

# Integer Programming

## Sheet 1

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November 7, 2023

### Exercise B.1

- (i) My interpretation of this exercise is to take this extension to mean switching out  $a_k^x x \leq b_k$  in the original formulation for  $A^k x \leq b^k$ . The extended alternative disjunction model then becomes

$$\begin{aligned} \min_{x \in \mathbb{R}^n} \quad & c^T x \\ \text{s.t.} \quad & 0 \leq x \leq u \\ & A^1 x \leq b^1 \\ & A^2 x \leq b^2 \end{aligned}$$

where we require that at least one of  $A^1 x \leq b^1$  or  $A^2 x \leq b^2$  is satisfied.

Letting  $M \geq \max_k \max_i \{a_i^k x - b_i^k \mid 0 \leq x \leq u\}$  we can extend the big- $M$  formulation to become

$$\begin{aligned} \min_{(x, y^1, y^2) \in \mathbb{R}^{n+1+1}} \quad & c^T x \\ & A^k x - b^k \leq M(\mathbf{1} - \mathbf{1}y^k), \quad (k = 1, 2) \\ & y^1 + y^2 = 1 \\ & 0 \leq x \leq u \\ & y^k \in \mathbb{B}^1 \text{ for } k = 1, 2 \end{aligned}$$

where  $\mathbf{1} = (\underbrace{1, \dots, 1}_{m^k \text{ times}})$  and  $m^k$  is the number of rows in  $A^k$ .

- (ii) There are two directions to prove.

( $\Rightarrow$ ): Assume that  $x \in P_1 \cup P_2$ . We can without loss of generality assume that  $x \in P_1$ . We then let

$$\begin{aligned} y^1 &= 1 \\ y^2 &= 0 \\ z^1 &= x \\ z^2 &= 0. \end{aligned}$$

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The constraints then become

$$\begin{aligned}x &= x \\ A^1 x &\leq b^1 \\ 0 &\leq x \leq u \\ 1 &= 1\end{aligned}$$

which are obviously all true.

( $\Leftarrow$ ): Assume that there exists  $x, z^1, z^2, y^1, y^2$  so that the constraints are satisfied. We want to show that  $x \in P_1 \cup P_2$ . Hence, we must show that both  $0 \leq x \leq u$  and  $A^1 x \leq b^1$  or  $A^2 x \leq b^2$ .

For the first requirement, notice that

$$0 \leq x = z^1 + z^2 \leq u(y^1 + y^2) = u.$$

It thus remains to show either  $A^1 x \leq b^1$  or  $A^2 x \leq b^2$ . Since  $y^1 + y^2 = 1$  we have either  $y^1 = 1 \wedge y^2 = 0$  or  $y^1 = 0 \wedge y^2 = 1$ . We can without loss of generality assume that  $y^1 = 1$  and  $y^2 = 0$ . The constraints then say that  $A^1 z^1 \leq b^1$  and  $0 \leq z^2 \leq u \cdot 0 = 0$  so that

$$A^1 x = A^1 z^1 + A^1 0 \leq b^1$$

as desired. This concludes the proof.

## Exercise B.2

The standard form is given by

$$\begin{aligned}\max_{x \in \mathbb{R}^5} & -x_{11} + x_{12} - x_{21} + x_{22} - x_3 \\ \text{s.t.} & -2x_{11} + 2x_{12} + x_3 \leq 2, \\ & 3x_{11} - 3x_{12} + x_{21} - x_{22} + 2x_3 \leq 6 \\ & x_{11} - x_{12} - x_{21} + x_{22} + 3x_3 \leq 3, \\ & -x_{11} + x_{12} + x_{21} - x_{22} - 3x_3 \leq -3 \\ & -x_{21} + x_{22} + x_3 \leq 0\end{aligned}$$

where  $x_1 = x_{11} - x_{12}$  and  $x_2 = x_{21} - x_{22}$ . This has corresponding dual

$$\begin{aligned}\min_{y \in \mathbb{R}^5} & 2y_1 + 6y_2 + 3y_3 - 3y_4 \\ \text{s.t.} & -2y_1 + 3y_2 + y_3 - y_4 = -1 \\ & y_2 - y_3 + y_4 - y_5 = -1 \\ & y_1 + 2y_2 + 3y_3 - 3y_4 + y_5 = -1 \\ & y_1, y_2, y_3, y_4, y_5 \geq 0.\end{aligned}$$

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### Exercise B.3

First, to see that the new system doesn't involve  $x_k$  is fairly straightforward. As  $a_{ij} = 0$  for  $i \in M_0^k$  we have that

$$\sum_{j=1}^n a_{ij}x_j \leq b_i, \quad (i \in M_0^k)$$

clearly doesn't involve  $x_k$ . Moreover, if we look at the  $k$ th term of the sum

$$\sum_{j=1}^n (a_{ik}a_{lj} - a_{lk}a_{ij})x_j$$

we see that it equates to

$$a_{ik}a_{lk} - a_{lk}a_{ik} = 0$$

so that the coefficient of  $x_k$  is zero. Hence this sum also does not involve  $x_k$  and so the new system does not involve  $x_k$ .

