^{m} b_i y_i$  which we want to minimize. We know by assumption that there exists at least one feasible solution  $\tilde{y} \in \mathbb{R}^m$  for (D). Setting  $y = \tilde{y} + \lambda u$  with  $\lambda \in \mathbb{R}_+$  we have that

$$\sum_{i=1}^{m} y_i a_{ij} = \sum_{i=1}^{m} \tilde{y}_i a_{ij} + \lambda \sum_{i=1}^{m} u_i a_{ij}$$
$$= \sum_{i=1}^{m} y_i a_{ij}$$
$$= c_i, \quad (j = 1, \dots, n)$$

so that y is also a feasible solution of (D). Moreover, letting  $\tilde{d} := \sum_{i=1}^m b_i \tilde{y}_i$  we have that

$$d_{\lambda} \coloneqq \sum_{i=1}^{m} b_i y_i$$
$$= \sum_{i=1}^{m} b_i \tilde{y}_i + \lambda \sum_{i=1}^{m} u_i b_i$$
$$= \tilde{d} + \lambda d'$$

where  $d' = \sum_{i=1}^{m} u_i b_i < 0$ . Hence by letting  $\lambda \to \infty$  we get that  $d_{\lambda} \to -\infty$  showing that (D) is unbounded.