For the second statement there are two directions to be proven.

( $\Rightarrow$ ): Assume that  $(x_1, \dots, x_{k-1}, x_{k+1}, \dots, x_n)$  satisfies the new system. We then want to find a value of  $x_k$  so that  $(x_1, \dots, x_k, \dots, x_n)$  satisfies the original system.

Since all the  $x_j$ 's are given, with the exception of  $x_k$  of course, this places restrictions on what values  $x_k$  can have. Solving the original system with respect to  $x_k$  we get the following constraints:

$$\begin{aligned} x_k &\leq \frac{b_i - \sum_{j \neq k} a_{ij}x_j}{a_{ik}} & a_{ik} &\in M_+^k \\ x_k &\geq \frac{b_i - \sum_{j \neq k} a_{ij}x_j}{a_{ik}} & a_{ik} &\in M_-^k. \end{aligned}$$

Putting this information together tells us that a value  $x_k$  exists if and only if

$$\max_{i \in M_-^k} \left( \frac{b_i - \sum_{j \neq k} a_{ij}x_j}{a_{ik}} \right) \leq \min_{i \in M_+^k} \left( \frac{b_i - \sum_{j \neq k} a_{ij}x_j}{a_{ik}} \right).$$

Letting  $i^*$  and  $l^*$  be the corresponding maximizing and minimizing indices we get, after rearranging terms, the requirement

$$\sum_{j \neq k} (a_{i^*k}a_{l^*j} - a_{l^*k}a_{i^*j})x_j \leq a_{i^*k}b_{l^*} - a_{l^*k}b_{i^*}$$

which holds by assumption.

( $\Leftarrow$ ): Suppose  $(x_1, \dots, x_k, \dots, x_n)$  satisfies the original system. By assumption we then have

$$\sum_{j=1}^n a_{ij}x_j \leq b_i, \quad (i \in M_0^k).$$

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We also have for all  $(i, l) \in M_+^k \times M_-^k$  that

$$\begin{aligned} \sum_{j=1}^n (a_{ik}a_{lj} - a_{lk}a_{ij})x_j &= a_{ik} \sum_{j=1}^n a_{lj}x_j - a_{lk} \sum_{j=1}^n a_{ij}x_j \\ &\leq a_{ik}b_l - a_{lk}b_i \end{aligned}$$

which concludes the proof.

### Exercise B.4

If

$$\sum_{j=1}^n a_{ij}x_j \leq b_i, \quad (i = 1, \dots, m)$$

is infeasible, then by the theorem of alternatives for linear inequalities (Farkas' lemma) we get that there exist a  $u \in \mathbb{R}^m$  such that

$$\begin{aligned} \sum_{i=1}^m u_i a_{ij} &= 0, \quad (j = 1, \dots, n) \\ u_i &\geq 0, \quad (i = 1, \dots, m) \\ \sum_{i=1}^m u_i b_i &< 0. \end{aligned}$$

The objective function of (D) is  $\sum_{i=1}^m b_i y_i$  which we want to minimize. We know by assumption that there exists at least one feasible solution  $\tilde{y} \in \mathbb{R}^m$  for (D). Setting  $y = \tilde{y} + \lambda u$  with  $\lambda \in \mathbb{R}_+$  we have that

$$\begin{aligned} \sum_{i=1}^m y_i a_{ij} &= \sum_{i=1}^m \tilde{y}_i a_{ij} + \lambda \sum_{i=1}^m u_i a_{ij} \\ &= \sum_{i=1}^m \tilde{y}_i a_{ij} \\ &= c_j, \quad (j = 1, \dots, n) \end{aligned}$$

so that  $y$  is also a feasible solution of (D). Moreover, letting  $\tilde{d} := \sum_{i=1}^m b_i \tilde{y}_i$  we have that

$$\begin{aligned} d_\lambda &:= \sum_{i=1}^m b_i y_i \\ &= \sum_{i=1}^m b_i \tilde{y}_i + \lambda \sum_{i=1}^m u_i b_i \\ &= \tilde{d} + \lambda d' \end{aligned}$$

where  $d' = \sum_{i=1}^m u_i b_i < 0$ . Hence by letting  $\lambda \rightarrow \infty$  we get that  $d_\lambda \rightarrow -\infty$  showing that (D) is unbounded.